



**DOCTORATE
THESIS MONITORING COMMITTEE PERIODIC EVALUATION REPORT**

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NOTE: The thesis committee members who have additional personal remarks may specify them with a report.

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Skew Polynomial Rings

- The theory of noncommutative polynomial rings was first introduced by Oystein Ore (1933), Nathan Jacobson (1943) and Bernard R. McDonald (1974).
- Let F be a finite field with characteristic p and let θ be an automorphism of F with $|\langle \theta \rangle| = m$. If K is the fixed subfield of F under θ , then $[F:K] = m$.

Example. Consider $F_4 = \{0, 1, \alpha, \alpha^2\}$ with the Frobenius automorphism $\theta: F_4 \rightarrow F_4$ where $\theta(\alpha) = \alpha^2$. We have $\theta(0) = 0, \theta(1) = 1, \theta(\alpha) = \alpha^2, \theta(\alpha^2) = \alpha$. Therefore $|\langle \theta \rangle| = 2$ and F_2 is the fixed subfield of F_4 under θ .

Definition. (B.R. McDonald, 1974) The set of skew polynomials

$$F[x; \theta] = \{a_0 + a_1x + \cdots + a_nx^n : a_i \in F\}$$

becomes a ring with the usual polynomial addition and the multiplication defined below;

$$(ax^i) * (bx^j) = a\theta^i(b)x^{i+j}$$

Example. Let $F_4 = \{0, 1, \alpha, \alpha^2\}$ with the Frobenius automorphism $\theta: F_4 \rightarrow F_4$ where $\theta(\alpha) = \alpha^2$.

$$(\alpha x) * (\alpha^2 x) = \alpha \theta(\alpha^2) x^2 = \alpha \alpha x^2 = \alpha^2 x^2$$

$$(\alpha^2 x) * (\alpha x) = \alpha^2 \theta(\alpha) x^2 = \alpha^2 \alpha x^2 = \alpha x^2$$

Note that, $F_4[x; \theta]$ is a noncommutative ring.

Important Notes on Skew Polynomial Rings

- Skew polynomial rings are noncommutative and non UFD.
- Right division algorithm holds in skew polynomial rings.
- Every right/left ideal of $F[x; \theta]$ is a principal ideal.
- I is a two sided ideal of $F[x; \theta]$, if there exist polynomials $f, g \in F[x; \theta]$ such that $I = f * F_4[x; \theta]$ and $I = F_4[x; \theta] * g$.
- If $|\langle \theta \rangle| = m$ and K is the fixed subfield of F under θ , the center of $F[x; \theta]$ is defined as $Z(F[x; \theta]) = \{a_0 + a_1x^m + \cdots + a_rx^{mr} : a_i \in K\}$.
- Any two sided ideal of $F[x; \theta]$ is generated by a polynomial in $Z(F[x; \theta])$.

Skew Cyclic Codes

Definition. (Boucher et al., 2007) Let F be a finite field with automorphism θ . A subspace C of F^n is considered as a skew cyclic code if for any vector $c = (c_0, c_1, \dots, c_{n-1}) \in C$, $(\theta(c_{n-1}), \theta(c_0), \dots, \theta(c_{n-2})) \in C$.

Skew Pseudo-Cyclic Codes

Definition. Let F be a finite field with automorphism θ , and let τ_v be a pseudo-cyclic shift defined on F^n . A subspace C of F^n is considered as a skew pseudo-cyclic code, if for any $c = (c_0, c_1, \dots, c_{n-1}) \in C$, $(\theta \circ \tau_v(c_{n-1}), \theta \circ \tau_v(c_0), \dots, \theta \circ \tau_v(c_{n-2})) \in C$.

Quasi-cyclic Codes

Theorem. Let C be a 1-generator quasi-cyclic code. Then C is generated by an element $\mathcal{g}(x) = (p_1(x)g_1(x), \dots, p_r(x)g_r(x))$ where $p_i(x), g_i(x) \in F[x]/\langle(x^{m_i} - 1)\rangle$ and $g_i(x)|x^{m_i} - 1 \forall i$. Let $h_i(x) = (x^{m_i} - 1)/g_i(x)$ and $\gcd(p_i(x), h_i(x)) = 1 \forall i$. Then, $\dim C = \deg(\text{lcm}(h_1(x), \dots, h_r(x)))$ and $d(C) \geq \sum_i d(\langle g_i(x) \rangle)$.

Quasi-pseudo-cyclic Codes

Theorem. Let C be a 1-generator quasi pseudocyclic code. Then C is generated by an element $\mathcal{g}(x) = (p_1(x)g_1(x), \dots, p_r(x)g_r(x))$ where $p_i(x), g_i(x) \in F[x]/\langle f_i(x) \rangle$ and $g_i(x)|f_i(x) \forall i$. Let $h_i(x) = f_i(x)/g_i(x)$ and $\gcd(p_i(x), h_i(x)) = 1 \forall i$. Then, $\dim C = \deg(\text{lcm}(h_1(x), \dots, h_r(x)))$ and $d(C) \geq \sum_i d(\langle g_i(x) \rangle)$.

