VinUniversity ICPC Team Notebook (2025-26)

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Mathematics

Combinatorics

```
\sum_{k=0}^{n} {n \choose k} = 2^n(x+y)^n = \sum_{k=0}^{n} {n \choose k} x^{n-k} y^k
 \sum_{k=0}^{n} \binom{k}{r} = \binom{n+1}{r+1}
(1+x)^{\alpha} = \sum_{k=0}^{\infty} {n \choose k} x^kk {n \choose k} = n {n-1 \choose k-1}
```

1.2 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
  if(x<=1) return false;</pre>
  if(x<=3) return true;</pre>
  if (!(x%2) || !(x%3)) return false;
   LL s=(LL) (sqrt((double)(x))+EPS);
  for(LL i=5; i<=s; i+=6)
    if (!(x%i) || !(x%(i+2))) return false;
  return true:
// Primes less than 1000:
                  103 107
                               109
                                     113
      157 163 167 173
                               179
                  379
                               389
                                                  409
      439 443 449 457
                               461
                                      463
                                           467
                                                  479
                                      557
      509 521 523
                         541
                               547
                                            563
                                                  569
                                                        643
727
811
      599 601 607
                         613
                               617
                                      619
                                            631
                                                  641
     661 673 677 683
751 757 761 769
829 839 853 857
                               691
773
859
                                     701
787
                                           709
797
                                                  719
809
                                                               733
                                      863
                                           877
                                                  881
                                                                     907
                                                        883
                                                                            911
// Other primes:
// 1 The largest prime smaller than 10 is 7.
// 2 The largest prime smaller than 100 is 97.
// 3 The largest prime smaller than 1000 is 997
// 4 The largest prime smaller than 10000 is 9973
// 5 The largest prime smaller than 100000 is 99991
// 6 The largest prime smaller than 1000000 is 999983
// 7 The largest prime smaller than 10000000 is 9999991
//8 The largest prime smaller than 100000000 is 99999989.
//9 The largest prime smaller than 100000000 is 999999937.
//10 The largest prime smaller than 10000000000 is 9999999967.

//11 The largest prime smaller than 100000000000 is 9999999977.

//12 The largest prime smaller than 1000000000000 is 99999999989.

//13 The largest prime smaller than 10000000000000 is 999999999971.
//14 The largest prime smaller than 1000000000000 is 9999999999973.
//16 The largest prime smaller than 100000000000000 is 99999999999937.
//17 The largest prime smaller than 1000000000000000 is 9999999999999997
// The 20 primes past 1e9+7 are.
// 1000000007 1000000009 1000000021 1000000033 1000000087 1000000093 1000000097
 // 1000000103 1000000123 1000000181 1000000207 1000000223 1000000241 1000000271
 // 1000000289 1000000297 1000000321 1000000349 1000000363 1000000403
```

Highly Composite Numbers

```
\# This program prints all hcn (highly composite numbers) <= MAXN (=10**18)
# The value of MAXN can be changed arbitrarily. When MAXN = 10**100, the
 # program needs less than one second to generate the list of hcn.
from math import log
MAXN = 10 **18
# TODO: Generates a list of the first primes (with product > MAXN).
primes = gen_primes() # primes = [2, 3, 5, 7, 11, ...]
# Generates a list of the hcn <= MAXN
def gen hcn():
     # List of (number, number of divisors, exponents of the factorization)
```

```
hen = [(1, 1, [])]
        for i in range(len(primes)):
                new_hcn = []
                for el in hcn:
                        new_hcn.append(el)
                        if len(el[2]) < i: continue</pre>
                        e_{max} = e1[2][i-1] if i >= 1 else int(log(MAXN, 2))
                        n = e1[0]
                        for e in range(1, e_max+1):
                                 n *= primes[i]
                                 if n > MAXN: break
                                div = el[1] * (e+1)
exponents = el[2] + [e]
                                new_hcn.append((n, div, exponents))
                new_hcn.sort()
                hcn = [(1, 1, [])]
                for el in new_hcn:
                        if el[1] > hcn[-1][1]: hcn.append(el)
\# Biggest HCN smaller than 10^9, 10^12, 10^18, and their number of divisors:
# 735134400
                        1344
                                   2^6*3^3*5^2*7*11*13*17
# 963761198400
                        6720
                                    2^6*3^4*5^2*7*11*13*17*19*23
# 897612484786617600
                       103680
                                   2^8*3^4*5^2*7^2*11*13*17*19*23*29*31*37
```

1.4 Number theory (modular, linear Diophantine)

```
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
       return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1;
        return ret;
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
       int xx = y = 0;
int yy = x = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q * yy; y = t;
        return a;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
       int x, y;
       VI ret;
int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                x = mod(x*(b / q), n);
                for (int i = 0; i < g; i++)
                        ret.push_back(mod(x + i*(n / g), n));
```

```
return ret;
 // computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
         int x, y;
         int g = extended_euclid(a, n, x, y);
         if (q > 1) return -1;
         return mod(x, n);
// computes x and y such that ax + by = c
// Computes x and y such that ax r by - c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
         if (!a && !b)
                  if (c) return false;
                  return true;
         if (!a)
                  if (c % b) return false;
                  x = 0; y = c / b;
                  return true:
         if (!b)
                  if (c % a) return false;
                  x = c / a; y = 0;
                  return true;
         int g = gcd(a, b);
         if (c % g) return false;
         x = c / g * mod_inverse(a / g, b / g);
         y = (c - a*x) / b;
         return true;
int main() {
         // expected: 2
         cout << gcd(14, 30) << endl;
         // expected: 2 -2 1
         int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;
         // expected: 95 451
         VI sols = modular_linear_equation_solver(14, 30, 100);
         for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
         cout << endl:
         // expected: 8
         cout << mod_inverse(8, 9) << endl;</pre>
         // expected: 23 105
                       11 12
         PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
         cout << ret.first << " " << ret.second << endl;</pre>
         ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
         cout << ret.first << " " << ret.second << endl;</pre>
          // expected: 5 -15
         if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;</pre>
         return 0;
```

1.5 Chinese Remainder Theorem

