

A8.2)

generator:  $X_{n+1} = (aX_n) \bmod 2^4$

$$m = 2^4 = 16$$

$$c = 0$$

$$0 < a < 16$$

$$\text{taking } X_0 = 1$$

a) Maximum period =  $2^{4-2} = 4$

b) for  $a = 5$

$\{ 5, 9, 13, 1, 5, 9, 13, 1, \dots \}$   
 $\xleftarrow{\text{period} = 4}$

for  $a = 11$

$\{ 11, 9, 3, 1, 11, 9, 3, 1, \dots \}$   
 $\xleftarrow{\text{period} = 4}$

for maximum value of  $a$ ,  $a = 5$  or  $a = 11$

c) The seed must be odd.

A8.4) let initial seed,  $X_0 = 1$ ,

for  $X_{n+1} = (6X_n) \bmod 13$

sequence:  $\{ 6, 10, 8, 9, 2, 12, 7, 3, 5, 4, 11, 1, \dots \}$   
 $\xleftarrow{\text{period} = 12}$

for  $X_{n+1} = (7X_n) \bmod 13$

sequence:  $\{ 7, 10, 5, 9, 11, 12, 6, 3, 8, 4, 2, 1, \dots \}$   
 $\xleftarrow{\text{period} = 12}$

The first sequence appears to be more random. ~~Also, it is more~~  
 Because of the pattern seen in the second half of the sequence  
 generated by the second generator it appears to be less random, although  
 period of both sequences generated is of size 12.



A.8.6) If we use a key of length 255 Bytes.  
First two bytes are zero.

$$\text{ie. } k[0] = 0 \\ k[1] = 0$$

∴ we have

$$k[2] = 255 \\ k[3] = 254 \\ k[4] = 253 \\ k[5] = 252$$

$$\vdots \\ k[255] = 2.$$

A.8.7) a) To store  $i, j$  &  $S$  it requires

$$[8 + 8 + (256 \times 8)] \text{ bits} \\ = [16 + 2048] \text{ bits} \\ = 2064 \text{ bits.}$$

b) The number of states are -

$$[256! \times 256^2] \approx 2^{1700}$$

This is equal to approx  $2^{1700}$

∴ 1700 bits required to represent the state.

A.8.8)

a) If we take the first 80 bits of  $v||c$  we will get  $v$   
↑  
initialization vector.

Now, since  $v, c$  &  $k$  are known, the message can easily be  
decrypted by using the -

$$RC4(v||k) \oplus c$$

b) for distinct  $i \neq j$ ,  
 If adversary notices that  $v_j = v_i$ ,  
 They will know that same key stream used to encrypt  
 ~~$m_i$  &  $m_j$~~   $m_i$  &  $m_j$   
 So, the messages  $(m_i, m_j)$  may be vulnerable and  
 can be decrypted.

c) As the key is fixed, the key stream will vary with the  $g_{06}$   
 $v$  which will be selected at random.  
 $\therefore 2^{40}$  messages. (By Birthday Paradox)

d)  $2^{40}$  messages can be encrypted using the key  $k$ , after which,  
 key should be changed,



PDF Created Using



## Camera Scanner

Easily Scan documents & Generate PDF



<https://play.google.com/store/apps/details?id=photo.pdf.maker>