

<u>S.No.</u>			
	1 Algorithm	Program	
	Design time	Implementation time	
	Domain knowledge	Programmer	
	Any language even English and Maths	Programming language	
	Hardware and software independent	Hardware and operating system dependent	
	Analyze an algorithm	Testing of programs	
	2 Priori Analysis	Posterior Testing	
	Algorithm	Program	
	Independent of language	Language dependent	
	Hardware independent	Hardware dependent	
	Time and space function	watch time and bytes	
	3 Characteristics of algorithm		
	Zero or more inputs		
	Must generate atleast one output		
	Definiteness		
	Finiteness		
	Effectiveness		
	4 How to analyze an algorithm		
	Time		
	Space		
	Network consumption : Data transfer amount		
	Power consumption		
	CPU registers		
	5 Frequency Count Method	Used for time snalysis of an algorithm	
	Assign 1 unit of time for each statement		
	For any repitition, calculate the frequency of repetition		

	for(i = 0; i < n; i++) --> condition is checked for n+1 times	2n + 2 units of time ~ n+1 as we see condition i < n only for now	
	any statement within the loop will execute for n times		
	Space complexity depends upon number and kind of variables used		
	6 Algorithm : sum(A, n)		
	Single for loop -		
	Time complexity: O(N)		
	Space complexity: O(N)		
	7 Algorithm : Add(A, B, n)	Sum of two square matrices of dimensions nXn	
	Two nested for loops -		
	Time complexity: O(N^2)		
	Outer for loop executes for N+1 times		
	Inner for loop executes for N * (N+1) times		
	Any statement within inner for loop executes for (N + 1) * (N + 1) times		
	Space complexity: O(N^2)		
	8 Algorithm : Multiply(A, B, n)		
	Three nested for loops -		
	Time complexity: O(N^3)		
	Space complexity: O(N^2)		
	9 Different algorithm conditions		
	For loops		
	for(i = n; i > 0; i--)	n+1 times	
	for(i = 0; i < n; i = i + 2)	n/2 times	

	2 nested for loops where both i and j range from 0 to n	n^2 times	
	2 nested for loops where j ranges from 0 to i	when $i = 0$; j loop repeats 0 times; when $i = 1$; j loop repeats 1 times; and so on...total number of repetitions: $0 + 1 + 2 + 3 + 4 + \dots + n = O(n^2)$	
	$p = 0$; for($i = 1$; $p \leq n$; $i++$){ $p = p + i$; }	$p = k(k+1)/2 \rightarrow$ assuming that the loop exits when p is greater than n $\rightarrow k(k+1) / 2 > n$	$\sim k^2 > n \rightarrow O(\sqrt{n})$
	for($i = 1$; $i < n$; $i = i * 2$)	will execute for 2^k times	$O(\log n)$
		Assume $i \geq n$; $i = 2^k \geq n$	
		$k = \log n$ with base 2	
	for($i = n$; $i \geq 1$; $i = i/2$)	i	
		n	
		$n/2$	
		$n/2^2$	
		$n/2^3$	
		
		$n/2^k$	
		Assume $i < 1 \Rightarrow n / 2^k < 1$	$\sim O(\log n)$ with base 2
	for($i = 0$; $i * i < n$; $i++$)	$i^2 < n$	
		$i^2 > -n$	
		$i^2 = n \rightarrow i = \sqrt{n}$	$\sim O(\sqrt{n})$
	for($i = 0$; $i < n$; $i++$) {.....}for($j = 0$; $j < n$; $j++$){.....}	$O(n)$	
	$p = 0$; for($i = 1$; $i < n$; $i * 2$){.....} for($j = 1$; $j < p$; $j * 2$){.....}	log n times for upper loop; log p times for lower loop	$\sim O(\log(\log n))$
	for($i = 0$; $i < n$; $i++$) {.....for($j = 0$; $j < n$; $j * 2$){.....}}	Outer loop repeats n times; inner loop repeats logn times	$\sim O(n \log n)$
	for($i = 1$; $i < n$; $i = i * 3$)		$\sim O(\log n)$ with base 3
	While loops		
	while vs. do while	do while will execute for minimum one time	
	for and while are almost similar	do while will execute as long as the condition is true; for loop will execute until the condition is false	

	a = 1;		
	while(a < b){ a = a *2;}	1, 2, 2^2, 2^3.....2^k repetitions	~O(logb) with base 2
		assume a > b; 2^k > b ==> k = logb with base 2	
	i = n; while(i > 1) {...i = i/2;}		~O(logn) with base 2
	i = 1; k = 1; while(k < n){....k = k + i; i++;}		
		i k	
		1	1
		2 1 + 1	
		3 2 + 2	
		4 2 + 2 + 3	
		5 2 + 2 + 3 + 4	
		
	m	m(m + 1) /2	
	Assume, k >= n	m(m + 1)/ 2 >= n	~O(root(n))
	while(m != n) { if(m > n) m = m - n; else n = n - m;}		~O(n)
	10 Types of time functions		
	O(1) --- constant		
	O(logn) --- logarithmic		
	O(n) --- linear		
	O(n^2) --- quadratic		
	O(n^3) --- cubic		
	O(2^n) --- exponential		
	11 Order of complexity		
	1 < logn < root(n) < n < nlogn < n^2 < n^3 <.....< 2^n < 3^n< n^n		

12	Asymptotic Notations		
	Representation of time complexity in simple form which is understandable		
	Big O Notation - works as an upper bound	The function $f(n) = O(g(n))$ iff for all positive constants c and n_0 , such that $f(n) \leq c * g(n)$ for all $n \geq n_0$; here, $f(n) = O(n)$	e.g. $2n + 3 \leq 10n$; All those functions in time order complexity above n become upper bound; below n become lower bound and n is the average bound
	Big Omega Notation - works as a lower bound	The function $f(n) = \Omega(g(n))$ iff for all positive constants c and n_0 , such that $f(n) \geq c * g(n)$ for all $n \geq n_0$; here, $f(n) = \Omega(n)$	e.g. $2n + 3 \geq 1n$
	Theta Notation - works as an average bound	The function $f(n) = \theta(g(n))$ iff for all positive constants c_1, c_2 and n_0 such that $c_1 * g(n) \leq f(n) \leq c_2 * g(n)$	e.g. $f(n) 2n + 3; 1n \leq 2n + 3 \leq 5n$
	Most useful is theta notation, then why do we need the other two?	In case we are not able to get the average bound, then we point to its upper or lower bound	
13	Examples for asymptotic notations		
a	$f(n) = 2n^2 + 3n + 4$		
	$2n^2 + 3n + 4 \leq 2n^2 + 3n^2 + 4n^2$ i.e. $9n^2$	$O(n^2)$	
	$2n^2 + 3n + 4 \geq 1n^2$	$\Omega(n^2)$	
	$1n^2 \leq 2n^2 + 3n + 4 \leq 9n^2$	$\Theta(n^2)$	
b	$f(n) = n^2 \log n + n$		
	$n^2 \log n \leq n^2 \log n + n \leq 10n^2 \log n$	$O(n^2 \log n)$	
		$\Omega(n^2 \log n)$	
		$\Theta(n^2 \log n)$	
c	$f(n) = n!$		
	$1 \leq 1 * 2 * 3 * 4 * \dots * n - 1 * n \leq n * n * n * \dots * n$	$O(n^n)$	
		$\Omega(1)$	
		Cannot find theta for $n!$	

d	$f(n) = \log n!$		
	$1 \leq \log(1 \cdot 2 \cdot 3 \dots n) \leq \log(n \cdot n \cdot n \dots n)$	$O(\log^n n)$	
		$\Omega(1)$	
		Cannot find theta for $\log n!$	
	14 Properties of Asymptotic notations		
	General properties -		
	if $f(n)$ is $O(g(n))$ then $a \cdot f(n)$ is $O(g(n))$		
	e.g. $f(n) = 2n^2 + 5$ is $O(n^2)$, then $7f(n)$ i.e. $14n^2 + 35$ is also $O(n^2)$	This would be true for both Ω and θ as well	
	Reflexive property -		
	If $f(n)$ is given then $f(n)$ is $O(f(n))$		
	e.g. $f(n) = n^2$ then $O(n^2)$	A function is an upper bound of itself	
		Similarly, a function is a lower bound of itself	
	Transitive property -		
	If $f(n)$ is $O(g(n))$ and $g(n)$ is $O(h(n))$ then $f(n)$ is $O(h(n))$		
	e.g. $f(n) = n$; $g(n) = n^2$ and $h(n) = n^3$	True for all notations	
	n is $O(n^2)$ and n^2 is $O(n^3)$ then n is $O(n^3)$		
	Symmetric property -		
	If $f(n)$ is $\theta(g(n))$ then $g(n)$ is $\theta(f(n))$	True for only $\theta(n)$	
	e.g. $f(n) = n^2$ $g(n) = n^2$; $f(n) = \theta(n^2)$ and $g(n) = \theta(n^2)$		
	Transpose symmetric -	True for BigO and Ω notations	
	if $f(n) = O(g(n))$ then $g(n)$ is $\Omega(f(n))$		
	e.g. $f(n) = n$ and $g(n)$ is n^2 then n is $O(n^2)$ and n^2 is $\Omega(n)$		

	If $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$ then $g(n) \leq f(n) \leq g(n)$ therefore $f(n) = \theta(g(n))$		
	If $f(n) = O(g(n))$ and $d(n) = O(e(n))$ then $f(n) + d(n) = O(\max(g(n), e(n)))$		
	e.g. $f(n) = n = O(n)$, $d(n) = n^2 = O(n^2)$ then $f(n) + d(n) = n + n^2 = O(n^2)$		
	If $f(n) = O(g(n))$ and $d(n) = O(e(n))$ then $f(n) * d(n) = O(g(n) * e(n))$		
	15 Comparison of functions		
	First method is substituting values for n and comparing		
	Second method is applying log on both sides		
		Properties of log -	
	Example -	$\log ab = \log a + \log b$	
	$f(n) = n^2 \log n$; $g(n) = n(\log n)^{10}$	$\log a/b = \log a - \log b$	
	Apply log	$\log a^b = b \log a$	
	$\log(n^2 \log n)$; $\log(n(\log n)^{10})$	$a^{(\log cb)} = b^{(\log ca)}$	
	$\log(n^2) + \log \log n$; $\log n + \log \log^{10} n$	$a^b = n$ then $b = \log_a n$	
	$2 \log n + \log \log n$; $\log n + 10 \log \log n$		
	here; $2 \log n$ is greater than $\log n$ and $\log n$ is a bigger term than $\log \log n$		
	so, first term is greater than the second one		
	$f(n) = 3n^{(\sqrt{n})}$; $g(n) = 2^{(\sqrt{n} \log_2(n))}$		
	Applying log		
	$3n^{(\sqrt{n})}$; $(n^{\sqrt{n}}) \log_2(2)$		
	$3n^{(\sqrt{n})}$; $n^{\sqrt{n}}$		
	first term is greater than the second one value wise but asymptotically they are equal		

