

- You will *not* be able to use a calculator or computer for either the midterm or the final exam, so please do not use one for this assignment. You may use one to *check* your answer, but please do not use one to solve the problem.
- Only part of the problems may be graded. But, you have to submit all the problems.
- The deadline is \*by 6pm on Friday\*, Nov. 15. You have a grace period until 9am the next day. The grace period is to take care of any technical issues you have while submitting the file. The grace period should give ample time for handling any issues, so No late HW after the grace period will be accepted, **regardless** your technical issues.
- Submit only pdf files.

1. 5 marks Solve the following LP problem using the revised simplex method:

$$\begin{array}{llll}
 \text{Maximize} & 2x_1 & +3x_2 & +3x_3 \\
 \text{subject to} & 3x_1 & +x_2 & \leq 60 \\
 & -x_1 & +x_2 & +4x_3 \leq 10 \\
 & 2x_1 & -2x_2 & +5x_3 \leq 15 \\
 & x_1, & x_2, & x_3 \geq 0.
 \end{array}$$

*Hint:* requires  $\leq 3$  iterations.

### Solution: First iteration

- Form matrices:

$$\begin{array}{ll}
 \vec{b} = \begin{bmatrix} 60 \\ 10 \\ 15 \end{bmatrix} & \vec{x}_B^* = \begin{bmatrix} 60 \\ 10 \\ 15 \end{bmatrix} \\
 \vec{x}_B = \begin{bmatrix} x_4 \\ x_5 \\ x_6 \end{bmatrix} & \vec{x}_N = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} \\
 \mathbf{B} = \begin{bmatrix} x_4 & x_5 & x_6 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} & \mathbf{N} = \begin{bmatrix} x_1 & x_2 & x_3 \\ 3 & 1 & 0 \\ -1 & 1 & 4 \\ 2 & -2 & 5 \end{bmatrix} \\
 \vec{c}_B^T = [0 \ 0 \ 0] & \vec{c}_N^T = [2 \ 3 \ 3]
 \end{array}$$

- Solve  $\vec{y}^T \mathbf{B} = \vec{c}_B^T$

$$[y_1 \ y_2 \ y_3] \mathbf{B} = [0 \ 0 \ 0] \implies [y_1 \ y_2 \ y_3] = [0 \ 0 \ 0]$$

- Choose entering column:

$$\vec{c}_N^T - y\mathbf{N} = \vec{c}_N^T - 0 = [2 \quad 3 \quad 3]$$

We pick column 2. Hence  $x_2$  enters and

$$\vec{a} = \begin{bmatrix} 1 \\ 1 \\ -2 \end{bmatrix}$$

- Solve  $\mathbf{B}\vec{d} = \vec{a}$ :

$$\begin{array}{ccc} x_4 & x_5 & x_6 \\ \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \vec{d} = \begin{bmatrix} 1 \\ 1 \\ -2 \end{bmatrix} \\ \vec{d} = \begin{bmatrix} 1 \\ 1 \\ -2 \end{bmatrix} \end{array}$$

- Find largest  $t$  such that  $\vec{x}_B^* - t\vec{d} \geq 0$

$$\begin{bmatrix} 60 \\ 10 \\ 15 \end{bmatrix} - t \begin{bmatrix} 1 \\ 1 \\ -2 \end{bmatrix}$$

So  $t = 10$ . Since the second component  $= 0$ , it follows that  $x_5$  leaves.

- Update things

$$\vec{x}_B^* = \begin{bmatrix} 50 \\ 10 \\ 35 \end{bmatrix}$$

$$\vec{x}_B = \begin{bmatrix} x_4 \\ x_2 \\ x_6 \end{bmatrix}$$

$$\vec{x}_N = \begin{bmatrix} x_1 \\ x_5 \\ x_3 \end{bmatrix}$$

$$\mathbf{B} = \begin{array}{ccc} & x_4 & x_2 & x_6 \\ \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & -2 & 1 \end{bmatrix} \end{array}$$

$$\mathbf{N} = \begin{array}{ccc} & x_1 & x_5 & x_3 \\ \begin{bmatrix} 3 & 0 & 0 \\ -1 & 1 & 4 \\ 2 & 0 & 5 \end{bmatrix} \end{array}$$

$$\vec{c}_B^T = [0 \quad 3 \quad 0]$$

$$\vec{c}_N^T = [2 \quad 0 \quad 3]$$

**Second iteration**

- Solve  $\vec{y}^T \mathbf{B} = \vec{c}_B^T$ :

$$[y_1 \ y_2 \ y_3] \begin{array}{c} x_4 \quad x_2 \quad x_6 \\ \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & -2 & 1 \end{bmatrix} \end{array} = [0 \ 3 \ 0] \implies [y_1 \ y_2 \ y_3] = [0 \ 3 \ 0]$$

- Choose entering column:

$$\begin{aligned} \vec{c}_N^T - y\mathbf{N} &= [2 \ 0 \ 3] - [0 \ 3 \ 0] \begin{array}{c} x_1 \quad x_5 \quad x_3 \\ \begin{bmatrix} 3 & 0 & 0 \\ -1 & 1 & 4 \\ 2 & 0 & 5 \end{bmatrix} \end{array} \\ &= [2 \ 0 \ 3] - [-3 \ 3 \ 12] = [5 \ -3 \ -9] \end{aligned}$$

