

CS 3250 Algorithms **Lecture 5**

Graphs: BFS & Detecting Cycles



Announcements

- HW1B Due Tuesday, September 12th by 9AM.
 - Brightspace Quiz on Asymptotic Notation It's tricky by design. Be sure to think before answering. You get (2) tries at it. Your highest score is recorded.
 - Gradescope Written HW A few written questions on the topic of asymptotic analysis.
 Submitted to Gradescope as a PDF.





Data Structures for Graphs

1. Adjacency Matrix –

- **Space:** An array of n x n is needed $\Theta(n^2)$ regardless of a sparse or dense graph in terms of number of edges.
- Speed: Depends on the question being asked.
 - Is there an edge from vertex 4 to vertex 8? Fast. Θ(1)
 - List all the neighbors of every vertex? Slow. Θ(n²)
 - Delete vertex 4? Slow. O(n²)

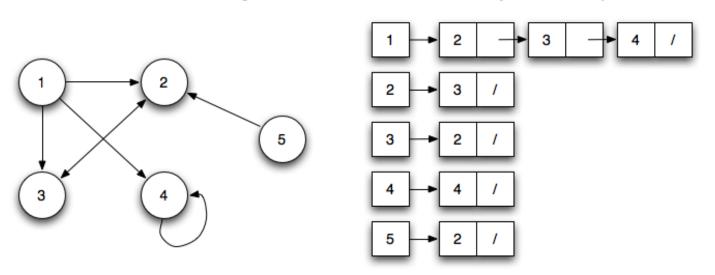
	1	2	3	4	5
1	0	1	1	1	0
2	0	0	1	0	0
3	0	1	0	0	0
4	0	0	0	0	0
5	0	1	0	0	0



Data Structures for Graphs

There are two main data structures used to represent graph G = (V, E) which contains n vertices and m edges.

2. Adjacency List - An adjacency list consists of an array of |V| pointers, where the ith element points to a linked list of all the edges incident to vertex i. For a weighted graph store the edge weight as an additional field in each edge node in the adjacency list.

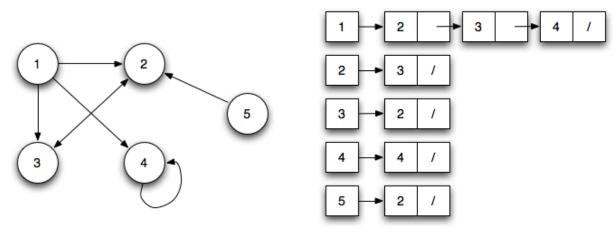




Data Structures for Graphs

2. Adjacency List –

- Space: O(|V| + |E|) = O(V + E) = O(max(V, E))
 - Cardinality symbol usually not included.
- Speed: Depends on the question being asked.
 - Is there an edge from vertex 4 to vertex 8? O(V) since a vertex might be adjacent to all other vertices.
 - List all the neighbors of every vertex? Θ(V + E)





Adjacency Matrix vs. Adjacency List

FACE OFF

Task	Winner	
Does edge (x, y) exist?	Adjacency Matrix	
What is the degree of vertex 4?	Adjacency List	
Add/Delete an edge	Adjacency Matrix	
Traverse the graph	Adjacency List	
Better for most problems	Adjacency List	



- Once you construct a graph, the next question becomes to do with the graph?
- Many questions are answered by exploring the graph.
- This idea of exploring a graph via its data structure is better known as a traversal.





- When traversing a graph, it's important to ensure the following:
 - 1. Efficiency We don't waste time visiting places we've already explored.
 - 2. Accuracy/Correctness We don't miss any places along the way (i.e., leave no stone unturned).



- One of the simplest graph traversals that can be used on a directed or undirected graph is known as **Breadth-First Search (BFS).** You saw it in 2201.
- Analogy: Imagine yourself as an explorer.
 - You gradually expand the frontier between discovered and undiscovered lands.
 - "lands" we discover and explore are the vertices
 - "paths" we take are the edges in the graph.





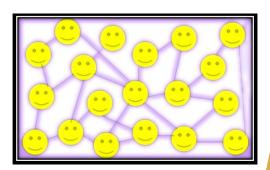
- Nobody (including your tech interviewer) will ever walk up to you and say, "I'll give you a bag of cash if you can write a breadth-first search for me."
- Breadth-first search shows up in numerous realworld applications (and job interview questions).
 - A user will describe a problem to you.
 - Your job is to recognize BFS is the right tool.