```
// Official version
// Source: https://cp-algorithms.com/math/chinese-remainder-theorem.html
struct Congruence {
    long long a, m;
};
long long chinese_remainder_theorem(vector<Congruence> const& congruences) {
    long long M = 1;
    for (auto const& congruence : congruences) {
        M *= congruence.m;
    }
    long long solution = 0;
    for (auto const& congruence : congruences) {
        long long a_i = congruence.a;
    }
```

```
long long M_i = M / congruence.m;
long long N_i = mod_inv(M_i, congruence.m);
solution = (solution + a_i * M_i % M * N_i) % M;
return solution;
```

2 Combinatorial optimization

2.1 Dinic max-flow for sparse graph

```
#include <bits/stdc++.h>
using namespace std;
const int N = 5001;
struct TEdge
    int v,rit; //rit: reverse edge
    long long cap, flow;
map<pair<int, int> , long long> ww;
int n,m;
void enter()
    cin >> n >> m;
    for (int i=0,a,b,c; i<m; i++)
        cin >> a >> b >> c;
        ww[{a,b}] += c;
        ww[{b,a}] += c;
vector< TEdge> g[N];
void init()
    int ru, rv;
   for (pair<pair<int, int>, int> p: ww)
        if (p.first.first < p.first.second)</pre>
            ru = g[p.first.first].size();
            rv = g[p.first.second].size();
            g[p.first.first].push_back({p.first.second, rv, p.second, 0});
            g[p.first.second].push_back({p.first.first, ru, p.second, 0});
int MF = 1;
int tt:
int d[N], ni[N];
bool bfs()
    for (int i=1; i<=n; i++)</pre>
       d[i] = 0;
    queue<int> qu;
    int u;
    qu.push(1);
    while (!qu.empty())
        u = qu.front();
        qu.pop();
        for (auto v : g[u])
            if (!d[v.v])
                if (v.flow + MF <= v.cap)</pre>
                    d[v.v] = d[u] + 1;
                    qu.push(v.v);
    return d[n];
long long dfs (int u, long long ff)
```

```
if (u ==n)
        return ff;
    for (;ni[u] < g[u].size(); ++ni[u])</pre>
        if (d[g[u][ni[u]].v] == d[u] + 1)
            int fff = dfs(g[u][ni[u]].v, min(g[u][ni[u]].cap - g[u][ni[u]].flow, ff));
            if (fff >= MF)
                 g[u][ni[u]].flow += fff;
                 g[g[u][ni[u]].v][g[u][ni[u]].rit].flow -= fff;\\
    return 0;
long long max_flow()
    long long res = 0,d;
    MF = 1 << 30;
    while (MF)
        while (bfs())
            for (int i=1; i<=n; i++)</pre>
                ni[i] = 0;
            do
                d = dfs(1, 1 << 30);
                 res += d;
            } while (d);
        MF >>= 1;
    return res;
int main()
    ios_base::sync_with_stdio(0);
    enter():
    init();
    cout << max_flow();</pre>
    return 0;
```

2.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
      max flow:
                          O(|V|^3) augmentations
      min cost max flow: O(|V|^4 * MAX EDGE COST) augmentations
      - graph, constructed using AddEdge()
      - source
// OUTPUT:
       - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric limits<L>::max() / 4;
struct MinCostMaxFlow {
 int N;
  VVL cap, flow, cost;
  VI found;
```

```
VL dist, pi, width;
  VPII dad;
  MinCostMaxFlow(int N) :
    N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
    this->cap[from][to] = cap;
this->cost[from][to] = cost;
  void Relax(int s, int k, L cap, L cost, int dir) {
  L val = dist[s] + pi[s] - pi[k] + cost;
  if (cap && val < dist[k]) {</pre>
      dist[k] = val;
      dad[k] = make_pair(s, dir);
      width[k] = min(cap, width[s]);
  L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0:
    width[s] = INF;
    while (s != -1) {
      int best = -1;
      found[s] = true;
      for (int k = 0; k < N; k++) {
        if (found[k]) continue;
        Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
        Relax(s, k, flow[k][s], -cost[k][s], -1);
        if (best == -1 || dist[k] < dist[best]) best = k;</pre>
      s = best;
    for (int k = 0; k < N; k++)
      pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
    L totflow = 0, totcost = 0;
    while (L amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
           totcost += amt * cost[dad[x].first][x];
          flow[x][dad[x].first] -= amt;
           totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
};
// BEGIN CUT
// The following code solves UVA problem #10594: Data Flow
int main() {
  int N, M;
  while (scanf("%d%d", &N, &M) == 2) {
    VVL v(M, VL(3));
for (int i = 0; i < M; i++)</pre>
      scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
    scanf("%Ld%Ld", &D, &K);
    MinCostMaxFlow mcmf(N+1);
    for (int i = 0; i < M; i++) {</pre>
      mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i][2]);
      mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i][2]);
    mcmf.AddEdge(0, 1, D, 0);
    pair<L, L> res = mcmf.GetMaxFlow(0, N);
    if (res.first == D) {
      printf("%Ld\n", res.second);
      printf("Impossible.\n");
```

```
return 0;
}
// END CUT
```

2.3 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// Running time:
       0(|V|^3)
       - graph, constructed using AddEdge()
// OUTPUT:
       - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std:
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
  int N = weights.size();
  VI used(N), cut, best_cut;
  int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
    VI w = weights[0]:
    VI added = used;
    int prev. last = 0:
    for (int i = 0; i < phase; i++) {</pre>
     prev = last;
       last = -1;
      for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j]; for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];
        used[last] = true;
        cut.push_back(last);
        if (best_weight == -1 || w[last] < best_weight) {</pre>
          best cut = cut;
          best_weight = w[last];
      } else {
        for (int j = 0; j < N; j++)
          w[j] += weights[last][j];
        added[last] = true;
  return make_pair(best_weight, best_cut);
// BEGIN CUT
// The following code solves UVA problem #10989: Bomb, Divide and Conquer
int main() {
 int N:
  cin >> N:
  for (int i = 0; i < N; i++) {
    int n, m;
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
      int a, b, c;
      cin >> a >> b >> c;
      weights[a-1][b-1] = weights[b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
    cout << "Case #" << i+1 << ": " << res.first << endl;
// END CUT
```

2.4 Graph cut inference