So we pick column 1. Hence  $x_1$  enters and

$$\vec{a} = \begin{bmatrix} 3 \\ -1 \\ 2 \end{bmatrix}$$

- Solve  $\mathbf{B}\vec{d} = \vec{a}$ :

$$\begin{array}{c} x_4 \quad x_2 \quad x_6 \\ \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & -2 & 1 \end{bmatrix} \end{array} \begin{bmatrix} d_1 \\ d_2 \\ d_3 \end{bmatrix} = \begin{bmatrix} 3 \\ -1 \\ 2 \end{bmatrix} \implies \vec{d} = \begin{bmatrix} 4 \\ -1 \\ 0 \end{bmatrix}$$

- Find largest  $t$  such that  $\vec{x}_B^* - t\vec{d} \geq 0$

$$\begin{bmatrix} 50 \\ 10 \\ 35 \end{bmatrix} - t \begin{bmatrix} 4 \\ -1 \\ 0 \end{bmatrix}$$

So  $t = 25/2$ . Since the first component is 0, it follows that  $x_4$  leaves.

- Update things

$$\vec{x}_B^* = \begin{bmatrix} 25/2 \\ 45/2 \\ 35 \end{bmatrix}$$

$$\vec{x}_B = \begin{bmatrix} x_1 \\ x_2 \\ x_6 \end{bmatrix}$$

$$\mathbf{B} = \begin{array}{c} \begin{matrix} x_1 & x_2 & x_6 \end{matrix} \\ \begin{bmatrix} 3 & 1 & 0 \\ -1 & 1 & 0 \\ 2 & -2 & 1 \end{bmatrix} \end{array}$$

$$\vec{c}_B^T = [2 \ 3 \ 0]$$

$$\vec{x}_N = \begin{bmatrix} x_4 \\ x_5 \\ x_3 \end{bmatrix}$$

$$\mathbf{N} = \begin{array}{c} \begin{matrix} x_4 & x_5 & x_3 \end{matrix} \\ \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 4 \\ 0 & 0 & 5 \end{bmatrix} \end{array}$$

$$\vec{c}_N^T = [0 \ 0 \ 3]$$

### Third iteration

- Solve  $\vec{y}^T \mathbf{B} = \vec{c}_B^T$ :

$$[y_1 \ y_2 \ y_3] \begin{array}{c} \begin{matrix} x_4 & x_2 & x_6 \end{matrix} \\ \begin{bmatrix} 3 & 1 & 0 \\ -1 & 1 & 0 \\ 2 & -2 & 1 \end{bmatrix} \end{array} = [2 \ 3 \ 0] \implies [y_1 \ y_2 \ y_3] = [5/4 \ 7/4 \ 0]$$

- Choose entering variable

$$\vec{c}_N^T - y\mathbf{N} = [0 \ 0 \ 3] - [5/4 \ 7/4 \ 0] \begin{array}{c} \begin{matrix} x_4 & x_5 & x_3 \end{matrix} \\ \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 4 \\ 0 & 0 & 5 \end{bmatrix} \end{array} = [-5/4 \ -7/4 \ -4]$$

Since all are negative, there are no entering columns. The current basic feasible solution is optimal.

- The current solution is

$$\vec{x}_B = \begin{bmatrix} x_1 \\ x_2 \\ x_6 \end{bmatrix} = \vec{x}_B^* = \begin{bmatrix} 25/2 \\ 45/2 \\ 35 \end{bmatrix}$$

$$\vec{x}_N = \begin{bmatrix} x_4 \\ x_5 \\ x_3 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

$$z = \vec{c}_B^T \vec{x}_B^* = 25 + 135/2 = 185/2$$

2. 5 marks Solve the following LP problem using the revised simplex method.

$$\begin{array}{llll}
 \text{Maximize} & 2x_1 & +3x_2 & \\
 \text{subject to} & x_1 & +2x_2 & \leq 6 \\
 & 2x_1 & +x_2 & \leq 8 \\
 & -x_1 & +x_2 & \leq 1 \\
 & & 2x_2 & \leq 1 \\
 & x_1, & x_2, & \geq 0.
 \end{array}$$

*Hint:* requires  $\leq 3$  iterations.

### Solution: First iteration

- Form matrices

$$\begin{array}{lll}
 \vec{x}_B^* = \begin{bmatrix} 6 \\ 8 \\ 1 \\ 1 \end{bmatrix} & \vec{x}_B = \begin{bmatrix} x_3 \\ x_4 \\ x_5 \\ x_6 \end{bmatrix} & \vec{x}_N = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \\
 \vec{c}_B^T = [0 \ 0 \ 0 \ 0] & \vec{c}_N^T = [2 \ 3] &
 \end{array}$$

$$\mathbf{N} = \begin{array}{c} x_1 \quad x_2 \\ \begin{bmatrix} 1 & 2 \\ 2 & 1 \\ -1 & 1 \\ 0 & 2 \end{bmatrix} \end{array} \quad \mathbf{B}_0 = \begin{array}{c} x_3 \quad x_4 \quad x_5 \quad x_6 \\ \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \end{array}$$