	$f(n) = n^{(\log n)}$; $g(n) = 2^{(\sqrt{n})}$		
	apply log,		
	$\log(n^{\log n})$; $\log(2^{\sqrt{n}})$		
	$\log n \cdot \log n$; $\sqrt{n} (\log_2 2)$		
	$\log^2 n$; \sqrt{n}		
	cannot judge, so apply log again		
	$2 \log \log n$; $1/2 \log n$		
	$\log \log n$ is smaller than $\log n$		
	thus, second term is greater		
	$f(n) = 2^{(\log n)}$; $g(n) = n^{(\sqrt{n})}$		
	$\log n \cdot \log_2 2$; $\sqrt{n} \cdot \log n$		
	$\log n$; $\sqrt{n} \cdot \log n$		
	second term is greater		
	$f(n) = 2n$; $g(n)$ is $3n$		
	both are equal asymptotically		
	$f(n) = 2^n$; $g(n) = 2^{(2n)}$		
	applying log		
	$\log(2^n)$; $\log(2^{2n})$		
	n ; $2n$	after applying log, do not cut coefficients	
	second function is greater		
	16 Best, worst and average case analysis		
	Example -		
a	Linear search		
	$A = \{8, 6, 12, 5, 9, 7, 4, 3, 16, 18\}$ key = 7		
	In linear search, it will start checking for the given key from left hand side		

	total in 6 comparisons, we would get our key		
	Best case - key element is present at first index		
	Best case time - 1 i.e. $B(n) = O(1)$; $\Omega(1)$; $\Theta(1)$		
	Worst case - key element is present at the last index		
	Worst case time - n i.e. $W(n) = O(n)$; $\Omega(n)$; $\Theta(n)$		
	Average case = all possible case time / no. of cases		
	average case analysis is very difficult for most of the cases		
	Here, average case time = $1 + 2 + 3 + \dots + n/2 = n(n+1)/2n = n+1/2$		
	$A(n) = n+1/2$		
b	Binary search tree		
	height = $\log n$		
	time taken for a particular key is $\log n$		
	Best case - element present in the root		
	Best case time - k i.e. $B(n) = O(1)$; $\Omega(1)$; $\Theta(1)$		
	Worst case - searching for a leaf element - depends upon the height of the tree		
	Worst case time - $\log n$ i.e. $O(\log n)$		
	min $w(n) = \log n$; max $w(n) = n$		
17	Disjoint sets		
	No common numbers between two sets - intersection is zero		
	Operations - find, union		
	Find - search or check membership		
	Union - Add an edge		
	<i>Kruskal algorithm: If you take an edge and both the vertices belong to the same set, then there is a cycle in the graph</i>		
	Weighted union is used while adding edges and detecting cycle		

	Collapsing find - process of directly linking node to a direct parent of a set is called collapsing find - reduces the time to find		
18	Divide and conquer - Strategy 1		
	Strategy - an approach for solving a problem		
	If a problem cannot be solved, divide it into sub-problems and find a solution for each sub problem, combine the solutions. <i>One point to note is that each sub problem should be similar to the original problem only.</i>		
	Recursive in nature		
	Should have one method to combine the solutions of each sub problem		
19	Problems under Divide and Conquer		
	Binary search		
	Finding maximum and minimum		
	MergeSort		
	QuickSort		
	Strassen's matrix multiplication		
20	Recurrence relation 1: $T(n) = T(n-1) + 1$		
	void test(int n)		
	{		
	if(n > 0){		
	printf("%d",n);		
	test(n-1)		
	}		
	}		
	test(3)		
	3. test(2)		

	2. test(1)		
	1 test(0)		
	each print statement takes constant time 1 and there are n+ 1 calls made to the function. we can ignore the last call when it is not printing		
	$f(n) = n + 1$ calls ; $O(n)$		
	$T(n) = T(n-1) + 1$; if we ignore if condition		
	Let us solve this relation;		
	if we know $T(n-1)$, we can get $T(n)$		
	$T(n-1) = T(n-2) + 1$		
	$T(n) = [T(n-2) + 1] + 1$		
	$T(n) = T(n-3) + 3$		
continue for k times		
	$T(n) = T(n-k) + k$		
	We would stop after k substitutions; now we need to find k		
	Assume $n - k = 0$; therefore $n = k$		
	$T(n) = T(n-n) + n$		
	$T(n) = T(0) + n$		
	$T(n) = n + 1$ i.e. $\theta(n)$		
21	Recurrence relation 2: $T(n) = T(n-1) + n$ (decreasing function)		
	void test(int n)	$T(n)$	
	{		
	if(n > 0)		1
	{		
	for(i = 0; i < n; i++)	$n+1$	
	{		
	printf("%d", n);	n	
	}		
	test(n-1);	$T(n-1)$	

	}		
	}		
		$T(n) = T(n-1) + 2n + 2$ i.e. $\theta(n)$	
	we can also write $T(n) = T(n-1) + n$ for $n > 0$		
	$T(n) = 1$ for $n = 0$		
	$T(n)$	n time	
	n $T(n-1)$	n-1 time	
	n-1. $T(n-2)$	n - 2 time	
	n-2 $T(n-3)$	n - 3 time	
		
	$T(2)$		
	2 $T(1)$	2 units of time	
	1 $T(0)$	1 unit of time	
	for $T(0)$ it does nothing	0 unit of time	
	time taken -		
		$0 + 1 + 2 + \dots + n-1 + n$	
	$\theta(n^2)$	$T(n) = n(n+1)/2$	
	$T(n) = T(n-1) + n$		
	$T(n-1) = T(n-2) + n-1$		
	thus, $T(n) = T(n-2) + (n-1) + n$	**remember, don't add the terms	
	$T(n) = T(n-3) + (n-2) + (n-1) + n$		
	$T(n) = T(n-k) + (n-(k-1)) + (n-(k-2)) \dots + (n-1) + n$	if we continue for k times	
	assume $n - k = 0$; $n = k$		
	Thus, $T(n) = T(n-n) + (n - n + 1) + (n - n + 2) \dots + (n-1) + n$		
	$T(n) = T(0) + n(n+1)/2$		
	$T(n) = 1 + n(n+1)/2$	$\theta(n^2)$; this extra 1 is owing to the calls	
22	Recurrence relation 3: $T(n) = T(n-1) + \log n$		
	void test(int n)	$T(n)$	

{		
if(n>0)		
{		
for(i = 1; i <n; i = i*2)		
{		
printf("%d", i);	log n times	
}		
test(n-1);	T(n-1)	
}		
}		
$T(n) = T(n-1) + \log n$ for $n > 0$		
$T(n) = 1$ for $n = 0$		
Solve using tree method,		
T(n)		
logn T(n-1)		
log(n-1) T(n-2)		
log(n-2) T(n-3)		
.....		
log2 T(!)		
log 1 T(0)		
logn + log(n-1) ++ log2 + log1		
$\log[n(n-1)(n-2)....2.1] = \log(n!)$	there is no tight bound for this function but there is an upper bound for it	
$O(n \log n)$		
Solving using induction method.		
$T(n) = T(n-1) + \log n$		
$T(n) = T(n-2) + \log(n-1) + \log(n)$		
$T(n) = T(n-3) + \log(n-2) + \log(n-1) + \log n$		

		
	$T(n) = T(n-k) + \log n + \log(n-1) + \dots \log 1$		
	Asume $n-k = 0$		
	$T(n) = T(0) + \log n!$		
	$T(n) = 1 + \log n!$		
	$O(n \log n)$		
23	How to get the direct answer for a recurrence relation?		
	$T(n) = T(n-1) + 1$	$O(n)$	
	$T(n) = T(n-1) + n$	$O(n^2)$	
	$T(n) = T(n-1) + \log n$	$O(n \log n)$	
	$T(n) = T(n-1) + n^2$	$O(n^3)$	
	$T(n) = T(n-2) + 1$	$O(n/2) \sim O(n)$	
	$T(n) = T(n-100) + n$	$O(n^2)$	
	$T(n) = 2T(n-1) + 1$???	
24	Recurrence relation 4: $T(n) = 2T(n-1) + 1$		
	Test(int n)	$T(n)$	
	{		
	if(n > 0)		1
	{		
	printf("%d", n);		1
	test(n-1);	$T(n-1)$	
	test(n-1);	$T(n-1)$	
	}		
	}		
		$T(n) = 2T(n-1) + 1$	
	$T(n) = 2T(n-1) + 1$ for $n > 0$		
	$T(n) = 1$ for $n = 0$		

	Solve using recursion tree method		
	1 T(n-1) T(n-1)		2
	1 T(n-2) T(n-2)	1 T(n-2) T(n-2)	4
	1 T(n-3) T(n-3) 1 T(n-3) T(n-3)	1 T(n-3) T(n-3) 1 T(n-3) T(n-3)	8
		
	T(0). T(0)		2 ^k
	1 + 2 + 2 ²+2 ^k = 2 ^(k+1) - 1		
	as, a + ar + ar ²ar ^k = a(r ^(k+1) - 1)/(r - 1)		
	Assume n - k = 0		
	thus, 2⁽ⁿ⁺¹⁾ - 1	O(2ⁿ)	
	Back substitution method		
	T(n) = 2T(n-1) + 1		
	T(n) = 4T(n-2) + 2 + 1		
	T(n) = 8T(n-3) + 4 + 2 + 1		
		
	T(n) = 2 ^k T(n - k) + 2 ^(k-1) + 2 ^(k)2 ³ + 2 ² + 1		
	Assume n - k = 0		
	n = k		
	T(n) = 2 ⁿ T(0) + 1 + 2 + 2 ² + 2 ⁿ⁻¹		
	T(n) = 2ⁿ + 2ⁿ - 1 i.e. 2⁽ⁿ⁺¹⁾ - 1		
25	Master theorem for decreasing function		
	T(n) = T(n-1) + 1	O(n)	
	T(n) = T(n-1) + n	O(n ²)	
	T(n) = T(n-1) + logn	O(nlogn)	
	T(n) = 2T(n-1) + 1	O(2 ⁿ)	
	T(n) = 3T(n-1) + 1	O(3 ⁿ)	