Graphs: Breadth-First Search in Action

- Breadth-First Search in Action
 - Shortest path in unweighted graph
 - Cycle detection in undirected graph



- Connected components
- Social Networks Who are Kristin's friends?
- GPS Navigation Systems What interesting places are near my current location?
- Search How many sites link to this site?



- The key to any graph traversal is keeping track of where you are.
- Question: Which of the following do you think is not a "state" that needs to be remembered during a BFS traversal?
 - Undiscovered
 - Discovered
 - Soon to be processed
 - Processed





- Question: Which of the following do you think is not a "state" that needs to be remembered during a BFS traversal?
 - Undiscovered
 - Discovered
 - Soon to be processed
 - Processed
- We don't need to worry about soon be processed.





- The key to any graph traversal is:
 - 1. Marking each vertex when first discovered
 - 2. Knowing when each vertex is **finished/processed**.
- We can also leave a breadcrumb to remember our "parent" (the person who discovered us). This is optional but can provide useful information.

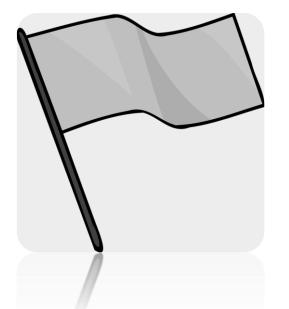


- We will label each vertex to be in one of three states (or colors).
 - 1. Undiscovered (white) I haven't yet discovered this vertex. It is unchartered territory (initially all vertices are undiscovered).



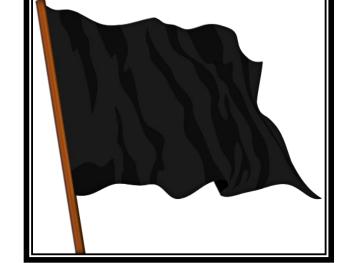


- We will label each vertex to be in one of three states (or colors).
 - 2. Discovered but unexplored (gray) I have discovered this vertex and planted my gray flag here. However, I haven't thoroughly explored the area yet.





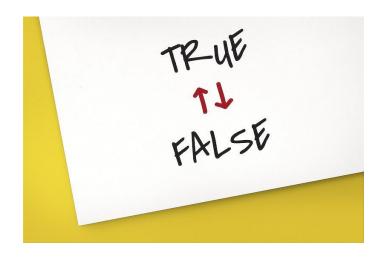
- We will label each vertex to be in one of three states (or colors).
 - 3. Processed/Explored (black) This vertex has not only been discovered, but I have also explored every aspect of this vertex, so I have finished processing it.





Breadth-First Search:

• Question: True/False. Every graph has a unique BFS traversal.

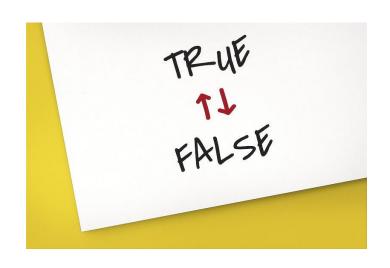






Breadth-First Search:

- Question: True/False. Every graph has a unique BFS traversal.
- Answer: False. It's quite possible to have multiple BFS traversals. That's why we often say on an exam or HW "when given a choice, choose the vertex that comes first alphabetically."

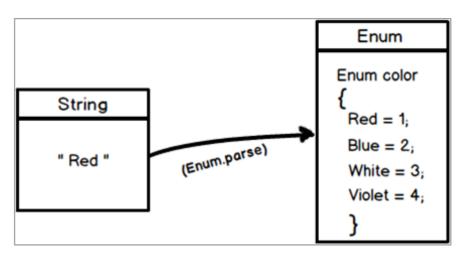






Breadth-First Search: Algorithm Details

- How should we implement these vertex states?
 - Some programmers use Enum types or Constants.
 - Some programmers use boolean arrays.
- The details do not matter provided they do not change the overall algorithm efficiency.



```
False False False False
True True True True
False True True True
False False True True
True True
True
```