- Solve  $\vec{y}^T \mathbf{B}_0 = \vec{c}_B^T$ . Since  $\mathbf{B}_0 = I$  and  $\vec{c}_B^T = 0$  we have  $\vec{y}^T = [0 \ 0 \ 0 \ 0]$ .
- Look at  $\vec{c}_N^T - \vec{y}^T \mathbf{N}$  to choose entering column:

$$\vec{c}_N^T - \vec{y}^T \mathbf{N} = [2 \ 3] \quad \text{since } \vec{y}^T = 0$$

We choose the second component — hence  $x_2$  enters and  $\vec{a} = \begin{bmatrix} 2 \\ 1 \\ 1 \\ 2 \end{bmatrix}$ .

- Solve  $\mathbf{B}_0 \vec{d} = \vec{a}$ . Since  $\mathbf{B}_0 = I$  we have  $\vec{d} = \vec{a}$ .
- Find largest  $t$  such that  $\vec{x}_B^* - t\vec{d} \geq 0$ :

$$\begin{bmatrix} 6 \\ 8 \\ 1 \\ 1 \end{bmatrix} - t \begin{bmatrix} 2 \\ 1 \\ 1 \\ 2 \end{bmatrix} \geq 0 \quad \implies \quad t = 1/2 \quad \text{and} \quad \begin{bmatrix} 6 \\ 8 \\ 1 \\ 1 \end{bmatrix} - (1/2) \begin{bmatrix} 2 \\ 1 \\ 1 \\ 2 \end{bmatrix} = \begin{bmatrix} 5 \\ 15/2 \\ 1/2 \\ 0 \end{bmatrix}$$

The fourth component is zero — hence  $x_6$  leaves.

- Update,  $x_2$  swaps with  $x_6$

$$\begin{aligned}\vec{x}_B^* &= \begin{bmatrix} 5 \\ 15/2 \\ 1/2 \\ 1/2 \end{bmatrix} & \vec{x}_B &= \begin{bmatrix} x_3 \\ x_4 \\ x_5 \\ x_2 \end{bmatrix} & \vec{x}_N &= \begin{bmatrix} x_1 \\ x_6 \end{bmatrix} \\ \vec{c}_B^T &= [0 \ 0 \ 0 \ 3] & \vec{c}_N^T &= [2 \ 0] \\ \mathbf{N} &= \begin{array}{cc} & \begin{matrix} x_1 & x_6 \end{matrix} \\ \begin{bmatrix} 1 & 0 \\ 2 & 0 \\ -1 & 0 \\ 0 & 1 \end{bmatrix} \end{array} & \mathbf{E}_1 &= \begin{bmatrix} 1 & 0 & 0 & 2 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 2 \end{bmatrix}\end{aligned}$$

### Second iteration

- Solve  $\vec{y}^T \mathbf{E}_1 = \vec{c}_B^T$ :

$$\begin{bmatrix} y_1 & y_2 & y_3 & y_4 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 2 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 2 \end{bmatrix} = [0 \ 0 \ 0 \ 3] \implies \vec{y}^T = [0 \ 0 \ 0 \ 3/2]$$

- Look at  $\vec{c}_N^T - \vec{y}^T \mathbf{N}$ :

$$[2 \ 0] - [0 \ 0 \ 0 \ 3/2] \begin{array}{cc} & \begin{matrix} x_1 & x_6 \end{matrix} \\ \begin{bmatrix} 1 & 0 \\ 2 & 0 \\ -1 & 0 \\ 0 & 1 \end{bmatrix} \end{array} = [2 \ -3/2]$$

We pick the first component —  $x_1$  enters and  $\vec{a} = \begin{bmatrix} 1 \\ 2 \\ -1 \\ 0 \end{bmatrix}$ .

- Solve  $\mathbf{E}_1 \vec{d} = \vec{a}$ :

$$\begin{bmatrix} 1 & 0 & 0 & 2 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 2 \end{bmatrix} \begin{bmatrix} d_1 \\ d_2 \\ d_3 \\ d_4 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \\ -1 \\ 0 \end{bmatrix} \implies \vec{d} = \begin{bmatrix} 1 \\ 2 \\ -1 \\ 0 \end{bmatrix}$$

- Find largest  $t$  such that  $\vec{x}_B^* - t\vec{d} \geq 0$ :

$$\begin{bmatrix} 5 \\ 15/2 \\ 1/2 \\ 1/2 \end{bmatrix} - t \begin{bmatrix} 1 \\ 2 \\ -1 \\ 0 \end{bmatrix} \geq 0 \implies t = 15/4 \text{ and } \begin{bmatrix} 5 \\ 15/2 \\ 1/2 \\ 1/2 \end{bmatrix} - (15/4) \begin{bmatrix} 1 \\ 2 \\ -1 \\ 0 \end{bmatrix} = \begin{bmatrix} 5/4 \\ 0 \\ 17/4 \\ 1/2 \end{bmatrix}$$

The second component is zero —  $x_4$  leaves.