	$T(n) = 2T(n-1) + n$	$O(2^{2n})$	
	$T(n) = 2T(n-2) + 1$	$O(2^{n/2})$	
	$T(n) = aT(n-b) + f(n)$		
	$a > 0, b > 0$ and $f(n) = O(n^k)$ where $k \geq 0$		
	if $a = 1, O(n^{k+1})$ or $O(n \cdot f(n))$		
	if $a > 1, O(n^k \cdot a^{n/b})$		
	if $a < 1, O(n^k)$ or $O(f(n))$		
26	Dividing functions		
	test(int n)	$T(n)$	
	{		
	if(n > 1)		
	{		
	printf("%d", n);		1
	test(n/2)	$T(n/2)$	
	}		
	}		
	$T(n) = T(n/2) + 1$ for $n > 1$		
	$T(n) = 1$ for $n = 1$		
	$T(n)$		
	1 $T(n/2)$		
	1 $T(n/2^2)$		
	1 $T(n/2^3)$		
continue for k times		
	1 $T(n/2^k)$		

	assume , $n/2^k = 1$		
	thus, we have taken k steps overall		
	since, $n/2^k = 1 \Rightarrow k = \log n$ with base 2	$O(\log n)$	
	Solving by substitution method		
	$T(n) = T(n/2) + 1$		
	$T(n) = T(n/2^2) + 2$		
	$T(n) = T(n/2^3) + 3$		
		
	$T(n) = T(n/2^k) + k$		
	assume $n/2^k = 1$		
	thus, $k = \log n$ with base 2		
	$T(n) = T(1) + \log n$		
	$O(\log n)$		
27	Recurrence relation: $T(n) = T(n/2) + n$		
	$T(n) = T(n/2) + n$ for $n > 1$		
	$T(n) = 1$ for $n=1$		
	$T(n)$		
	$T(n/2) \quad n$		
	$T(n/2^2) \quad n/2$		
	$T(n/2^3) \quad n/2^2$		
		
	$T(n/2^k). \quad n/2^{(k-1)}$		
	$T(n) = n + n/2 + n/2^2 + n/2^3..... + n/2^k$		
	$T(n) = n[1 + 1/2 + 1/2^2 + 1/2^3 +.....1/2^k]$		
	$T(n) = n \cdot 1 = n$		
	$O(n)$		

	Using substitution method		
	$T(n) = T(n/2) + n$		
		
	$T(n) = T(n/2^2) + n/2 + n$		
		
	$T(n) = T(n/2^3) + n/2^2 + n/2 + n$		
		
	$T(n) = T(n/2^k) + n/2^{k-1} + \dots + n/2^2 + n/2 + n$		
	Assume $n/2^k = 1$		
	$k = \log n$ with base 2		
	$T(n) = T(1) + n[1/2^{k-1} + \dots + 1/2^2 + \dots + 1]$		
	$T(n) = 1 + 2n \sim O(n)$		
28	Recurrence Relation: $T(n) = 2T(n/2) + n$		
	void test(int n)	$T(n)$	
	{		
	if(n > 1)		
	{		
	for(int i = 0; i < n ; i++)		
	{		
	stmt	n	
	}		
	test(n/2);	$T(n/2)$	
	test(n/2);	$T(n/2)$	
	$T(n) = 2T(n/2) + n$ for $n > 1$		
	$T(n) = 1$ for $n = 1$		
	Solve using recursion tree method,		

	$T(n)$		
	$T(n/2). \quad T(n/2) \quad n$	n	
	$T(n/2^2) \quad T(n/2^2) \quad T(n/2^2) \quad T(n/2^2) \quad n/2$	n	
	$T(n/2^3). \quad T(n/2^3). \quad T(n/2^3). \quad T(n/2^3). \quad T(n/2^3). \quad T(n/2^3). \quad T(n/2^3). \quad T(n/2^3).$	n	
	n	
	$T(n/2^k).....$		
		n	
	assume $n / 2^k = 1$		
	$k = \log n$ with base 2		
	$T(n) = nk \sim O(n \log n)$		
	Using backsubstitution method;		
	$T(n) = 2T(n/2) + n$		
	$T(n/2) = 2T(n/2^2) + n/2$		
	$T(n) = 2[2T(n/2^2) + n/2] + n$		
	$T(n) = 2^2T(n/2^2) + n + n$		
	$T(n) = 2^3T(n/2^3) + 3n$		
continue for k times		
	$T(n) = 2^kT(n/2^k) + kn$		
	Asume $T(n/2^k) = T(1)$		
	$k = \log n$ with base 2		
	Thus, $T(n) = n + n \log n \sim O(n \log n)$		
	29 Masters Theorem for dividing functions		
	$T(n) = aT(n/b) + f(n)$	$\log a$ with b	
	$a \geq 1; b > 1; f(n) = \theta(n^k \log^p n)$	k	
	case 1: if $\log a$ with base b $> k$ then $\theta(n^{\log a \text{ with base b}})$		
	case 2: if $\log a$ with base b $= k$ then		

	if $p > -1$ $\theta(n^k \log^{(p+1)} n)$		
	if $p = -1$ $\theta(n^k \log \log n)$		
	if $p < -1$ then $\theta(n^k)$		
	case 3: if \log_a with base $b < k$		
	then, if $p \geq 0$, $\theta(n^k \log^p n)$		
	if $p < 0$, $\theta(n^k)$		
	$T(n) = 2T(n/2) + 1$		
	$a = 2$		
	$b = 2$		
	$f(n) = \theta(n^0) * \log n^0$		
	$k = 0$; $p = 0$		
	here, \log_a with base $b > k$		
	$\theta(n^1)$ where \log_a with base b is 1		
	$T(n) = 4T(n/2) + n$		
	\log_a with base $b = 2$		
	$k = 1$		
	$p = 0$		
	this is an example of case 1		
	$\theta(n^2)$		
	$T(n) = 8T(n/2) + n$		
	\log_8 with base $2 = 3 > k = 1$		
	$\theta(n^3)$		
	$T(n) = 9T(n/3) + 1$		
	\log_a with base $b = 2 > k$		
	$\theta(n^2)$		

	$T(n) = 9T(n/3) + n^2$		
	loga with base $b = 2 = k$	case 2	
	$\theta(n^2)$		
	$T(n) = 8T(n/2) + n$		
	$\theta(n^3)$		
	$T(n) = 2T(n/2) + n$		
	loga with base $b = k = 1$; $p = 0$		
	case 2		
	$\theta(n \log n)$		
	$T(n) = 4T(n/2) + n^2$		
	$\theta(n^2 \log n)$		
	$T(n) = 4T(n/2) + n^2 \log n$		
	$\theta(n^2 \log n^2)$		
	$T(n) = 8T(n/2) + n^3$		
	$\theta(n^3 \log n)$		
	$T(n) = 2T(n/2) + n/\log n$		
	log a with base $b = k = 1$		
	$p = -1$		
	$\theta(n \log \log n)$		
	$T(n) = 2T(n/2) + n/\log n^2$		
	$p = -2$		
	$\theta(n)$		

	$T(n) = 2T(n/2) + n^2$		
	loga with base $b < k$		
	$\theta(n^2)$		
	$T(n) = 2T(n/2) + n^2$		
	$\theta(n^2 \log n)$		
	$T(n) = 2T(n/2) + n^3$		
	loga with base $b < k$		
	$\theta(n^3)$		
30	$T(n) = 2T(n/2) + 1$		
	loga with base $b = 1$		
	$k = 0$		
	loga with base $b > k$		
	$\theta(n^1)$		
	$T(n) = 4T(n/2) + 1$		
	loga with base $b = 2$		
	$k = 0$		
	$\theta(n^2)$		
	$T(n) = 4T(n/2) + n$		
	loga with base $b = 2$		
	$k = 1$		
	$\theta(n^2)$		
	$T(n) = 8T(n/2) + n^2$		
	loga with base $b = 3$		
	$k = 2$		

	$\theta(n^3)$		
	$T(n) = 16T(n/2) + n^2$		
	log a with base b = 4		
	k = 2		
	$\theta(n^4)$		
	$T(n) = T(n/2) + n$		
	log a with base b = 0		
	k = 1		
	$\theta(n)$		
	$T(n) = 2T(n/2) + n^2$		
	log a with base b = 1		
	k = 2		
	$\theta(n^2)$		
	$T(n) = 2T(n/2) + n^2 \log n$		
	log a with base b = 1		
	k = 2		
	$\theta(n^2 \log n)$		
	$T(n) = 4T(n/2) + n^3 \log^2 n$		
	log a with base b = 2		
	k = 3		
	$\theta(n^3 \log^2 n)$		
	$T(n) = 2T(n/2) + n^2 / \log n$		
	log a with base b = 1		
	k = 2		

	$\theta(n^2)$		
	$T(n) = T(n/2) + 1$		
	log a with base b = 0		
	k = 0		
	$\theta(\log n)$		
	$T(n) = 2T(n/2) + n$		
	log a with base b = 1		
	k = 1		
	p = 0		
	$\theta(n \log n)$		
	$T(n) = 2T(n/2) + n \log n$		
	log a with base b = 1		
	k = 1		
	p = 1		
	$\theta(n \log^2 n)$		
	$T(n) = 4T(n/2) + n^2$		
	log a with base b = 2		
	k = 2; p = 0		
	$\theta(n^2 \log n)$		
	$T(n) = 4T(n/2) + (n \log n)^2$		
	log a with base b = 2		
	k = 2, p = 2		
	$\theta(n^2 \log^3 n)$		
	$T(n) = 2T(n/2) + n/\log n$		

	log a with base b = 1		
	k = 1; p = -1		
	theta(nloglogn)		
	$T(n) = 2T(n/2) + n/\log^2 n$		
	log a with base b = 1		
	k = 1; p = -2		
	theta(n)		
31	Root function Recurrence relation		
	$T(n) = T(\sqrt{n}) + 1$ for $n > 2$		
	$T(n) = 1$ for $n = 2$		
	$T(n) = T(\sqrt{n}) + 1$		
	$T(n) = T(n^{1/2}) + 1$equation 1		
	using substitution		
	$T(n) = T(n^{1/2^2}) + 2$equation 2		
	$T(n) = T(n^{1/2^3}) + 3$equation 3		
	$T(n) = T(n^{1/2^k}) + k$equation 4		
	assume, $n = 2^m$		
	$T(2^m) = T(2^{m/2^k}) + k$		
	assume $T(2^{m/2^k}) = T(2)$		
	thus, $m/2^k = 1$		
	$m = 2^k$		
	$k = \log m$ with base 2		
	substituting value of n		
	$m = \log n$ with base 2		
	therefore, $k = \log \log n$ with base 2		
	theta(loglogn with base 2)		

