# A GREEDY CONSTRUCTION IN SKOLEM'S 1920 PAPER

#### 1. Introduction

In the not so well known second section<sup>1</sup> of Skolem's 1920 paper [1] he gives an algorithm to determine the universally quantified first-order sentences that hold in Schröder's Gruppenkalkül. (The Gruppenkalkül was introduced in the first volume of Schröder's famous three-volume *Algebra der Logik*.) From the modern viewpoint the Gruppenkalkül is just a version of what is now called lattice theory. Skolem's version of the Gruppenkalkül is yet another version of modern lattice theory.

### 2. SKOLEM'S GRUPPENKALKÜL

Skolem axiomatized Schröder's Gruppenkalkül (lattice theory) in a relational language with two ternary relations and one binary relation, but without equality as follows:

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\begin{array}{lll} \mathrm{II.} & x \leqslant x \text{ for every } x \\ \mathrm{III.} & x \leqslant y \text{ and } y \leqslant z \text{ imply } x \leqslant z \\ \mathrm{III}_{\times}. & Mxyz \text{ implies } z \leqslant x \text{ and } z \leqslant y \\ \mathrm{IV}_{\times}. & Mxyz \text{ and } u \leqslant x \text{ and } u \leqslant y \text{ imply } u \leqslant z \\ \mathrm{V}_{\times}. & Mxyz \text{ and } x \leqslant x' \text{ and } x' \leqslant x \text{ and } y \leqslant y' \text{ and } \\ & y' \leqslant y \text{ and } z \leqslant z' \text{ and } z' \leqslant z \text{ imply } Mx'y'z' \\ \mathrm{VI}_{\times}. & \text{for all } x \text{ and } y \text{ there is a } z \text{ such that } Mxyz \\ \end{array} \qquad \begin{array}{ll} \mathrm{III}_{+}. & Jxyz \text{ implies } x \leqslant z \text{ and } y \leqslant z \\ \mathrm{IV}_{+}. & Jxyz \text{ and } x \leqslant u \text{ and } y \leqslant u \text{ imply } z \leqslant u \\ \mathrm{V}_{+}. & Jxyz \text{ and } x \leqslant x' \text{ and } x' \leqslant x \text{ and } y \leqslant y' \text{ and } \\ & \text{and } y' \leqslant y \text{ and } z \leqslant z' \text{ and } z' \leqslant z \text{ imply } Jx'y'z' \\ \end{array}
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**Definition 2.1.** A structure  $S = (S, M, J, \leq)$ , where M and J are ternary relations on S, and  $\leq$  is a binary relation on S, will be called a **Skolem structure**. A Skolem structure which satisfies Axioms I–VI will be called a **Skolem lattice**.

## 3. SKOLEM'S ALGORITHM

Skolem views axioms I–V as *production* rules that can only augment the three relations. Given a finite Skolem structure  $\mathbf{S} = (S, M, J, \leqslant)$  he first applies I–V repeatedly until the three relations on S no longer increase. (This a polynomial time procedure.)

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<sup>&</sup>lt;sup>1</sup>Lösung des Problems zu entscheiden, ob eine gegebene Aussage des Gruppenkalkuls beweisbar ist oder nicht.

<sup>&</sup>lt;sup>2</sup>Skolem used the notations (ab), abc and abc instead of  $a \le b$ , Mabc and Jabc used here. In Schröder's Algebra der Logik one finds the corresponding notations to be  $a \le b$ , ab = c and a + b = c.

Suppose at this point VI does not hold, say  $a,b \in S$  but there is no  $c \in S$  such that Mabc. He adds a new element, say  $\alpha$ , to S, adds  $Mab\alpha$ , adds  $\alpha \leqslant \alpha$ , and describes which  $a \leqslant \alpha$  and which  $\alpha \leqslant a$  to add, where  $a \in S$ . The bulk of his argument is to show that I–V still hold, and that  $\leqslant$  has not increased on the set S.

A similar 1-element extension is used if there are  $a, b \in S$  but there is no  $c \in S$  such that Jabc. By repeated application of these procedures a tower of 1-element extensions can be created whose union satisfies I–VI, with no further augmentation of  $\leq$  on S since the initial closure under I–V.

The augmentation of  $\leq$  on S by I–V was Skolem's algorithm to determine for  $a, b \in S$  if the information in S and the lattice axioms force  $a \leq b$ .

For an example (now using modern notation) he applies this to the two halves

$$(3.1) a \wedge (b \vee c) \leq (a \wedge b) \vee (a \wedge c)$$

$$(3.2) (a \wedge b) \vee (a \wedge c) \leqslant a \wedge (b \vee c)$$

of the distributive law. The associated Skolem structure S on the set  $S = \{a, b, c, d, e, f, g, h\}$  of eight elements is the following:

$$\begin{array}{ccc} M & J & \leqslant \\ abd & bcf \\ ace & deg \\ afh & & \end{array}$$

Closing S under the production rules I–V gives 21 pairs in  $\leq$  which include gh but not hg. Thus 3.2 holds in lattices, but not 3.1. Consequently the distributive law

$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

does not hold in lattices.<sup>3</sup>

## REFERENCE

(1) Skolem, Thoralf. Logisch-kombinatorische Untersuchungen über die Erfüllbarkeit oder Beweisbarkeit mathematischer Sätze nebst einem Theoreme über dichte Mengen. Skrifter utgitt av Det Norske Videnskaps-Akademi i Oslo, I. Matematisk-naturvidenskabelig Klasse, 1920, no. 4.

<sup>&</sup>lt;sup>3</sup>Originally proved by Schröder in Vol. I of *Algebra der Logik* in a rather tedious manner, using 990 quasi-group equations, to refute a claim of C.S. Peirce a decade earlier that the distributive law held in lattices. Soon after this Dedekind gave a simple proof in his paper on Dualgruppen, a very modern version of lattices, by exhibiting a 5-element lattice that did not satisfy the distributive law. It is not clear if Skolem was aware of Dedekind's proof from the 1890s.