

Lab 3 Report

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Changes

proc.h

- Line 52: Created uint pages

```
52     uint pages;                // CS153 Lab 3: Number of pages
```

memlayout.h

Defined top of stack address

- Line 18: Defined memory address for top of stack.

```
17     // Lab 3: Top of Stack Address
18     #define TOPSTACK (KERNBASE-4)    // The first address is one word
19                                     // under the kernel base
```

proc.c

Copied pages to child processes

- Line 204: Copied pages to child process as well.

```
203     // CS153 Lab 3: copy number of pages
204     np->pages = curproc->pages;
```

syscall.c

Implemented Lab 3 by replacing curproc->sz with TOPSTACK throughout the file.

- Line 24, 42, 45, 74: Use TOPSTACK instead of curproc->sz.

```

17  int
18  fetchint(uint addr, int *ip)
19  {
20      // CS153 Lab 3: Change checks to top of the stack rather than sz
21      //      Removes need to myproc as we use the top of stack
22      // struct proc *curproc = myproc();
23
24      if(addr >= TOPSTACK || addr+4 > TOPSTACK)
25          return -1;
26      *ip = *(int*)(addr);
27      return 0;
28  }

33  int
34  fetchstr(uint addr, char **pp)
35  {
36      char *s, *ep;
37
38      // CS153 Lab 3: Change checks to top of the stack rather than sz
39      //      Removes need to myproc as we use the top of stack
40      // struct proc *curproc = myproc();
41
42      if(addr >= TOPSTACK)
43          return -1;
44      *pp = (char*)addr;
45      ep = (char*)TOPSTACK;
46      for(s = *pp; s < ep; s++){
47          if(*s == 0)
48              return s - *pp;
49      }
50      return -1;
51  }
--

```

```

68 // CS153 Lab 3: Change checks to top of the stack rather than sz
69 // Removes need to myproc as we use the top of stack
70 // struct proc *curproc = myproc();
71
72 if(argint(n, &i) < 0)
73     return -1;
74 if(size < 0 || (uint)i >= TOPSTACK || (uint)i+size > TOPSTACK)
75     return -1;
76 *pp = (char*)i;
77 return 0;
78 }

```

vm.c

Copies over information into virtual memory.

- Lines 340 - 352: Loop in order to copy information onto the new stack.

```

339 // CS153 Lab 3: Loop to copy the new stack
340 for(i = PGROUNDUP(TOPSTACK - (PGSIZE * myproc()->pages)); i < TOPSTACK; i += PGSIZE) {
341     if((pte = walkpgdir(pgdir, (void *) i, 0)) == 0)
342         panic("copyuvm: pte should exist");
343     if(!(*pte & PTE_P))
344         panic("copyuvm: page not present");
345     pa = PTE_ADDR(*pte);
346     flags = PTE_FLAGS(*pte);
347     if((mem = kalloc()) == 0)
348         goto bad;
349     memmove(mem, (char*)P2V(pa), PGSIZE);
350     if(mappages(d, (void*)i, PGSIZE, V2P(mem), flags) < 0)
351         goto bad;
352 }

```

exec.c

Allocates all pages starting from the top of the stack. Replaces all sz with TOPSTACK.

- Line 65, 67: Replace sz with TOPSTACK.
- Line 102: Set number of pages to 1.

```

63 // CS153 Lab 3: Allocate page at the top of the stack
64 sz = PGROUNDUP(sz);
65 if((sp = allocvm(pgdir, TOPSTACK - PGSIZE, TOPSTACK)) == 0)
66     goto bad;
67 sp = TOPSTACK;
101 // CS153 Lab 3: set number of pages to 1
102 curproc->pages = 1;
103 cprintf("Initial number of pages by the process: %d\n", curproc->pages); //cs153 - Lab3

```

trap.c

Implemented page faults

- Lines 81 - 93: Implemented page faults.

```

81 // CS153 Lab 3: Page Fault
82 case T_PGFLT:
83     // Check if the address in rcr2 is in the next page
84     if ((rcr2() < PGROUNDUP(TOPSTACK - myproc()->pages * PGSIZE)) && (rcr2() > PGROUNDUP(TOPSTACK - (myproc()->pages+1)*PGSIZE))) {
85         myproc()->pages++;
86         if (allocvm(myproc()->pgdir, (TOPSTACK - myproc()->pages * PGSIZE), (TOPSTACK - (myproc()->pages - 1) * PGSIZE)) == 0) {
87             cprintf("case T_PGFLT from trap.c: allocvm failed. Number of current allocated pages: %d\n", myproc()->pages);
88             exit();
89         }
90         cprintf("case T_PGFLT from trap.c: allocvm succeeded. Number of pages allocated: %d\n", myproc()->pages);
91     }
92     break;

```

Extra Credit

We were unable to implement a stack that grows into the heap because this would cause the stack to overlap the heap. As a result, the stack would overwrite information contained within the heap. The heap is allocated within vm.c and shown below. Within the I/O space is the heap. By growing the stack towards the heap, we effectively reach outside of the size limitations of the stack and cause an overflow, which would break the kernel and trigger an overflow error.

```

103 // This table defines the kernel's mappings, which are present in
104 // every process's page table.
105 static struct kmap {
106     void *virt;
107     uint phys_start;
108     uint phys_end;
109     int perm;
110 } kmap[] = {
111     { (void*)KERNBASE, 0,          EXTMEM,    PTE_W}, // I/O space
112     { (void*)KERNLINK, V2P(KERNLINK), V2P(data), 0},    // kern text+rodata
113     { (void*)data,     V2P(data),    PHYSTOP,   PTE_W}, // kern data+memory
114     { (void*)DEVSPACE, DEVSPACE,     0,        PTE_W}, // more devices
115 };

```