32	Binary Search Iterative Method		
	To perform binary search, the prerequisite is that the list must be in sorted order	A = {3, 6, 8, 12, 14, 17, 25, 29, 31, 36, 42, 47, 53, 55, 62}	
	we need two index pointers, one is low at the starting point and the other is high at the end point	l = 1, h = 15 (lowest and highest index); mid = 8	
	mid = low + high / 2 and we take the floor value	key value = 42; A[mid] = 29 --> key > A[mid]	
	the key value is on the right hand side as key value is greater than A[mid]		
	we will change low to mid + 1	l = 9, h = 15; mid = 9 + 15 / 2 = 12	
		A[mid] = 47 > key	
	we will change high to mid - 1 as key < A[mid]		
		h = 11, l = 9, mid = 10; A[mid] = 36	
		A[mid] < key	
	we will change low to mid + 1	l = 11; h = 11; mid = 11; A[mid] = 42	
	we can return the index as we have found the key value	A[mid] = key	
	therefore, binary search looks faster than linear search. It just took 4 comparisons		
	int BinSearch(A, n, key)		
	{		
	l = 1, h = n		
	mid = l + h / 2 - take floor value		
	while(l <= h){		
	if(key == A[mid])		
	{ return index i.e.element is found}		
	else if(key < A[mid])		
	{h= mid-1;}		
	else {		
	l = mid + 1;}		
	}		

	return 0;		
	}		
	Time taken for binary search = $\log n$		
	min time: $O(1)$		
	max time: $O(\log n)$		
	avg time = add time for each element and divide by number of elements		
33	Binarysearch Recursive method		
	Algorithm RBinarySearch(l,h,key)	$T(n)$	
	{		
	if(l==h)		1
	{		
	if(A[l]== key)		
	{		
	return l;		
	}		
	else		
	{		
	return 0;		
	}		
	else		
	{		
	mid = $l + h / 2$ //taking floor value		1
	if(key == A[mid])		1
	{return mid;}		
	if(key < A[mid])		1
	{		
	return RBinarySearch(l, mid - 1, key)	$T(n/2)$	
	}		

	else		
	{		
	return RBinarySearch(mid+1, h, key)	$T(n/2)$	
	}		
	}		
		$T(n) = 1; n = 1$	
		$T(n) = T(n/2) + 1 \text{ for } n > 1$	
		$\theta(\log n)$	
	34 Heaps		
a	Representation of a binary tree using an array		
	T {A, B, C, D, E, F, G}		
	if a node is at index i;		
	its left child is at node $2*i$		
	its right child is at node $2*i + 1$		
	its parent is at node $i/2$		
	if there are missing nodes, we leave a blank in its place in the array		
b	Full binary tree		
	In its height, it has maximum number of nodes and if we wish to add a node, height would increase		
	Max no. of nodes = $2^h - 1$		
c	Complete binary tree		
	there is no missing element from first element to the last element in array representation of the binary tree		
	Every full binary tree is also a complete binary tree		
	A complete binary tree is a full binary tree until height $h - 1$		
	Height of a complete binary tree would be minimum i.e. $\log n$		
d	Heap		

	Heap is a complete binary tree		
	Max Heap: every node has value greater than all its descendants {50, 30, 20, 15, 10, 8, 16}		
	Min Heap: every node has value smaller or equal to than all its descendants {10, 30, 20, 35, 40, 32, 25}		
35	Insert operation in a max heap		
	Insert 60 in the above given max heap		
	this value should be inserted in the last free space in the array		
	i.e. left child of the left most leaf node		
	Then, adjust the elements to make it as a heap		
	So, compare and move 60 up the levels and in the array check at $i/2$ indices where initially i would be the last empty index where 60 was inserted		
	Time taken would be equal to the number of swaps		
	this depends upon the height of the tree i.e. $\log n$, hence $O(\log n)$		
	minimum time is of no swaps $O(1)$; max would be $O(\log n)$		
36	Delete operation in a max heap		
	From the heap, we need to remove the root / top most element only		
	The last element in the complete binary tree would come in its place		
	Adjust the elements to maintain heap order		
	From the root towards the leaf, adjust		
	Compare the children ($2i$ and $2i + 1$) and whichever child is greater than compare with the parent		
	Time taken depends upon the height; max could be $O(\log n)$		
	Whenever you delete from max heap, you get the next max element and in case of min heap, it would be the next min element		

37 HeapSort		
For a given set of numbers, create a heap		
Delete all the elements from the heap		
Total N elements we have inserted; each element we assume is moved up to the root; so time taken $O(N\log N)$		
Then we delete the elements		
Store deleted elements in the array in free space in the end		
Deletion also takes $O(N\log N)$ time		
Thus, heapsort takes $O(N\log N)$		
38 Heapify		
The process of creating heap but direction is opposite than creating a heap		
$O(N)$		
39 Priority Queue		
elements will have priority and they would be inserted and deleted as per the priority order		
For min heap, smaller the no. higher the priority		
For max heap, greater the no. higher the priority		
$O(\log N)$ for insertion and/or deletion		
40 TwoWay MergeSort - Iterative method	Algorithm Merge(A, B, m, n)	
merging two sorted lists to get a sorted result	{i = 1, j = 1, k = 1;	
A = {2, 8, 15, 18} i	while(i <= m && j <= n){	
B = { 5, 9, 12, 17} j	if(A[i] < B[j])	
Compare A(i) with B(j) to get C(k) and move to next location	{	
m + n elements are obtained , thus theta(m + n)	C[k++] = A[i++];	
	}	
	else {	
	C[k++] = B[j++];	

		}	
		for(; i <=m; i++){	
		C[k++]=A[i];	
		}	
		for(;j<= n;j++){	
		C[k++] = B[j];	
		}	
		}	
41	Merging more than two lists		
	M-way merging		
	A = {4, 6, 12}		
	B = {3, 5, 9}		
	C = {8, 10, 16}		
	D = {2, 4, 18}		
	One way is that we merge A and B; C and D and then finally merge the two resulting lists --> so we perform merge three times here		
	Another way is that we first merge A and B; then we merge resulting list with C ; and the resulting list with D		
	Two-way mergesort is an iterative process whereas mergeSort is a recursive process		
	A = {9, 3, 7, 5, 6, 4, 8, 2} - given an array and we have to sort them using 2-way mergesort		
1st pass	We would consider each element as a sorted list and merge	merged n elements in this pass	
	First select two lists 3 and 9; then merge them - 3, 9		
	Similarly, we select two lists 7 and 5 , merge them - 5 and 7		
	Another lists we get are {4, 6} and {2, 8}		
	Now, we have 4 lists with two elements each		
2nd pass	When we merged we kept the resulting 4 lists in another array B; B = {{3, 9}, {5, 7}, {4, 6}, {2, 8}}	merged n elements in this pass	

	We merge two lists each		
3rd pass	$C = \{\{3, 5, 7, 9\}, \{2, 4, 6, 8\}\}$	merged n elements in this pass	
	we merge the above two lists to get a single sorted list		
	$D = \{2, 3, 4, 5, 6, 7, 8, 9\}$		
	$\log(\text{no of elements}) = \text{no. of passes}$		
	Time complexity: $O(n \log n)$		
42 MergeSort			
	$A = \{9, 3, 7, 5, 6, 4, 8, 2\}$	Algorithm MergeSort(l, h){	$T(n)$
	If there is a single element, we can consider it as a base or small problem {Divide and conquer}		
		if(l < h){	
		mid = (l + h) / 2;	1
		MergeSort(l, mid);	$T(n/2)$
		MergeSort(mid + 1, h);	$T(n/2)$
		Merge(l, mid, h);	n
		}	$T(n) = 2T(n/2) + n$ for $n > 1$
		}	$T(n) = 1$ for $n = 1$
	time complexity: $\theta(n \log n)$		using master's theorem, $a = 2, b = 2, k = 1$
	merging is done in post order traversal		\log_a with base $b = 1 = k$
			thus, it is case 2
			$\theta(n \log n)$
43 Pros of MergeSort		Cons of MergeSort	
	works great for Large size lists	Extra space (not inplace sort)	
	suitable for Linked List	no small problem	
	supports external sorting	recursive and uses a stack (need $n + \log n$ space) i.e. space complexity: $O(n + \log n)$ where n is the extra space and $\log n$ is the stack space	
	stable: the order of duplicates is maintained		
		insertion sort ($O(n^2)$)	

		mergesort $O(n \log n)$	
		for small problems, $n \leq 15$; insertionsort works better --> use insertion sort	
43 QuickSort			
	students arranging themselves in increasing order of heights		
	10 80 90 60 30 20		
	5 6 3 4 2 1 9		
	4 6 7 10 16 12 13 14		
	A = {10, 16, 8, 12, 15, 6, 3, 9, 5, INFINITY}	partition(l, h){	
	select first element as a pivot	pivot = A[l];	
	pivot = 10	i = l; j = h;	
	we need to find the sorted position for 10	while(i < j){do	
	i starting from pivot and j starting from infinity	{	
	i would check for elements greater than 10; j would check for elements smaller than pivot	i++;	
	we are using the partitioning procedure	} while(A[i] <= pivot);	
	increment i until next value is greater than 10 and decrement j until next value is smaller than pivot; stop and swap	do	
	{10, 5, 8, 9, 3, 6, 15, 12, 16}	{	
	send pivot element at j position	j--;	
	now, we can sort the two lists around the partitioning position by performing quicksort recursively	}while(A[j] > pivot);	
		if(i < j){	
	QuickSort(l, h)	swap(A[i], A[j]);	
	{	}	
	if(l < h)	swap(A[l], A[j]);	
	{	return j;	
	j = partition(l, h);	}	
	QuickSort(l, j);		
	QuickSort(j+1, h);		

	}		
	}		
44 QuickSort Analysis			
suppose it is partitioning in the middle of 1 and 15th index			
then, two partitions: [1, 7] ; [9, 15]			
further partitions: [1, 3]; [5, 7]; [9, 11]; [13, 15]			
at each level , n elements are being handled			
and there are logn levels			
thus time complexity for best case: $O(n \log n)$			
median : middle element of a sorted list			
best case of quicksort is that the partitioning occurs exactly at the middle			
worstcase: if we have an already sorted list			
time complexity for worstcase: $O(n^2)$			
to handle this, try taking middle element as a pivot			
2. select random element as a pivot			
45 Strassen's matrix multiplication			
A = [a11 a12			
a21. a22]			
B = [b11 b12			
b21 b22]			
Cij = Summing up Aik*Bkj			
for(i = 0; i < n ; i++){			
for(j = 0; i < n ; j++){			
C[i,j]= 0;			
for(k=0;k<n;k++){			
C[i,j] += A[i, k]*B[k, i];			

