Revision

1. Dinect Methods

- Gauss eliminahm 0 (2/3h3)
- Gauss Jordon O(n3)
- LU Decomposition
 - o GE
 - o Doolttle

- o Crout (Toidiag mal) · Thomson algorithm - O(n)
- · Cholesky decomposition D(1/3 n3) (Symmetric poeter definle)

Forward Error Analysis

Recall
$$Cp = \left(\frac{\Delta f}{f}\right) \qquad \chi + D\pi$$

$$|n|$$

$$|\Delta n|$$

$$|Df| = Cp \left(\frac{\Delta n}{n}\right)$$

$$A \times = b$$
$$\times = A^{-1}b$$

What would be the relative charges in X, ie $\frac{\Delta X}{X}$, due to small perturbations in A, $\frac{\Delta A}{A}$, or b, $\frac{\Delta b}{b}$.

for a vector $X = [n_1 n_2 - n_n]^T$ room is a generalized length of the vector. It is denoted by IIXII and it saturates the forlowing properties || X || is a non negative real number $|| \times || = 0 \qquad \text{iff} \qquad \times = 0 \qquad \text{2i = 0} \qquad \text{i=1,--b}$ [| K X]] = (K/ || X/ $|| \times + Y || \le || \times || + || Y || - || Toisngular || i-equality$

Matrix Norm

 $f \mid_{x \in x} f$

AX -> vector of sine nx1

The matrix norm is the maximum extension sato achievable

 $\|A\| = \max_{X\neq 0} \frac{\|Ax\|}{\|x\|}$ for all non-zero

Maximum is taken over all possible non-zew X

Another supresentative

 $||A|| = \max_{\|A\|=1} \|AX\|$

Satisfy the following properties 1 A 1 > 0 || KA || = | K | || A || $\|A + B\| \leq \|A\| + \|B\| = \begin{cases} a_{11}n_{1} + a_{12}n_{2} \\ a_{21}n_{1} + a_{22}n_{2} \end{cases}$ $\|A\chi\| \leq \|A\|\|\chi\|$ $\begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix} \begin{bmatrix} n_1 \\ n_2 \end{bmatrix}$ [A B] (B) 1 All = max 11Ax1/0

 $\|A\|_{\infty} = \max \left(\sum_{j=1}^{\infty} \beta_{ij} \right)^{-1}$ $\|A\|_{1} = man \left(\sum_{i=1}^{n} |a_{ij}| \right) \cdot coumn$

Frankle
$$A = \begin{bmatrix} -2 & 3 \\ 4 & -5 \end{bmatrix}$$

$$\begin{bmatrix} 5 \\ 9 \end{bmatrix}$$

$$\begin{bmatrix} A \\ b \end{bmatrix} = 9$$

$$\begin{bmatrix} 6 & 8 \end{bmatrix}$$

Spectral radius

$$\begin{cases}
Aai = \lambda_i a_i \\
\lambda_i - egen \\
\lambda_i - e$$

Example
$$x_1 + 2x_2 = 10$$
 (443)
 $x_1 + 2x_2 = 10.4$
 $x_1 + 2x_2 = 10.4$
 $x_2 = 10.4$
 $x_3 = 4$
 $x_4 = 11.1$
 $x_4 = 11.1$

$$\|A\|_{\infty} = 3.$$

$$A^{-1} = \begin{bmatrix} -10 & 10 \\ 5.5 & -5 \end{bmatrix}$$

$$\|A^{-1}\|_{\infty} = 20$$

$$C(A) = ||A^{-1}||_{1} ||A||_{2}$$

= $3 \cdot | \times 20 = 62$
 $i(||-condutated|)$

Smallest condular number $A = I \quad C(A) = 1$

$$\frac{\|\Delta n\|}{\|x\|} \leq C_{p} \frac{\|\Delta n\|}{\|A\|}$$

$$\leq 62. \frac{0.05}{3.1}$$

$$\frac{\|\Delta n\|}{\|n\|} \leq 1$$

$$\|\Delta n\| \leq \|x\|_{\infty}$$

$$\|\Delta n\| \leq 4$$

Determinant is not a good measure

of the ill or well conditioning of

the media: $A = \begin{bmatrix} 1 & 2 \\ 1 & 1 \end{bmatrix} \quad \text{det}(A) = -0.2.$ $A = \begin{bmatrix} 1000 & 2010 \\ 1100 & 2010 \end{bmatrix} \quad \text{det}(A) = 2 \times 10^{5}$