Skew Quasi-pseudo-cyclic Codes

Theorem. Let C be a 1-generator skew quasi pseudocyclic code over F . Then C is generated by an element $\mathcal{g}(x) = (p_1(x)g_1(x), \dots, p_r(x)g_r(x))$ where $p_i(x), g_i(x) \in F[x; \theta]/\langle f_i(x) \rangle$ and $g_i(x)$ is a right divisor of $f_i(x) \forall i$. Let $h_i(x) = f_i(x)/g_i(x)$ and $\text{rgcd}(p_i(x), h_i(x)) = 1 \forall i$. If $f_i(x) \in Z(F[x; \theta])$, then $\dim C = \deg(\text{llcm}(h_1(x), \dots, h_r(x)))$ and $d(C) \geq \sum_i d(\langle g_i(x) \rangle)$.

Code Snippet and an example of Skew quasi-pseudo-cyclic codes

```
R1<t>:=PolynomialRing(GF(2));
p:=t^2+t+1;
F<a>:=ext< GF(2) | p >;
R2<y>:=PolynomialRing(F);
R<X>:=TwistedPolynomials(F;q:= 2); // teta maps a to a^t
teta:= hom< F -> F | a^2 >;

n1:=29;
V1:=VectorSpace(F,n1);

g1s:=[a^2,a,a^2,0,a^2,a,1,a,1,1,1,a,1,1,1,0,0,0,0,0,0,0,0,0,0,0,0]; // coefficient sequence for g1
g1v:=V1!g1s; // coefficient vector for g1
g1:=R!g1s; // g1 as a skew polynomail

h1s:=[1,1,a,0,1,0,0,0,0,0,0,0,0,0,0,1,0,0,0,0,0,0,0,0,0,0,0]; // coefficient sequence for h1
h1:=R!h1s; // h1 as a skew polynomial

f1:=h1*g1;
f1s:=ElementToSequence(f1);

//p1
p1s:=[1,0,0,0,1,0,0,0,1,0,0,0,1,0,0,0];
p1:=R!p1s;
q1,r1:=Quotrem(p1*g1,f1); "r1=", r1; u1s:=ElementToSequence(r1) cat [0:x in [1..n1-1-Degree(r1)]];

//p2
p2s:=[1,0,1,0,0,0,0,0];
p2:=R!p2s;
q2,r2:=Quotrem(p2*g1,f1); "r2=",r2; u2s:=ElementToSequence(r2) cat [0:x in [1..n1-1-Degree(r2)]];

T1:=CompanionMatrix(R2!f1s);
"T1=",T1;

// 1st part
G1:=Matrix(F,n1,n1,[]);
G1[1]:=V1!u1s;
for i in [1..n1-1] do
G1[i+1]:=(V1!([teta(G1[i,j]): j in [1..n1]]))*T1;
end for;
"G1=",G1;
MinimumDistance(LinearCode(EchelonForm(G1)));

// 2nd part
G2:=Matrix(F,n1,n1,[]);
G2[1]:=V1!u2s;
for i in [1..n1-1] do
G2[i+1]:=(V1!([teta(G2[i,j]): j in [1..n1]]))*T1;
end for;
"G2=",G2;
MinimumDistance(LinearCode(G2));

G:=HorizontalJoin(G1,G2);G;
C:=LinearCode(G); C;
MinimumDistance(C);
```

1st Part;

$$f1 = X^{29} + X^{28} + X^{27} + a^2X^{26} + X^{25} + X^{24} + X^{23} + a^2X^{22} + X^{21} + \\ a^2X^{20} + aX^{19} + X^{18} + a^2X^{17} + a^2X^{15} + a^2X^{14} + aX^{13} + X^{12} + \\ a^2X^{11} + a^2X^{10} + X^9 + a^2X^8 + aX^5 + X^4 + X^3 + X^2 + a^2$$

$$g1 = X^{14} + X^{13} + X^{12} + aX^{11} + X^{10} + X^9 + X^8 + aX^7 + X^6 + aX^5 + \\ a^2X^4 + a^2X^2 + aX + a^2$$

$$h1 = X^{15} + X^4 + aX^2 + X + 1$$

$$p1 = X^{12} + X^8 + X^4 + 1$$

$$f1 = h1 * g1, \text{rgcd}(p1, h1) = 1,$$

$$p1 * g1 \bmod f1 = X^{26} + X^{25} + X^{24} + aX^{23} + X^{18} + aX^{17} + \\ a^2X^{16} + aX^{14} + a^2X^{10} + X^9 + X^8 + aX^7 + aX^6 + \\ a^2X^2 + aX + a^2$$

1st Part \rightarrow [29,15,8] code

2nd Part; $f1, g1, h1$ same with part 1

$$p2 = X^2 + 1$$

$$p2 * g1 \bmod f1 = X^{16} + X^{15} + a^2X^{13} + a^2X^{11} + a^2X^9 + aX^6 + \\ aX^5 + aX^3 + aX + a^2$$

2nd Part \rightarrow [29,15,8] code

$C1|C2 \rightarrow$ [58,15,23] code