```
// Special-purpose {0,1} combinatorial optimization solver for
// problems of the following by a reduction to graph cuts:
                             sum_i psi_i(x[i])
// x[1]...x[n] in {0,1} + sum_{i < j} phi_{ij}(x[i], x[j])
// where
        psi_i : {0, 1} --> R
    phi_{ij} : {0, 1} x {0, 1} --> R
// \quad phi_{ij}(0,0) \ + \ phi_{ij}(1,1) \ <= \ phi_{ij}(0,1) \ + \ phi_{ij}(1,0) \quad (\star)
// This can also be used to solve maximization problems where the // direction of the inequality in (*) is reversed.
// INPUT: phi -- a matrix such that phi[i][j][u][v] = phi_{ij}(u, v)
           psi -- a matrix such that psi[i][u] = psi_i(u)
           x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution
// To use this code, create a GraphCutInference object, and call the
// DoInference() method. To perform maximization instead of minimization,
// ensure that #define MAXIMIZATION is enabled.
#include <vector>
#include <iostream>
using namespace std:
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef vector<VVI> VVVI;
typedef vector<VVVI> VVVVI;
const int INF = 1000000000;
// comment out following line for minimization
#define MAXIMIZATION
struct GraphCutInference {
  int N:
  VVI cap, flow;
  VI reached:
  int Augment(int s, int t, int a) {
    reached[s] = 1;
    if (s == t) return a;
    for (int k = 0; k < N; k++) {
      if (reached[k]) continue;
      if (int aa = min(a, cap[s][k] - flow[s][k])) {
         if (int b = Augment(k, t, aa)) {
           flow[s][k] += b;
           flow[k][s] = b;
           return b:
    return 0;
  int GetMaxFlow(int s, int t) {
    N = cap.size();
    flow = VVI(N, VI(N));
    reached = VI(N);
    int totflow = 0;
    while (int amt = Augment(s, t, INF)) {
      totflow += amt:
      fill(reached.begin(), reached.end(), 0);
    return totflow;
  int DoInference(const VVVVI &phi, const VVI &psi, VI &x) {
    int M = phi.size();
    cap = VVI(M+2, VI(M+2));
    VI b (M);
    int c = 0;
    for (int i = 0; i < M; i++) {
  b[i] += psi[i][1] - psi[i][0];
  c += psi[i][0];</pre>
      for (int j = 0; j < i; j++)
  b[i] += phi[i][j][1][1] - phi[i][j][0][1];
for (int j = i+i; j < M; j++) {
    cap[i][j] = phi[i][j][0][1] + phi[i][j][1][0] - phi[i][j][0][0] - phi[i][j][1][1];</pre>
         b[i] += phi[i][j][1][0] - phi[i][j][0][0];
         c += phi[i][j][0][0];
```

```
#ifdef MAXIMIZATION
    for (int i = 0; i < M; i++) {
  for (int j = i+1; j < M; j++)</pre>
        cap[i][j] *= -1;
    c *= -1;
#endif
    for (int i = 0; i < M; i++) {</pre>
      if (b[i] >= 0) {
        cap[M][i] = b[i];
      } else {
        cap[i][M+1] = -b[i];
        c += b[i];
    int score = GetMaxFlow(M, M+1);
    fill(reached.begin(), reached.end(), 0);
    Augment (M, M+1, INF);
    x = VI(M);
    for (int i = 0; i < M; i++) x[i] = reached[i] ? 0 : 1;
     score += c:
#ifdef MAXIMIZATION
    score \star = -1:
#endif
    return score:
};
int main() {
  // solver for "Cat vs. Dog" from NWERC 2008
  int numcases;
  cin >> numcases;
  for (int caseno = 0; caseno < numcases; caseno++) {</pre>
    int c, d, v;
    cin >> c >> d >> v;
    VVVVI phi(c+d, VVVI(c+d, VVI(2, VI(2))));
    VVI psi(c+d, VI(2));
for (int i = 0; i < v; i++) {</pre>
      char p, q;
      int u, v;
      cin >> p >> u >> q >> v;
      u--; v--;
if (p == 'C') {
        phi[u][c+v][0][0]++;
         phi[c+v][u][0][0]++;
      } else {
        phi[v][c+u][1][1]++;
        phi[c+u][v][1][1]++;
    GraphCutInference graph;
    cout << graph.DoInference(phi, psi, x) << endl;</pre>
  return 0;
```

3 Geometry

3.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
//
// Running time: O(n log n)
//
// INPUT: a vector of input points, unordered.
// OUTPUT: a vector of points in the convex hull, counterclockwise, starting
// with bottommost/leftmost point
#include <cstdio>
#include <cstdio>
#include <vector>
```

```
#include <algorithm>
#include <cmath>
#include <map>
using namespace std;
#define REMOVE REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
  T x, y;
PT() {}
  PT(T x, T y) : x(x), y(y) {}
  bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }
  bool operator == (const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
 return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 && (a.y-b.y)*(c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
   vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  pts = dn:
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {
   if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();</pre>
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
  pts = dn:
#endif
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
int main() {
 int t;
scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    int n;
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);
vector<PT> h(v);
    map<PT, int> index;
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h);
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {
  if (i > 0) printf(" ");
     printf("%d", index[h[i]]);
    printf("\n");
```

3.2 Miscellaneous geometry

// END CUT

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std:
double INF = 1e100;
double EPS = 1e-12;
  double x, y;
   PT() {}
  PT(double x, double y) : x(x), y(y) {}
  PT(const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y);
                                  const { return PT(x*c, y*c );
  PT operator * (double c)
  PT operator / (double c)
                                  const { return PT(x/c, y/c ); }
double dot(PT p, PT q) { return p.x*q.x*p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
    return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT (-p.y,p.x); }
PT RotateCW90(PT p)
                          { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
 // project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                             double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
 // determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
       && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
  if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
    dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
```

