

Types, Expressions, and States

CS 536: Science of Programming, Fall 2021

1. Which of the following expressions are legal or illegal according to the syntax we're using? Assume x, y, z are integer variables and b is an array name.
 - a. $(\text{if } x > y \text{ then } x \text{ else } y)$
 - b. $(\text{if } x < y \text{ then } -1 \text{ else } (\text{if } x = y \text{ then } 0 \text{ else } 1))$
 - c. $(\text{if } y = 0 \text{ then } f \text{ else } g)(17)$
 - d. $b[0][1]$ /* What type must b have for this to be legal? */
 - e. b /* Remember we're given that b is an array */
2. Which of the following are legal ways to write out a state? (And if not, why not?)
 - a. $\{x = 5, y = 2\}$
 - b. $\{x = \text{five}, y = \text{one plus one}\}$
 - c. $\{x = 5, y = x \text{ minus } 3\}$
 - d. $\{x = 5, y = \alpha - 3\}$ where $\alpha = 5$
 - e. $\{x = 5, y = (\text{the value of } x \text{ in this environment minus } 3)\}$
 - f. $\{ \}$
3. Let $e_4 \equiv x = y + 1 \wedge y = z^2 - 3 \wedge z = 6$. Write out the textual definition of a state σ_4 in which e_4 evaluates to true. Use only bindings that map variables to constants. $\sigma_4 = \{ x = 34, y = 33, z = 6 \}$
5. Which of the following states are well-formed and also proper for the expression $b[i] + 0 * y$? If ill-formed, why? If taking the value might cause a runtime error, why?
 - a. $\{i = 0, b = (3, 4, 8), y = 3, z = 5\}$
 - b. $\{i = 0, b = (6), y = 5\}$
 - c. $\{i = 0, b = 6, y = 5\}$
 - d. $\{i = 1, b = (3, 4, 8)\}$
 - e. $\{i = 1, i = 2, y = 0, b = (2, 6)\}$
 - f. $\{i = 5, b = (1, 2), y = 4\}$

CS 536 Solution to Practice 3 (Types, Expressions, and States)

1. (Legal and illegal expressions)
 - a. legal
 - b. legal
 - c. illegal because the conditional expression can't yield a function/operator
 - d. $b[0][1]$ is legal (b must be a 2-dimensional array)
 - e. b (all by itself) is illegal, since b we've assumed is an array
2. (Legal ways to represent states)
 - a. $\{x = 5, y = 2\}$ is legal
 - b. $\{x = \text{five}, y = \text{one plus one}\}$ is legal because "five" and "one" etc. refer to semantic objects.
 - c. $\{x = 5, y = x \text{ minus } 3\}$ is illegal: To be legal, " x minus 3" has to be a value, so " x " has to be a value (it has to be the name of a mathematical object like 5). But the binding $x = 5$ tells us " x " is a variable that can appear in an expression, so " x " is a syntactic object. It can't be syntactic and semantic at the same time.
Also, in this nicely word-processed document, " x " is presented in *this font*, so we know it's supposed to be a syntactic object. On paper, you see the difference between " x " and " x ". Even so, if someone wrote $\{ x = 5, y = x \text{ minus } 1 \}$ on the blackboard, it would have to be illegal because of using X in two incompatible ways.
 - d. $\{x = 5, y = \alpha - 3\}$ where $\alpha = 5$ — is legal. We infer that symbols x and y are syntactic objects and α is the name of the semantic object 5.
 - e. $\{x = 5, y = (\text{the value of } x \text{ in this environment, minus } 3)\}$ is legal. Since "the value of x in this environment" is just another name (albeit complicated) for the mathematical object 5, it's legal to use here.
 - f. $\{ \}$ is legal, since it's just another way to write \emptyset , the empty state.
3. $\sigma_4 = \{z = 6, y = 33, x = 34\}$
4. (Proper states)
 - a. (Well-formed and) Proper: The extra binding for z isn't a problem
 - b. (Well-formed and) Proper: The value of b is an array of length 1.
 - c. (Well-formed but) Improper: The value of b can't be an integer.

- d. (Well-formed but) Improper: We need a binding for y even though we're multiplying it by zero. [So our semantics uses eager evaluation, not lazy evaluation.]
- e. Ill-formed: We have two bindings for i .
- f. (Well-formed and) Proper but causes a runtime error, since b has size 2.