

Sorting

① Bubble Sort:

Question:

12	23	35	15	11
----	----	----	----	----

1st Iteration:

12	23	35	35	11
----	----	----	----	----

12	23	15	11	35
----	----	----	----	----

2nd Iteration:

12	23	15	11	35
----	----	----	----	----

12	25	23	11	35
----	----	----	----	----

12	15	11	23	35
----	----	----	----	----

3rd Iteration:

12	15	11	23	35
----	----	----	----	----

12	11	15	23	35
----	----	----	----	----

4th Iteration:

12	11	15	23	35
----	----	----	----	----

11	12	15	23	35
----	----	----	----	----

for (a=0; a < n-1; a++)

for (i=0; i < n-1; i++)

{

if (arr[i] > arr[i+1])

{

temp = arr[i]

arr[i] = arr[i+1]

arr[i+1] = temp

}

Algorithm

Step 01: Start the program

Step 02: Initialize an unsorted list.

Step 03: Using for loop from 0 to no. of element-1. (Outer loop)

Step 04: Apply inner for loop from 0 till n-1.

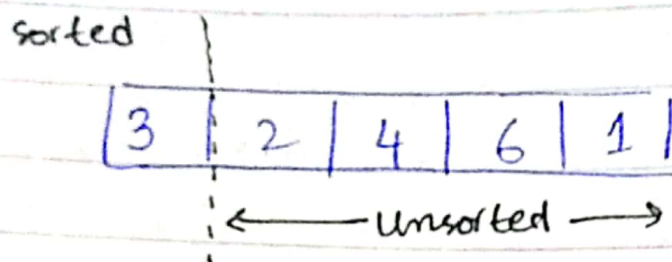
Step 05: Using if to check if the first number is greater than the second or not.

Step 06: If condition gets false swap these numbers.

Step 07: End the program.

② Insertion Sort:

Jenny's Lecture



Code:

```
for (int i = 1; i <= 4; i++)
```

```
{
```

```
    temp = arr[i];
```

```
    j = i - 1;
```

```
    while (j >= 0 && arr[j] > temp)
```

```
    {
```

```
        arr[j + 1] = arr[j];
```

```
        j--;
```

```
    }
```

```
    arr[j + 1] = temp;
```

```
}
```


Algorithm:

Step 01: Start the program.

Step 02: Initialize an array.

Step 03: Start a for loop from 1 to length of array - 1.

Step 04: Assign ~~the~~ value of i in temporary variable.

Step 05: Initialize j as one less i .

Step 06: Apply a while loop with two conditions:

1. $j > 0$

and

$arr[j] > temp$.

Step 07: assign value of $arr[j]$ into $arr[j+1]$.

Step 08: Decrement j by 1.

Step 09: assign value of $temp$ into $arr[j+1]$.

Sorted

3 | 2 | 4 | 6 | 1 | 1
Unsorted

2 | 3 | 4 | 6 | 1 | 1
S U

2 | 3 | 4 | 6 | 1 | 1
S U

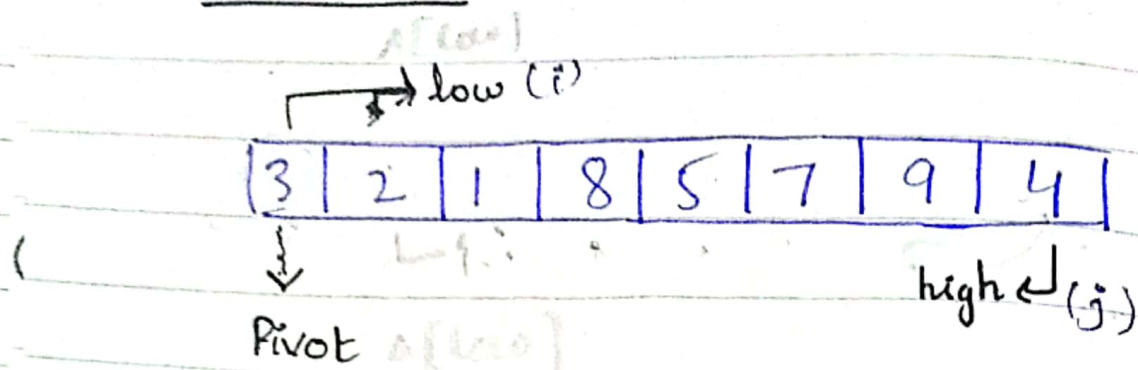
6 | 2 | 3 | 4 | 6 | 1
S U

2 | 1 | 2 | 3 | 4 | 6

Sorted

Contk with Harry.

