CS 453/698: Software and Systems Security

Module: Bug Finding Tools and Practices

Lecture: Fuzz testing (a.k.a., fuzzing)

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Winter 2025

Outline

- Introduction
- Testing with concrete inputs
- 3 Fuzzing overview
- 4 Program state coverage: "natural selection" in the fuzzing world
- Conclusion

Program assurance

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Intro

- Existing practice: testing with manual effort
 - a.k.a., unit tests, E2E tests, quality assurance, etc.

Overview 00000

Program assurance

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 - i.e., automated, evolutionary, and random generation of test cases

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- In research pipeline: symbolic execution
 - i.e., automated, systematic, and deterministic exploration of search space

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 - i.e., automated, systematic, and deterministic exploration of search space
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 - i.e., automated, efficient, and practical exploration of search space

```
1 // implementation of `calc`
2 fn calc(
3     x: u64, y: u64, n: u64
4 ) -> (u64, u64, u64) {
5     /*
6     *
7     * code omitted for now
8     *
9     */
10 }
```

calc is the function we want to check for vulnerabilities.

let (a, b, i) = calc(x, y, n);

// use the results

//

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We know how to use calc (as shown in the main function).

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    x: u64, y: u64, n: u64
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     use the `calc` function
  pub fn main() {
    let (x, y, n) = /* input */;
    let (a, b, i) = calc(x, y, n);
    assert!(n-a-b+i != 42);
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More importantly, we know when calc is wrong (by the assertion the main function).

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We know how to use calc (as shown in the main function).

More importantly, we know when calc is wrong (by the assertion the main function).

Real-world software is of course more complicated than this simple example.

Running example: open-box analysis

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     x: u64, y: u64, n: u64
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     return (a, b, i);
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```

Q: Now we have the source code of the calc function, how do we know it is correct?

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   // use the `calc` function
  pub fn main() {
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Q: Now we have the source code of the calc function, how do we know it is correct?

A: Just try it with some concrete inputs and see if the results match with our expectations.

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\bullet x=0, y=1, n=2 \rightarrow a=2, b=2, i=3
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A: The *de facto* answer is: when achieved 100% code coverage.

CFG and code coverage

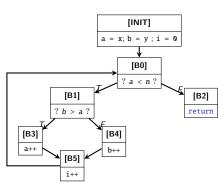
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Figure: the control-flow graph (CFG) of function calc(..)



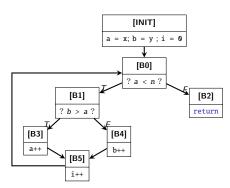
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Figure: the control-flow graph (CFG) of function calc(...)



100% code coverage usually means:

- all nodes in the CFG, or
- all edges in the CFG

100% coverage does not imply a worry-free program

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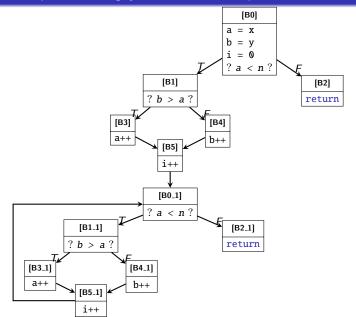
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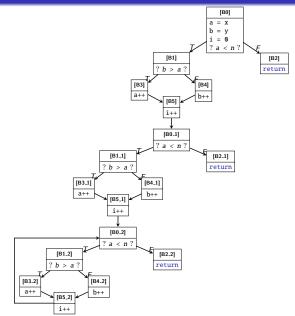
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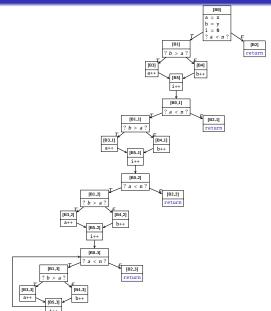
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In reality, it is infeasible to manually construct test cases that cover every new node / edge in the loop-unrolled CFG.

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A genetic programming solution: coverage-guided fuzzing.

- lacktriangle Randomly selects a seed s from a pool of test cases (seed pool)
 - e.g., seed test case $s = \{ x=0, y=1, n=2 \}$
- $oldsymbol{0}$ Mutate seed s to produce a new test case t
 - e.g., test case $t = \{ x=0, y=42, n=2 \}$
- ullet Execute t and collect coverage on the CFG during execution
 - if t yields new coverage, save t to the seed pool
 - otherwise, discard t
- Go back to step 1 completing one round of evolution

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History: why do we call it "fuzzing"?

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In 80's, someone remotely logged into a unix system over a dial-up network link during a storm.

The rain caused a lot of random noise on the dial-up link.

And these noise caused applications that were using data off the dial-up network line to crash.

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Gist of the story? — The rain tests the program way better than human beings.

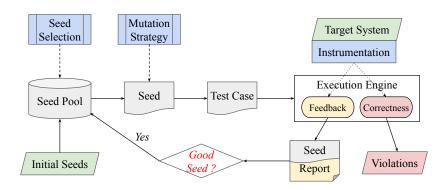
Evolution: from the rain-fuzzer to modern fuzzing

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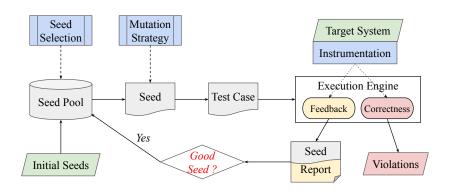
The key is **genetic algorithm**.