- Update —  $x_1$  swaps with  $x_4$ :

$$\begin{aligned} \vec{x}_B^* &= \begin{bmatrix} 5/4 \\ 15/4 \\ 17/4 \\ 1/2 \end{bmatrix} & \vec{x}_B &= \begin{bmatrix} x_3 \\ x_1 \\ x_5 \\ x_2 \end{bmatrix} & \vec{x}_N &= \begin{bmatrix} x_4 \\ x_6 \end{bmatrix} \\ \vec{c}_B^T &= [0 \ 2 \ 0 \ 3] & \vec{c}_N^T &= [0 \ 0] \\ \mathbf{N} &= \begin{array}{cc} & \begin{matrix} x_4 & x_6 \end{matrix} \\ \begin{bmatrix} 0 & 0 \\ 1 & 0 \\ 0 & 0 \\ 0 & 1 \end{bmatrix} & \end{array} & \mathbf{E}_1 &= \begin{bmatrix} 1 & 0 & 0 & 2 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 2 \end{bmatrix} & \mathbf{E}_2 &= \begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & 2 & 0 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \end{aligned}$$

### Third iteration

- Solve  $\vec{y}^T \mathbf{E}_1 \mathbf{E}_2 = \vec{c}_B^T$ :

$$\begin{aligned} [u_1 \ u_2 \ u_3 \ u_4] \begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & 2 & 0 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} &= [0 \ 2 \ 0 \ 3] \implies \vec{u}^T = [0 \ 1 \ 0 \ 3] \\ [y_1 \ y_2 \ y_3 \ y_4] \begin{bmatrix} 1 & 0 & 0 & 2 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 2 \end{bmatrix} &= [0 \ 1 \ 0 \ 3] \implies \vec{y}^T = [0 \ 1 \ 0 \ 1] \end{aligned}$$

- Look at  $\vec{c}_N^T - \vec{y}^T \mathbf{N}$ :

$$[0 \ 0] - [0 \ 1 \ 0 \ 1] \begin{array}{cc} & \begin{matrix} x_4 & x_6 \end{matrix} \\ \begin{bmatrix} 0 & 0 \\ 1 & 0 \\ 0 & 0 \\ 0 & 1 \end{bmatrix} & \end{array} = [-1, -1]$$

There are no entering variables or columns, so current solution is optimal.

- Optimal solution is

$$\vec{x}_B = \begin{bmatrix} x_3 \\ x_1 \\ x_5 \\ x_2 \end{bmatrix} = \vec{x}_B^* = \begin{bmatrix} 5/4 \\ 15/4 \\ 17/4 \\ 1/2 \end{bmatrix} \quad \vec{x}_N = \begin{bmatrix} x_4 \\ x_6 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

$$z = \vec{c}_B^T \vec{x}_B = 2 \times (15/4) + 3 \times (1/2) = 9$$

3. 5 marks Your colleague is working on the following LP problem:

$$\begin{array}{llll} \text{Maximize } z = & -x_1 & +x_2 & +2x_3 \\ \text{subject to} & x_1 & +2x_2 & -x_3 \leq 2 \\ & 2x_1 & +3x_2 & +x_3 \leq 6 \\ & -2x_1 & +4x_2 & +2x_3 \leq 8 \\ & x_1, & x_2, & x_3 \geq 0. \end{array}$$

He tells you that he has found an optimal solution when  $x_2, x_5, x_6$  are non-basic and  $x_1, x_3, x_4$  are basic.

- (a) Write down the original dictionary and the dictionary that is thought to be optimal: Give  $\mathbf{B}, \mathbf{N}, \vec{c}_N^T, \vec{c}_B^T, \vec{x}_B, \vec{x}_N, \vec{b}$ , and remember to write in the basis headings (that is, label the columns of matrices with the corresponding variables).

**Solution:**

- The original dictionary is given by

$$\begin{array}{lll} & x_4 & x_5 & x_6 \\ \mathbf{B} = & \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} & \mathbf{N} = & \begin{bmatrix} x_1 & x_2 & x_3 \\ 1 & 2 & -1 \\ 2 & 3 & 1 \\ -2 & 4 & 2 \end{bmatrix} \\ \vec{x}_B = & \begin{bmatrix} x_4 \\ x_5 \\ x_6 \end{bmatrix} & \vec{x}_N = & \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} & \vec{b} = & \begin{bmatrix} 2 \\ 6 \\ 8 \end{bmatrix} \\ \vec{c}_B^T = & [0 & 0 & 0] & \vec{c}_N^T = & [-1 & 1 & 2] \end{array}$$



- So if we make  $x_2, x_5, x_6$  non-basic and the other variables basic, we get

$$\mathbf{B} = \begin{array}{c} x_1 \quad x_3 \quad x_4 \\ \begin{bmatrix} 1 & -1 & 1 \\ 2 & 1 & 0 \\ -2 & 2 & 0 \end{bmatrix} \end{array} \quad \mathbf{N} = \begin{array}{c} x_2 \quad x_5 \quad x_6 \\ \begin{bmatrix} 2 & 0 & 0 \\ 3 & 1 & 0 \\ 4 & 0 & 1 \end{bmatrix} \end{array}$$

$$\vec{x}_B = \begin{bmatrix} x_1 \\ x_3 \\ x_4 \end{bmatrix} \quad \vec{x}_N = \begin{bmatrix} x_2 \\ x_5 \\ x_6 \end{bmatrix}$$

$$\vec{c}_B^T = [-1 \quad 2 \quad 0] \quad \vec{c}_N^T = [1 \quad 0 \quad 0]$$

- (b) What are the values of the basic and non-basic variables in your colleague's dictionary? Is this a feasible solution?