```
return false;
    return true:
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
 b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
 // tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++){</pre>
    int j = (i+1)%p.size();
    if ((p[i].y \le q.y \&\& q.y < p[j].y ||
      p[j].y \le q.y && q.y < p[i].y) &&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
      c = !c;
  return c:
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
     \textbf{if} \ (\texttt{dist2}(\texttt{ProjectPointSegment}(\texttt{p[i]},\ \texttt{p[(i+1)\$p.size()]},\ \texttt{q)},\ \texttt{q)} \ < \ \texttt{EPS)} 
      return true;
    return false:
\ensuremath{//} compute intersection of line through points a and b with
// circle centered at c with radius r >
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret:
  b = b-a:
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R | | d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
  the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
```

```
area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs (ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {
  int j = (i+1) % p.size();</pre>
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale:
 // tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++)</pre>
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
int l = (k+1) % p.size();
if (i == 1 || j == k) continue;
if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
         return false:
  return true:
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "
        << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "</pre>
        << ProjectPointSegment (PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
  // expected: 1 0 1
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
        << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
        << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 0 0 1
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
        << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
        << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 1 1 1 0
  cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
<< SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "</pre>
        << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
  // expected: (1,2)
  cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
  // expected: (1,1)
  cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
  v.push_back(PT(0,0));
  v.push_back(PT(5,0));
  v.push_back(PT(5,5));
  v.push_back(PT(0,5));
  // expected: 1 1 1 0 0
  cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
       << PointInPolygon(v, PT(2,0)) << " "
        << PointInPolygon(v, PT(0,2)) << " "
        << PointInPolygon(v, PT(5,2)) << " "
        << PointInPolygon(v, PT(2,5)) << endl;
  // expected: 0 1 1 1 1
```

cerr << PointOnPolygon(v, PT(2,2)) << " "

```
<< PointOnPolygon(v, PT(2,0)) << " "
      << PointOnPolygon(v, PT(0,2)) << " "
      << PointOnPolygon(v, PT(5,2)) << " "
      << PointOnPolygon(v, PT(2,5)) << endl;
               (5,4) (4,5)
               blank line
               (4,5) (5,4)
               blank line
               (4,5) (5,4)
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
 u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0); \\  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl; 
// area should be 5.0
// centroid should be (1.166666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
```

3.3 Slow Delaunay triangulation

```
// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
// Running time: O(n^4)
// INPUT:
              x[] = x-coordinates
              v[] = v-coordinates
// OUTPUT: triples = a vector containing m triples of indices
                         corresponding to triangle vertices
#include < vector >
using namespace std;
typedef double T;
struct triple {
    int i, j, k;
    triple() {}
    triple(\textbf{int } i, \textbf{ int } j, \textbf{ int } k) \ : \ i(i), \ j(j), \ k(k) \ \{\}
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
        int n = x.size();
         vector<T> z(n);
         vector<triple> ret;
        for (int i = 0; i < n; i++)
z[i] = x[i] * x[i] + y[i] * y[i];</pre>
        for (int i = 0; i < n-2; i++) {</pre>
             for (int j = i+1; j < n; j++) {
    for (int k = i+1; k < n; k++) {
                      if (j == k) continue;
                      double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
                      double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                      double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
                      bool flag = zn < 0;
                      for (int m = 0; flag && m < n; m++)</pre>
                           flag = flag && ((x[m]-x[i])*xn +
                                             (y[m]-y[i])*yn +
                                             (z[m]-z[i])*zn <= 0);
                      if (flag) ret.push_back(triple(i, j, k));
        return ret;
int main()
```

```
T xs[]={0, 0, 1, 0.9};
T ys[]={0, 1, 0, 0.9};
vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
vector<triple> tri = delaunayTriangulation(x, y);

//expected: 0 1 3
// 0 3 2

int i;
for(i = 0; i < tri.size(); i++)
    printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
return 0;</pre>
```

4 Numerical algorithms

4.1 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
    (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
             a[][] = an nxn matrix
              b[][] = an nxm matrix
                    = an nxm matrix (stored in b[][]
             A^{-1} = an nxn matrix (stored in a[][])
              returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T det = 1;
  for (int i = 0; i < n; i++) {</pre>
    int pj = -1, pk = -1;
for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
      for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
        if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
if (pj != pk) det *= -1;
irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;
    for (int p = 0; p < m; p++) b[pk][p] *= c;
    for (int p = 0; p < n; p++) if (p != pk) {
      c = a[p][pk];
      for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;</pre>
      for (int \hat{q} = 0; \hat{q} < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
```

```
return det;
int main() {
  const int n = 4;
  const int m = 2;
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
  a[i] = VT(A[i], A[i] + n);</pre>
    b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
  // expected: 60
  cout << "Determinant: " << det << endl;</pre>
  // expected: -0.233333 0.166667 0.133333 0.0666667
                  0.166667 0.166667 0.333333 -0.333333
                   0.233333 0.833333 -0.133333 -0.0666667
                  0.05 -0.75 -0.1 0.2
  cout << "Inverse: " << endl;
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)
    cout << a[i][j] << ' ';</pre>
    cout << endl;
  // expected: 1.63333 1.3
                   -0.166667 0.5
                   2.36667 1.7
                   -1.85 -1.35
  cout << "Solution: " << endl;
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < m; j++)
    cout << b[i][j] << ' ';</pre>
     cout << endl;
```

4.2 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT: a[][] = an nxm matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
             returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
 int n = a.size();
  int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m && r < n; c++) {
    int j = r;
    for (int i = r + 1; i < n; i++)
      if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s; for (int i = 0; i < n; i++) if (i != r) {
      T t = a[i][c];
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];</pre>
    r++;
```