	}		
	}		
	}		
	$C_{11} = a_{11} \cdot b_{11} + a_{12} \cdot b_{21}$		
	$C_{21} = a_{11} \cdot b_{12} + a_{12} \cdot b_{22}$	$A = [a_{11}]$	
	$C_{21} = a_{21} \cdot b_{11} + a_{22} \cdot b_{21}$	$B = [b_{11}]$	
	$c_{22} = a_{21} \cdot b_{12} + a_{22} \cdot b_{22}$	$C = [a_{11} \cdot b_{11}]$	
	for $[2 \times 2]$ matrix, we would use above formula	for $[1 \times 1]$ matrix, use above formula	
	we assume that the matrix has dimensions of power of 2	Algorithm MM(A, B, n)	
		{	
		if(n <= 2	
	8 times the function is calling itself	{	
	$T(n) = 8T(n/2) + n^2$ for $n > 1$	C = 4 formula stated above;	
	$a = 8, b = 2, \log a$ with base $b = 3$	}	
	$k = 2$	else	
	it is case 1 of master's theorem	{	
	$\theta(n^3)$	mid = $n/2$	
		MM(A ₁₁ , B ₁₁ , $n/2$) + MM(A ₁₂ , B ₂₁ , $n/2$);	
		MM(A ₁₁ , B ₁₂ , $n/2$) + MM(A ₁₂ , B ₂₂ , $n/2$);	
		MM(A ₂₁ , B ₁₁ , $n/2$) + MM(A ₂₂ , B ₂₁ , $n/2$);	
		MM(A ₂₂ , B ₂₂ , $n/2$) + MM(A ₂₁ , B ₁₂ , $n/2$);	
		}	
		}	
	Strassen's approach -		
	has given 4 different formulas with 7 multiplications	$P = (A_{11} + A_{22})(B_{11} + B_{22})$	
	$C_{11} = A_{11} \cdot B_{11} + A_{12} \cdot B_{21}$	$Q = (A_{21} + A_{22}) B_{11}$	
	$C_{21} = A_{11} \cdot B_{12} + A_{12} \cdot B_{22}$	$R = A_{11}(B_{12} - B_{22})$	
	$C_{21} = A_{21} \cdot B_{11} + A_{22} \cdot B_{21}$	$S = A_{22}(B_{21} - B_{11})$	

	$C22 = A21 \cdot B12 + A22 \cdot B22$	$T = (A11 + A12)B22$	
		$U = (A21 - A11)(B11 + B12)$	
		$V = (A12 - A22)(B21 + B22)$	
		$C11 = P + S - T + V$	
		$C12 = R + T$	
		$C13 = Q + S$	
		$C22 = P + R - Q + U$	
		$T(n) = 7T(n/2) + n^2$ for $n > 2$	
		$T(n) = 1$ for $n \leq 2$	
		using master's theorem,	
		$O(n^{\log_7 \text{ with base } 2}) = O(n^{2.81})$	
	Strategies used for solving optimization problems - Greedy Method, Dynamic programming, branch and bound		
46	Greedy method		
	Design with which we can adopt to solve similar problems		Greedy method says that each problem should be solved in stages - each stage we give an input, check if the solution is feasible then we pick it up and move to next stage
	Solving optimization problems		Algorithm Greedy(a, n) $a = \{a1, a2, a3, a4, a5\}$; $n = 5$
	Optimization problem: Problems which require either minimum or maximum result		{
	Suppose we have a problem P where we need to travel from source A to destination B, we can have several solutions such as walking on foot, travel by an airplane, ride on a bike, travel by a bus, drive on a car, go by a train and so on..... Now, we notice that we also have some constraints. The solutions that satisfy the conditions given in a problem are called feasible solutions	Minimum cost journey - " Minimization problem "; then feasible solutions giving minimum cost are called optimal solutions . There can be many feasible solutions but only one optimal solution	for $i = 1$ to n do { $x = \text{select}(a)$:
	Example: selecting a car to purchase	Example: Hire a person for your company	if feasible(x) then

	Method 1: Looking at all the models available in the city	Method 2: Conduct an assessment center to filter people at each stage and get the best person	{
	Method 2: Checking for the features of the cars and filtering and selecting based on your preferences - greedy method	So, the person may not be the best but the approach is greedy here as we are using our criteria and constraints to choose the best person	solution = solution + x;
47 Knapsack problem			}
	n = 7; m = 15	Bag capacity is 15 kgs and we have been given 7 objects. we have to fill this bag with these objects. Profit is the gain we get by transferring this object. Problem is a container loading problem. Problem is filling the container with the objects as the capacity of container is limited	}
Objects	{0 1 2 3 4 5 6 7 }	Optimization and maximization problem	}
Profits	{P 10 5. 15 7 6 18 3 }	Constraints : Bag weight limit	
Weights	{W 2 3 5 7 1 4 1 }		
Profit by weight	{P/W 5 1.3 3 1 6 4.5 3 }		
0 ≤ x ≤ 1	Objects are divisible i.e. we can take just half kg of object 1 and may be 2 kgs of object 2 and so on		
x	()		
	x1 x2 x3 x4 x5 x6 x7		
Method 1	Take the thing that have maximum profit		
Method 2	Take things with smaller weight so that you can put in more things		
Method 3	Take things that have highest profit by weight		
	Let's use method 3		

	First, I include object 5 that has maximum profit by weight. Then I check remaining weight I can put in. We can still put in 14 kgs. Then we select all the quantity of object 1. Remaining weight limit 12 kgs. Add all of object 6. Remaining weight limit 8 kgs. Add all of object 3. Remaining weight limit 3 kgs. Add all of object 7. Remaining weight limit is 2 kgs. Add 2/3 of object 2 as we have only 2 kgs limit remaining.						
x	(1 2/3 1 0 1 1 1)						
	Calculate total profit and verify weight						
	Total weight = $1*2 + 2/3*3 + 1*5 + 0*7 + 1*1 + 1*4 + 1*1 = 15$					//Multiplying x elements by Weight w for each object	
	Total profits = $1*10 + 2/3*5 + 1*15 + 1*6 + 1*18 + 1*3 = 54.6$					//Multiplying x elements by Profit P for each object	
48	0/1 Knapsack problem						
	Objects are indivisible and fractions are not allowed i.e. either you include the whole thing or you do not include it at all						
49	Job sequencing with deadlines					n = 5 (tasks)	
Jobs	J1	J2	J3	J4	J5		
Profits	20	15	10	5	1		
Deadlines	2	2	1	3	3		
	Assume that there is a machine, on which each job has to be processed and each job takes 1 unit of time (hour) for completion						
	Set of the jobs which can be completed within their deadlines such that profit is maximized					Constraints: deadlines must be met	
deadlines	0-----1-----2-----3					maximum 3 slots / jobs	
time slots	9am-----10am-----11am-----12am						
Jobs chosen	J2 J1 J4						
Profits	$15 + 20 + 5 = 40$						
Sequence	J1 --> J2 ---> J4 J2 --> J1 --> J4						

Job consider	Slot assign	Solution	Profit
J1	[1,2]	J1	20
J2	[0,1][1,2]	J1J2	20 + 15
J3	[0,1][1,2]	J1J2	20 + 15
J4	[0,1][1,2][2,3]	J1J2J4	20 + 15 + 5
J5	[0,1][1,2][2,3]	J1J2J4	20 + 15 + 5
50	Job sequencing with deadlines another example	n = 7 (jobs)	
Jobs	J1 J2 J3 J4 J5 J6 J7		
Profits	35 30 25 20 15 12 5		
Deadlines	3 4 4 2 3 1 2		
deadlines	0.....1.....2.....3.....4	4 SLOTS AVAILABLE	
Jobs chosen	J4 J3 J1 J2		
Profits	20 25 35 30	110	
51	Optimal Merge Patern		
	A = {3, 8, 12, 20}		
	B = {5, 9, 11, 16}		
	C = {3, 5, 8, 9, 11, 12, 16 20}	How merging works for two sorted lists ; time = $\theta(m + n)$	
	what happens if we have 4 lists?		
List	A B C D		
Sizes	6 5 2 3		
Choice 1	We can merge at a time two lists - first A and B; the merge it with C and finally with D. total cost= 11 + 13 + 16 = 40		

Choice 2	Merge A and C; D and E that gives cost 11 and 5 respectively. Merge resulting two lists which will cost 16. Total cost would be $11 + 5 + 16 = 32$						
Choice 3	Merge C and E, resulting list is merged with B and then finally with A; total cost = $5 + 10 + 16 = 31$						
optimal method	Always merge two small sized lists, then combined time would be reduced						
Example:							
List	x1	x2	x3	x4	x5		
Sizes	20	30	10	5	30		
Increasing order of sizes	5	10	20	30	30		
Lists	x4	x3	x1	x2	x5		
	First x4 and x3 are merged, cost = 15; then result is merged with x1; cost = 35; x2 and x5 are merged with cost = 60; the two resulting lists are merged with cost = 95						
Total cost	$15 + 35 + 60 + 95 = 205$						
	$3*5 + 3*10 + 2*20 + 2*30 + 2*30 = 205$					//multiplying distance of each node and size of each node	
52 Huffman Coding							
	Compression technique used to reduce the size of data or message						
Message	BCCABBDDEAECBBAEDDCC						
	Length = 20						
	it has to be sent using ASCII codes (8-bit)						
	A	65	01000001			Size = $8*20 = 160$ bits	
	B	66	01000010				