Breadth-First Search: Pseudocode

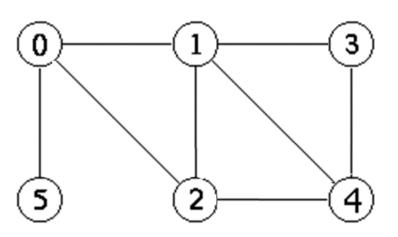
```
BFS(G, s) [Initially all vertices undiscovered]
 Set start vertex s discovered and set parent to nil,
 Add s to the TO-DO list (enqueue)
 WHILE (there are vertices on the TO-DO list)
    v = take a vertex off the To-Do list (dequeue)
    Optional: Do early vertex processing
    FOR (all of v's neighbors w)
          IF (state[w] == UNDISCOVERED/white) THEN
                 Change status of vertex w to DISCOVERED/gray)
                 Add w to the TO-DO list (enqueue)
                 Optional: Record the level where w was discovered
                 Optional: Record the parent/predecessor of w, which is v
          Move to the next neighbor w of v
    LOOP
 Change status of vertex v to PROCESSED/black
 Optional: Do late vertex processing
```



LO_OP

Example: Breadth-First Search

• Example: Perform a BFS on the undirected graph below starting at vertex 0. When given a choice between two vertices, choose the one that comes first numerically.



TO DO LIST

STATE

PROCESSED:



Breadth-First Search:

- Question: Does our breadth-first search algorithm work correctly for any undirected graph G?
 - A. Yes of course. We're experienced programmers.
 - B. Definitely not. We're experienced programmers.
 - C. I thought it did, but now I'm not sure.



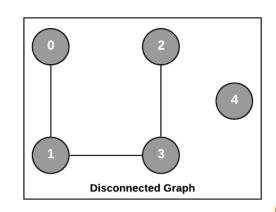




Breadth-First Search:

- Answer: Definitely not. Our current BFS algorithm does not work correctly for all types of undirected graphs. Consider the disconnected graph G shown below and a BFS traversal that starts at vertex 0. Our current algorithm would never find vertex 4.
- We need to add a controlling loop to ensure we hit all vertices in the graph.







Breadth-First Search: Optional Tasks

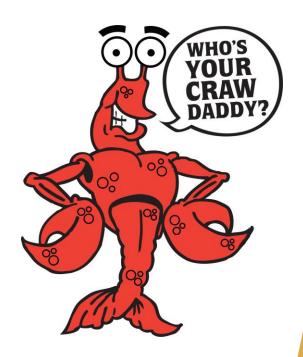
- You may decide to perform additional work during your graph traversal depending on your needs.
- This is fine provided you do not sacrifice overall algorithm efficiency.
- Here are some tasks commonly performed with BFS.





Breadth-First Search: Optional Tasks

- Task: Who's Your Daddy?
 - What: Who discovered us (i.e., our "parent")?
 - When: Helpful if you need to report the shortest path from one vertex to another.





Breadth-First Search: Optional Tasks

- Task: Level Up!
 - What: Remember every vertex's distance from the starting vertex (e.g., "six degrees of separation").
 - When: Helpful if you need to determine the shortest path from one vertex to another.





Breadth-First Search: Optional Details

Placeholders for tasks in our BFS algorithm:

- processVertexEarly(v) Optional method you can include if you need to do prework before exploring a vertex.
- 2. processVertexLate(v) Optional method you can include if you need to do postwork after exploring a vertex.
- 3. processEdge (v, w) Optional method you can include if you need to do any work while processing an edge.





Now, let's THINK about the algorithm as an algorist might.

 Question: Given an undirected graph, G, how many times can we discover a particular vertex?





Let's THINK about the algorithm as an algorist might.

- Question: Given an undirected graph, G, how many times can we discover a particular vertex?
- Answer: Once. I can only discover it once at which point, I plant my flag in it. If I (or anyone else) sees that land again, they will see my flag and know it's already been discovered.





Questions: Now, THINK about the algorithm as an algorist might. Given an undirected graph, G:

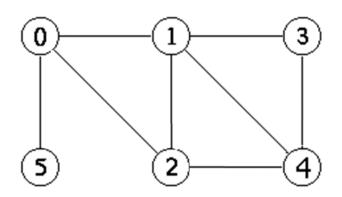
- 1. How many times can we discover a vertex?
- 2. Could a neighbor of the active vertex already be...
 - a. Discovered but not processed (gray)?
 - b. Discovered and processed (black)?
- 3. What color are the vertices on the TO-DO queue?
- 4. Can there be more than one copy of a vertex on the queue?
- 5. How can you tell if there's a cycle in an undirected graph?



Questions: Now, THINK about the algorithm as an algorist might. Given an undirected graph, G:

- 2. Could a neighbor of the active vertex be...
 - a. Discovered but not processed (gray)?

Answer: Yes. In the graph below, consider vertex 1 which we discover while actively working on vertex 0, but will not be processed until later in the BFS.