③ Quick Sort =



- ⇒ Take first num as Pivot.
- ⇒ i will be incremented and j will be decremented.
- ⇒ i will ~~find~~^{stop} at pivot se bada num
- j will stop at pivot se chota num.
- ⇒ agr i or j ne ek dosre ko cross hui kro
hoga to dono ki values swap else
Code. i and pivot swap.

```
void main ()  
{
```

```
    int A[] = {9, 4, 4, 8, 7, 5, 6};  
    int n = 9;
```

```
    printArray(A, n);
```

```
    * quickSort(A, 0, n-1);
```

```
    printArray(A, n);
```

```
    return 0;
```

```
}
```

```
void quickSort (int A[], int low, int high)
{
```

```
    int partitionIndex;
```

```
    if (low < high)
    {
```

```
        partitionIndex = partition(A, low, high);
```

```
        quickSort(A, low, partitionIndex - 1);
```

```
        quickSort(A, partitionIndex + 1, high);
```

```
    }
```

```
}
```

```
int partition (int A[], int low, int high)
```

```
{
```

```
    int pivot = A[low];
```

```
    int i = low + 1;
```

```
    int j = high;
```

```
    int temp;
```

```
    do
```

```
    {
```

```
        while (A[i] <= pivot)
```

```
        {
```

```
            i++;
```

```
        }
```

```
        while (A[j] > pivot)
```

```
        {
```

```
            j--;
```

```
        }
```


if ($i < j$)

{

temp = A[i];

A[i] = A[j];

A[j] = temp;

}

}

while ($i < j$);

temp = A[low];

A[low] = A[j];

A[j] = temp;

return j;

}

Algorithm:

Main Function

Step 01: Start the program.

Step 02: Initialize an array as well as its length (n).

Step 03: Call the function quickSort(A, 0, n-1).

Step 04: Print that array.

Step 05: End the program.

quicksort (arr, low, high) function.

- Step 01: Initialize low = first ^{index} element (0th index) of array, it is also called Pivot.
- Step 02: Initialize high as last index of arr.
- Step 03: If $low < high$, call the partition(arr, low, high) function.
- Step 04: Using recursion properly, again call the quicksort(arr, low, p.index-1) but with diff. arguments.
- Step 05: Partition index is basically the variable which consist the value of Pivot's position.
- Step 06: Again call quicksort(arr, p.index+1, high).
- Step 07: End this function.

partition(arr, low, high) function.

- Step 01: Initialize pivot = arr[low].
 $i = low + 1$.
 $j = high$.
- Step 02: Applying do while loop with the condition ($i < j$).
- Step 03: Inside loop apply while (arr[i] <= pivot)
- ```
{
 i++;
}
```



```

while (arr[j] > pivot)
{
 j--;
}

```

Applying if condition:

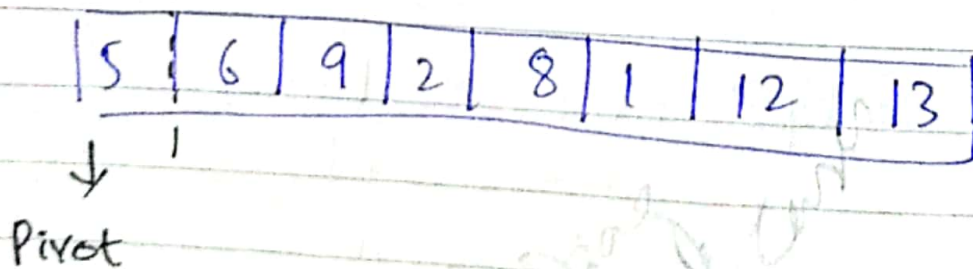
```

if (i < j)
{
 swap arr[i] with arr[j]
}

```

Step 04: If loop condition gets set to false, swap  $arr[j]$  by pivot ( $arr[low]$ ).

