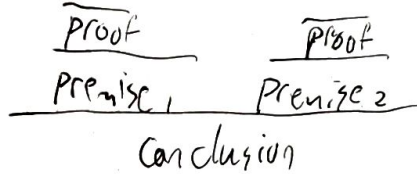


# Writing proofs

axioms

Proof trees:



Ex.  $x=y$

$$\frac{\frac{\frac{x=y}{x+1=y+1} \quad x:=x+1 \quad \{x=y+1\}}{\vdash \{x+1=y+1\} \quad x:=x+1 \quad \{x=y+1\}} \quad \frac{\frac{\frac{y+1}{x=y+2-1} \quad y:=y+2 \quad \{x=y-1\}}{\vdash \{x=y+2-1\} \quad y:=y+2 \quad \{x=y-1\}}}{\vdash \{x=y\} \quad x:=x+1; y:=y+2 \quad \{x=y-1\}}$$

$$\vdash \{x>0\} \text{ if } (x>0) \{z:=1\} \text{ else } \{z:=0\} \{z=1\}$$

$$\frac{\vdash \{T\} x:=1 \{x=1\} \text{ Assign}}{\vdash \{T\} x:=1 \{x=1\}}$$

$$\frac{\vdash \{T\} x:=1 \{x=1\} \neq \vdash \{x>0\} \text{ if } \dots \{z=1\}}{\vdash \{T\} x:=1; \text{ if } x>0 \{z:=1\} \text{ else } \{z:=0\} \{z=1\}}$$

$$\frac{\vdash \{p\} \leq \{q\} \quad p' \Rightarrow p \quad q \Rightarrow q'}{\vdash \{p'\} \leq \{q'\}} \text{ weaken}$$

Can always make precondition stronger      Can always make postcond weaker

$$\frac{\vdash \{1=1\} x:=1 \{x=1\} \text{ Assign} \quad T \Rightarrow 1=1 \quad x=1 \Rightarrow x>0}{\vdash \{T\} x:=1 \{x>0\}} \text{ weaken}$$

$$\frac{\text{Assign} \quad \frac{\vdash \{1=1\} z:=1 \{z=1\}}{\vdash \{x>0 \wedge x>0\} z:=1 \{z=1\}} \quad \text{weaken} \quad \frac{\vdash \{x>0 \wedge x \leq 0\} z:=0 \{F\}}{\vdash \{x>0\} z:=0 \{F\}}}{\vdash \{x>0\} z:=0 \{F\}} \text{ weaken}$$

$$\vdash \{T\} x:=1; \text{ if } x>0 \{z:=1\} \text{ else } \{z:=0\} \{z=1\}$$

## Alternative Proof styles

### Hilbert-style

1.  $\{l=1\} x:=1 \{x=1\}$  Assign
2.  $\{T\} x:=1 \{x>0\}$  Weaken 1
3.  $\{l=1\} z:=1 \{z=1\}$  Assign
4.  $\{x>0 \wedge x>0\} z:=1 \{z=1\}$  Weaken 3
5.  $\{F\} z:=0 \{F\}$  Assign
6.  $\{x>0 \wedge x \leq 0\} z:=0 \{F\}$  Weaken 5
7.  $\{x>0\}$  if  $(x>0)$  then  $\{z:=1\}$  else  $\{z:=0\} \{z=1\}$  If 4, 5
8.  $\{T\} x:=1; \text{ if } \dots \{z=1\}$  Sequence 2, 7

Proof Outlines - Annotate the program

$x := 1;$	$\{T\} \Rightarrow \{l=1\}$
if $(x>0)$ then {	$\{x=1\} \Rightarrow \{x>0\}$
$z:=1$	$\{x>0 \wedge x>0\} \Rightarrow \{l=1\}$
} else {	$\{z=1\}$
$z:=0$	$\{x>0 \wedge x \leq 0\} \Rightarrow \{F\}$
}	$\{F\}$
	$\{z=1 \vee F\} \Rightarrow \{z=1\}$

$\{T\}$  if  $(x > 0)$  <sup>then</sup>  $\{x := 1\}$  else  $\{x := 2\}$  ; if  $(y > 0)$  <sup>then</sup>  $\{z := 1\}$  else  $\{z := 0\}$   $\{z = 1\}$

1.  $\{1 > 0\} x := 1 \{x > 0\}$  Assign

2.  $\{2 > 0\} x := 2 \{x > 0\}$  Assign

3.  $\{T \wedge x > 0\} x := 1 \{x > 0\}$  Weaken 1

4.  $\{T \wedge x > 0\} x := 2 \{x > 0\}$  Weaken 2

5.  $\{T\}$  if  $x > 0$  then  $\{x := 1\}$  else  $\{x := 2\}$   $\{x > 0 \vee x > 0\}$  If 3, 4

6.  $\{1 = 1\} z := 1 \{z = 1\}$  Assign

7.  $\{x > 0 \wedge y > 0\} z := 1 \{z = 1\}$  Weaken 6

8.  $\{F\} z := 0 \{F\}$  Assign

9.  $\{x > 0 \wedge y \leq 0\} z := 0 \{F\}$  Weaken 8

10.  $\{x > 0\}$  if  $(x > 0)$  then  $\{z := 1\}$  else  $\{z := 0\}$   $\{z = 1 \vee F\}$  If 7, 9

11.  $\{x > 0\}$  if ...  $\{z = 1\}$  Weaken 10

12.  $\{T\}$  if  $(x > 0)$  ...  $\{x > 0\}$  Weaken 5

13.  $\{T\}$  if  $(x > 0)$  ... ; if  $(y > 0)$  ...  $\{z = 1\}$  Sequence 12, 11

Alternate rule for if: 
$$\frac{\{p \wedge e\} s_1 \{q_1\} \quad \{p \wedge \neg e\} s_2 \{q_2\}}{\{p\} \text{ if } e \text{ then } \{s_1\} \text{ else } \{s_2\} \{q\}}$$

Would let us conclude 12 w/o weakening

What if postconditions of  $s_1$  and  $s_2$  don't match? Weaken!

$\{x > 0 \wedge y \leq 0\} z := 0 \{z = 1\}$  Weaken 9 (b.c.  $F \Rightarrow z = 1$ )

No loss of generality since  $q_1 \Rightarrow q_1 \vee q_2$  and  $q_2 \Rightarrow q_1 \vee q_2$