

- Original Heuristic:

Instead of simply calculating *Manhattan distance (MD)*, which is the sum of the absolute differences between the two vectors, we can find a closer cost added to MD.

$$H'(n) = \begin{cases} MD & , \text{if } MD \leq 1 \\ MD + 3 & , \text{o. w.} \end{cases}$$

Relax the problem by simplifying the 1*2 pieces to two 1*1 pieces, there are only one 2*2 pieces, and all rest pieces are 1*1 (labeled 'x').

- Say MD=0, we are at goal state, actual cost is 0.
- Say MD =1, the best case, we can move to goal state at cost of 1.

○ Ex.

xxxx		xxxx
xxxx		xxxx
x11x	----down-->	x00x
x11x	cost: 1	x11x
x00x		x11x

- Say MD=2, the best case, we can move to goal state in 5 moves

○ Ex1.

xxxx		xxxx		xxxx		xxxx
xxxx		xxxx		xxxx		xxxx
110x.	---right-->	011x	-move empty-->	x11x	--down-->	x00x
110x.	cost: 1	011x	cost: 3	x11x	cost: 1	x11x
xxxx		xxxx		x00x		x11x

○ Ex2.

xxxx		xxxx		xxxx		xxxx
x11x		x00x		xxxx		xxxx
x11x.	---right-->	x11x	-move empty-->	x11x	--down-->	x00x
x00x.	cost: 1	x11x	cost: 6	x11x	cost: 1	x11x
xxxx		xxxx		x00x		x11x

- In best case for MD>= 2, 2*2 piece need **at least cost of 3** to swap the empty pieces with 1*1 pieces and prepare the empty pieces for the next move.

- Why admissible?

$H'(n)$ is admissible as it never overestimates the cheapest cost from state n to a goal state.

In the relaxed rules, for MD>=2, the 2*2 piece needs to move at least **MD** moves, and between (**MD-1**) pairs of consecutive moves we need to move the empty pieces, with a **min cost of 3** each.

- Why does it dominate *Manhattan distance heuristic*?

- For all possible number of MD numbers, $H'(n) \geq H(n)$; and for MD>=2, $H'(n) > H(n)$.

- Implementation:

- get_original_heuristic(board)
- A* with original heuristic for the classic HRD config still runs at costs of 116.