

Syntactic Substitution

CS 536: Science of Programming, Fall 2021

1. Calculate $[b+c/c][i+1/i](x+i*b+c = 0)$.
2. Let p be $\exists x. x < y \wedge x^2 \geq y+k$
 - a. What is $[5/x]p$?
 - b. What is $[5/y]p$?
 - c. What is $[5/z]p$?
 - d. What is $[y^2/y]p$?
 - e. What is $[y^k/y]p$?
 - f. What is $[(x + y) \div 2/y]p$?
3. Give an example where $[e/v][e'/w](v * w)$ and $[e'/w][e/v](v * w)$ are
 - a. Syntactically equal (\equiv)
 - b. Syntactically unequal ($\not\equiv$).
4. In the predicate $(\exists x . x < y \wedge x^2 \geq y+k)$, x is bound, but in $(x < y \wedge x^2 \geq y+k)$, x is free
— is this a contradiction?
6. Let $p \equiv (\forall x. \exists y. R(x, y, z)) \wedge (\exists z. R(x, y, z))$ where R is a boolean function over three arguments.
 - a. What is $[17/w]p$?
 - b. What is $[17/x]p$?
 - c. What is $[y^2/y]p$?
 - d. What is $[y^2/z]p$?
 - e. What is $[a+b/z][a*z/y]p$?

Solution to Practice 12 (Syntactic Substitution)

1. $[b+c/c][i+1/i](x+i*b+c = 0) \equiv [b+c/c](x + (i+1)*b + c = 0)$
 $\equiv x+(i+1)*b+(b+c) = 0$
2. Let $p \equiv \exists x . x < y \wedge x^2 \geq y+k$
 - 2a. $[5/x]p \equiv p$ unchanged
 - 2b. $[5/y]p \equiv [5/y](\exists x . x < y \wedge x^2 \geq y+k) \equiv \exists x . x < 5 \wedge x^2 \geq 5+k$
 - 2c. $[5/z]p \equiv p$ unchanged because z doesn't occur in p
 - 2d. $[y*2/y]p \equiv [y*2/y](\exists x . x < y \wedge x^2 \geq y+k) \equiv \exists x . x < y*2 \wedge x^2 \geq y*2+k$
 - 2e. $[y*k/y]p \equiv [y*k/y](\exists x . x < y \wedge x^2 \geq y+k) \equiv \exists x . x < y*k \wedge x^2 \geq y*k+k$
 - 2f. $[(x+y) \div 2/y]p \equiv [(x+y) \div 2/y](\exists x . x < y \wedge x^2 \geq y+k)$
 $\equiv [(x+y) \div 2/y] \exists v.[v/x](x < y \wedge x^2 \geq y+k)$ (note renaming of x to v)
 $\equiv [(x+y) \div 2/y] \exists v.(v < y \wedge v^2 \geq y+k)$
 $\equiv \exists v . v < (x+y) \div 2 \wedge v^2 \geq (x+y) \div 2 + k$
3. (Cases where $[e'/w][e/v](v * w)$ and $[e/v][e'/w](v * w)$ are \equiv and $\not\equiv$.)
 - 3a. One case is when v doesn't occur in e' and w doesn't occur in e .
Example: $[a*w/w][v*2/v](v * w) \equiv [a*w/w](v*2 * w)$
 $\equiv v*2 * (a*w) \equiv [v*2/v](v * (a*w))$
 $\equiv [v*2/v][a*w/w](v * w)$
 - 3b. One case is when w appears in e and v appears in e' , at least, for certain e and e' .
Example: $[a*v/w][w-3/v](v * w) \equiv [a*v/w]((w-3) * w) \equiv (w-3) * (a * v)$
but $[w-3/v][a*v/w](v * w) \equiv [w-3/v](v * (a*v)) \equiv (w-3) * (a * (w-3))$
4. No, this is exactly what a quantifier does: It captures the x 's that are free in its body and makes them bound with respect to any context that includes the quantified predicate.
6. Substitutions with $p \equiv (\forall x . \exists y . R(x, y, z)) \wedge \exists z . R(x, y, z)$:
 - 6a. $[17/w]p \equiv p$ (because w doesn't occur in p)
 - 6b. $[17/x]p \equiv (\forall x . \exists y . R(x, y, z)) \wedge \exists z . R(17, y, z)$
 - 6c. $[y*2/y]p \equiv (\forall x . \exists y . R(x, y, z)) \wedge \exists z . R(x, y*2, z)$
 - 6d. $[y*2/z]p \equiv (\forall x . \exists v . [y*2/z][v/y]R(x, y, z)) \wedge \exists z . R(x, y, z)$ (using v as a fresh variable)
 $\equiv (\forall x . \exists v . R(x, v, y*2)) \wedge \exists z . R(x, y, z)$
 - 6e. $[a+b/z][a*z/y]p$
 $\equiv [a+b/z](\forall x . \exists y . R(x, y, z)) \wedge \exists v . [a*z/y][v/z]R(x, y, z)$ (using v as a fresh variable)
 $\equiv [a+b/z](\forall x . \exists y . R(x, y, z)) \wedge \exists v . [a*z/y]R(x, y, v)$ (only the first y is quantified)
 $\equiv [a+b/z](\forall x . \exists y . R(x, y, z)) \wedge \exists v . R(x, a*z, v)$
 $\equiv ((\forall x . \exists y . R(x, y, a+b)) \wedge \exists v . R(x, a*(a+b), v))$ (parens around $a+b$ are required)
(No renaming necessary because we have no quantification of a or b .)