

Control over CAN and Flexray  
**Embedded Control Systems**

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# Contents

<b>1</b>	<b>Part 1</b>	<b>2</b>
1.1	Introduction . . . . .	2
1.2	Response Time analysis . . . . .	2
1.2.1	Fixed-Priority Non-Preemptive . . . . .	2
1.2.2	Fixed-Priority Preemptive . . . . .	3
1.2.3	Response time analysis per processing unit . . . . .	3
1.2.4	Response time analysis for the CAN bus messages . . . . .	3
1.3	System Model Derivation and Control Parameter Design . . . . .	3
1.3.1	Controller Structure . . . . .	4
1.3.2	Controller Design . . . . .	5
1.4	Design decision . . . . .	6
1.5	Results . . . . .	6
<b>2</b>	<b>Part 2</b>	<b>7</b>
2.1	Introduction . . . . .	7
2.2	Response Time analysis . . . . .	7
2.2.1	Response time analysis per processing unit . . . . .	7
2.2.2	Response time analysis for the CAN bus messages . . . . .	9
2.3	Optimization for sensor-to-actuator delay . . . . .	9
2.4	Design decision . . . . .	11
2.5	System model and Results . . . . .	11
2.6	Conclusions . . . . .	11
<b>3</b>	<b>Part 3</b>	<b>13</b>
3.1	Introduction . . . . .	13
3.2	Answer all the questions . . . . .	13
3.2.1	Theoretical analysis versus actual implementation . . . . .	14
3.3	Design decision . . . . .	14
3.4	Results . . . . .	14
3.5	Conclusions . . . . .	14
3.6	Results . . . . .	14
3.7	Conclusion . . . . .	14

# Chapter 1

## Part 1

### 1.1 Introduction

### 1.2 Response Time analysis

#### 1.2.1 Fixed-Priority Non-Preemptive

CAN messages follow fixed-priority non-preemptive protocol. The response time analysis for fixed-priority non-preemptive messages can be computed as follows.

##### **BlockingTime**

Blocking time of a message is defined as the time period for which the message can be blocked by a lower priority message. Therefore blocking time can be computed by finding the maximum of execution of times of all the messages having lower priority than the current message. It can be computed as shown in equation 1.1.

$$B_{mi} = \max_{k \in lp(m_i)} (c_k) \quad (1.1)$$

##### **BusyPeriod**

Busy Period is defined as the maximum time at which the execution of the message gets completed. It includes the blocking time, delay due to execution of higher priority messages and the message execution time. Busy period of message  $m_i$  with priority  $i$  can be computed from the equation 1.4 with initialization condition shown in equation 1.2 and termination condition as shown in equation 1.3.

$$t_{mi}^{k+1} = B_{mi} + \sum_{\forall m \in hp(m_i) \cup m_i} \left\lceil \frac{t_{mi}^k}{p_m} \right\rceil c_m \quad (1.2)$$

$$t_{mi}^0 = c_{mi} \quad (1.3)$$

$$t_{mi}^{k+1} = t_{mi}^k \quad (1.4)$$

##### **ResponseTime**

If the busy period of the message is less than the period of the message, then the response time of the message is equal to busy period. But if busy period is greater than the period of the message, multiple messages arrive in this instance and response time is computed as the maximum of response times of each individual instances. The number of instances can be computed from equation 1.5. The waiting time and response time of instance  $q$  can be computed using equations 1.6 and equation 1.7 respectively. The initialization

and termination conditions for each instance are governed by equations 1.8 and 1.9 respectively. The total response time can be computed using equation 1.10.

$$Q_{mi} = \lceil \frac{t_{mi}}{p_{mi}} \rceil \quad (1.5)$$

$$w_{mi}^{k+1}(q) = B_{mi} + qc_{mi} + \sum_{\forall m \in hp(m_i)} \lceil \frac{w_{mi}^k(q) + \tau_{bit}}{p_m} c_m \rceil \quad (1.6)$$

$$R_{mi}(q) = w_{mi}(q) - qp_{mi} + c_{mi} \quad (1.7)$$

Initialization and Termination:

$$w_{mi}^0(q) = B_{mi} + qc_{mi} \quad (1.8)$$

$$w_{mi}^{k+1}(q) = w_{mi}^k(q) \quad (1.9)$$

Total Response time of the message:

$$R_{mi} = \max_{q=0 to Q_{mi}} (R_{mi}(q)) \quad (1.10)$$

## 1.2.2 Fixed-Priority Preemptive

Unlike the above case messages or tasks can be preempted in Fixed-Priority preemptive case. The processors PU1 and PU2 employ FPP for their task management. The response time of FPP can be calculated as below.