```
return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
     { 5, 11, 10, 8},
    { 9, 7, 6, 12},
    { 4, 14, 15, 1},
    {13, 21, 21, 13}};
  VVT a(n);
for (int i = 0; i < n; i++)</pre>
    a[i] = VT(A[i], A[i] + m);
  int rank = rref(a);
  // expected: 3
  cout << "Rank: " << rank << endl;
  // expected: 1 0 0 1
                 0 1 0 3
                  0 0 1 -3
                  0 0 0 3.10862e-15
                 0 0 0 2.22045e-15
  cout << "rref: " << endl;
 cout << "rrer: " << end1;
for (int i = 0; i < 5; i++) {
  for (int j = 0; j < 4; j++)
    cout << a[i][j] << ' ';</pre>
    cout << endl:
```

4.3 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear programs of the form
         maximize c^T x
         subject to Ax <= b
                         x >= 0
// INPUT: A -- an m x n matrix
            b -- an m-dimensional vector
            c -- an n-dimensional vector
            x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
             above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9:
struct LPSolver {
  int m. n:
  VI B N:
  VVD D;
   LPSolver (const VVD &A, const VD &b, const VD &c) :
     m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
     for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] = A[i][j]; for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + 1] = b[i]; }
    for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
N[n] = -1; D[m + 1][n] = 1;</pre>
  void Pivot(int r, int s) {
    double inv = 1.0 / D[r][s];
for (int i = 0; i < m + 2; i++) if (i != r)
    for (int j = 0; j < n + 2; j++) if (j != s)
    D[i][j] -= D[r][j] * D[i][s] * inv;
    for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
    for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;</pre>
```

```
D[r][s] = inv;
    swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
    while (true) {
      int s = -1;
      for (int j = 0; j <= n; j++) {
  if (phase == 2 && N[j] == -1) continue;</pre>
        if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j] < N[s]) s = j;
      if (D[x][s] > -EPS) return true;
      int r = -1;
for (int i = 0; i < m; i++) {
        if (D[i][s] < EPS) continue;</pre>
        if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||</pre>
           (D[i][n+1] / D[i][s]) == (D[r][n+1] / D[r][s]) && B[i] < B[r]) r = i;
      if (r == -1) return false;
      Pivot(r, s);
  DOUBLE Solve (VD &x) {
    int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -numeric_limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1;
          if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s] && N[j] < N[s]) s = j;
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return D[m][n + 1];
1:
int main() {
  const int m = 4;
  const int n = 3;
  DOUBLE A[m][n] = {
    { 6, -1, 0 },
    \{-1, -5, 0\},
    { 1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE _b[m] = { 10, -4, 5, -5 };

DOUBLE _c[n] = { 1, -1, 0 };
  VVD A(m);
  VD b(\underline{b}, \underline{b} + m);
  VD c(_c, _c + n);
  for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
  LPSolver solver(A, b, c);
  DOUBLE value = solver.Solve(x);
  cerr << "VALUE: " << value << endl; // VALUE: 1.29032
  cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
  cerr << endl:
  return 0:
```

5 Graph algorithms

5.1 Fast Dijkstra's algorithm

```
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
// Running time: O(|E| log |V|)

#include <queue>
#include <cstdio>
```

```
using namespace std;
const int INF = 2000000000;
typedef pair<int, int> PII;
int main() {
        scanf("%d%d%d", &N, &s, &t);
        vector<vector<PII> > edges(N);
        for (int i = 0; i < N; i++) {</pre>
                int M;
                scanf("%d", &M);
for (int j = 0; j < M; j++) {
                        int vertex, dist;
                        scanf("%d%d", &vertex, &dist);
                         edges[i] push_back(make_pair(dist, vertex)); // note order of arguments here
        // use priority queue in which top element has the "smallest" priority
        priority_queue<PII, vector<PII>, greater<PII> > Q;
        vector<int> dist(N, INF), dad(N, -1);
        Q.push(make_pair(0, s));
        dist[s] = 0:
        while (!Q.empty()) {
                PII p = Q.top();
                Q.pop();
int here = p.second;
                if (here == t) break;
                if (dist[here] != p.first) continue;
                for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++) {
                         if (dist[here] + it->first < dist[it->second]) {
                                 dist[it->second] = dist[here] + it->first;
                                 dad[it->second] = here;
                                 Q.push(make_pair(dist[it->second], it->second));
        printf("%d\n", dist[t]);
        if (dist[t] < INF)</pre>
                for (int i = t; i != -1; i = dad[i])
                         printf("%d%c", i, (i == s ? '\n' : ' '));
        return 0;
Sample input:
5 0 4
2 1 2 3 1
2 2 4 4 5
3 1 4 3 3 4 1
2 0 1 2 3
2 1 5 2 1
Expected:
4 2 3 0
```

5.2 Strongly connected components

```
#include <vector>
#include <algorithm>
using namespace std;
vector < bool > visited; // keeps track of which vertices are already visited
// runs depth first search starting at vertex v.
// each visited vertex is appended to the output vector when dfs leaves it.
void dfs(int v, vector<vector<int> > const& adj, vector<int> &output) {
    visited[v] = true;
    for (auto u : adj[v])
       if (!visited[u])
           dfs(u, adj, output);
    output.push_back(v); // This is used to record the t_out of each vertices
// input: adj -- adjacency list of G
// output: components -- the strongy connected components in G
// output: adj_cond -- adjacency list of G^SCC (by root vertices)
void strongly_connected_components(vector<vector<int> > const& adj,
                                  vector<vector<int> > &components,
                                  vector<vector<int> > &adj cond) {
    int n = adj.size();
    components.clear(), adj_cond.clear();
```

```
vector<int> order; // will be a sorted list of G's vertices by exit time
visited.assign(n, false);
// first series of depth first searches
for (int i = 0; i < n; i++)
    if (!visited[i])
        dfs(i, adj, order);
// create adjacency list of G^T
vector<vector<int> > adj_rev(n);
for (int v = 0; v < n; v++)
    for (int u : adj[v])</pre>
        adj_rev[u].push_back(v);
visited.assign(n, false);
reverse (order.begin(), order.end());
vector<int> roots(n, 0); // gives the root vertex of a vertex's SCC
// second series of depth first searches
for (auto v : order)
    if (!visited[v]) {
        std::vector<int> component;
        dfs(v, adj_rev, component);
        components.push_back(component);
        int root = *min_element(begin(component), end(component));
        // actually, we can choose any element in the component!!!.
        for (auto u : component)
            roots[u] = root;
// add edges to condensation graph
adj_cond.assign(n, {});
for (int v = 0; v < n; v++)
    for (auto u : adj[v])
        if (roots[v] != roots[u])
            adj_cond[roots[v]].push_back(roots[u]);
```

