## **Assignment 3**

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<ul><li>1. Prove or disprove the following statement. If all capacities in a network are distinct, then there exists a unique flow function that gives the maximum flow.</li><li>Solution:</li></ul>
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2. An edge of a flow network is called <b>critical</b> if decreasing the capacity of this edge results in a decrease in the maximum flow value. Present an efficient algorithm that, given an s-t network G finds any critical edge in a network(assuming one exists).  Solution:
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3. Let $G=(V,E)$ be an undirected weighted graph with two distinguished vertices $s,t\in V$ . Give an efficient algorithm to find a minimum weight cut that separates $s$ from $t$ . Solution:
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<b>4.</b> You are given a matrix with fractional elements between 0 and 1. The sum of all numbers in each row and in each column is integer. Prove that we can always round each element to 0 or 1 so that the sum of each row and each column remains unchanged and design a polynomial time algorithm to find such a rounding result. <b>Solution:</b>
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5. Suppose that, in addition to edge capacities, a flow network has **vertex capacities**. That is each vertex has a limit on how much flow can pass though. Show how to transform a flow network G = (V, E) with vertex capacities into an equivalent flow network G' = (V', E') without vertex capacities, such that a maximum flow in G' has the same value as a maximum flow in G. How many

**Solution:** 

vertices and edges does G' have?

<b>6.</b> Consider a bipartite graph $G=(X\cup Y,E)$ with parts $X$ and $Y$ . Each part contains $2k$ vertices (i.e. $ X = Y =2k$ ). Suppose that $deg(u)\geq k$ for every $u\in X\cup Y$ . Prove that $G$ has a perfect matching. Solution:
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7. You are designing a experiment in which you want to measure certain properties $p_1, \ldots, p_n$ of a yeast culture. You have a set of tools $t_1, \ldots, t_m$ that can each measure a subset $S_i$ of the properties. For example, tool $t_i$ measures $S_i$ may equal $\{p_7, p_8\}$ . To be sure that your results are not due to noise or other artifact, you must measure every property at least $k$ times using $k$ different tools.
• Give a polynomial-time algorithm that decides whether the tools you have are sufficient to measure the desired properties the desired number of times.
• Suppose each tool $t_i$ comes from manufacturer $M_i$ and we have the additional constraint that the tools to test any property $p_i$ can't all come from the same manufacturer. Give a polynomial-time algorithm to solve this problem.
Solution:
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8. Consider a flow network $G=(V,E)$ with positive edge capacities $\{c(e)\}$ . Let $f:E\to\mathbb{R}_{\geq 0}$ be a maximum flow in $G$ , and $G_f$ be the residual graph. Denote by $S$ the set of nodes reachable from $S$ in $G_f$ and by $T$ the set of nodes from which $t$ is reachable in $G_f$ . Prove that $V=S\cup T$ if and only if $G$ has a <b>unique</b> $S$ - $t$ minimum cut. Solution:
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