### ResponseTime

The waiting time and response time of each instance can be calculated from equations 1.11 and 1.12. The initialization and termination conditions for each task can be obtained from equations 1.13 and 1.14. The response time is nothing but the settling value of the response time.

$$w_{mi}^{k+1}(q) = w_{mi}^k(q) + \sum_{\forall m \in hp(m_i)} \lceil \frac{R_{mi}^k(q)}{p_m} c_m \rceil \quad (1.11)$$

$$R_{mi}(q) = w_{mi}(q) + c_{mi} \quad (1.12)$$

Initialization and Termination:

$$w_{mi}^0(q) = 0; \quad (1.13)$$

$$R_{mi}^{k+1}(q) = R_{mi}^k(q) \quad (1.14)$$

### ResponseTime

## 1.2.3 Response time analysis per processing unit

## 1.2.4 Response time analysis for the CAN bus messages

# 1.3 System Model Derivation and Control Parameter Design

The objectives for the controller is to properly control the given system with a set of design parameters. These parameters, shown in Table

There are namely a number of steps taken until all parameters can be considered to satisfy the performance constraints of the controller. These basic steps can be seen in the following list.

Table 1.1: By running the Matlab script ResponsetimeAnylysis\_FPP.m with the different parameters given for PU1 and PU2 these response times are obtained for each of the tasks. These files are then delivered as PU1.m PU2.m

PU1	$T_1$	$T_2$	$T_3$	$T_4 (T_s)$
Matlab (ms)	0.1	2.1	4.1	7.2

PU2	$T_5$	$T_6$	$T_7$	$T_8$
Matlab (ms)	6	3	9	5

Table 1.2: Response times for the CAN bus messages

CAN	$m_2$	$m_1$	$m_3$	$m_8$
Matlab (ms)	2	3	4	4

1. Step 1: Derive the system model. Determine the A, B and C matrices in relation to all constants by hand calculations. We are already provided with the matrices in assignment description, hence no need to derive. Insert this into Matlab.
2. Step 2: Derive the values for sampling period and sensor-to-actuator delay based on our system design.
3. Step 3: Choosing desired poles according to the given requirements. Verifying the controllability of the system for the choosen values. Computing the controller Feed-Forward and Feed-Back gains  $F$  and  $K$ . Inserting these calculations into MATLAB. More about this step can be found in Section
4. Step4: Design K and F values, compute current input activation(u) and outputs(x) from equations, insert these calculations to MATLAB and simulate the system.
5. Step5: Apply multi-objective genetic algorithm on the above system, with pole positions as parameters, and settling-time and maximum input voltage as our objectives. Obtain pareto optimal points satisfying design constraints.

Table 1.3: Constants and design parameters referenced to in calculations

Symbol	Description	Value	Unit
$\theta$	Rotor position	-	rad
$i_m$	Motor current	-	A
J	Inertia	$3.2284 \cdot 10^{-6}$	$Kgm^2$
b	Friction coefficient	$3.5077 \cdot 10^{-6}$	Nms/rad
R	Armature Resistance	4	Ohm
L	Inductance of Motor Winding	$2.75 \cdot 10^{-6}$	H
$K_m$	Motor constant	0.0274	Nm/A
K	Feed-BackGain	-	-
F	Feed-Forward Gain	-	-

### 1.3.1 Controller Structure

The system can be described with the second order differential equation shown in Equation 1.15 as given in the assignment description. Equation 1.16 and Equation 1.17 describe the statespace representation of our system. From these equations the Feedback- and Feedforward gain can be determined in relations to the constants in the system.

