W1 1-4,5 Discrete Probability

1、一些概念

U: finite set (e.g. $U = \{0,1\}^n$)

Def: Probability distribution P over U is a function P: $U \rightarrow [0,1]$ such that $\Sigma P(x) = 1$

Uniform distribution: for all $x \in U$: P(x) = 1/|U|

Point distribution at x0 : P(x0) = 1, $\forall x \neq x0$: P(x) = 0

Distribution vector: (P(000), P(001), P(010), ..., P(111))

2. Events

For a set $A \subseteq U$: $Pr[A] = \Sigma P(x) \in [0,1]$ (Pr[U]=1)

The set A is called an event

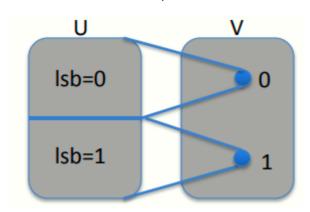
3、The union bound

For events A_1 and A_2 , $Pr[A_1 \cup A_2] \le Pr[A_1] + Pr[A_2]$

不等式恒成立, 当事件A1和A2相互独立时取等号

4. Random Variables

Def: a random variable X is a function X:U \rightarrow V Example: X: $\{0,1\}$ n $^{\wedge}$ \rightarrow $\{0,1\}$; X(y) = lsb(y) \in $\{0,1\}$



More generally: rand. var. X induces a distribution on V: $Pr[X=v] := Pr[X^{-1}(v)]$

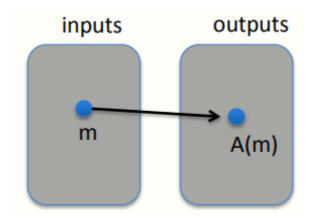
5. The uniform random variable

Let U be some set, e.g. $U = \{0,1\}^n$

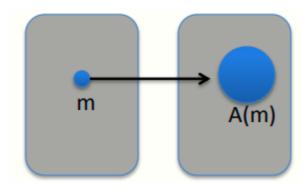
We write $r \leftarrow U$ to denote a uniform random variable over U for all $a \in U$: Pr[r = a] = 1/|U|

6. Randomized algorithms

Deterministic algorithm: y ← A(m),对于每次相同的输入,确定性算法总能得到相同的输出



Randomized algorithm y \leftarrow A(m ; r) where r \leftarrow {0,1}ⁿ , output is a random variable , 随机化算法每一次输入往往得到不同的输出



7. Independence

Def: events A and B are independent if $Pr[A \text{ and B}] = Pr[A] \cdot Pr[B]$

random variables X,Y taking values in V are independent if $\forall a,b \in V$: $Pr[X=a \text{ and } Y=b] = Pr[X=a] \cdot Pr[Y=b]$

8, XOR

XOR of two strings in {0,1}ⁿ is their bit-wise addition mod 2

Thm:Y a rand. var. over $\{0,1\}^n$, X an indep. uniform var. on $\{0,1\}^n$ Then Z := Y \oplus X is uniform var. on $\{0,1\}^n$

9. The birthday paradox

Let r1 , ..., rn \in U be indep. identically distributed random vars. Thm: when n= 1.2 × $|U|^{(1/2)}$ then Pr[$\exists i\neq j$: ri = rj] \geq 1/2