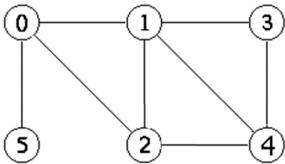




Questions: Now, THINK about the algorithm as an algorist might. Given an undirected graph, G:

- 2. Could a neighbor of the active vertex be...
 - b. Discovered and processed (black)?

Answer: Yes, of course. At some point during the BFS, the vertex 0 will be flagged as both discovered and then processed.

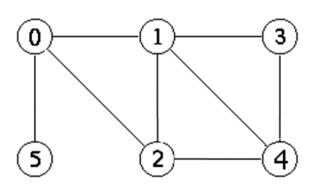




Questions: Now, THINK about the algorithm as an algorist might. Given an undirected graph, G:

3. What color are the vertices on the TO-DO queue? **Answer: They are gray** because they are discovered but not processed.





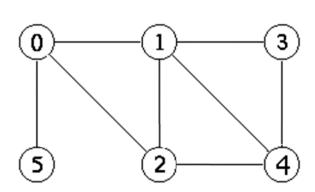


Questions: Now, THINK about the algorithm as an algorist might. Given an undirected graph, G:

4. Can there be more than one copy of a vertex on the queue?

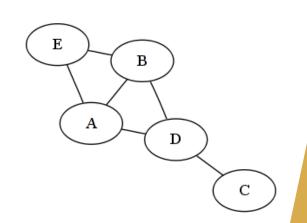
Answer: No. Once a vertex is placed on the queue, it is marked gray/discovered. As we've previously noted, a vertex can only be discovered once.







- Question: 5. How can we detect a cycle in an undirected graph?
- **Answer:** If we encounter a vertex during traversal that is in the discovered state, but "we" didn't discover it, we've found a cycle (i.e., got to same vertex via a different edge).





- Let's expand on this idea and modify our BFS to detect cycles in a connected, undirected graph.
- Hmm...what do we need to do and where?
 - If we are looking at an edge that has not been processed, we should ask whether this edge leads to a vertex that has already been discovered by someone else.
 - If that vertex has already been discovered by someone else, then there must be two ways to get to it.
 - That means we just found a cycle.



```
BFS(G, s) [Initially all vertices undiscovered]
 Set start vertex s discovered and set parent to nil,
 Add s to the TO-DO list (enqueue)
 WHILE (there are vertices on the TO-DO list)
    v = take a vertex off the To-Do list (dequeue)
    FOR (all of v's neighbors w)
          IF (state[w] != PROCESSED) THEN
                 processEdge(v, w)
          IF (state[w] == UNDISCOVERED) THEN
                 Change status of vertex w to DISCOVERED)
                 Add w to the TO-DO list (enqueue)
                 Record the parent/predecessor of w, which is v
          Move to the next neighbor w of v
    LOOP
 Change status of vertex v to PROCESSED
 L00P
```





END BFS

Now for the important piece...

```
processEdge(x, y)
    IF ((state[y] == DISCOVERED) THEN
        PRINT "Cycle found:" + x + "," + y)
END Method
```



 Question: Would our current code also work for determining cycles in a directed graph?





- Question: Would our current code also work for determining cycles in a directed graph?
- Answer: No, it would not work to determine cycles in a directed graph. For a path to be a cycle you must be able to get back to the first vertex. The fact that a vertex has already been discovered by a someone else, does not mean you can get from that vertex back to whoever discovered it. The edge might only be directed toward the vertex and not back.



Breadth-First Search: Manager Loop

BFSManager(G)

END FOR

```
//housekeeping
FOR each vertex u in V(G)
 state[u] = undiscovered
 parent[u] = nil
time = 0
FND FOR
//catch all vertices
FOR each vertex u in V(G)
  IF (state[u] != discovered) THEN
       BFS(G, u)
```

The manager loop for BFS ensures that if the graph is not connected, we still discover all vertices





Breadth-First Search Summary

- When to use BFS: BFS is a useful traversal method when you know the solution isn't far from the source vertex (e.g., friend suggestions on Facebook that are one-degree away from you).
- Where BFS Shines: Shortest Path (unweighted), Spanning Tree (undirected), Social-Networks.





That's All For Now...

- Coming to a Slideshow Near You Soon...
 - 1. BFS Tree of Discovery
 - 2. Analysis of the BFS algorithm
 - 3. Applications of BFS algorithm

That's All For Now



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