$$\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} \theta \\ \dot{\theta} \\ i \end{bmatrix} = \begin{bmatrix} \theta \\ \omega \\ i \end{bmatrix} \quad \text{and} \quad \dot{\mathbf{x}} = \begin{bmatrix} \dot{\theta} \\ \ddot{\theta} \\ \dot{i} \end{bmatrix} \quad (1.15)$$

$$\dot{\mathbf{x}} = \mathbf{A}\mathbf{x} + \mathbf{B}u \quad \text{and} \quad y = \mathbf{C}\mathbf{x} \quad \text{where} \quad \text{input: } u = V, \quad \text{output: } y = i \quad (1.16)$$

$$\dot{\mathbf{x}} = \begin{bmatrix} 0 & 1 & 0 \\ 0 & \frac{-b}{J} & \frac{K_m}{J} \\ 0 & \frac{-K_m}{L} & \frac{-R}{L} \end{bmatrix} \begin{bmatrix} \theta \\ \omega \\ i \end{bmatrix} + \begin{bmatrix} 0 \\ 0 \\ \frac{1}{L} \end{bmatrix} V \quad \text{and} \quad y = \begin{bmatrix} 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} \theta \\ \omega \\ i \end{bmatrix} \quad (1.17)$$

### 1.3.2 Controller Design

In this subsection we discuss the method for designing the controller design parameters FeedbackGain- $K$  and FeedForwardGain- $F$ . In this problem we have the scenario where  $D_c < h$ . Since our controller operates in discrete sample time, we start with converting our continuous system as described in equation 1.18 into discrete domain. Therefore on applying ZOH sampling with period  $h_c$  and constant sensor-actuator Delay  $D_c$  we achieve the equation 1.19 for discrete sample time system. From 1.19 we can notice that the next output not only depends on current input but also on previous input. Hence, we simplify the system and get it into standard form representing equation 1.20 by invoking equation 1.21. After applying above simplification input  $u[k]$  can be represented in terms of controller gains using equation 1.22. and matrices  $\Phi, C$  are converted to corresponding augmented matrices  $\Phi_{aug}, C_{aug}$ , whereas  $\Gamma_1, \Gamma_2$  are converted to single augmented matrix  $\Gamma_{aug}$ .

$$\dot{\mathbf{x}} = \mathbf{A}\mathbf{x} + \mathbf{B}u \quad \text{and} \quad y = \mathbf{C}\mathbf{x} \quad (1.18)$$

$$x[k+1] = \phi x[k] + \Gamma_1(D_c)u[k-1] + \Gamma_0(D_c)u[k] \quad \text{and} \quad y[k] = Cx[k] \quad (1.19)$$

$$x[k+1] = \phi z[k] + \Gamma_{aug}u[k] \quad \text{and} \quad y[k] = C_{aug}z[k] \quad (1.20)$$

$$z[k] = \begin{bmatrix} x[k] \\ u[k-1] \end{bmatrix} \quad (1.21)$$

$$u[k] = Kz[k] + Fr \quad (1.22)$$

$$K = -[0 \ 0 \ \dots \ 1] \gamma_{aug}^{-1} H(\phi_{aug}) \quad (1.23)$$

$$F = \frac{1}{C_{aug}(I - \phi_{aug} - \Gamma_{aug}K)^{-1} \Gamma_{aug}} \quad (1.24)$$

Before deriving the controller gains it is important to verify if the system is controllable under given configuration. This can be verified by calculating  $\det(\gamma_{aug})$ . If the determinant is not equal to 0, then system is controllable. After verifying the controllability of the system  $K$  can be derived by applying Ackermanns formula described in equation 1.23 and  $F$  can be derived from equation 1.24. However, one can notice the matrix  $H(\Phi_{aug})$  depends on the desired poles which in turn depends on the design requirements. The desired poles alter the frequency spectrum of the transfer function of the system and play a significant role in controlling the behaviour of output parameters of the system. Therefore various design requirements such as Overshoot, settling time, boundary ranges of the parameters in the system depends on the desired poles and thus in turn influence controller gains  $K$  and  $F$ . Following design, work flow has been developed in MATLAB to derive pareto optimal points (pole locations) satisfying design constraints using multi-objective genetic algorithm and one of the optimal point is selected for simulation.

## **1.4 Design decision**

## **1.5 Results**

Firstly: Response time analysis

Secondly: Control system input and output

# Chapter 2

## Part 2

### 2.1 Introduction

Analyzing and confirming broad analysis done on paper and in Matlab allows us to confirm the hypothesis made about the behavior of the CAN bus and its specific tasks from Part 1. This analysis is done in Inchron, which allows us to explore the real-time behavior of this embedded system in full detail. Some exploration on how the system should work has already been done in Part 1, [Add Table reference to periods and priorities](#), this is used here by feeding the settings of the embedded system to the tool. The settings included is the hierarchy of the Processing Units (PUs) and their tasks which each have different periods, execution time and priority. The tree view, Figure 2.1, shows the details of the hierarchy from the Inchron's perspective.