	C	67			
	D	68			
	E	69			
	Can we use our own codes instead of ASCII codes?				
	Fixed size method				
Character	A	B	C	D	E
Count	3	5	6	4	2
Code	000	001	010	011	100
message	BCCABBDDAECCBBAEDDCC				
bit code	001010....				
					size = $20 \times 3 = 60$ bits
					$5 \times 8 = 40$ bits for ASCII code translations
					$5 \times 3 = 15$ bits --> our assigned codes
					$40 + 15 = 55$ bits
					message: 60 bits
					char: 55 bits
					total message size: 115 bits
					so, the message size reduced from 160 bits to 115 bits
					thus, 40% reduction in size with fixed sized code
	Huffman coding - variable sized code				
					element that appears more / often should have a smaller sized code
character	A	B	C	D	E
count	3	5	6	4	2
code					
	first, arrange the letters with increasing count / frequency				
character	E	A	D	B	C
count	2	3	4	5	6

code	000 001 01 10 11	Merge two smaller ones, we get 5, then combine with D, we get 9. Combine B and C, we get 11. Finally, combine two resulting lists, we get 20.	
bit count	6 9 8 10 12	On left side paths, mark as 0 and on right side mark as 1	
total bits for message	45 bits	Bit count for message can also be obtained from the tree, by counting number of edges for a letter and multiplying by the number of occurrences for that letter in the message i.e. summation of distance and frequency of a letter	
ASCII codes for chart	5*8 = 40 bits		
assigned codes	12 bits		
total bits for tree/table	52 bits		
Size of total msg	52 + 45 = 97 bits		
Message transferred	001111101101111001011000111110100010100000110	A tree or a table would be needed along with it	
Decoding	BCCD...		
53	Minimum Cost Spanning Tree		
	$G = (V, E)$		
	$V = \{1, 2, 3, 4, 5, 6\}$	$ V = n = 6$	
	$E = \{(1,2), (2,3), (3,4), (4,5), (5, 6), (6,1)\}$	$ V - 1 = 5$ edges	
	the tree should not have a cycle		
	S is a subset of G, WHERE IN $S = (V', E')$	$V' = V; E' = V - 1$	
	Number of edges in graph = 6 out of which I have to select 5 edges for spanning tree - thus i can select in 6C5 ways	Suppose we have 7 edges, out of which the seventh edge (3,5) divides the graph into two cycles of less tha 6 vertices, then we can select 5 edges for spanning tree in 7C5 - 2 ways	

General formula	$ E C(V -1)$ - no. of cycles		
	Now, if we have a weighted graph, I wish to know the number of possible spanning tree		
	Vertices = 4		
	Edges = 3		
	cost = 14		
	similarly , depending upon the edges we select, cost may vary each time		
	Can I found the minimum cost spanning tree?		
Method 1	Try all possible spanning trees and get the minimum cost spanning tree		
Method 2	Prim's algorithm (Greedy method)		
Method 3	Krskal's algorithm (Greedy method)		
Method 2:	Prim's algorithm		
	Select the minimum cost edge from the graph first	(6,1) ; w = 10	
	Then, following this select minimum cost edge but make sure it is connected to previously chosen vertices	(5, 6); w = 25	
		(5, 4); w = 22	
		(4, 3); w = 12	
		(3, 2); w = 16	
		(2, 7); w = 14	
	Now, if we add costs of all the chosen edges, total cost = 99		
	For non connected graphs we cannot find the minimum cost spanning tree or spanning tree		
Method 3	Kruskal's method		
	Always select smallest cost edge		
	(1,6); w = 10		

	(3, 4); w = 12		
	(2, 7); w = 14		
	(2, 3); w = 16		
	(4, 5); w = 22		
	(5, 6); w = 25		
	total cost = 99		
	vertices count : $ V $	To get a minimum cost edge each time, min heap can be used	
	edges count: $ V - 1$	$\theta(n \log n)$	
	$\theta(V E)$		
	$\theta(n \cdot e) = \theta(n^2)$		
	for non-connected graphs, spanning tree cannot be found		
	Kruskal algo may give spanning tree for those non connected components but not for the graph as a whole		
	if in a certain graph, certain edges' weights are missing, then use the given weights of remaining edges to guess the weight		
	54 Dijkstra algorithm		
	Single source shortest path to all the vertices		
	find the shortest path to a vertex and update it to other vertices. this updation is called relaxation		
	Relaxation		
	$\text{if}(d[u] + c(u,v) < d[v]) \{ d[v] = d[u] + c(u,v) \}$		
	no of vertices = $ V $		
	at most no. of vertices relaxing = $ V $		
	worst case time of Dijkstra algorithm: $\theta(n^2)$		

	Example - starting vertex is 1						
selected vertex	2	3	4	5	6		
	4	50	45	10	infinity	infinity	
	5	50	45	10	25	infinity	
	2	45	45	10	25	infinity	
	3	45	45	10	25	infinity	
	6	45	45	10	25	infinity	
	Another example - starting vertex is 1						
selected vertex	{2,	3,	4}				
	2	{3,	infinity,	5}			
	4	{3,	infinity,	5}			
	3	{3,	7,	5}			
		{3,	7,	5}			
	Another example - starting vertex is 1						
selected vertex	{2,	3,	4}				
	2	{3,	infinity,	5}			
	4	{3,	infinity,	5}			
	3	{3,	7,	5}			
		{-3,	7,	5}			
	Dijkstra algorithm might work or might not work in case of an edge having negative weightage						
55	Dynamic programming						
	Dynamic programming vs greedy method						

	In Greedy method, we try to follow a predefined procedure that gives us the best / optimal result. The procedure is already known for optimization. But, in dynamic programming, we try to get all the solutions and then decide the best solution. Mostly dynamic programming questions are solved using recursive procedures. They follow a principle of optimality. In greedy method, decision is taken just once and followed through whereas in dynamic programming, decision is taken at each step		
	Example:		
	<i>Fibonacci series</i>		
	fib(n) = 0 if n = 0	$T(n) = 2T(n-1) + 1$ {Approximating $T(n-2) \sim T(n-1)$ here}	
	fib(n) = 1 if n = 1	Time taken would be $O(2^n)$ by using Master's theorem	
	fib(n) = fib(n-2) + fib(n-1) if n > 1	Why can't we reduce the function calls to reduce the time taken?	
		For this, we would take a global array and initially fill it with -1	
	int fib(n) {	$F = \{-1, -1, -1, -1, -1\}$	
	if(n <= 1){	Then, as the function calls f(1), mark it as 1	
	return n;}	Then f(0) is marked as 1	
	return fib(n-2) + fib(n-1);	Then use the stored result to get the rest.	
	}	Finally, F would get updated as we solve: $F = \{0, 1, 1, 2, 3, 5\}$	
		Total 6 calls are made then i.e. n+1 calls i.e. $O(n)$	
		This is called result of memorization	
	From the above example, we can see reduction in number of calls from $O(2^n)$ to $O(n)$ using memorization. It follows top down approach. The same problem can be solved using tabular method (iterative process) as shown below:		
	int fib(int n) {		
	if(n <= 1) {		
	return n;}		
	F[0] = 0; F[1] = 1;		

	for(int i = 2; i <= n; i++){		
	F[i] = F[i-2] + F[i - 1];		
	}		
	return F[n];		
	}		
	F= {0, 1, 1, 2, 3, 5}		
	This is a bottom - up approach i.e. starting from F[0] and moving to F[n]		
56	Multistage Graph		
	<i>A multistage graph is a directed weighted graph. The vertices are divided into stages such that the edges are connecting vertices from one stage to next stage only. First stage and last stage will have only one vertex to represent start and end point. This is usually used to represent resource allocation.</i>		
	The objective of the problem is that I have to select a path which gives me minimum cost.	//it is a minimization or optimization problem	
	Dynamic programming works on principle of optimality. Principle of optimality says that a problem can be solved in a sequence of decisions.		
	From first stage I have to select one optimal vertex that leads to minimum cost and I have to take this decision at each stage. Thus, I can apply dynamic programming here.		
V	1 2 3 4 5 6 7 8 9 10 11 12		
Cost	16 7 9 18 15 7 5 7 4 2 5 0	cost(5, 12) = 0 ; here 5 is the stage and 12 is the vertex	
d	2/3 7 6 8 8 10 10 10 12 12 12 12	cost(4, 9) = 4	
		cost(4, 10) = 2	
	Formula for multistage graph:	cost(4, 11) = 5	
	cost(lth stage, jth vertex no.) = cost(i, j) = min{C(j, l) + cost(i+1, l)}	cost(3, 6) = min{ C(6, 9) + cost(4, 9) , C(6, 10) + cost(4, 10)} = min{6 + 4, 5 + 2} = 7	
		Similarly, cost(3, 7) = min{8, 5} = 5	
	Now, we will solve it by going in forward direction and taking decisions based on above data;	cost(3, 8) = min{7, 11} = 7	

	d(1,1) = 2	Similarly, cost(2, 2) = min{C(2, 6) + cost(3, 6), C(2, 7) + cost(3, 7), C(2, 8) + cost(3, 8)} = min{11, 7, 8} = 7	
	d(2,2) = 7	cost(2, 3) = min{9, 12} = 9	
	d(3, 7) = 10	cost(2, 4) = min{18} = 18	
	d(4, 10) = 12	cost(2, 5) = min{16, 15} = 15	
	Path: 2---> 7---> 10----> 12	cost(1,1) = min{16, 16, 21, 17} = 16	
	d(1, 1) = 3		
	d(2, 3) = 6		
	d(3, 6) = 10		
	d(4, 10) = 12		
	Path: 3---> 6---> 10----> 12		
	So, we have two paths with same cost.		
57	Multistage Graph (Program)		
	Cost adjacency Matrix	main(){	
	0 1 2 3 4 5 6 7 8	int stages = 4, min;	
0	0 0 0 0 0 0 0 0 0	int n = 8;	
1	0 0 2 1 3 0 0 0 0	int cost[9], d[9], path[9];	
2	0 0 0 0 0 2 3 0 0	int c[9][9] = {{0,0,0,0,0,0,0,0,0}, {0,0,2,1,3,0,0,0,0}, {0,0,0,0,0,2,3,0,0},}	
3	0 0 0 0 0 6 7 0 0	cost[n] = 0;	
4	0 0 0 0 0 6 8 9 0	for(int i = n-1; i >=1; i--){	
5	0 0 0 0 0 0 0 0 6	min = 32767;	
6	0 0 0 0 0 0 0 0 4	for(int k = i + 1; k <= n; k++){	
7	0 0 0 0 0 0 0 0 5	if(C[i][k] != 0 && C[i][k] + C[k] < min){	
8	0 0 0 0 0 0 0 0 0	min = C[i][k] + C[k] ;	
		d[i] = k;	
	0 1 2 3 4 5 6 7 8		
cost	-- 9 7 11 12 6 4 5 0	}	
d	-- 2 6 6 5 8 8 8 --	}	

path	-- 1 2 6 8	cost[i] = min;	
		}	
	Path is calculated using the following formula: $p[i] \ d[p[i-1]]$; $p[2] = d[p[2-1]] = d[1] = 2$	$p[1] = 1$; $p[\text{stages}] = n$;	
	time complexity: $O(n^2)$	for($i = 2$; $i < \text{stages}$; $i++$){	
		$p[i] = p[d[i-1]]$;	
	58 All Pairs Shortest Path		
	If I use Dijkstra algorithm on each vertex to find the shortest path of all, the time complexity would be $O(n^3)$.		
A0 =	1 2 3 4		
	1 0 3 INF 7		
	2 8 0 2 INF		
	3 5 INF 0 1		
	4 2 INF INF 0		
	Considering vertex 1 as intermediate vertex	$A0[2,3] \quad A0[2,1] + A0[1,3]$	
A1 =	1 2 3 4	2 < 8 + INF	
	1 0 3 INF 7		
	2 8 0 2 15	$A0[2,4] \quad A0[2,1] + A0[1,4]$	
	3 5 8 0 1	INF > 8 + 7 = 15	
	4 2 8 INF 0		
		$A0[3,2] \quad A0[3,1] + A0[1,2]$	
	Considering vertex 2 as intermediate vertex	INF > 5 + 3 = 8	
A2 =	1 2 3 4		
	1 0 3 5 7	$A[1,3] \quad A[1,2] + A[2,3]$	
	2 8 0 2 15	INF < 3 + 2 = 5	
	3 5 8 0 1		
	4 2 5 7 0	for($k = 1$; $k \leq n$; $k++$){	
		for($i = 1$; $i \leq n$; $i++$){	
	Considering vertex 3 as intermediate matrix	for($j = 1$; $j \leq n$; $j++$){	