**Solution:**

- So the current basic feasible solution is given by  $\vec{x}_N = 0$  and  $\vec{x}_B^* = \mathbf{B}^{-1}\vec{b}$ . To find the latter we solve  $\mathbf{B}\vec{x}_B^* = \vec{b}$

$$\begin{array}{c} x_1 \quad x_3 \quad x_4 \\ \begin{bmatrix} 1 & -1 & 1 \\ 2 & 1 & 0 \\ -2 & 2 & 0 \end{bmatrix} \begin{bmatrix} x_1^* \\ x_3^* \\ x_4^* \end{bmatrix} = \begin{bmatrix} 2 \\ 6 \\ 8 \end{bmatrix} \end{array}$$

which gives  $x_4^* = 6, x_3^* = \frac{14}{3}, x_1^* = \frac{2}{3}$ . The other variables are zero.

- Since all the variables are non-negative, it follows that the dictionary is feasible.

- (c) Is your colleague's dictionary optimal?

**Solution:**

- To check optimality we need to verify that  $\vec{c}_N^T - \vec{c}_B^T \mathbf{B}^{-1} \mathbf{N} \leq 0$ .
- First compute  $\vec{y}^T = \vec{c}_B^T \mathbf{B}^{-1}$  by solving  $\vec{y}^T \mathbf{B} = \vec{c}_B^T$ :

$$[y_1 \quad y_2 \quad y_3] \begin{array}{c} x_1 \quad x_3 \quad x_4 \\ \begin{bmatrix} 1 & -1 & 1 \\ 2 & 1 & 0 \\ -2 & 2 & 0 \end{bmatrix} \end{array} = [-1 \quad 2 \quad 0]$$

This gives  $y_1 = 0$ ,  $2y_2 - 2y_3 = -1$ ,  $y_2 + 2y_3 = 2$ . Hence  $\bar{y}^T = [0, \frac{1}{3}, \frac{5}{6}]$

- So  $\bar{c}_N^T - \bar{y}^T \mathbf{N}$  is:

$$[1 \quad 0 \quad 0] - [0 \quad \frac{1}{3} \quad \frac{5}{6}] \begin{matrix} & x_2 & x_5 & x_6 \\ \begin{bmatrix} 2 & 0 & 0 \\ 3 & 1 & 0 \\ 4 & 0 & 1 \end{bmatrix} & = & [-\frac{10}{3} & -\frac{1}{3} & -\frac{5}{6}] \end{matrix}$$

Hence the current dictionary is optimal.

(d) What is the optimal value?

**Solution:**

- The objective function is therefore  $z = -1 \times \frac{2}{3} + 0 + 2 \times \frac{14}{3} = \frac{26}{3}$ .

4. 5 marks [This is continuation of the problem 3] Your colleague is working on the following LP problem:

$$\begin{array}{llll} \text{Maximise } z = & -x_1 & +x_2 & +2x_3 \\ \text{subject to} & x_1 & +2x_2 & -x_3 \leq 2 \\ & 2x_1 & +3x_2 & +x_3 \leq 6 \\ & -2x_1 & +4x_2 & +2x_3 \leq 8 \\ & x_1, & x_2, & x_3 \geq 0. \end{array}$$

He tells you that he has found an optimal solution when  $x_2, x_5, x_6$  are non-basic and  $x_1, x_3, x_4$  are basic. **[You can use the results of Problem 3.]**

(a) Your colleague then decides to change the third constraint to

$$-2x_1 + 4x_2 + 2x_3 \leq 8 + \beta$$

Over what range of  $\beta$  is your colleague's dictionary feasible and optimal. What is the corresponding solution + optimal value?

**Solution:**

- Changing  $\vec{b}$  does not effect the coefficients of the non-basic variables in  $z$ . Hence they all remain negative.
- So we only need to check feasibility, we do this by seeing if  $\vec{x}_B^* \geq 0$

- To do this we solve  $\mathbf{B}\vec{x}_B^* = \vec{b}$ :

$$\begin{array}{ccc} x_1 & x_3 & x_4 \\ \begin{bmatrix} 1 & -1 & 1 \\ 2 & 1 & 0 \\ -2 & 2 & 0 \end{bmatrix} & \begin{bmatrix} x_1^* \\ x_3^* \\ x_4^* \end{bmatrix} & = \begin{bmatrix} 2 \\ 6 \\ 8 + \beta \end{bmatrix} \end{array}$$

which gives

$$\vec{x}_B^* = \begin{bmatrix} \frac{4-\beta}{6} \\ \frac{14+\beta}{3} \\ \frac{12+\beta}{2} \end{bmatrix}$$