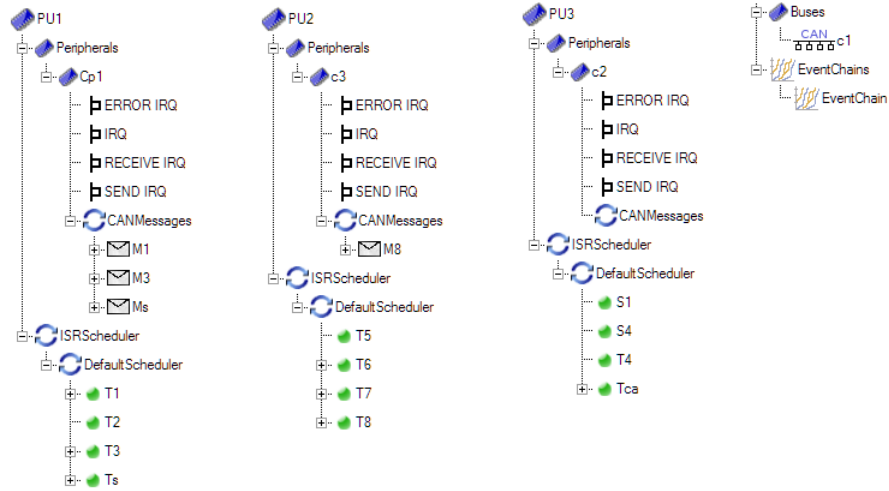


Figure 2.1: This shows the coarse grain hierarchy of the system ported into Inchron for verification and simulation of the real time components of our embedded system.

### 2.2 Response Time analysis

#### 2.2.1 Response time analysis per processing unit

For this analysis all details have to be imported to Inchron as stated earlier. Now paying special attention to the messages which transfer the packets between controllers which make a full system.



When validating and simulating the model with the settings mentioned in Part 1 and Figure 2.1 and 2.5 we can observe the response times for each task when inserting the Worst Case Execution Time into the tool. Next each processing unit and their Worst Case Response Time (WCRT) will be investigated and then compared to the Matlab model (Response Time analysis). The resulting figures were obtained after several tries with various settings until the correct configuration of the tool was found, e.g. setting the schedule as preemptive was not set as it was not found in the first try.

## PU1

For PU1 a fixed preemptive priority scheme with four tasks T1,T2,T3 and Ts. Now when validating the model, Figure 2.2 is obtained showing the results and they are compared to the Matlab response times in Table 2.1.

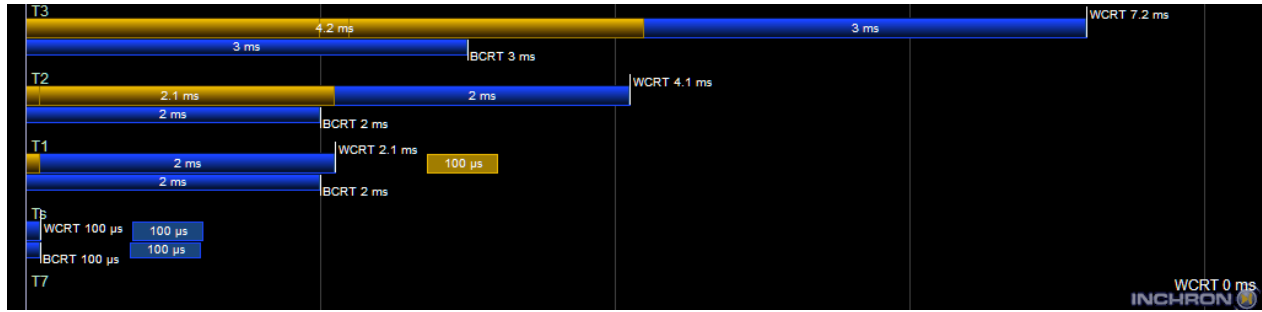


Figure 2.2: Showing the WCRT analysis for PU1. The Results can be read from each horizontal bar.

## PU2

Now PU2, has the same sheme as PU1, fixed preemptive priority, with four tasks T5,T6,T7 and T8. Now when validating the model, Figure 2.3 is obtained showing the results and they are compared to the Matlab response times in Table 2.1.



