

Chapter 6: Architecture

Logical / Shift Instructions

Programming

- **High-level languages:**
 - e.g., C, Java, Python
 - Written at higher level of abstraction
- **High-level constructs:** loops, conditional statements, arrays, function calls
- **First, introduce instructions that support these:**
 - Logical operations
 - Shift instructions
 - Multiplication & division
 - Branches & Jumps

Ada Lovelace, 1815-1852

- Wrote the first computer program
- Her program calculated the Bernoulli numbers on Charles Babbage's Analytical Engine
- She was the daughter of the poet Lord Byron



Logical Instructions

- **and, or, xor**

- and: useful for **masking** bits
 - Masking all but the least significant byte of a value:
 $0xF234012F \text{ AND } 0x000000FF = 0x0000002F$
- or: useful for **combining** bit fields
 - Combine $0xF2340000$ with $0x000012BC$:
 $0xF2340000 \text{ OR } 0x000012BC = 0xF23412BC$
- xor: useful for **inverting** bits:
 - $A \text{ XOR } -1 = \text{NOT } A$ (remember that $-1 = 0xFFFFFFFF$)

Logical Instructions: Example 1

Source Registers

s1	0100 0110	1010 0001	1111 0001	1011 0111
s2	1111 1111	1111 1111	0000 0000	0000 0000

Assembly Code

and s3, s1, s2

or s4, s1, s2


xor s5, s1, s2

Result

s3	0100 0110	1010 0001	0000 0000	0000 0000
s4	1111 1111	1111 1111	1111 0001	1011 0111
s5	1011 1001	0101 1110	1111 0001	1011 0111

Logical Instructions: Example 2

Source Values

t3	0011	1010	0111	0101	0000	1101	0110	1111
imm	1111	1111	1111	1111	1111	1010	0011	0100
								

Assembly Code

```
andi s5, t3, -1484
ori  s6, t3, -1484
xori s7, t3, -1484
```

Result

s5							
s6							
s7							

-1484 = **0xA34** in 12-bit 2's complement representation.

Shift Instructions

Shift amount is in (lower 5 bits of) a register

- `sll`: shift left logical
 - **Example:** `sll t0, t1, t2 # t0 = t1 << t2`
- `srl`: shift right logical
 - **Example:** `srl t0, t1, t2 # t0 = t1 >> t2`
- `sra`: shift right arithmetic
 - **Example:** `sra t0, t1, t2 # t0 = t1 >>> t2`

Immediate Shift Instructions

Shift amount is an immediate between 0 to 31

- `slli`: shift left logical immediate
 - **Example:** `slli t0, t1, 23 # t0 = t1 << 23`
- `srlr`: shift right logical immediate
 - **Example:** `srlr t0, t1, 18 # t0 = t1 >> 18`
- `srai`: shift right arithmetic immediate
 - **Example:** `srai t0, t1, 5 # t0 = t1 >>> 5`

Chapter 6: Architecture

Multiplication and Division

Multiplication

32 × 32 multiplication → 64 bit result

mul s3, s1, s2

s3 = lower 32 bits of result

mulh s4, s1, s2

s4 = upper 32 bits of result, treats operands as signed

$\{s4, s3\} = s1 \times s2$

Example: $s1 = 0x40000000 = 2^{30}$; $s2 = 0x80000000 = -2^{31}$

$s1 \times s2 = -2^{61} = 0xE0000000\ 00000000$

$s4 = 0xE0000000$; $s3 = 0x00000000$

Division

32-bit division → 32-bit quotient & remainder

- `div s3, s1, s2` # $s3 = s1 / s2$
- `rem s4, s1, s2` # $s4 = s1 \% s2$

Example: $s1 = 0x00000011 = 17$; $s2 = 0x00000003 = 3$

$$s1 / s2 = 5$$

$$s1 \% s2 = 2$$


$$s3 = 0x00000005; s4 = 0x00000002$$

Chapter 6: Architecture

Branches & Jumps

Branching

- Execute instructions out of sequence
- Types of branches:
 - **Conditional**
 - branch if equal (`beq`)
 - branch if not equal (`bne`)
 - branch if less than (`blt`)
 - branch if greater than or equal (`bge`)
 - **Unconditional**
 - jump (`j`)
 - jump register (`jr`)
 - jump and link (`jal`)
 - jump and link register (`jalr`)



