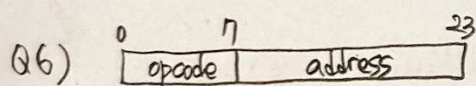


I honestly will not do any cheating. 송민정 ~~Real~~

- Q1) Because Cache can help speed up by lower gap between CPU and main memory's speed. It has small capacity but fast.
- Q2) It deals with the potential speed up of a program using multiple processors compared to a single processor. The law indicates the amount of speedup as a function of the fraction of code that can be executed in parallel.
- Q3) SAAS is Software as a Service. It provides service to customers in form of a software, specifically application software, running on and accessible in the cloud.
- Q4) There are cell phone, tooth brushes, video cameras, calculator, washing machines.
- Q5) It means that split caches into 2 caches; one dedicated to instructions and one dedicated to data.



1) IR: 24 bits, PC: 16 bits

2) 01110000 1010 1010 1111 1111

3) 10101010 0011 1111 1111 0011

$$\begin{aligned} Q1) \text{ CPI} &= \frac{60000 \times 1 + 40000 \times 2 + 10000 \times 3 + 30000 \times 2}{200000} \\ &= \frac{(6+8+2+6) \frac{10000}{20}}{\frac{200000}{20}} = \frac{41}{20} = 2.05 \end{aligned}$$

$$\begin{aligned} \text{MIPS} &= \frac{I_c}{T \times 10^6}, [T = I_c \times \text{CPI} \times t] \\ &= \frac{f}{\text{CPI} \times 10^6} = \frac{100 \times 10^6}{2.05 \times 10^6} = \frac{20}{100} \times \frac{100}{2.05} = \frac{2000}{41} \end{aligned}$$

$$\begin{aligned} T &= I_c \times \text{CPI} \times t = I_c \times \text{CPI} \times \frac{1}{f} \\ &= 266666 \times 2.05 \times \frac{1}{100 \times 10^6} = 0.2 \times 2.05 \times 10^{-2} = 41 \times 10^{-5} \end{aligned}$$

$$Q4) 1) \frac{1}{100 \text{ MHz}} = \frac{1}{10^8} = 10^{-8}$$

\therefore 70 ns are needed

$$2) 100/3 = 33.3 \text{ MB/s}$$

$$3) 100/5 = 20 \text{ MB/s}$$

Q9) 1 block = 4 words

$$2^{23} \text{ byte} / 2^3 \text{ byte} = 2^{20} \text{ byte} \rightarrow 2^{20} / 2^2 = 2^{18} \text{ block}$$

1 word = 64 bit = 8 Byte

$$2^{18} / 2^2 = 2^{16}$$

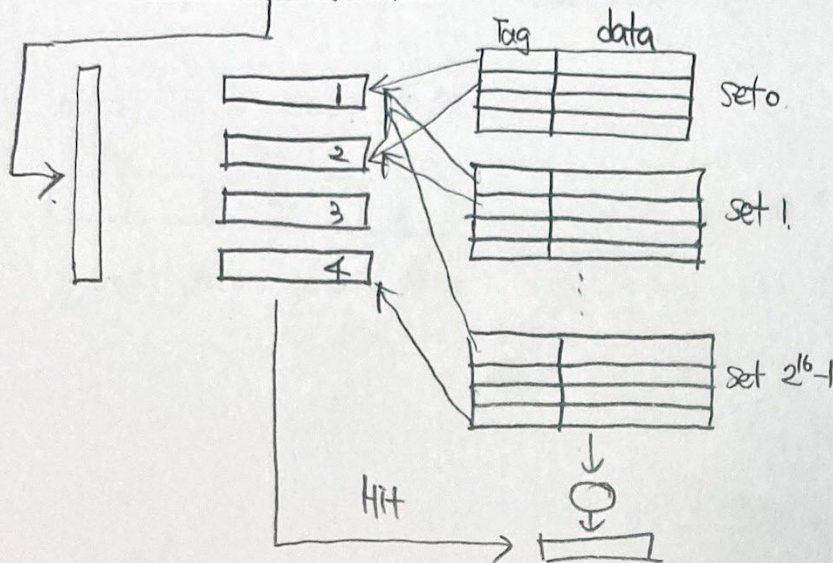
a)

Tag	Set	Word
4	16	3

$$23 - 16 - 3 = 4$$

b)

Tag	Set	Word
4	16	3



c) When we add 2 bits LRU, if the data hits, it changes to "1".

And others are zero or 1. LRU means "least recently used". So, least recently used, (LRU 2 bits are "00") will be replaced first.

Q10) 1 word = 1 Byte, 1 block = 8 Byte = 8 word

[Direct-map] 32 lines = 2^5 64K = 2^{16}

a)

Tag	line	word
8	5	3

b) 0001 0001 / 0001 1011 (slot 3)

1100 0011 / 0011 0100 (slot 6)

c) 0001 1010 0001 1010 ~ 0001 1010 0001 1111 are stored.

Q10) d) 256 bytes

e) 0001 1010 0001 1010

TAG	word
13	3

0/001 1/010 0/001 1/010 \Rightarrow 0343 : Tag

f) 0001 1010 0001 1010

Tag	Set	word
9	4	3

0/001 1/010 0/001 1/010 Tag: 034 Set: 3

Q11)	cache	time	probability
		9 ns	0.9
	main 0, cache X	80 ns + 9 ns	(0.1)(0.6)
	main X, cache X	12 ms + 80 ns + 9 ns	(0.1)(0.4)

$$\therefore 9 \text{ ns}(0.9) + 89 \text{ ns}(0.06) + (12,000,089) \text{ ns}(0.04)$$

$$= (811534 + 48,000,356) \text{ ns} \times 10^{-2}$$

$$= (48,001,700) \times 10^{-2} \text{ ns} = 480017 \text{ ns}$$

$$\begin{array}{r} 7889 \\ \times 6 \\ \hline 534 \end{array}$$

$$\begin{array}{r} + 480 \\ 7344 \\ 356 \\ \hline 1700 \end{array}$$