VinUniversity ICPC Team Notebook (2025-26)

Contents

1	Mathematics			
	1.1	Combinatorics	1	
	1.2	Prime numbers	1	
	1.3	Highly Composite Numbers	1	
	1.4	Number theory (modular, linear Diophantine)	2	
	1.5	Chinese Remainder Theorem	2	
	1.6	Discrete Log	3	
2	Geo	metry	3	
	2.1	Convex hull	3	
	2.2	Miscellaneous geometry	3	
	2.3	Slow Delaunay triangulation	5	
	2.4	Point in Polygon	5	
3	Nun	nerical algorithms	6	
	3.1	Systems of linear equations, matrix inverse, determinant	6	
	3.2	Reduced row echelon form, matrix rank	6	
	3.3	Simplex algorithm	7	
4	Gra	ph algorithms	7	
	4.1	Fast Dijkstra's algorithm	7	
	4.2	Bellman Ford's shortest path for negative cycle detection	8	
	4.3	Strongly connected components	8	
	4.4	Bridges And Articulation Points	8	
	4.5	Eulerian path	9	
5	Con	abinatorial optimization	9	
	5.1	Dinic max-flow for sparse graph	9	
	5.2	Min-cost max-flow	LO	
	5.3	Min Cost Matching	Ι1	
	5.4	Push-relabel max-flow	Ι1	
	5.5	Unweighted Bipartite Matching	12	
	5.6	Global min-cut	13	
	5.7	Graph cut inference	13	
6	Data	a structures 1	4	
	6.1	Binary Indexed Tree	14	
	6.2	DSU rollback	14	
7	Stri	ng algorithms	.5	
	7.1	Suffix array	15	
	7.2	Knuth-Morris-Pratt	15	
8	Miso	cellaneous 1	6	
	8.1	Longest increasing subsequence	16	
	8.2	Dates	16	
	8.3	C++ input/output	16	
	8.4	Latitude/longitude	١7	
	8.5	Random STL stuff	١7	
	8.6	Longest common subsequence	١7	
	8.7		١7	
	8.8		18	
	8.9	FFT	18	

1 Mathematics

1.1 Combinatorics

$$\binom{n}{k} = \frac{n!}{k!(n-k)!}$$

```
 \binom{n}{k} = \binom{n}{n-k} 
 \binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} 
 \sum_{k=0}^{r} \binom{m}{k} \binom{n}{r-k} = \binom{m+n}{r} 
 \sum_{i=r}^{n} \binom{i}{r} = \binom{n+1}{r+1} 
 \sum_{k=0}^{n} \binom{n}{k} = 2^{n} 
 (x+y)^{n} = \sum_{k=0}^{n} \binom{n}{k} x^{n-k} y^{k} 
 \sum_{k=0}^{n} \binom{k}{r} = \binom{n+1}{r+1} 
 (1+x)^{\alpha} = \sum_{k=0}^{\infty} \binom{\alpha}{k} x^{k} 
 k\binom{n}{k} = n\binom{n-1}{k-1}
```

1.2 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
   if(x<=1) return false;</pre>
   if(x<=3) return true;</pre>
   if (!(x%2) || !(x%3)) return false;
    LL s=(LL) (sqrt((double)(x))+EPS);
   for(LL i=5;i<=s;i+=6)
      if (!(x%i) || !(x%(i+2))) return false;
   return true:
 // Primes less than 1000:
                                           179
        227 229 233 239
       283 293 307 311 313 317 331 337 347 349 349 347 349 443 449 457 461 463 467 479 487 491 509 521 523 541 547 557 563 569 571 577 599 601 607 613 617 619 631 641 643 647 661 673 677 683 691 701 709 719 727 733 751 757 761 769 773 787 797 809 811 821
                                                                                               823
                                  857 859 863 877
941 947 953 967
                839
                                 857
                                                                     881
                        853
                                                                                                        911
// 1 The largest prime smaller than 10 is 7.
// 2 The largest prime smaller than 100 is 97.
// 3 The largest prime smaller than 1000 is 997.
// 4 The largest prime smaller than 10000 is 9973.
// 5 The largest prime smaller than 100000 is 99991.
// 6 The largest prime smaller than 1000000 is 999983.
// 6 The largest prime smaller than 10000000 is 9999991.

// 7 The largest prime smaller than 10000000 is 9999991.

// 8 The largest prime smaller than 10000000 is 9999998.

// 9 The largest prime smaller than 1000000000 is 999999997.

//11 The largest prime smaller than 100000000000 is 9999999997.

//12 The largest prime smaller than 100000000000 is 99999999998.