5.3 Bridges And Articulation Points

```
// Official version
#include <bits/stdc++.h>
using namespace std;
const int maxN = 10010;
int n, m;
bool joint[maxN];
int timeDfs = 0, bridge = 0;
int low[maxN], num[maxN];
vector <int> g[maxN];
void dfs(int u, int pre) {
   int child = 0;
    num[u] = low[u] = ++timeDfs;
    for (int v : g[u]) {
        if (v == pre) continue;
         if (!num[v]) {
             dfs(v, u);
             low[u] = min(low[u], low[v]);
             if (low[v] == num[v]) bridge++;
             if (pre != -1 && low[v] >= num[u]) joint[u] = true;
        else low[u] = min(low[u], num[v]);
    if (pre == -1) {
        if (child > 1) joint[u] = true;
     cin >> n >> m;
    for (int i = 1; i <= m; i++) {</pre>
        int u, v;
        cin >> u >> v;
         g[u].push_back(v);
         g[v].push_back(u);
    for (int i = 1; i <= n; i++)
        if (!num[i]) dfs(i, -1);
    int cntJoint = 0;
```

```
for (int i = 1; i <= n; i++) cntJoint += joint[i];
cout << cntJoint << ' ' << bridge;</pre>
```

5.4 Eulerian path

```
struct Edge;
typedef list<Edge>::iterator iter;
struct Edge
        int next_vertex;
        iter reverse_edge;
        Edge(int next_vertex)
                :next_vertex(next_vertex)
1:
const int max vertices = ;
list<Edge> adj[max_vertices];
                                        // adjacency list
vector<int> path;
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                adj[v].pop_front();
                find path(vn):
        path.push_back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
```

5.5 Unweighted Bipartite Matching

```
// Max matching for unweighted bipartie graph
// Kuhn's algorithm O(n^2)
Given a **bipartite graph** SG = (X \setminus y, E). The vertices of XX are denoted x_1, x_2, \cdot x_3.
      x_m, and the vertices of $Y$ are denoted $y_1, y_2, \ldots, y_n$.
A **matching** on G is a set of edges E \subseteq E such that no two edges in E A share a common
**Requirement:** Find a **maximum matching** (having the most edges) on $G$.
* **Line 1:** Contains two integers, $m$ and $n$ ($1 \le m, n \le 100$).
* **Subsequent lines:** Each line contains two positive integers, $i$ and $j$, representing an edge $(
     x_i, y_j) \in E.
* **Line 1:** The number of edges in the maximum matching found ($K$).
* **$K$ subsequent lines: ** Each line contains two numbers, $u$ and $v$, representing the edge $(x_u,
      y_v)$ chosen for the maximum matching.
#include <bits/stdc++.h>
using namespace std;
const int N = 102:
int n, m, Assigned[N];
int Visited[N], t = 0;
vector<int> a[N];
bool visit(int u) {
```

```
if (Visited[u] != t)
         Visited[u] = t;
    else
         return false;
    for (int i = 0; i < a[u].size(); i++) {</pre>
         int v = a[u][i];
         if (!Assigned[v] || visit(Assigned[v])) {
              Assigned[v] = u;
              return true;
    return false;
int main() {
     scanf("%d%d", &m, &n);
    int x, y;
    while (scanf("%d%d", &x, &y) > 0)
         a[x].push_back(y);
    int Count = 0;
    for (int i = 1; i <= m; i++) {</pre>
         t++;
         Count += visit(i);
    printf("%d\n", Count);
    for (int i = 1; i <= n; i++)
    if (int j = Assigned[i])
        printf("%d %d\n", j, i);</pre>
```

6 Data structures

6.1 Binary Indexed Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while(x <= N) {
    tree[x] += v;
    x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
  int res = 0:
  while(x) {
    res += tree[x]:
    x -= (x & -x);
  return res:
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
  int idx = 0, mask = N;
  \textbf{while} \, (\texttt{mask \&\& idx} \, < \, \texttt{N}) \quad \{
    int t = idx + mask;
    if(x >= tree[t]) {
      idx = t:
      x -= tree[t];
    mask >>= 1;
  return idx;
```

6.2 DSU rollback

#include <bits/stdc++.h>

```
using namespace std;
using 11 = long long;
// BeginCodeSnip{DSU}
class DSU {
 private:
        vector<11> p, sz, sum;
        // stores all history info related to merges
        vector<pair<ll &, ll>> history;
        DSU(int n) : p(n), sz(n, 1), sum(n) { iota(p.begin(), p.end(), 0); }
        void init sum(const vector<11> a) {
                for (int i = 0; i < (int)a.size(); i++) { sum[i] = a[i]; }</pre>
        int get(int x) \{ return (p[x] == x) ? x : get(p[x]); \}
        11 get_sum(int x) { return sum[get(x)]; }
        void unite(int a, int b) {
                a = get(a);
                b = get(b);
                if (a == b) { return; }
                if (sz[a] < sz[b]) { swap(a, b); }</pre>
                // add to history
history.push_back({p[b], p[b]});
                history.push_back({sz[a], sz[a]});
                history.push_back({sum[a], sum[a]});
                sz[a] += sz[b];
                sum[a] += sum[b];
        void add(int x, ll v) {
                x = get(x);
                history.push_back({sum[x], sum[x]});
sum[x] += v;
        int snapshot() { return history.size(); }
        void rollback(int until) {
                while (snapshot() > until) {
                        history.back().first = history.back().second;
                        history.pop_back();
// EndCodeSnip
const int MAXN = 3e5:
DSU dsu (MAXN);
struct Ouerv {
        int t, u, v, x;
vector<Query> tree[MAXN * 4];
void update(Query &q, int v, int query_l, int query_r, int tree_l, int tree_r) {
        return:
        int m = (tree_1 + tree_r) / 2;
        update(q, v * 2, query_1, query_r, tree_1, m);
update(q, v * 2 + 1, query_1, query_r, m + 1, tree_r);
void dfs(int v, int 1, int r, vector<11> &ans) {
        int snapshot = dsu.snapshot();
        // perform all available operations upon entering
        for (Query &q : tree[v]) {
                if (q.t == 1) { dsu.unite(q.u, q.v); }
                if (q.t == 2) { dsu.add(q.v, q.x); }
        if (1 == r) {
                // answer type 3 query if we have one
                for (Query &q : tree[v]) {
                        if (q.t == 3) { ans[1] = dsu.get_sum(q.v); }
                 // go deeper into the tree
                int m = (1 + r) / 2;
                dfs(2 * v, 1, m, ans);
                dfs(2 * v + 1, m + 1, r, ans);
```