- In order for this to be non-negative, we need

$$-12 \leq \beta \leq 4$$

- And in this range the optimal solution is

$$\begin{array}{lll} x_1^* = \frac{4-\beta}{6} & x_2^* = 0 & x_3^* = \frac{14+\beta}{3} \\ x_4^* = \frac{12+\beta}{2} & x_5^* = x_6^* = 0 & \end{array}$$

and the objective function is

$$\begin{aligned} z &= -x_1 + x_2 + 2x_3 = -\frac{4-\beta}{6} + 2\frac{14+\beta}{3} \\ &= \frac{26}{3} + \frac{5\beta}{6} \end{aligned}$$

- (b) Your (increasingly indecisive) colleague then decides to leave the third constraint as it was originally, but instead changes the objective function to

$$z = (\gamma - 1)x_1 + x_2 + 2x_3$$

Over what range of  $\gamma$  is the dictionary feasible and optimal. What is the corresponding solution and optimal value?

**Solution:**

- Changing  $\vec{c}$  does not change the current value of  $\vec{x}_B^*$  and so the dictionary remains feasible.

- Hence we only need to check optimality, which we do by checking that  $\bar{c}_N^T - \bar{c}_B^T \mathbf{B}^{-1} \mathbf{N} \leq 0$ .
- First solve  $\bar{y}^T \mathbf{B} = \bar{c}_B^T$ :

$$[y_1 \ y_2 \ y_3] \begin{matrix} & x_1 & x_3 & x_4 \\ \begin{bmatrix} 1 & -1 & 1 \\ 2 & 1 & 0 \\ -2 & 2 & 0 \end{bmatrix} & = & [\gamma - 1 & 2 & 0] \end{matrix}$$

$$\bar{y}^T = [0 \quad \frac{1+\gamma}{3} \quad \frac{5-\gamma}{6}]$$

- Now check  $\bar{c}_N^T - \bar{y}^T \mathbf{A}_N$ :

$$\begin{aligned} \bar{c}_N^T - \bar{y}^T \mathbf{A}_N &= [1 \ 0 \ 0] - [0 \ \frac{1+\gamma}{3} \ \frac{5-\gamma}{6}] \begin{matrix} & x_2 & x_5 & x_6 \\ \begin{bmatrix} 2 & 0 & 0 \\ 3 & 1 & 0 \\ 4 & 0 & 1 \end{bmatrix} \end{matrix} \\ &= [\frac{-10-\gamma}{3} \quad \frac{-1-\gamma}{3} \quad \frac{\gamma-5}{6}] \end{aligned}$$

- Now in order for  $\bar{c}_N^T$  to be non-positive we require

$$-1 \leq \gamma \leq 5$$

- In this range the optimal solution is

$$x_1^* = \frac{2}{3} \quad x_2^* = 0 \quad x_3^* = \frac{14}{3} \quad x_4^* = 6 \quad x_5^* = x_6^* = 0$$

and the optimal value is

$$\begin{aligned} z &= (\gamma - 1)x_1^* + x_2^* + 2x_3^* = \frac{2\gamma - 2}{3} + 0 + \frac{28}{3} \\ &= \frac{26}{3} + \frac{2}{3}\gamma \end{aligned}$$

5. 5 marks Consider the following LP problem in which the objective function depends

on a parameter  $\gamma$ :

$$\begin{array}{rcll} \max & (4 + \gamma)x_1 + 5x_2 & & \\ \text{such that} & x_1 + x_2 & \leq & 5 \\ & x_1 + 2x_2 & \leq & 8 \\ & 2x_1 + x_2 & \leq & 8 \\ & x_1, x_2 & \geq & 0 \end{array}$$

When  $-\frac{3}{2} \leq \gamma \leq 1$  this has optimal dictionary

$$\begin{array}{rcll} x_1 & = & 2 & -2x_3 + x_4 \\ x_2 & = & 3 & +x_3 - x_4 \\ x_5 & = & 1 & +3x_3 - x_4 \\ \hline z & = & (23 + 2\gamma) & -(3 + 2\gamma)x_3 + (\gamma - 1)x_4 \end{array}$$

Starting from this dictionary find the optimal solution and the optimal value as a function of  $\gamma$ ,

(a) when  $-\frac{3}{2} \leq \gamma \leq 1$ ,

**Solution:** The initial dictionary is

$$\begin{array}{rcll} x_1 & = & 2 & -2x_3 + x_4 \\ x_2 & = & 3 & +x_3 - x_4 \\ x_5 & = & 1 & +3x_3 - x_4 \\ \hline z & = & (23 + 2\gamma) & -(3 + 2\gamma)x_3 - (1 - \gamma)x_4 \end{array}$$

This is feasible and optimal while  $-\frac{3}{2} \leq \gamma \leq 1$ . This gives

$$x_1 = 2 \quad x_2 = 3 \quad x_3 = x_4 = 0 \quad x_5 = 1 \quad z = 23 + 2\gamma$$

