

# Syntactic Substitution

## CS 536: Science of Programming, Fall 2021

1. Calculate  $[b+c/c][i+1/i](x+i*b+c=0)$ .
2. Let  $p$  be  $\exists x. x < y \wedge x^2 \geq y+k$ 
  - a. What is  $[5/x]p$ ?
  - b. What is  $[5/y]p$ ?
  - c. What is  $[5/z]p$ ?
  - d. What is  $[y*2/y]p$ ?
  - e. What is  $[y*k/y]p$ ?
  - f. What is  $[(x+y) \div 2/y]p$ ?
3. Give an example where  $[e/v][e'/w](v * w)$  and  $[e'/w][e/v](v * w)$  are
  - a. Syntactically equal ( $\equiv$ )
  - b. Syntactically unequal ( $\not\equiv$ ).
4. In the predicate  $(\exists x. x < y \wedge x^2 \geq y+k)$ ,  $x$  is bound, but in  $(x < y \wedge x^2 \geq y+k)$ ,  $x$  is free — is this a contradiction?
6. Let  $p \equiv (\forall x. \exists y. R(x, y, z)) \wedge (\exists z. R(x, y, z))$  where  $R$  is a boolean function over three arguments.
  - a. What is  $[17/w]p$ ?
  - b. What is  $[17/x]p$ ?
  - c. What is  $[y*2/y]p$ ?
  - d. What is  $[y*2/z]p$ ?
  - e. What is  $[a+b/z][a*z/y]p$ ?

*Solution to Practice 12 (Syntactic Substitution)*

1.  $[b+c/c][i+1/i](x+i*b+c = 0) \equiv [b+c/c](x + (i+1)*b + c = 0)$   
 $\equiv x+(i+1)*b+(b+c) = 0$
2. Let  $p \equiv \exists x. x < y \wedge x^2 \geq y+k$ 
  - 2a.  $[5/x]p \equiv p$  unchanged
  - 2b.  $[5/y]p \equiv [5/y](\exists x. x < y \wedge x^2 \geq y+k) \equiv \exists x. x < 5 \wedge x^2 \geq 5+k$
  - 2c.  $[5/z]p \equiv p$  unchanged because  $z$  doesn't occur in  $p$
  - 2d.  $[y^2/y]p \equiv [y^2/y](\exists x. x < y \wedge x^2 \geq y+k) \equiv \exists x. x < y^2 \wedge x^2 \geq y^2+k$
  - 2e.  $[y*k/y]p \equiv [y*k/y](\exists x. x < y \wedge x^2 \geq y+k) \equiv \exists x. x < y*k \wedge x^2 \geq y*k+k$
  - 2f.  $[(x+y) \div 2/y]p \equiv [(x+y) \div 2/y](\exists x. x < y \wedge x^2 \geq y+k)$   
 $\equiv [(x+y) \div 2/y]\exists v.[v/x](x < y \wedge x^2 \geq y+k)$  (note renaming of  $x$  to  $v$ )  
 $\equiv [(x+y) \div 2/y] \exists v.(v < y \wedge v^2 \geq y+k)$   
 $\equiv \exists v. v < (x+y) \div 2 \wedge v^2 \geq (x+y) \div 2 + k$
3. (Cases where  $[e'/w][e/v](v * w)$  and  $[e/v][e'/w](v * w)$  are  $\equiv$  and  $\neq$ .)
  - 3a. One case is when  $v$  doesn't occur in  $e'$  and  $w$  doesn't occur in  $e$ .  
Example:  $[a*w/w][v^2/v](v * w) \equiv [a*w/w](v^2 * w)$   
 $\equiv v^2 * (a*w) \equiv [v^2/v](v * (a*w))$   
 $\equiv [v^2/v][a*w/w](v * w)$
  - 3b. One case is when  $w$  appears in  $e$  and  $v$  appears in  $e'$ , at least, for certain  $e$  and  $e'$ .  
Example:  $[a*v/w][w-3/v](v * w) \equiv [a*v/w]((w-3) * w) \equiv (w-3) * (a * v)$   
but  $[w-3/v][a*v/w](v * w) \equiv [w-3/v](v * (a*v)) \equiv (w-3) * (a * (w-3))$
4. No, this is exactly what a quantifier does: It captures the  $x$ 's that are free in its body and makes them bound with respect to any context that includes the quantified predicate.
6. Substitutions with  $p \equiv (\forall x. \exists y. R(x, y, z)) \wedge \exists z. R(x, y, z)$ :
  - 6a.  $[17/w]p \equiv p$  (because  $w$  doesn't occur in  $p$ )
  - 6b.  $[17/x]p \equiv (\forall x. \exists y. R(x, y, z)) \wedge \exists z. R(17, y, z)$
  - 6c.  $[y^2/y]p \equiv (\forall x. \exists y. R(x, y, z)) \wedge \exists z. R(x, y^2, z)$
  - 6d.  $[y^2/z]p \equiv (\forall x. \exists v. [y^2/z][v/y]R(x, y, z)) \wedge \exists z. R(x, y, z)$  (using  $v$  as a fresh variable)  
 $\equiv (\forall x. \exists v. R(x, v, y^2)) \wedge \exists z. R(x, y, z)$
  - 6e.  $[a+b/z][a*z/y]p$   
 $\equiv [a+b/z](\forall x. \exists y. R(x, y, z)) \wedge \exists v. [a*z/y][v/z]R(x, y, z)$  (using  $v$  as a fresh variable)  
 $\equiv [a+b/z](\forall x. \exists y. R(x, y, z)) \wedge \exists v. [a*z/y]R(x, y, v)$  (only the first  $y$  is quantified)  
 $\equiv [a+b/z](\forall x. \exists y. R(x, y, z)) \wedge \exists v. R(x, a*z, v)$   
 $\equiv ((\forall x. \exists y. R(x, y, a+b)) \wedge \exists v. R(x, a*(a+b), v))$  (parens around  $a+b$  are required)  
(No renaming necessary because we have no quantification of  $a$  or  $b$ .)