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On The Construction of Codes over Term Rank Metric Spaces

Sümeýra Bedir

1 Introduction

The vector space of $m \times n$ matrices over a fixed finite field F_q of q elements become a metric space denoted by M_{TR} , with the term rank metric derived from the term rank weight defined below, given A as an $m \times n$ matrix with $\mathcal{I}(A)$ being the set of rows/columns of A which contains all the nonzero entries of A ;

$$\|A\|_{TR} = \min |\mathcal{I}(A)|$$

If A and B are two $m \times n$ matrices, the term rank distance is defined as;

$$d_{TR} = \|A - B\|_{TR}$$

Codes over term rank metric spaces have applications in information transmission via memoryless matrix channels which appear in the data storage systems, memory cards and some wireless communication systems.[8,9] These codes are considered as k -dimensional vector subspaces of $F_q^{m \times n}$. The minimum distance of a code over term rank metric space, denoted by D_{TR} , should clearly be less than or equal to the minimum of $\{m, n\}$ and assuming without the loss of generality that $m \leq n$, we have;

$$D_{TR} = \min_{A \in C - \{0\}} \|A\|_{TR} \leq m$$

The only known bound for optimality of codes over M_{TR} is the Singleton bound, which is expressed in the following version;

$$k \leq n(m - D_{TR} + 1)$$

If we have the equality, the code is considered to be optimal.

In [1], authors have introduced a construction method for optimal codes over M_{TR} , by using the correspondence between polynomials and $p(x)$ -circulants. With this method, they construct $[n \times n, n]$ -codes, and for the construction of $[m \times n, n]$ -codes, they address the shortening method. In this study, we introduce a new construction, which will guide as an analogue to the usual construction of codes over vector spaces, and will give a way for direct construction of $[m \times n, n]$ -codes over term rank spaces in general. We also propose a script for finding the minimum term rank distance of a given code using Python packages.

2 Code Construction and Examples

Let $p(x) = a_0 + a_1x + \dots + x^m$ be a monic divisor of degree m of a polynomial $f = x^n - 1$ of degree n and consider the following matrix P obtained from the companion matrix of p horizontally joined with an $m \times (n - m)$ block matrix of zeroes.

$$P = \begin{bmatrix} 0 & 1 & 0 & \dots & 0 & \dots & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & \dots & 0 & 1 & 0 & \dots & 0 \\ -a_0 - a_1 \dots - a_{m-1} & 0 & \dots & 0 \end{bmatrix}$$

We define a cyclic shift by vertically shifting the columns of P to the right hand side. We can obtain this shift by successive operations of the companion matrix of f to the matrix P .

The F_q sub matrix space spanned by n matrices of cyclic shifts of P construct a form of a cyclic code over M_{TR} .

We generalize this to the pseudo-cyclic case as follows;

Theorem 1 *Let F_q be a finite field with q elements and let f be a monic polynomial and p a divisor polynomial of f over $F_q[x]$, with $\deg f = n$ and $\deg p = m$. Let P be the matrix obtained from the companion matrix of p horizontally joined with an $m \times (n - m)$ block matrix of zeroes, and T be the companion matrix of f . The F_q -sub matrix space spanned by the following set of $m \times n$ matrices;*

$$\{P, PT, PT^2, \dots, PT^{n-1}\}$$

constructs a pseudo-cyclic code over the term rank metric space M_{TR} .

Example 2 *Let F_4 be the finite field with 4 elements; $F_4 = \{0, 1, a, a^2\}$. Consider $f(x) = x^9 - 1$ and take $p(x) = x^3 + a^2$ as a divisor of f .*

Therefore we have $m = 3$, $n = 9$, and

$$P = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ a^2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}, T = \begin{bmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \ddots & \vdots \\ 0 & \vdots & \ddots & \ddots & 0 \\ 0 & 0 & \dots & 0 & 1 \\ 1 & 0 & \dots & \dots & 0 \end{bmatrix}_{9 \times 9}$$

Applying T to P , constructs the desired cyclic shift;

$$PT = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & a^2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix},$$

$$PT^2 = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & a^2 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix},$$

$$\vdots$$

$$PT^8 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & a^2 \end{bmatrix}$$

And the subspace generated by the spanning set $\{P, PT, PT^2, \dots, PT^8\}$ becomes a cyclic $[3 \times 9, 9]$ -code over the F_4 -matrix space of 3×9 matrices.