A3 =	1 2 3 4	$A[i,j] = \min(A[i,j], A[i,k] + A[k,j])$	
1	0 3 5 6	}	
2	7 0 2 3	}	
3	5 8 0 1	}	
4	2 5 7 0		
	Considering vertex 4 as intermediate matrix		
A4 =	1 2 3 4		
1	0 3 5 6		
2	5 0 2 3		
3	3 6 0 1		
4	2 5 7 0		
	Formula used to make the above matrices:		
	$A[i,j] = \min\{A[i,j], A[i,k] + A[k,j]\}$		
	time complexity = $O(n^3)$		
59	Matrix chain multiplication		
	$A1 \cdot A2 \cdot A3 \cdot A4$		
	(5X4) (4X6) (6X2) (2X7)		
	For generating a single matrix C after single multiplication of 2 matrices of order (5X4) and (5X3) , we would need to do 60 multiplications		
	$((A1.A2).A3).A4$		
	or $(A1.A2).(A3.A4)$		
	orthere could be several ways		
	then, how to choose theright way?		
	$T(n) = 2nCn/n+1$ trees are possible		

	thus, with 3 nodes ; $T(3) = 5$					
	Using tabular approach (bottom up approach),					
m	1	2	3	4	in $m[1,1]$ i.e. A_1 , nothing is multiplied, hence it can be taken as zero	
	1	0	120	88	158	$m[1,2] = A_1.A_2$
	2	-	0	48	104	$(5 \times 4) \quad (4 \times 6)$
	3	-	-	0	84	Total cost of multiplication = $5 * 4 * 6 = 120$
	4	-	-	-	0	
					$m[1,3] = A_1.A_2.A_3$	
s	1	2	3	4	Two possibilities: $A_1.(A_2.A_3)$ or $(A_1.A_2).A_3$	
	1	-	1	1	3	$(5 \times 4) \quad (4 \times 6) \quad (6 \times 2)$
	2	-	-	2	3	for $A_1.(A_2.A_3) \rightarrow m[1,1] + m[2,3] + (5*4*2)$
	3	-	-	-	3	for $(A_1.A_2).A_3 \rightarrow m[1,2] + m[3,3] + (5*6*2)$
	4	-	-	-	-	$0 + 48 + 40$ $120 + 0 + 60$
					88	180
					Similarly, $m[2,4]$	
	formula:				$A_2.(A_3.A_4) \quad (A_2.A_3).A_4$	
	$m[i,j] = m[i,k] + m[k+1, j] + d_{i-1} * d_k * d_j$				$(4 \times 6) \quad (6 \times 2)(2 \times 7) \quad (4 \times 6) \quad (6 \times 2)(2 \times 7)$	
	we are generating $n(n-1)/2$ elements				for $A_2.(A_3.A_4) \rightarrow m[2,2] + m[3,4] + (4*6*7)$	for $(A_2.A_3).A_4 \rightarrow m[2,3] + m[4,4] + (4*2*7)$
	time complexity = $O(n^3)$				$0 + 84 + 168$	$48 + 0 + 56$
					252	104
					$m[1,4]$	
					$\min\{m[1,1] + m[2,4] + (5*4*7), m[1,2] + m[3,4] + (5*6*7), m[1,3] + m[4,4] + (5*2*7)\}$	
					$\min\{0+104+140, 120 + 84+210, 88+70\}$	
					$\min\{244, 414, 158\}$	
60	Matrix chain multiplication - A few pointers					
	Condition of the multiplication: The number of columns in the first matrix involved in the multiplication must be equal to the number of rows in the second matrix					
A=	a11	a12	a13		2X3 dimension	

	a21 a22 a23		
B=	b11 b12	3X2 dimension	
	b21 b22		
	b31 b32		
A*B =	a11*b11 + a12*b21 + a13*b31 a11*b12 + a12*b22 + a13*b32	12 multiplications (2*3*2)	
	a21*b11 + a22*b21 + a23*b31 a21*b12 + a22*b22 + a31*b32	2X2 dimensions of the resultant matrix	
	A1 X A2 X A3 {Multiplication of more than two matrices}		
	2X3 3X4 4X2		
	d0 d1. d1 d2 d2 d3		
	Same answer by the two following methods (Associative property)		
	Method 1: (A1 X A2) X A3	Method 2: A1 X (A2 X A3)	
	2X3 3X4 i.e (2*3*4) = 24 multiplications for A1 X A2	2X3 3X4 4X2	
	Now, (A1 X A2) X A3 would require (2*4*2) = 16 multiplications	A2 X A3 requires (3*4*2) = 24 multiplications	
	Thus, altogether 40 multiplications are required	A1 X (A2 X A3) requires (2*3*2) = 12 multiplications	
		Altogether, 36 multiplications are needed here.	
	Now, dynamic programming asks us to find all the possible methods for matrix multiplication and check which one costs the minimum --> this implies that for 10 matrices, there would be numerous methods and we would have to check for all before proceeding with any one of them. Thus, we need a formula to check all that		
	We need to find C[1,3]		
	Method 1: (A1 X A2) X A3	Method 2: A1 X (A2 X A3)	
	C[1,2] = 24; C[3,3] = 0	C[1,1] = 0; C[2,3] = 24	
	C[1,2] + C[3,3] + d0*d2*d3 = 40	C[1,1] + C[2,3] + d0*d1*d3 = 36	
	C[i,j] = C[i, k] + C[k+1, j] + di-1 * dk * dj		
	After generalization,		

	$C[i,j] = \min \{ C[i,k] + C[k+1,j] + d_{i-1} * d_k * d_j \}$		
	where $i \leq k \leq j$		
	$A_1 \times A_2 \times A_3 \times A_4$		
	$d_0 \ d_1 \ d_1 \ d_2 \ d_2 \ d_3 \ d_3 \ d_4$		
	Check which method works the best for the above matrix chain multiplication	$2n \ C \ n / n + 1$ multiplications are possible where n = no. of matrices - 1	
	1. $A_1 (A_2 (A_3 A_4))$	Modified formula: $2(n-1) \ C \ (n-1) / n$	
	2. $A_1 ((A_2 A_3) A_4)$	Now, for $n = 4$, $2*3C_3 / 4 = 6*5*4/3*2*1 / 4 = 5$	
	3. $(A_1 A_2)(A_3 A_4)$	$n = 5$, 14 multiplications	
	4. $(A_1 (A_2 A_3)) A_4$		
	5. $((A_1 A_2) A_3) A_4$		
	Applying the formula:		
4-1 = 3 values	$C[1,4] = \min \{ k = 1; C[1,1] + C[2,4] + d_0 * d_1 * d_4,$		
	$k = 2; C[1,2] + C[3,4] + d_0 * d_2 * d_4,$		
	$k = 3; C[1,3] + C[4,4] + d_0 * d_3 * d_4 \}$		
	$1 \leq k \leq 4$		
	1. $A_1 (A_2 A_3 A_4)$		
	2. $(A_1 A_2) (A_3 A_4)$		
	3. $(A_1 A_2 A_3) A_4$		
	here, $C[1,1] = 0$; $C[4,4] = 0$		
4-2 = 2 values	$C[2,4] = \min \{ k = 2; C[2,2] + C[3,4] + d_1 * d_2 * d_4$		
	$k = 3; C[2,3] + C[4,4] + d_1 * d_3 * d_4 \}$		
	$2 \leq k \leq 4$		

4-3 = 1 value	$C[3,4] = C[3,3] + C[4,4] + d_2 \cdot d_3 \cdot d_4$					
	A1 X A2 X A3 X A4					
	3X2 2X4 4X2 2X5					
	d0 d1 d1 d2 d2 d3 d3 d4					
	As we notice that there is a repetition of the values needed , such as C[3,4] or C[4,4], there is unnecessary calculation, we are repeating, thus we should use a table (4X4)					
C table	1	2	3	4	$C[1,2] = \min\{k=1; C[1,1] + C[2,2] + d_0 \cdot d_1 \cdot d_2\}$	
1	0	24	28	58	Thus, $C[1,2] = 3 \cdot 2 \cdot 4 = 24$	
2	-	0	16	36	$C[2,3] = \min\{k = 2; C[2,2] + C[3,3] + d_1 \cdot d_2 \cdot d_3\}$	
3	-	-	0	40	Thus, $C[2,3] = 2 \cdot 4 \cdot 2 = 16$	
4	-	-	-	0	$C[3,4] = d_2 \cdot d_3 \cdot d_4 = 4 \cdot 2 \cdot 5 = 40$	
					$C[1,3] = \min\{k=1; C[1,1] + C[2,3] + d_0 \cdot d_1 \cdot d_3\}$	
k table	1	2	3	4	$k = 2; C[1,2] + C[3,3] + d_0 \cdot d_2 \cdot d_3\}$	
1	-	1	1	3	$C[1,3] = \min\{16 + 3 \cdot 2 \cdot 2, 24 + 3 \cdot 4 \cdot 2\}$	
2	-	-	2	3	$C[1,3] = \min\{28, 48\} = 28$	
3	-	-	-	3		
4	-	-	-	-	$C[2,4] = \min\{C[2,2] + C[3,4] + d_1 \cdot d_2 \cdot d_4, C[2,3] + C[4,4] + d_1 \cdot d_3 \cdot d_4\}$	
					$C[2,4] = \min\{40 + 2 \cdot 4 \cdot 5, 16 + 2 \cdot 2 \cdot 5\}$	
	The result shows that we need to do minimum of 58 multiplications to get the result of A1 X A2 X A3 X A4				$C[2,4] = \min\{80, 36\} = 36$	
	The k table will give the paranthesization					
	((A1) (A2 A3))(A4)					
	How much time it has taken ? $1+2+3+4 = 4(5) / 2$ i.e. $n(n+1)/2 \sim n^2$				$C[1,4] = \min\{k = 1; C[1,1] + C[2, 4] + d_0 \cdot d_1 \cdot d_4,$	
	we also tried n posisble values of k to compute this value, thus time taken is $n^2 \cdot n$ i.e. $O(n^3)$				$k = 2; C[1,2] + C[3,4] + d_0 \cdot d_2 \cdot d_4,$	
					$k=3; C[1,3] + C[4,4] + d_0 \cdot d_3 \cdot d_4\}$	
					$C[1,4] = \min\{36 + 3 \cdot 2 \cdot 5, 24 + 40 + 3 \cdot 4 \cdot 5, 28 + 3 \cdot 2 \cdot 5\}$	