Figure 2.3: Showing the WCRT analysis for PU2. The Results can be read from each horizontal bar.

### PU3

Although PU3 has a time division multiplexing (TDM) its tasks will still be analyzed here, shown in Figure 2.4. It is important to state that  $T_4$  is not a real-time task as it is sending a message *out to the blue* but  $T_{ca}$  is an important task, actuating on the sensor value and completing the sensor to actuator delay.

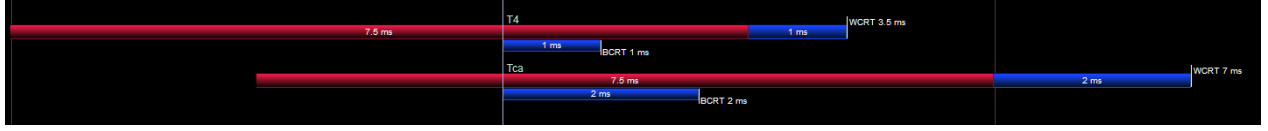


Figure 2.4: Showing the WCRT analysis for PU3 (TDM). The Results can be read from each horizontal bar

Table 2.1: By running the Matlab script `ResponsetimeAnylysis_FPP.m` with the different parameters given for PU1 and PU2 these response times are obtained. These files are then delivered as PU1.m and PU2.m

	PU1	$T_1$	$T_2$	$T_3$	$T_4 (T_s)$
Matlab (ms)		0.1	2.1	4.1	7.2
Inchron (ms)		0.1	2.1	4.1	7.2

	PU2	$T_5$	$T_6$	$T_7$	$T_8$
Matlab (ms)		6	3	9	5
Inchron (ms)		6	3	9	5

### 2.2.2 Response time analysis for the CAN bus messages

PU1 and PU2 are the only units within the system that are sending messages, shown for clarity in Figure 2.5, but PU3 contains the computing and actuating task which will receive the  $m_s$  message and mark the end of the sensor to actuator delay.

	Name	ECU	Peripheral	CAN_ID (Decimal)	Size	Format
1	M1	PU1	Cp1	0	8	Standard
2	M3	PU1	Cp1	1	8	Standard
3	Ms	PU1	Cp1	2	8	Standard
4	M8	PU2	c3	3	8	Standard

Figure 2.5: All can messages in the system, indicating PU source, individual CAN id and message size in bytes

## 2.3 Optimization for sensor-to-actuator delay

Now optimizing the sensor-to-actuator delay requires running the `CrhonOpt` tool within `Inchron` and setting up objectives for the real time requirements. That will be set to target: event-chain, which is set to be smaller or equal to 5ms.

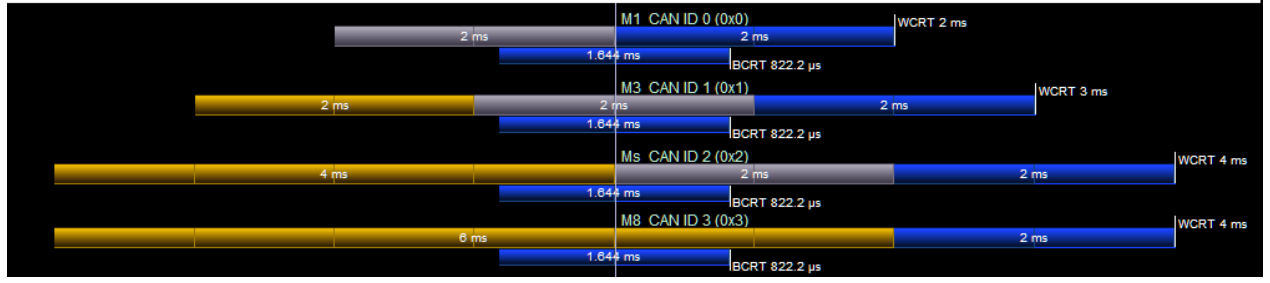


Figure 2.6: Showing the WCRT analysis for the CAN messages. The Results can be read from each horizontal bar and is compared in Table 2.2.

Table 2.2: CAN messages of the system compared from the Inchron tool suite and Matlab. Showing identical results

CAN	$m_2$	$m_1$	$m_3$	$m_8$
Matlab ms	2	3	4	4
Inchron	2	3	4	4

The initial priorities were selected according to period and execution time, shown in Figure 2.7 on the left, but when running the simulation to enhance the sensor to actuator delay, the tool changed the priorities in PU2, shown in Figure 2.7 on the right. This showed to be an improvement although this did not change the sensor to actuator delay.

