

Simply Typed Lambda Calculus (STLC)

$\tau ::= \text{mit } (\tau \rightarrow \tau) | \tau \times \tau | \tau + \tau | \text{void}$

$e ::= x \mid () \mid e \cdot e \mid \lambda x : t \cdot e \mid (e, e) \mid \text{fst} e \mid \text{snd} e \mid \text{inr } e \mid \text{inl } e$

Case e of Exercised aborte

Pre:c

$$\frac{\Gamma(x) = \tau \quad (\text{I-1})}{\Gamma \vdash x : \tau} \quad \frac{}{\Gamma \vdash () : \text{unit}} \quad (\text{I-2}) \quad \frac{\Gamma \vdash e_1 : \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2} \quad (\text{I-3})$$

$$\frac{r_1 x - r_2 + e = c'}{r_1 x - r_2 + e = c} \quad (T4)$$

$$\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2 \text{ (T-5)}}{\Gamma \vdash (e_1, e_2) : \tau_1 \times \tau_2} \quad \frac{\Gamma \vdash e : \tau_1 \times \tau_2 \text{ (T-6)}}{\Gamma \vdash \text{fst } e : \tau_1} \quad \frac{\Gamma \vdash e : \tau_1 \times \tau_2 \text{ (T-7)}}{\Gamma \vdash \text{snd } e : \tau_2}$$

$\text{dim} \mathbb{S}^2 = \int_x \times \int_y$

Person = string \times int
 ↑ ↑
 name age

test \in negative / positive \leftarrow enum (e.g. Java)

lolla: proof \rightarrow bool

proof: Vaccine of string \times int
or test of test

proof \in (string \times int) $^+$ test

test \equiv unit + unit

neg test \Leftarrow in negative

negative \in in 1 ()

pfizer-vax ≡ in | ("Pfizer", 2)

positive \geq inc ()

$\text{tolla} \in \lambda p:\text{proof. case } p \text{ of}$

{ vox. sind vox = 2. 11 test vox = "J & V";
test. test = negative }

$$\frac{\Gamma \vdash e : \tau_1}{\Gamma \vdash \text{inl } e : \tau_1 + \tau_2} \quad (T-8) \quad \frac{\Gamma \vdash e : \tau_2}{\Gamma \vdash \text{inr } e : \tau_1 + \tau_2} \quad (T-9)$$

$$\frac{\Gamma \vdash e : \tau_1 + \tau_2 \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma \vdash y : \tau_2 + e_2 : \tau}{\Gamma \vdash \text{case } e \text{ of } \{x.e_1, y. e_2\} : \tau} \quad (T-10)$$

Notice: Binary product $\tau_1 \times \tau_2$

- 1 way to create (pair)
- 2 things to do with it (fst, snd)

n-ary product $\tau_1 \times \dots \times \tau_n$

- 1 way to create (tuple)
- n things to do (projections)

Binary sum $\tau_1 + \tau_2$

- 2 ways to create (inl, inr)
- 1 thing to do (case)

n-ary sum $\tau_1 + \dots + \tau_n$

- n ways to create
- 1 thing to do

0-ary (nullary) product

- 1 way to create (supply nothing) = Unit!
- 0 things to do

nullary sum: 0 ways to create, 1 thing to do "void"

$$\frac{\Gamma \vdash e : \text{void}}{\Gamma \vdash \text{abort } e : \tau} \quad (T-11)$$

$$\frac{v \text{ val}}{(\lambda x : \tau e) v \mapsto [v/x]e}$$

$$(\lambda x. e) \text{ val} \quad (\lambda x. e) \text{ val}$$

$$c \text{ val}$$

$$\frac{v_1 \text{ val} \quad v_2 \text{ val}}{(v_1, v_2) \text{ val}}$$

$$\frac{e_1 \mapsto e'_1}{(e_1, e_2) \mapsto (e'_1, e_2)}$$

$$\frac{v_1 \text{ val} \quad e_2 \mapsto e'_2}{(v_1, e_2) \mapsto (v_1, e'_2)}$$

$$\frac{\dots e \mapsto e'}{f_3 t : e \mapsto f_3 t e'}$$

$$\frac{v_1 \text{ val} \quad v_2 \text{ val}}{f_3 t (v_1, v_2) \mapsto v_1}$$

$$\frac{e \mapsto e'}{\text{inl } e \mapsto \text{inl } e'}$$

$$\frac{v \text{ val}}{\text{inr } v \text{ val}}$$

$$\frac{e \mapsto e'}{\text{case } e \text{ of } \{x. e_1, y. e_2\} \mapsto \text{case } e' \text{ of } \{x. e'_1, y. e'_2\}}$$

$$\text{case inl } v \text{ of } \{x. e_1, y. e_2\} \mapsto [v/x]e_1$$

$$\text{case inr } v \text{ of } \{x. e_1, y. e_2\} \mapsto [v/y]e_2$$

Progress: If $\Gamma \vdash e : \tau$ then $e \text{ val}$ or $e \mapsto e'$

Preservation: If $\Gamma \vdash e : \tau$ and $e \mapsto e'$ then $\Gamma \vdash e' : \tau$.

Fun fact: If $\vdash e : \tau$ then $\exists v \text{ st. } v \text{ val and } e \mapsto^* v$
 (remember: not true for untyped calculus)

$$\frac{\Gamma, x : \tau \vdash x : \tau \quad \Gamma, x : \tau \vdash x : \tau}{\Gamma, x : \tau \vdash x x : \tau}$$

$$\tau = \tau \rightarrow \tau' \text{ what?}$$

$$\frac{}{\Gamma \vdash x. xx : \tau \rightarrow \tau}$$

$$\Gamma \vdash (\lambda x. xx) (\lambda x. xx) :$$