Training a program to play the snake game with genetic algorithm

Feedback-guided evolution process



Feedback-guided evolution process



Natural selection — survival of the fittest

Demo with AFL++

Acknowledgement: this demo is based on one of the examples used in the "Fuzzing with AFL" workshop by Michael Macnair.

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pub fn foo(a: num, b: num) {
       let c = if (a >= \emptyset) {
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       } else {
       };
7
       // irrelevant operations
9
10
       let d = if (b >= \emptyset) {
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       } else {
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       };
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17
       assert!(c != d);
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19 }
```

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⇒ if the fuzzer generates an input that expands the coverage, that input is a good seed.

Illustration of different coverage metrics

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       } else {
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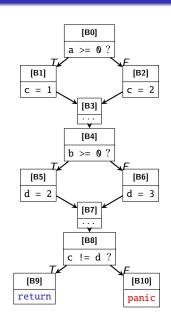


Illustration of different coverage metrics

- Cover every line?
 - Block coverage
- Cover every if-else branch?
 - Branch coverage
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 - Return coverage
- Cover every path?
 - Path coverage

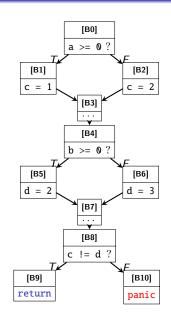
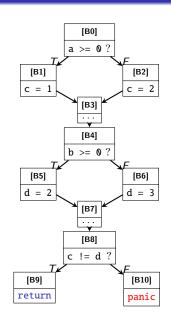


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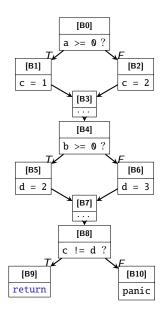
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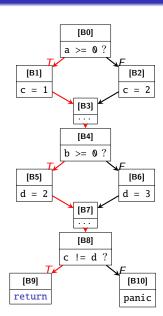
Path coverage: a theoretical optimum

Claim: A program is saturately tested if we obtain a set of inputs that covers every feasible path of the program CFG.

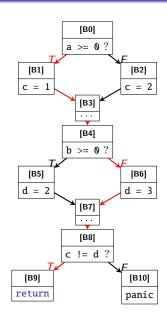
NOTE: feasible paths include paths that leads to explicit and implicit panics.



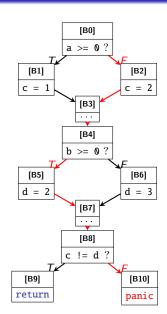
 \bullet a = 1, b = 1



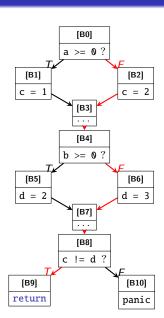
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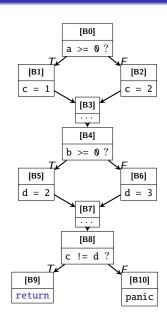


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No new program behaviors can be discovered \implies the program is saturately tested



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Short answer: I don't know... AFL (American Fuzzy Lop) didn't adopt path coverage, so everyone follows suite...

Long answer:

- tracking block / branch coverage is stateless while tracking path coverage requires stateful instrumentations.
- different parts of the execution are not necessarily related, i.e., a new path does not necessarily mean interesting findings.
- it is hard to quantitatively measure the completeness of path coverage (because of infeasible paths). But by default, all branches should be somewhat feasible.

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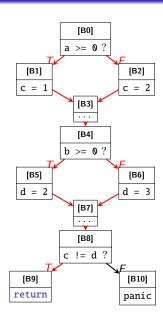
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In practice, branch coverage hits a nice balance between effectiveness and easiness of instrumentation.

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Two seeds already covered most of the branches.

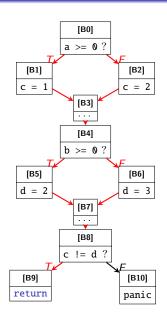


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A seed that yields new path but is considered as a bad seed as it yields no new branch coverage.



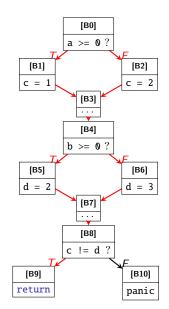
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Two seeds already covered most of the branches.

• a = 1, b = -1

A seed that yields new path but is considered as a bad seed as it yields no new branch coverage.

⇒ fuzzer is not rewarded by mutating a and b, hence, lowering their priorities and the panic case may never be found,



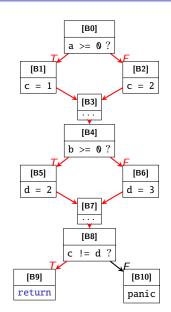
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A seed that yields new path but is considered as a bad seed as it yields no new branch coverage.

⇒ fuzzer is not rewarded by mutating a and b, hence, lowering their priorities and the panic case may never be found, especially when fuzzing complex CFGs



Outline

- Introduction
- 2 Testing with concrete inputs
- 3 Fuzzing overview
- 4 Program state coverage: "natural selection" in the fuzzing world
- Conclusion

The goal of fuzzing

Q: What is fuzzing doing essentially? Try to describe it in a way that is as abstract/general as possible.

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A: To drive the execution of a system into desired states.

Elaborating on the definition

- What is special about the target system?
 - Do we know the source code?
 - Do we know the input format?
 - What are the challenges when executing the "system"?
- What do we mean by a state?
 - How can we tell that one state is different from another?
- What do we mean by desired?
 - New/unseen behavior?
 - Closeness to targeted execution points?
- What do we mean by driving the execution?
 - What can possibly be one mutation?
 - How do you select the next mutation?

Conclusion ○○○●

 \langle End \rangle