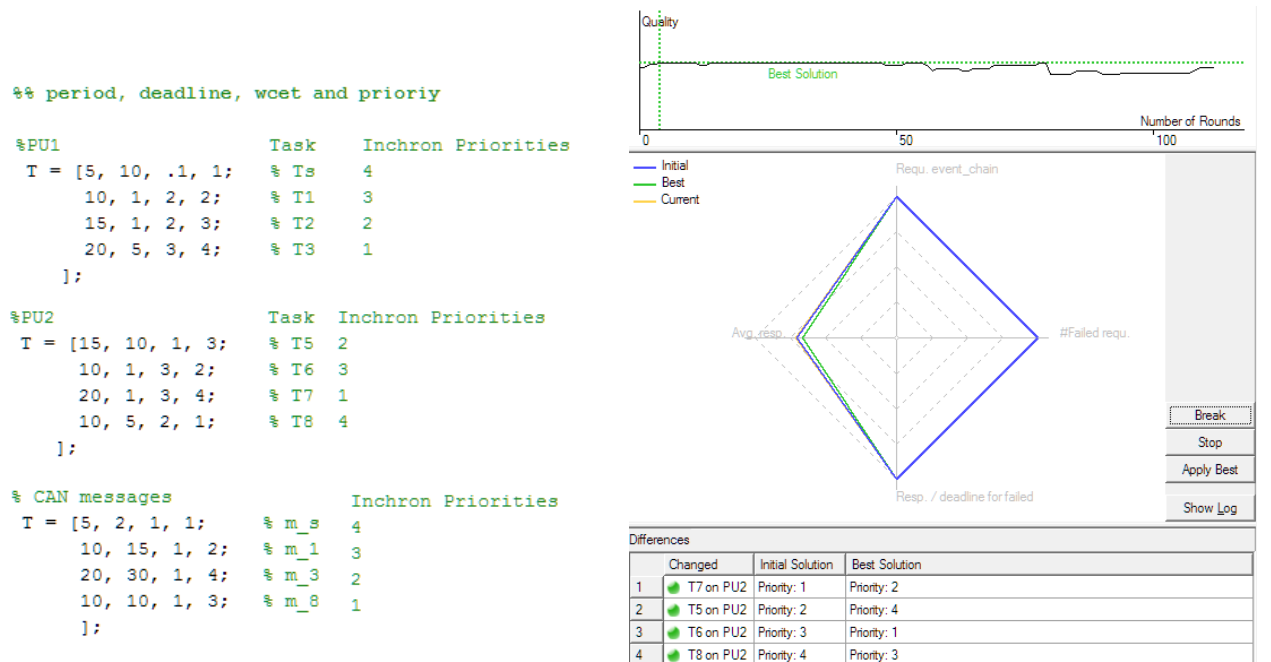


Figure 2.7: On the left, initial priorities used in the Matlab script and on the right the ChronOPT optimization of these parameters.

## 2.4 Design decision

Design decision are used from Part 1, seen in Section 1.4. To satisfy the design decision made in Part 1 saying that the sensor to actuator delay should be 8.1 ms as mentioned in Section 1.4 as well, some changes have to be made to the Inchron project. Without any changes to the project the Sensor to actuator delay will be 12 ms seen in Figure 2.8, as well notice the blue area which is 4.9 ms which should be 1 ms. So what we did was add a offset of  $4.9 \text{ ms} - 1 \text{ ms} = 3.9 \text{ ms}$  to the  $T_1$  task. This allows the model to act correctly and show the correct sensor to actuator delay of 8.1 ms and also lowering the sampling period to 10 ms.