11 A 11 = 3100

 $A^{-1} = \begin{bmatrix} -0.01 & 0.01 \\ 0.0055 & -0.005 \end{bmatrix} \begin{bmatrix} ||A^{-1}||_{\infty} = 0.02 \\ ||A^{-1}||_{\infty} = 0.02 \end{bmatrix}$ $C(A) = ||A^{-1}||_{\infty} ||A||_{\infty}$ $= 62 \qquad ||A^{-1}||_{\infty}$

Recommendation that offer estimating X, substitute it in the equation and see whether the equation is satisfied a not

AX = b

Instead of X me estimated of

An = 5

9.56-5 is' small then the estandi his good)

$$An = b$$
 $||b|| = ||An||$
 $||b|| \le ||A|| ||a||$

$$||A^{-1}|| ||A|| ||S||$$

$$||B^{-1}|| ||A|| ||S||$$

Example
$$n_1 + 2n_2 = 10 \rightarrow 10$$

 $m_1 + 2n_2 = 10.9 \rightarrow 10.2$

Matrix Norm

The l_{∞} norm of a matrix can be easily computed from the entries of the matrix.

If $A = (a_{ij})$ is an $n \times n$ matrix, then

$$||A||_{\infty} = \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|.$$

Proof First we show that $||A||_{\infty} \le \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|$. Let \mathbf{x} be an n-dimensional vector with $1 = ||\mathbf{x}||_{\infty} = \max_{1 \le i \le n} |x_i|$. Since $A\mathbf{x}$ is also an n-dimensional vector,

$$||A\mathbf{x}||_{\infty} = \max_{1 \le i \le n} |(A\mathbf{x})_i| = \max_{1 \le i \le n} \left| \sum_{j=1}^n a_{ij} x_j \right| \le \max_{1 \le i \le n} \sum_{j=1}^n |a_{ij}| \max_{1 \le j \le n} |x_j|.$$

But $\max_{1 \le j \le n} |x_j| = ||\mathbf{x}||_{\infty} = 1$, so

$$||A\mathbf{x}||_{\infty} \le \max_{1 \le i \le n} \sum_{i=1}^{n} |a_{ij}|.$$

Consequently,

$$||A||_{\infty} = \max_{\|\mathbf{x}\|_{\infty} = 1} ||A\mathbf{x}||_{\infty} \le \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|.$$
 (7.3)

Now we will show the opposite inequality, that $||A||_{\infty} \ge \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|$. Let p be an integer with

$$\sum_{j=1}^{n} |a_{pj}| = \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|,$$

and \mathbf{x} be the vector with components

$$x_j = \begin{cases} 1, & \text{if } a_{pj} \ge 0, \\ -1, & \text{if } a_{pj} < 0. \end{cases}$$

Then $\|\mathbf{x}\|_{\infty} = 1$ and $a_{pj}x_j = |a_{pj}|$, for all j = 1, 2, ..., n, so

$$||A\mathbf{x}||_{\infty} = \max_{1 \le i \le n} \left| \sum_{j=1}^{n} a_{ij} x_j \right| \ge \left| \sum_{j=1}^{n} a_{pj} x_j \right| = \left| \sum_{j=1}^{n} |a_{pj}| \right| = \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|.$$

This result implies that

$$||A||_{\infty} = \max_{\|\mathbf{x}\|_{\infty} = 1} ||A\mathbf{x}||_{\infty} \ge \max_{1 \le i \le n} \sum_{i=1}^{n} |a_{ij}|,$$

which, together with Inequality (7.3), gives

$$\|A\|_{\infty} = \max_{1 \le i \le n} \sum_{j=1}^{n} |a_{ij}|.$$
 From: Numerical Analysis by Burden & Faires