```
// undo operations upon exiting
        dsu.rollback(snapshot);
int main() {
        int n, q;
        cin >> n >> q;
        vector<11> a(n);
        for (int i = 0; i < n; i++) { cin >> a[i]; }
        dsu.init_sum(a);
        map<pair<int, int>, int> index_added;
        for (int i = 0; i < q; i++) {
                 int t:
                 cin >> t;
                 if (t == 0) {
                          int u, v;
                          if (u > v) swap(u, v);
                          // store index this edge is added, marks beginning of interval
                          index_added[{u, v}] = i;
                 } else if (t == 1) {
                          int u, v;
                          cin >> u >> v;
                          if (u > v) swap(u, v);
                          Query cur_q = \{1, u, v\};
                          // add all edges that are deleted to interval [index added, i - 1]
                          update(cur_q, 1, index_added[{u, v}], i - 1, 0, q - 1); index_added[{u, v}] = -1;
                 } else if (t == 2) {
                          int v, x;
                          cin >> v >> x;
                          Query cur_q = \{2, -1, v, x\};
                          // add all sum queries to interval [i, q-1]
                          update(cur_q, 1, i, q - 1, 0, q - 1);
                 } else if (t == 3) {
                          int v;
                          cin >> v;
                          Query cur_q = \{3, -1, v\};
                          // add all output queries to interval [i, i]
                          update(cur_q, 1, i, i, 0, q - 1);
         // add all edges that are not deleted to interval [index added, q-1]
        for (auto [edge, index] : index_added) {
                 if (index != -1) {
                          Query cur_q = {1, edge.first, edge.second};
                          update(cur_q, 1, index, q - 1, 0, q - 1);
        vector<11> ans (q, -1);
        descript disc(q, c<sub>j</sub>,
dfs(1, 0, q - 1, ans);
for (int i = 0; i < q; i++) {
    if (ans[i] != -1) { cout << ans[i] << "\n"; }</pre>
```

7 String algorithms

7.1 Suffix array

```
#include tots/stdc++.h>
using namespace std;

struct SA {
    string s;
    vector<int> p;
    int n;

SA (string s) : s(s) {
        s = s + "$";
        n = s.size();
        p.resize(n);

    for (int i=0; i<n; ++i)
        p[i] = i;

    sort (p.begin(), p.end(), [&] (int a, int b) {
        return s[a] < s[b];
    });
    vector<int> rank(n, 0);
    for (int i=0; i<n; ++i) {
</pre>
```

```
rank[i] = lower_bound(p.begin(), p.end(), i, [&] (int a, int b) {
                 return s[a] < s[b];
             }) - p.begin();
        vector<int> rank_new(n), p_new(n), cnt(n);
        for (int k=1; k<n; k*=2) {
             for (int i = 0; i < n; i++) {
                 p_new[i] = p[i] - k;
                 if (p_new[i] < 0) p_new[i] += n;</pre>
             cnt.assign(n, 0); rank_new.assign(n, 0);
             for (int i = 0; i < n; i++)
                 cnt[rank[p_new[i]]]++;
             for (int i = 1; i < n; i++)
    cnt[i] += cnt[i-1];</pre>
             for (int i = n-1; i >= 0; i--)
                p[--cnt[rank[p_new[i]]]] = p_new[i];
             rank_new[p[0]] = 0;
             int classes = 0;
             for (int i = 1; i < n; i++) {
                 pair<int, int> cur = {rank[p[i]], rank[(p[i] + k) % n]};
                 pair<int, int> prev = {rank[p[i-1]], rank[(p[i-1] + k) % n]};
                 if (cur != prev)
                     ++classes;
                 rank_new[p[i]] = classes;
             rank.swap(rank new);
};
// Input "ppppplppp" -> Output "9 5 8 4 7 3 6 2 1 0"
// "ababba" -> "6 5 0 2 4 1 3"
int main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);
    string s; cin >> s;
    SA sa(s):
    for (auto x : sa.p) cout << x << " ";
```

7.2 Knuth-Morris-Pratt

```
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respecitvely.
#include <iostream>
#include <string>
#include <vector>
using namespace std:
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
  pi = VI(p.length());
  int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != p[i])
      k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP(string& t, string& p)
  VI pi;
  buildPi(p, pi);
  int k = -1;
  for(int i = 0; i < t.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != t[i])
      k = (k == -1) ? -2 : pi[k];
    if(k == p.length() - 1) {
     // p matches t[i-m+1, ..., i]
cout << "matched at index " << i-k << ": ";</pre>
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
 return 0;
```

```
int main()
{
    string a = "AABAACAADAABAABA", b = "AABA";
    KMP(a, b); // expected matches at: 0, 9, 12
    return 0;
}
```

8 Miscellaneous

8.1 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
     INPUT: a vector of integers
    OUTPUT: a vector containing the longest increasing subsequence
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
  VI dad(v.size(), -1);
  for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASNG
    PII item = make_pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.end(), item);
    item.second = i;
#else
    PII item = make_pair(v[i], i);
VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
    if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().second);
      best.push_back(item);
      dad[i] = it == best.begin() ? -1 : prev(it)->second;
      *it = item;
  for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push back(v[i]);
  reverse(ret.begin(), ret.end());
  return ret:
```