## 2.5 System model and Results

Firstly: The Response time analysis for each of the Processing units was performed in Inchron and the timings match the timings obtained from the Matlab model as seen in Table 2.1 as for the CAN messages, seen in Table 2.2. Secondly: Plots from chronVIEW (before and after optimization)

Last: Control system input and output

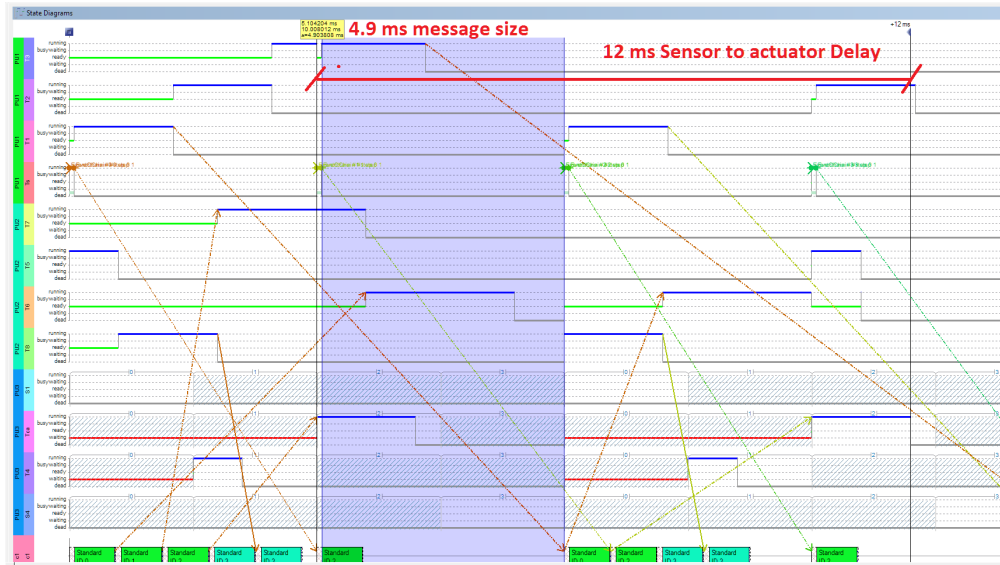


Figure 2.8: The first 20ms (equal to largest period) of the **optimized state diagram** for the CAN model. Showing the sensor to actuator delay and the wrong message time of  $m_1$

## 2.6 Conclusions

To conclude, the theoretical/Matlab analysis goes hand in hand with the Inchron modeling when used correctly. The tool is very powerful and can count for all design decision made in Part 1. The tool allows for a broad overview of the complicated system as well as troubleshooting more difficult sections of the analysis to maximize the performance of the CAN bus according to the theoretical analysis.