We'll talk
about these
when discuss
function calls

Conditional Branching

RISC-V assembly

```
addi s0, zero, 4      # s0 = 0 + 4 = 4
addi s1, zero, 1      # s1 = 0 + 1 = 1
slli s1, s1, 2        # s1 = 1 << 2 = 4
beq  s0, s1, target   # branch is taken
addi s1, s1, 1        # not executed
sub  s1, s1, s0        # not executed

target:               # label
add  s1, s1, s0        # s1 = 4 + 4 = 8
```

Labels indicate instruction location. They can't be reserved words and must be followed by a colon (:)

The Branch Not Taken (bne)

RISC-V assembly

```
addi    s0, zero, 4           # s0 = 0 + 4 = 4
addi    s1, zero, 1           # s1 = 0 + 1 = 1
slli    s1, s1, 2              # s1 = 1 << 2 = 4
bne     s0, s1, target        # branch not taken
addi    s1, s1, 1              # s1 = 4 + 1 = 5
sub     s1, s1, s0             # s1 = 5 - 4 = 1

target:
add     s1, s1, s0             # s1 = 1 + 4 = 5
```

Unconditional Branching (j)

RISC-V assembly

```
j            target            # jump to target
srai         s1, s1, 2         # not executed
addi         s1, s1, 1         # not executed
sub          s1, s1, s0        # not executed

target:
add          s1, s1, s0        # s1 = 1 + 4 = 5
```


Chapter 6: Architecture

Conditional Statements & Loops

Conditional Statements & Loops

- **Conditional Statements**

- `if` statements
- `if/else` statements

- **Loops**

- `while` loops
- `for` loops

If Statement

C Code

```
if (i == j)
    f = g + h;
```

```
f = f - i;
```

RISC-V assembly code

```
# s0 = f, s1 = g, s2 = h
# s3 = i, s4 = j
```

Assembly tests opposite case ($i \neq j$) of high-level code ($i == j$)

If/Else Statement

C Code

```
if (i == j)
    f = g + h;
```

```
else
    f = f - i;
```

RISC-V assembly code

```
# s0 = f, s1 = g, s2 = h
# s3 = i, s4 = j
```

Assembly tests opposite case ($i \neq j$) of high-level code ($i == j$)

While Loops

C Code

```
// determines the power
// of x such that 2x = 128
int pow = 1;
int x    = 0;

while (pow != 128) {
    pow = pow * 2;
    x = x + 1;
}
```

RISC-V assembly code

```
# s0 = pow, s1 = x
```

Assembly tests opposite case (`pow == 128`) of high-level code
(`pow != 128`)

For Loops

```
for (initialization; condition; loop operation)  
    statement
```

- **initialization:** executes **before** the loop begins
- **condition:** is tested **at the beginning** of each iteration
- **loop operation:** executes at the **end** of each iteration
- **statement:** executes **each time** the condition is met

For Loops

C Code

```
// add the numbers from 0 to 9
int sum = 0;
int i;

for (i=0; i!=10; i = i+1) {
    sum = sum + i;
}
```

RISC-V assembly code

```
# s0 = i, s1 = sum
```

Less Than Comparison

C Code

```
// add the powers of 2 from 1
// to 100
int sum = 0;
int i;

for (i=1; i < 101; i = i*2) {
    sum = sum + i;
}
```

RISC-V assembly code

```
# s0 = i, s1 = sum
```


Less Than Comparison: Version 2

C Code

```
// add the powers of 2 from 1
// to 100
int sum = 0;
int i;

for (i=1; i < 101; i = i*2) {
    sum = sum + i;
}
```