3 Computing The Minimum Term Rank Distance

As in the case in general coding theory, computing the minimum distance and obtaining optimal codes is an important issue also for codes over term rank metric spaces. It is already known that, the term rank weight of a matrix A is equal to the maximum size of a matching of the bipartite graph for which A is the bi-adjacency matrix[9]. Currently, there is not any in-built function for computing the term rank of a matrix in commonly used computer algebra systems. As an example for codes over F_4 -matrix spaces, we used the following Python packages and scripts for computing the minimum term rank distance of a code over M_{TR} , where we print the list of all matrices in the code from MAGMA computer algebra system to a text file as follows;

```
//Magma Code for writing matrices to a file:
K<a>:=GF(2^2); F<x>:= PolynomialRing(K);
p:=; f:=;
m:=Degree(p);n:=Degree(f);
```

```

T:= CompanionMatrix(f); V:= KMatrixSpace(K,m,n);
M:=MatrixRing(K,n);
Z1 := [0: x in [1..m*(n-m)]];
P:= HorizontalJoin(CompanionMatrix(p),Matrix(K, m,
n-m, Z1));
B := { V!P*T^i : i in [0..n-1]};
S:=sub< V | B >;
SetOutputFile("mxn.txt");
for s in S do
for i in [1..m] do
for j in [1..n] do
print s[i,j];end for;
print "$";end for;
print "@";end for;
UnsetOutputFile();

##Python script for computing minimum term rank distance:
import numpy as np
import networkx as nx
from networkx.algorithms import bipartite
import itertools
from networkx.convert import _prep_create_using
from networkx.convert_matrix import _generate_weighted_edges
import scipy
from scipy import linalg
n = 11;m = 5
fname = "mxn.txt"
fhand = open(fname)
L = list();S = str()
for line in fhand:
    line = line.strip()
    if "a^2" in line:
        line = line.replace("a^2","1")
    elif "a" in line:
        line = line.replace("a","1")
    S = S + line
M = S.strip().split("@")
for s in M:
    s = s.split("$")
    L.append(s)
for l in L:
    if len(l) < m+1 :

```

```

L.remove(l)
else:
    l.remove("")
K = list()
for item in L:
    M = list()
    for i in range(m):
        M.append([int(r) for r in item[i]])
    A = scipy.sparse.csr_matrix(M)
    G = nx.bipartite.from_biadjacency_matrix(A)
    D = nx.bipartite.maximum_matching(G)
    termrank = int(len(D.items())/2)
    if termrank != 0:
        K.append(termrank)
print "D_tr = ", min(K)

```

The cyclic code in the above example has a minimum term rank distance of 3, and therefore it is optimal.

Example 3 Let F_q be the finite field with 4 elements; $F_4 = \{0, 1, a, a^2\}$. Consider $f(x) = x^6 + a^2x^2 + a$ and take $p(x) = x^4 + x^2 + a$ as a divisor of f . We have $m = 4$, $n = 6$, and

$$P = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ a & 0 & 1 & 0 & 0 & 0 \end{bmatrix}, T = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ a & 0 & a^2 & 0 & 0 & 0 \end{bmatrix}$$

Applying T to P , constructs a pseudo-cyclic shift as follows;

$$PT = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & a & 0 & 1 & 0 & 0 \end{bmatrix}, PT^2 = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & a & 0 & 1 & 0 \end{bmatrix}, PT^3 = \begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ a & 0 & a^2 & 0 & 0 & 0 \\ 0 & 0 & 0 & a & 0 & 1 \end{bmatrix},$$

$$PT^4 = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 1 \\ a & 0 & a^2 & 0 & 0 & 0 \\ 0 & a & 0 & a^2 & 0 & 0 \\ a & 0 & a^2 & 0 & a & 0 \end{bmatrix}, PT^5 = \begin{bmatrix} a & 0 & a^2 & 0 & 0 & 0 \\ 0 & a & 0 & a^2 & 0 & 0 \\ 0 & 0 & a & 0 & a^2 & 0 \\ 0 & a & 0 & a^2 & 0 & a \end{bmatrix}$$

And the subspace generated by the spanning set $\{P, PT, PT^2, PT^3, PT^4, PT^5\}$ becomes a pseudo-cyclic $[4 \times 6, 6]$ -code over the F_4 -matrix space of 4×6 matrices.

4 Future Studies

- We aim to find the restrictions and conditions for constructing optimal codes with the proposed construction method.
- New bounds (implementations of sphere packing-covering bounds) are to be explored as in the usual vector space case.

5 References

References

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