		$C[1,4] = \min\{66, 124, 58\} = 58$	
61	Matrix chain multiplication Program		
	A1 X A2 X A3 X A4	main{	
	5X4 4X6 6X2 2X7	int n = 5;	
		int P[] = {5, 4, 6, 2, 7};	
	P = { 5,4,6,2,7}	int m[5][5] = {0};	
		int s[5][5] = {0};	
		int j, min, q;	
		for(int d = 1; d < n - 1; d++)	
		{	
		for(int i = 1; i < n - d; i++)	
		{	
		j = i + d;	
		min = 32767;	
		for(int k = 1; k <= j - 1; k++)	
		{	
		q = m[i][k] + m[k + 1][j] + P[i	
		- 1] * P[k] * P[j];	
		if(q < min)	
		{	
		min = q;	
		s[i][j] = k;	
		}	
		}	
		m[i][j] = min;	
		}	
		}	
		count << m[1][n - 1];	
		}	

62	Single source shortest path (Bellman Ford Algorithm)	Example of Bellman Ford algorithm:	
	For doing this, we already have Dijkstra algorithm but it may not work correctly if we have negative weights, thus we need some other method that works with negative weights	edges --> (3,2)(4,3)(1,4)(1,2)	
	Bellman Ford algorithm says that we should relax the edges N-1 times where the number of vertices is equal to N	mark source vertex 1 as 0 and rest all as infinity	
	$ V = N = 7$	there are 4 vertices, so we should relax all the edges for 3 times	
	So, we should relax them for $ V - 1$ times	First iteration:	
	so, we would cover all possible paths even the longest path	for (3,2) --> infinity - 10 is infinity only, so no change	
	Relaxation means between a pair of vertices u and v if there is an edge, then check if:	for (4,3) --> infinity + 3 is infinity only, so no change	
	if($d[u] + C(u,v) < d[v]$) {	for (1,4), $0 + 5 < \text{infinity}$, thus vertex 4 is updated to 5	
	$d[v] = d[u] + C(u,v)$ }	for (1,2) $0 + 4 < \text{infinity}$, thus vertex 2 is updated to 4	
	edgeList --> (1,2)(1,3)(1,4)(2,5)(3,2)(3,5)(4,3)(4,6)(5,7)(6,7)	Second iteration:	
	Now, I have to relax these edges for $ V - 1$ i.e. 6 times	for (3,2), vertex 2 is already 4, which is less than $d[u] + C(u,v)$ in this case	
	Initially, mark the distance for source vertex as 0 and for the rest of the vertices as infinity	for (4,3) vertex 3 is updated to $5+3 = 8$	
	Now, let's relax edge (1,2)	for (1,4)--> no change	
	here, $d[u] = 0$; $d[v] = \text{infinity}$; $C(u,v) = 6$	for(1,2) --> no change	
	$0 + 6 < \text{infinity}$; thus $d[v] = 6$	Third iteration	
	for vertex 2, distance is 6;	for(3,2), $8 - 10 = -2 < d[v]$ which is 4 right now; thus updated for vertex 2 as -2	
	similarly, relaxing (1,3); thus its distance is updated to 5 from infinity	for the rest of the edges there won't be any change	
	in similar way, (1,4) is relaxed, following that the distance of vertex 4 is updated to 5 from infinity	results obtained :	
	Now, relaxing (2,5); $d[u] = 6$; $C(u,v) = -1$; $d[v] = \text{infinity}$	vertex 1 --> 0	
	the distance (2,5) is updated to $6 - 1 = 5$	vertex 2 --> -2	
	similarly, relaxing (3,2); $d[u] = 5$, $C(u,v) = -2$, $d[v] = 6$	vertex 3 --> 8	

5 - 2 = 3 < 6 thus, d[v] is updated here to 3 i.e. at vertex 2 , distance is updated to 3	vertex 4 --> 5	
Now, relaxing (3,5), d[u] = 5 , C(u,v) = 1, d[v] = 5 here d[u] + C(u,v) is not smaller than d[v] hence the distance of vertex 5 is not modified	Now, if I relax one more time extra, there's no change	
Moving to (4,3), relaxing it --> d[u] = 5, C(u,v) = -2; d[v] = 5, d[u] + C(u,v) = 3 < d[v] hence distance of vertex 3 is updated to 3	Drawback of Bellman Ford algorithm:	
following this (4,6) is relaxed again and checked, d[u] = 5, C(u,v) = -1; thus distance of vertex 6 is updated to 5-1 = 4	let us an edge(2,4) in the above example	
Now, relaxing (5,7), d[u] = 5; C(u,v) = 3; d[v] = infinity	we see even after N-1 iterations, there is one vertex changing, we note that there's a problem	
thus, d[v] is updated to 8 for edge (5,7)	the reason is that there is a cycle of edges where total weight of edges is negative i.e. 5 + 3+ (-10) = -2 , thus graph cannot be solved	
moving to (6,7), d[u] = 4, C(u,v) = 3; d[v] = 8	hence, for a negative weighted cycle , the bellman ford algorithm fails	
d[v] is updated to 7 here for edge (6,7)		
let us continue second time;		
there won't be any change in (1,2), (1,3), (1,4)		
when we check for (2,5); d[u] = 3 C(u,v) = -1 d[v] = 5; d[v] for edge (2,5) is updated to 2		
similarly, for edge (3,2), the value is change to 3 - 2 = 1; earlier it was 3		
(4,3) and (4,6), there's no change		
for (5,7) d[u] has changed from 5 to 2; thus d[v] changes to 2 + 3 = 5		
for (6,7) there won't be any change		
let us continue third time;		
there won't be any change in (1,2), (1,3), (1,4),		
when we check for (2,5); d[u] = 1 C(u,v) = -1 d[v] = 2; d[v] for edge (2,5) is updated to 0		
for (3,2) there own't be any change		
for (3,5) again thee won't be any change		

	for (4,3), (4,6) --> no change		
	for (5,7) d[v] gets updated to $0 + 3 = 3$		
	for (6,7) --> no change		
	let us check for fourth time,		
	we notice for all edges --> no change		
	results obtained:		
	vertex 1 --> 0		
	vertex 2 --> 1		
	vertex 3 --> 3		
	vertex 4 --> 5		
	vertex 5 --> 0		
	vertex 6 --> 4		
	vertex 7 --> 3		
	so, finally these are the shortest paths		
	time complexity: $O(E (V - 1)) \sim O(V E) \sim O(N^2)$		
	If it is a complete graph, that is between every two vertex there is an edge, then number of edges is $N(N - 1) / 2$		
	i.e. $ E = N(N - 1) / 2$		
	then time complexity = $O(E V) = O(N((N - 1) / 2)(N - 1)) \sim O(N^3)$		
	63 0/1 Knapsack Problem		
	$m = 8$; $n = 4$ where m is the capacity of the bag		
	$P = \{1, 2, 5, 6\}$		
	$w = \{2, 3, 4, 5\}$		
	Have to give a solution of what all should be included to maximize profit		
	x can be either 0 or 1; we cannot take fraction of an object		
	x would be like $\{1, 0, 1, \dots\}$		
	sum of profits is maximized i.e. $\max\{\sum(P_i x_i)\}$		

now, check for object 3, if there is a value 2; check if the value 2 is there in for object 2 as well at the same place, if yes, then object 3 was not taken i.e. $x_3 = 0$; if the same 2 is there for object 1 then $x_2 = 0$ else $x_2 = 1$		
here, $x_3 = 0$		
$x_2 = 1$		
$x_1 = 0$ as the remaining profit is 0		
$x = \{0 \ 1 \ 0 \ 1\}$ for maximizing the profit		
let us use sets method		
we will prepare sets of P and w		
$s_0 = \{(0,0)\}$		
$s_0(1) = \{(1,2)\}$ here, we added first object to elements of set s_0		
$s_1 = \{(0,0), (1,2)\}$ merged the above two sets to get s_1		
$s_1(1) = \{(2,3), (3,5)\}$, added second object to elements of s_1		
$s_2 = \{(0,0), (1,2), (2,3), (3,5)\}$ merged the above two sets to get s_2		
$s_2(1) = \{(5,4), (6,6), (7,7), (8,9)\}$		
$s_3 = \{(0,0), (1,2), (2,3), (3,5), (5,4), (6,6), (7,7)\}$ removed (8,9) as it exceeded the permitted limit		
in the above set, we notice that as profit increases , weight increases, but at (5,4) profit has increased from (3,5) whereas weight has decreased		
thus, we discard (3,5) with lesser profit {dominance rule}		
so, $s_3 = \{(0,0), (1,2), (2,3), (5,4), (6,6), (7,7)\}$		
now, considering the fourth object,		
$s_3(1) = \{(6,5), (7,7), (8,8), (11,9), (12,11), (13,12)\}$		
$s_4 = \{(0,0), (1,2), (2,3), (5,4), (6,6), (7,7), (8,8)\}$ merged above two sets and removed elements using dominance rule and those exceeding the weight limits		
time taken is almost (2^n)		
maximum order is (8,8)		
(8,8) belongs to s_4		

	check whether it belongs to s3 ---> it does not, hence object 4 is included		
	now $(8,8) - (6,5) = (2,3)$		
	$(2,3)$ --> check if it belongs to s3 --> yes, check whether it belongs to s2 -> yes, thus object 3 is not included --> now, check if it belongs to s1 --> no, that means object 2 is included and object 1 is not		
	$2-2, 3-3 = (0,0)$		
	$(0,0)$ belongs to set 1 and set 0 as well therefore object 1 is not included		
	$x = \{0 \ 1 \ 0 \ 1\}$		