

## Assignment 2

8.2

(a)

$$x_{n+1} = (ax_n) \bmod 2^4$$

$$\begin{aligned} \text{Maximum period} &= 2^4 - 2 \\ &= 4 \end{aligned}$$

$$\begin{aligned} &5x_n \bmod 2^4 \\ &5 \end{aligned}$$

(b)  $a = 5, 11$

(c) Seed must be odd:

8.4

$$x_{n+1} = (6x_n) \bmod 13$$

Sequence : 1, 6, 10, 8, 9, 2, 12, 7, 3, 5, 4, 11, 1, ...

$$x_{n+1} = (7x_n) \bmod 13$$

Seq : 1, 7, 10, 5, 9, 11, 12, 6, 3, 8, 4, 2, 1, ...

Because of the patterns seen in 2<sup>nd</sup> half of the latter sequence, most people would consider it to be less random than the first sequence.

8.6

Use a key of length 255 bytes.  
The first two bytes are zero,  
 $K[0] = K[1] = 0$

Thereafter we have

$$K[2] = 255, K[3] = 254, \dots, \\ K[255] = 2.$$

This key will leave  $S$  unchanged.

consider

8.7 (a) Let us ~~state~~  $i, j, S$  which  
requires

$$8 + 8 + (256 * 8) = 2064 \text{ bits.}$$

(b) The number of states is  
 $(256! \times 2562) \times 21700$

Therefore 1700 bits are required.

8.8 (a) By taking the first 80 bits of  $v \parallel c$   
we obtain the initialization vector  $v$ .  
Since  $v, c, k$  are known,  
the message can be recovered  
by computing  $RC4(v \parallel k) * c$ .

(b) If adversary observes the  $v_i = v_j$   
for distinct  $i, j$  then he/she

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knows that the same key stream was used to encrypt both  $m_i$  and  $m_j$ .

In this case, the messages  $m_i$  and  $m_j$  may be vulnerable to the type of cryptanalysis carried out in above part (a).

(c) Since the Key is fixed, the key stream varies with the choice of the 80 bit  $V$ , which is selected randomly. Thus after approximately messages are sent, we expect the same  $V$  and hence the same Key stream, to be used more than once.

(d) The key  $K$  should be changed sometime before 240 messages are sent.

### 8.5 Algorithm

- (1) Generate ' $n$ ' number of any random numbers —
- (2) Finding GCD for pairs in random numbers using Euclid's algorithm. (for ' $n$ ' numbers, we ~~have~~ <sup>will see</sup>  $n/2$  pairs)
- (3) Find probability of  $\text{GCD} = 1$  (wherever this occurs)  
→ calculated in (3)
- (4) Estimate  $\pi = \sqrt{6 / (\text{probability gcd as } 1)}$



## Program:

```
#include <random>
```

```
#include <iostream>
```

```
#include <chrono>
```

```
int main() number_generator()
```

```
{
```

```
    unsigned int num;
```

```
    unsigned int seed = chrono::system_clock::now().time_since_epoch().count();
```

```
    minstd_rand0 generator(seed);
```

```
    // standard linear congruential engine
```

```
    num = generator();
```

```
    return num;
```

```
}
```

```
int Euclid's_algo(int n, int m)
```

```
{
```

```
return (m != 0) ? gcd(m, n % m) : n;
```

```
    if (n == 0)
```

```
        return m;
```

```
    return euclid's_algo(n % m, n);
```

```
}
```

```
int square root ( int n)
```

```
{  
    int sq = sqrt(n)  
    return sq;  
}
```

```
int main()
```

```
{  
    float phi, int val;  
int number-generator();  
    int probability;  
    int n;  
    int random[100]; // array of random  
    int GCD[50];      numbers.  
    cin >> n; // number of nos.
```

```
    for (int i=0; i<n; i++)  
        random[i] = number-generator();
```

```
    for (int i=0; i<n; i+=2)  
    {  
        GCD[i] = euclids_algo(random[i],  
                                random[i+1]);
```

```
    }
```

```
    for (int i=0; i<n/2; i++)
```

```
    {  
        if (GCD[i] == 1)
```

```
            num_1++;  
    }
```

$$\text{probability} = \frac{\text{num}-1}{\binom{n}{2}};$$

$$\text{val} = \frac{6}{\text{probability}};$$

$$\text{phi} = \text{sqrt}(\text{val});$$

~~if (phi != 0)~~

cout << "PI value is" << phi

return 0;