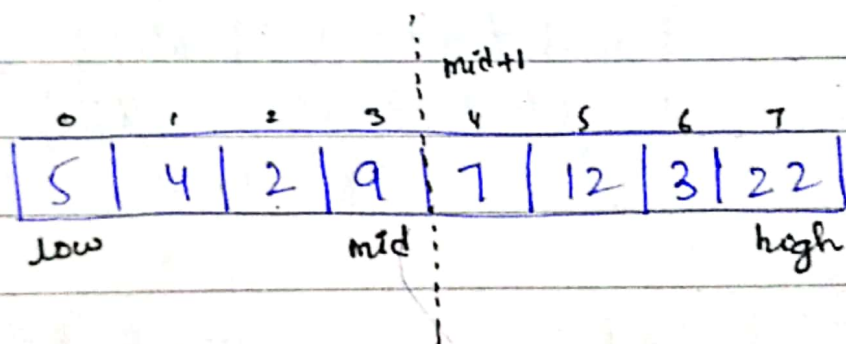
Step 05: Return  $j$ .



## Merge Sort:

→ First we will break an array into two equal parts. We consider first element as low and last as high. The breakpoint will be mid. We do this again and again by recursiveness and when only one element is left we will apply merge() function.

In merge(), we have two parts low  $\rightarrow$  mid and mid + 1  $\rightarrow$  high, we will apply ~~three~~ <sup>two</sup> diff. loops on each and start comparison and store the values in new array B[]



### Algorithm:

Recursive Merge Sort Function.

~~void ms(int A[]).~~

Step 01: Start the program

Step 02: Apply if to check low < high.

Sub-step (2a): Find mid point by using formula  $mid = \frac{low + high}{2}$ .

Sub-step (2b): Again call this function from low to mid.

Sub-step (2c): Call this function again from mid+1 to high.

Sub-step (2d): Call merge() function.

Step 03: End the function.

PseudoCode:

```
void MS (int A[], int low, int high)
{
 if (low < high)
 {
 mid = (low+high)/2;
 MS (A, low, mid);
 MS (A, mid+1, high);
 Merge (A, low, mid, high);
 }
}
```

Algorithm for merge() function.

Step 01: Start the function.

Step 02: Initialize  $i = \text{low}$ ,  $j = \text{mid} + 1$  and  $k = \text{low}$ .

Step 03: Apply while loop with the condition  $(i < \text{mid} \ \&\& \ j < \text{high})$

Sub-step (3a): Check if  $(A[i] < A[j])$

Sub-step-3a(i): Copy  $A[i]$  in new array  $B[k]$  and increment  $i$  and  $k$ .



Sub-step(3b): Else copy  $A[j]$  to  $B[k]$  and increment  $j$  and  $k$  by one.

Step 04: Now for copying remaining elements, ~~from  $A$  to  $B$~~ , <sup>low to mid</sup> apply while loop with condition  $(i \leq \text{mid})$

Sub-step(4a):  $B[k] = A[i]$  and  $i++$ ,  $k++$

Step 05: Copying from  ~~$A$~~  <sup>mid+1</sup> to  ~~$B$~~  <sup>high</sup>, while  $(j < \text{high})$

Sub-step(5a):  $B[k] = A[j]$  and  $k++$  &  $j++$

Step 06: End the function.

### Pseudocode:

```
void merge (A[], mid, low, high) {
```

```
 int i = low, j = mid + 1, k = low;
```

```
 while (i < mid && j < high)
```

```
 {
```

```
 if ($A[i] < A[j]$) {
```

```
 $B[k] = A[i]$
```

```
 $i++$; $k++$;
```

```
 }
```

```
 else
```

```
 { A $B[k] = A[j]$;
```

```
 $j++$; $k++$;
```

```
 }
```

```
 }
```

```
 while (i < mid)
```

```
 {
```

```
 $B[k] = A[i]$;
```

```
 $k++$; $i++$;
```

```
 }
```

```
 while (j < high) {
```

```
 $B[k] = A[j]$;
```

```
 $k++$; $j++$;
```

# Diagram of Merge Sort