(b) when  $-\infty \leq \gamma \leq -\frac{3}{2}$ ,

**Solution:** When  $\gamma \leq -3/2$  then the coefficient of  $x_3$  becomes positive and so  $x_3$  must enter, and  $x_1$  must leave. Pivoting gives:

$$\begin{array}{rcll} x_3 & = & 1 & -(1/2)x_1 + (1/2)x_4 \\ x_2 & = & 4 & -(1/2)x_1 - (1/2)x_4 \\ x_5 & = & 4 & -(3/2)x_1 + (1/2)x_4 \\ \hline z & = & 20 & +(3/2 + \gamma)x_1 - (5/2)x_4 \end{array}$$

We see that this is feasible optimal for all  $\gamma < -3/2$ . This gives

$$x_1 = 0 \quad x_2 = 4 \quad x_3 = 1 \quad x_4 = 0 \quad x_5 = 4 \quad z = 20$$

(c) when  $1 \leq \gamma \leq 6$  and

**Solution:** If  $\gamma > 1$  then the coefficient of  $x_4$  becomes positive and so it must enter the basis. The leaving variable is  $x_5$ . Pivoting gives:

$$\begin{array}{rclcl} x_1 & = & 3 & +x_3 & -x_5 \\ x_2 & = & 2 & -2x_3 & +x_5 \\ x_4 & = & 1 & +3x_3 & -x_5 \\ \hline z & = & (22+3\gamma) & +(\gamma-6)x_3 & +(1-\gamma)x_5 \end{array}$$

this is clearly feasible + optimal while  $1 \leq \gamma \leq 6$ . This gives

$$x_1 = 3 \quad x_2 = 2 \quad x_3 = 0 \quad x_4 = 1 \quad x_5 = 0 \quad z = 22 + 3\gamma$$

(d) when  $6 \leq \gamma \leq \infty$ .

**Solution:** If  $\gamma > 6$  the coefficient of  $x_3$  becomes positive and it must enter the basis, and we see that  $x_2$  leaves. Pivoting gives

$$\begin{array}{rclcl} x_1 & = & 4 & -(1/2)x_2 & -(1/2)x_5 \\ x_3 & = & 1 & -(1/2)x_2 & +(1/2)x_5 \\ x_4 & = & 4 & -(3/2)x_2 & +(1/2)x_5 \\ \hline z & = & (16+4\gamma) & +(3-\gamma/2)x_2 & +(2-\gamma/2)x_5 \end{array}$$

This is feasible + optimal for all  $\gamma \geq 6$ .

$$x_1 = 4 \quad x_2 = 0 \quad x_3 = 1 \quad x_4 = 4 \quad x_5 = 0 \quad z = 16 + 4\gamma$$

Do not solve the problem from scratch each time. You will find it easier to do this problem using dictionaries rather than the revised simplex method.

**Solution:** So to summarise we have

$$\begin{array}{llll} \gamma \leq -3/2 & x_1^* = 0 & x_2^* = 4 & z = 20 + 0\gamma \\ -3/2 \leq \gamma \leq 1 & x_1^* = 2 & x_2^* = 3 & z = 23 + 2\gamma \\ 1 \leq \gamma \leq 6 & x_1^* = 3 & x_2^* = 2 & z = 22 + 3\gamma \\ 6 \leq \gamma & x_1^* = 4 & x_2^* = 0 & z = 16 + 4\gamma \end{array}$$

6. 5 marks Consider the following LP problem in which one of the constraints bounds

depends on a parameter  $\beta$ :

$$\begin{array}{rcll} \max & 4x_1 + 5x_2 & & \\ \text{such that} & x_1 + x_2 & \leq & 5 \\ & x_1 + 2x_2 & \leq & 8 + \beta \\ & 2x_1 + x_2 & \leq & 8 \\ & x_1, x_2 & \geq & 0 \end{array}$$

When  $-1 \leq \beta \leq 2$  the optimal dictionary is

$$\begin{array}{rcll} x_1 & = & 2 - \beta & -2x_3 + x_4 \\ x_2 & = & 3 + \beta & +x_3 - x_4 \\ x_5 & = & 1 + \beta & +3x_3 - x_4 \\ \hline z & = & 23 + \beta & -3x_3 - x_4 \end{array}$$

Starting from this dictionary find the optimal solution and the optimal value as a function of  $\beta$ ,

(a) when  $-1 \leq \beta \leq 2$ ,

**Solution:** The initial dictionary is

$$\begin{array}{rcll} x_1 & = & 2 - \beta & -2x_3 + x_4 \\ x_2 & = & 3 + \beta & +x_3 - x_4 \\ x_5 & = & 1 + \beta & +3x_3 - x_4 \\ \hline z & = & 23 + \beta & -3x_3 - x_4 \end{array}$$

This is dual-feasible and feasible when  $-1 \leq \beta \leq 2$ . This gives

$$x_1 = 2 - \beta \quad x_2 = 3 + \beta \quad x_3 = x_4 = 0 \quad x_5 = 1 + \beta \quad z = 23 + \beta$$