RISC-V assembly code

```
# s0 = i, s1 = sum
    addi    s1, zero, 0
    addi    s0, zero, 1
    addi    t0, zero, 101

loop:
    slt     t2, s0, t0
    beq     t2, zero, done
    add     s1, s1, s0
    slli    s0, s0, 1
    j       loop

done:
```

slt: set if less than instruction

```
slt t2, s0, t0 # if s0 < t0, t2 = 1
               # otherwise t2 = 0
```

Chapter 6: Architecture

Arrays

Arrays

- Access large amounts of similar data
- **Index:** access each element
- **Size:** number of elements

Arrays

- 5-element array
- **Base address** = 0x123B4780 (address of first element, `array[0]`)
- First step in accessing an array: load base address into a register

Address	Data
123B4790	<code>array[4]</code>
123B478C	<code>array[3]</code>
123B4788	<code>array[2]</code>
123B4784	<code>array[1]</code>
123B4780	<code>array[0]</code>

Main Memory

Accessing Arrays

// C Code

```
int array[5];  
array[0] = array[0] * 2;  
array[1] = array[1] * 2;
```

RISC-V assembly code

```
# s0 = array base address
```

Address	Data
123B4790	array[4]
123B478C	array[3]
123B4788	array[2]
123B4784	array[1]
123B4780	array[0]

Main Memory

Accessing Arrays Using For Loops

// C Code

```
int array[1000];  
int i;  
  
for (i=0; i < 1000; i = i + 1)  
    array[i] = array[i] * 8;
```

RISC-V assembly code

```
# s0 = array base address, s1 = i
```

Accessing Arrays Using For Loops

RISC-V assembly code

```
# s0 = array base address, s1 = i
# initialization code
lui    s0, 0x23B8F          # s0 = 0x23B8F000
ori    s0, s0, 0x400        # s0 = 0x23B8F400
addi   s1, zero, 0         # i = 0
addi   t2, zero, 1000      # t2 = 1000

loop:
    bge    s1, t2, done     # if not then done
    slli   t0, s1, 2        # t0 = i * 4 (byte offset)
    add    t0, t0, s0       # address of array[i]
    lw     t1, 0(t0)        # t1 = array[i]
    slli   t1, t1, 3        # t1 = array[i] * 8
    sw     t1, 0(t0)        # array[i] = array[i] * 8
    addi   s1, s1, 1        # i = i + 1
    j      loop            # repeat
done:
```

ASCII Code

- **ASCII:** *American Standard Code for Information Interchange*
- Each text character has unique byte value
 - For example, S = 0x53, a = 0x61, A = 0x41
 - Lower-case and upper-case differ by 0x20 (32)

Cast of Characters: ASCII Encodings

#	Char	#	Char	#	Char	#	Char	#	Char	#	Char
20	space	30	0	40	@	50	P	60	`	70	p
21	!	31	1	41	A	51	Q	61	a	71	q
22	“	32	2	42	B	52	R	62	b	72	r
23	#	33	3	43	C	53	S	63	c	73	s
24	\$	34	4	44	D	54	T	64	d	74	t
25	%	35	5	45	E	55	U	65	e	75	u
26	&	36	6	46	F	56	V	66	f	76	v
27	'	37	7	47	G	57	W	67	g	77	w
28	(38	8	48	H	58	X	68	h	78	x
29)	39	9	49	I	59	Y	69	i	79	y
2A	*	3A	:	4A	J	5A	Z	6A	j	7A	z
2B	+	3B	;	4B	K	5B	[6B	k	7B	{
2C	,	3C	<	4C	L	5C	\	6C	l	7C	
2D	-	3D	=	4D	M	5D]	6D	m	7D	}
2E	.	3E	>	4E	N	5E	^	6E	n	7E	~
2F	/	3F	?	4F	O	5F	_	6F	o		

Accessing Arrays of Characters

// C Code

```
char str[80] = "CAT";
```

```
int len = 0;
```

```
// compute length of string
```

```
while (str[len]) len++;
```

RISC-V assembly code

```
# s0 = array base address, s1 = len
```