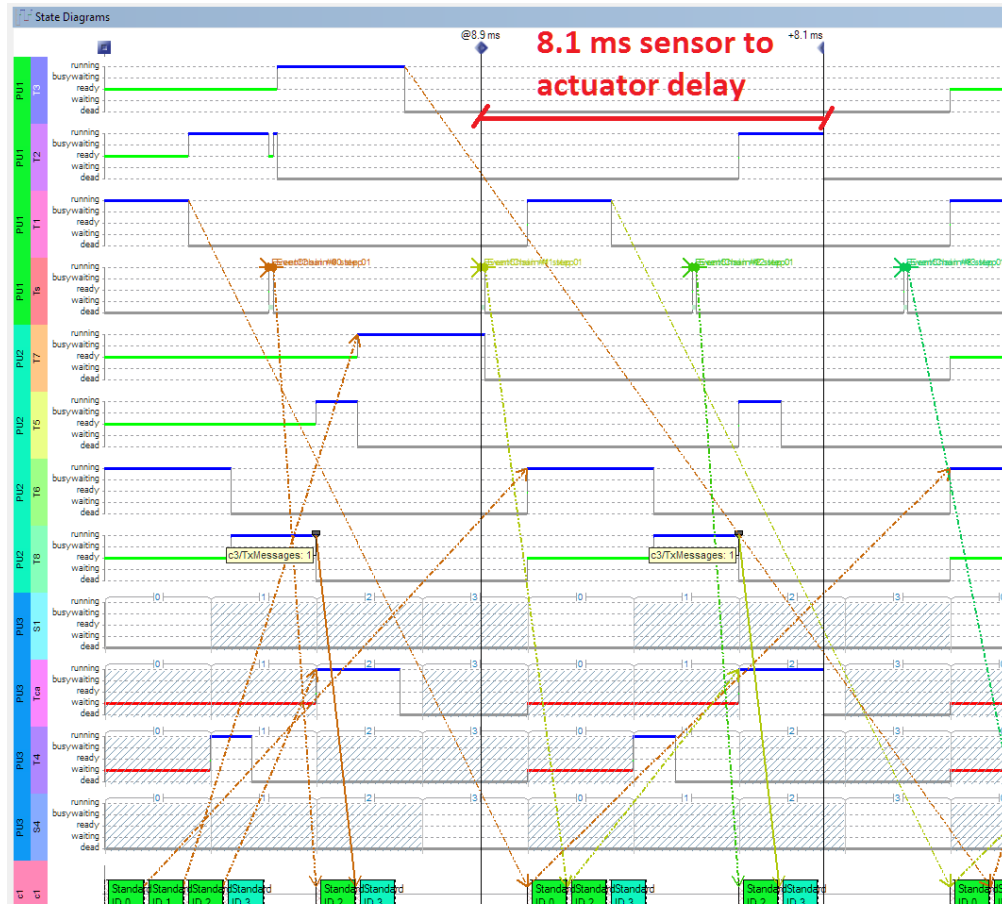


Figure 2.9: The first 20ms (equal to largest period) of the **non optimized state diagram** for the CAN model.

# Chapter 3

## Part 3

### 3.1 Introduction

Now describe the setup for the FlexRay modeling part, Figure 3.1 shows the most important given parts and also the tree view in Inchron along with the FlexRay bus connections.

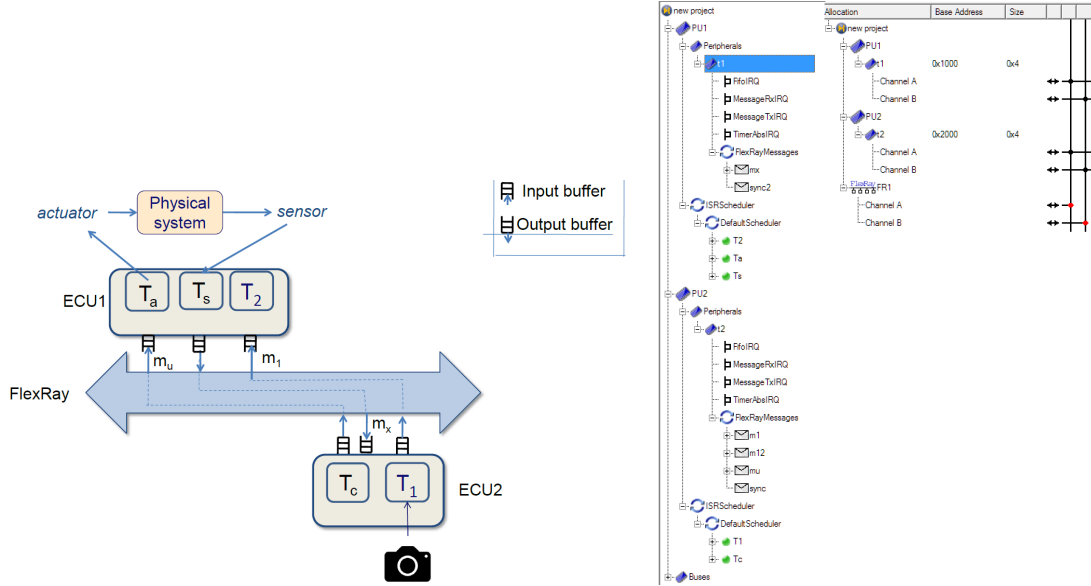


Figure 3.1: On the left, system diagram and on the right the tree view and architecture of the Inchron project.

### 3.2 Answer all the questions

1. What is the maximum cycle length possible?
2. What are the maximum possible durations for:
  - (a) static slot?
  - (b) dynamic slot?
  - (c) Network idle time?

3. Derive the implicitly defined parameters (e.g., no. of minislots).
4. Why do you need idle phase within a dynamic slot?
5. Does the above parameters conform to FlexRay specification?
6. Answer the above questions in your report Electrical Engineering

What is the maximum cycle length possible?

What are the maximum possible durations for

static slot?

dynamic slot?

Network idle time?

Derive the implicitly defined parameters (e.g., no. of minislots).

Why do you need idle phase within a dynamic slot?

Does the above parameters conform to FlexRay specification?

Answer the above questions in your report Electrical Engineering

### **3.2.1 Theoretical analysis versus actual implementation**

## **3.3 Design decision**

## **3.4 Results**

Firstly: Solution to the design problem. (Include the parameters you have chosen)

Secondly: from chronVIEW for your design

## **3.5 Conclusions**

## **3.6 Results**

## **3.7 Conclusion**