(b) when  $2 \leq \beta \leq \infty$ ,

**Solution:** When  $\beta \geq 2$  the dictionary remains dual feasible, but stops being feasible. In particular,  $x_1 < 0$ . Hence  $x_1$  is the leaving variable. Rewriting the pivot row gives  $0 = (2 - \beta) - x_1 - 2x_3 + x_4$ . Adding  $q \times$  this equation to  $z$  gives

$$z = 23 + 2q - 2\beta q - qx_1 + x_3(-3 - 2q) + x_4(-1 + q)$$

The largest  $q$  that keeps dual-feasibility is  $q = 1$ . Hence  $x_4$  is the entering variable. Pivoting gives

$$\begin{array}{rcll} x_4 & = & -2 + \beta & +x_1 + 2x_3 \\ x_2 & = & 5 & -x_1 - x_3 \\ x_5 & = & 3 & -x_1 + x_3 \\ \hline z & = & 25 & -x_1 - 5x_3 \end{array}$$

This is feasible and dual-feasible for all  $\beta \geq 2$ , and gives

$$x_1 = 0 \quad x_2 = 5 \quad x_3 = 0 \quad x_4 = -2 + \beta \quad x_5 = 3 \quad z = 25$$

(c) when  $-4 \leq \beta \leq -1$ ,

**Solution:** In this range  $x_5$  becomes negative — so  $x_5$  is the leaving variable. Rewriting the pivot row gives  $0 = 1 + \beta - x_5 + 3x_3 - x_4$ . Adding  $q \times$  this equation to  $z$  gives

$$z = (23 + \beta) + q(1 + \beta) - x_5q + x_3(-3 + 3q) + x_4(-1 - q)$$

Hence  $q = 1$  and  $x_3$  is the leaving variable. Pivoting gives

$$\begin{array}{rcl} x_1 & = & 8/3 - (1/3)\beta + (1/3)x_4 - (2/3)x_5 \\ x_2 & = & 8/3 + (2/3)\beta - (2/3)x_4 + (1/3)x_5 \\ x_3 & = & -1/3 - (1/3)\beta + (1/3)x_4 + (1/3)x_5 \\ \hline z & = & 24 + 2\beta - 2x_4 - x_5 \end{array}$$

This is feasible and optimal for  $-4 \leq \beta \leq -1$ , and gives

$$x_1 = \frac{8 - \beta}{3} \quad x_2 = \frac{8 + 2\beta}{3} \quad x_3 = -\frac{1 + \beta}{3} \quad x_4 = 0 \quad x_5 = 0$$

$$z = 24 + 2\beta$$

(d) and when  $-8 \leq \beta \leq -4$ .

**Solution:** In this range  $x_2$  becomes negative — so  $x_2$  is the leaving variable. Rewrite the pivot row to get  $0 = (8/3) + (2/3)\beta - x_2 - (2/3)x_4 + (1/3)x_5$ . Adding  $q \times$  this row to  $z$  gives:

$$z = 24 + 2\beta + q((8/3) + (2/3)\beta) - qx_2 + x_4(-2 - (2/3)q) + x_5((-1 + (1/3)q)).$$

Hence  $q = 3$  and  $x_5$  is the entering variable.

$$\begin{array}{rcl} x_1 & = & 8 + \beta - 2x_2 - x_4 \\ x_3 & = & -3 - \beta + x_2 + x_4 \\ x_5 & = & -8 - 2\beta + 3x_2 + 2x_4 \\ \hline z & = & 32 + 4\beta - 3x_2 - 4x_4 \end{array}$$

This is feasible and optimal for  $-8 \leq \beta \leq -4$ , and gives

$$x_1^* = 8 + \beta \quad x_2^* = 0 \quad x_3 = -3 - \beta \quad x_4 = 0 \quad x_5 = -8 - 2\beta$$

$$z = 32 + 4\beta$$



(e) What happens to this problem when  $\beta < -8$ ?

**Solution:** It suffices to see that the relevant constraint becomes

$$x_1 + 2x_2 \leq 8 + \beta < 0$$

Since  $x_1, x_2$  are non-negative, this constraint cannot be satisfied and the problem becomes infeasible.

If we instead try to apply the dual-simplex method we find that  $x_1$  becomes negative, and so the pivot row is  $0 = 8 + \beta - x_1 - 2x_2 - x_4$ . Adding  $q \times$  this row to  $z$  gives:

$$z = 32 + 4\beta + q(8 + \beta) - qx_1 + x_2(-3 - 2q) + x_4(-4 - q)$$

The coefficients of the variables stay negative for all  $q > 0$ . Hence there is no entering variable. So we cannot pivot. Hence the problem is infeasible (the dual problem is unbounded).

Do not solve the problem from scratch each time. You will find it easier to do this problem using dictionaries and the dual-simplex method.

**Solution:** In summary

$\beta < -8$	infeasible		
$-8 \leq \beta \leq -4$	$x_1^* = 8 + \beta$	$x_2^* = 0$	$z^* = 32 + 4\beta$
$-4 \leq \beta \leq -1$	$x_1^* = \frac{8 - \beta}{3}$	$x_2^* = \frac{8 + 2\beta}{3}$	$z^* = 24 + 2\beta$
$-1 \leq \beta \leq 2$	$x_1^* = 2 - \beta$	$x_2^* = 3 + \beta$	$z^* = 23 + \beta$
$2 \leq \beta \leq \infty$	$x_1^* = 0$	$x_2^* = 5$	$z^* = 25$