8.2 Dates

```
// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.

#include <iostream>
#include <string>
using namespace std;

string dayofWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};

// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y){
    return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 + 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
```

```
d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
 int x, n, i, j;
  x = jd + 68569;
 n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;

x = 1461 * i / 4 - 31;
  i = 80 * x / 2447;
  d = x - 2447 * j / 80;
 x = j / 11;

m = j + 2 - 12 * x;
 y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
 return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
 int jd = dateToInt (3, 24, 2004);
  int m, d, y;
  intToDate (jd, m, d, y);
  string day = intToDay (jd);
  // expected output:
  // 2453089
// 3/24/2004
// Wed
  << day << endl;
```

8.3 Regular expressions

```
// Code which demonstrates the use of Java's regular expression libraries.
// This is a solution for
     Loglan: a logical language
     http://acm.uva.es/p/v1/134.html
// In this problem, we are given a regular language, whose rules can be
// inferred directly from the code. For each sentence in the input, we must
// determine whether the sentence matches the regular expression or not. The
// code consists of (1) building the regular expression (which is fairly
// complex) and (2) using the regex to match sentences.
import java.util.*;
import java.util.regex.*;
public class LogLan {
    public static String BuildRegex () {
         String space = " +";
         String A = "([aeiou])";
String C = "([a-z&&[^aeiou]])";
        String MOD = "(g" + A + ")";

String BA = "(b" + A + ")";

String DA = "(d" + A + ")";
         String LA = "(1" + A + ")";
         String NAM = "([a-z]*" + C + ")";
         String PREDA = "(" + C + C + A + C + A + "|" + C + A + C + C + A + ")";
         String predstring = "(" + PREDA + "(" + space + PREDA + ")*)";
        String predname = "(" + LA + space + predstring + "|" + NAM + ")";
String preds = "(" + predstring + "(" + space + A + space + predstring + ")*)";
         String predclaim = "(" + predname + space + BA + space + preds + "|" + DA + space +
         String verbpred = "(" + MOD + space + predstring + ")";
         String statement = "(" + predname + space + verbpred + space + predname + "|" + predname + space + verbpred + ")";
         String sentence = "(" + statement + "|" + predclaim + ")";
         return "^" + sentence + "$";
    public static void main (String args[]) {
         String regex = BuildRegex();
         Pattern pattern = Pattern.compile (regex);
```

```
Scanner s = new Scanner(System.in);
while (true) {
    // In this problem, each sentence consists of multiple lines, where the last
    // line is terminated by a period. The code below reads lines until
    // encountering a line whose final character is a '.'. Note the use of
          s.length() to get length of string
          s.charAt() to extract characters from a Java string
         s.trim() to remove whitespace from the beginning and end of Java string
    // Other useful String manipulation methods include
         s.compareTo(t) < 0 if s < t, lexicographically
s.indexOf("apple") returns index of first occurrence of "apple" in s</pre>
         s.lastIndexOf("apple") returns index of last occurrence of "apple" in s
         s.replace(c,d) replaces occurrences of character c with d
         s.startsWith("apple) returns (s.indexOf("apple") == 0)
          s.toLowerCase() / s.toUpperCase() returns a new lower/uppercased string
          Integer.parseInt(s) converts s to an integer (32-bit)
          Long.parseLong(s) converts s to a long (64-bit)
         Double.parseDouble(s) converts s to a double
    String sentence = "":
    while (true) {
        sentence = (sentence + " " + s.nextLine()).trim();
        if (sentence.equals("#")) return;
        if (sentence.charAt(sentence.length()-1) == '.') break;
    // now, we remove the period, and match the regular expression
    String removed_period = sentence.substring(0, sentence.length()-1).trim();
    if (pattern.matcher (removed_period).find()){
        System.out.println ("Good");
    } else {
        System.out.println ("Bad!");
```

8.4 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
     // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
    cout << 100.0/7.0 << endl;</pre>
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal cout << hex << 100 << " " << 1000 << ^{\circ} " " << 10000 << dec << endl;
```

8.5 Latitude/longitude

```
/* Converts from rectangular coordinates to latitude/longitude and vice versa. Uses degrees (not radians). \ast
```

```
#include <iostream>
#include <cmath>
using namespace std;
struct 11
  double r, lat, lon;
};
struct rect
  double x, y, z;
1:
11 convert (rect& P)
  11 Q;
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
  Q.lat = 180/M_PI*asin(P.z/Q.r);
  Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
  return 0:
rect convert(11& 0)
  rect P:
 P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
 P.z = Q.r*sin(Q.lat*M_PI/180);
int main()
  rect A:
 A.x = -1.0; A.y = 2.0; A.z = -3.0;
 B = convert(A);
cout << B.r << " " << B.lat << " " << B.lon << endl;</pre>
  A = convert(B);
  cout << A.x << " " << A.y << " " << A.z << endl;
```

8.6 Longest common subsequence

```
#include <bits/stdc++.h>
using namespace std;
int main() {
    int n; cin >> n; int m; cin >> m;
    vector<int> a(n); for (int &x:a) cin >> x;
    vector<int> b(m); for (int &x:b) cin >> x;
    vector< vector< int > > f(n+1, vector<int> (m+1, 0));
    for (int i=0; i<n; ++i) {</pre>
        for (int j=0; j<m; ++j) {
            if (a[i] == b[j]) +
                f[i+1][j+1] = f[i][j] + 1;
            } else {
                f[i+1][j+1] = max(f[i][j+1], f[i+1][j]);
    cout << f[n][m] << endl;
    int x=n,y=m; vector<int> trace;
    while (x>0&&y>0) {
        if (a[x-1] == b[y-1]) {
            trace.push_back(a[x-1]);
        } else if (f[x][y] == f[x-1][y]) x--;
        else y--;
    for (int i=trace.size()-1; i>=0; --i) cout << trace[i] << " ";</pre>
```