1. Problem 1: Fibonacci Series

The function calls itself recursively. When the punction encounters O/1 value, it is relationed to the function that added it \sim so on.

```
Function that make the call value thats sectivened
 gib(3) n=3 \Rightarrow gib(2) + gib(1)
 gib(2) n=2 =) gib(1) + gib(0)
 glb(1) n=1 =) & (n=1) =) suction I
 Value 1 seaturned to fib(2)
 f(b(a))  n=a =) I + f(b(a))
 glb(o) n=0 => if (n==0) => section o
 (26) n=2 => (+0=I
 pib(3)  n=3 =) pib(a) \leftarrow pib(a) = 1 + pib(a)
 fib (i) n=1 => septerin 1
 Alb (3) n=3 => 1+1=2
 gzb(4) n=4 => 2 + gzbG)
 pib(a) p=a \Rightarrow pib(i) + pib(o) = 1
 fib(4) n=4 => 2+1=3
 qzb(s) n=5 => qzb(4) + qzb(3)
                   = 3 + fib(3)
 400C3) n=3 2
 f(b(s)) n=s \Rightarrow 3+2=5
```

2. Poroblem 2: Merge Asonays

- a) Code uploaded in Github
- b) Time complexity of the function Souting:

 Since there is I for loop that suns for the length of averay 'n'

 1 while loop that suns also suns 'n' times,

 The time complexity is ni. (1)

But the perogenam auto contains the part of building the average using user input.

If the number of array 'k' is longth of array 'n' is too large, the array building part has to be taxon into contain alreation.

.. Food lange K x n: ((k x n + n2)

- c) Ways to imposove my implementation:
- 1. I could have used 'menge look' to soort the unified array, it would have reduced the time complexity
- 2. If the input arrays are larger in Size, compate nating them would take longer. So instead of compating them, a recursive sorting algorithm for each array would be better.

3. Paroblem 3: Remove duplicates

- a) code uploaded in Github.
- b) Time Complexity: The function has 2 for loops

 1st for loop pick one comment & the 2nd for loop

 is used for comparing the 1st comment to all

 the other complexity is \oplus (n²)
- c) ways to imposore my implementation:

This would have produced a much simpler code.