//13 The largest prime smaller than 1000000000000 is 99999999999.
//14 The largest prime smaller than 1000000000000 is 9999999999973.
//16 The largest prime smaller than 100000000000000 is 999999999999973.
//17 The largest prime smaller than 100000000000000 is 9999999999999999.
//18 The largest prime smaller than 1000000000000000 is 99999999999999999.
// The 20 primes past 1e9+7 are.
// 1000000007 1000000009 1000000021 1000000033 1000000087 1000000093 1000000097 // 1000000103 1000000123 1000000181 1000000207 1000000223 1000000241 1000000271
 // 1000000289 1000000297 1000000321 1000000349 1000000363 1000000403
```

1.3 Highly Composite Numbers

```
\# This program prints all hcn (highly composite numbers) <= MAXN (=10++18) \# The value of MAXN can be changed arbitrarily. When MAXN = 10++100, the \# program needs less than one second to generate the list of hcn.
```

```
from math import log
MAXN = 10**18
# TODO: Generates a list of the first primes (with product > MAXN).
primes = gen_primes() # primes = [2, 3, 5, 7, 11, ...]
# Generates a list of the hcn <= MAXN.
def gen_hcn():
    # List of (number, number of divisors, exponents of the factorization)
        hcn = [(1, 1, [])]
        for i in range(len(primes)):
                new\_hcn = []
                for el in hcn:
                         new_hcn.append(el)
                         if len(e1[2]) < i: continue
e_max = e1[2][i-1] if i >= 1 else int(log(MAXN, 2))
                         for e in range(1, e_max+1):
                                 n *= primes[i]
                                 if n > MAXN: break
                                 div = el[1] * (e+1)
                                 exponents = e1[2] + [e]
                                 new_hcn.append((n, div, exponents))
                new hcn.sort()
                hcn = [(1, 1, [])]
                for el in new_hcn:
     if el[1] > hcn[-1][1]: hcn.append(el)
        return hon
# Biggest HCN smaller than 10^9, 10^12, 10^18, and their number of divisors:
                         1344
                                     2^6*3^3*5^2*7*11*13*17
# 963761198400
                         6720
                                     2^6*3^4*5^2*7*11*13*17*19*23
# 897612484786617600
                                     2^8*3^4*5^2*7^2*11*13*17*19*23*29*31*37
```

1.4 Number theory (modular, linear Diophantine)

```
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a:
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1:
        return ret;
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t; 
 t = yy; yy = y - q*yy; y = t;
        return a;
```

#include <iostream>

```
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
         int x, y;
         VI ret;
         int g = extended_euclid(a, n, x, y);
         if (!(b%g)) {
                  x = mod(x*(b / g), n);
                  for (int i = 0; i < g; i++)
                           ret.push_back(mod(x + i*(n / g), n));
         return ret:
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
         int x, y;
         int g = extended_euclid(a, n, x, y);
         if (g > 1) return -1;
         return mod(x, n);
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
         if (!a && !b)
                  if (c) return false;
                  x = 0; v = 0;
                  return true;
         if (!a)
                  if (c % b) return false;
                  return true;
         if (!b)
                  if (c % a) return false;
                  x = c / a; y = 0;
                  return true:
         int g = gcd(a, b):
         if (c % g) return false;
         x = c / g * mod_inverse(a / g, b / g);
         y = (c - a*x) / b;
         return true;
int main() {
         // expected: 2
         cout << gcd(14, 30) << endl;
         // expected: 2 -2 1
        int x, y;
int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
         // expected: 95 451
         VI sols = modular_linear_equation_solver(14, 30, 100);
         for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
         cout << mod_inverse(8, 9) << endl;</pre>
         // expected: 23 105
                     11 12
        PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
cout << ret.first << " " << ret.second << endl;</pre>
         ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << end;
            expected: 5 -15
         if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;</pre>
```

1.5 Chinese Remainder Theorem

```
// Official version
// Source: https://cp-algorithms.com/math/chinese-remainder-theorem.html
struct Congruence {
    long long a, m;
};
```

```
long long chinese_remainder_theorem(vector<Congruence> const& congruences) {
   long long M = 1;
   for (auto const& congruence : congruences) {
        M *= congruence.m;
   }

   long long solution = 0;
   for (auto const& congruence : congruences) {
      long long a_i = congruence.a;
      long long M_i = M / congruence.m;
      long long N_i = mod_inv(M_i, congruence.m);
      solution = (solution + a_i * M_i % M * N_i) % M;
   }
   return solution;
```

1.6 Discrete Log

```
#include <bits/stdc++.h>
using namespace std;
// returns any x such that a^x = b \pmod{m}
// O(m^0.5) complexity
int discrete_log(int a, int b, int m) {
    assert(gcd(a, m) == 1);
    int n = (int) sqrt(m) + 1;
    for (int i = 0; i < n; ++i)
        an = ((long long) an * a) % m;
    unordered_map<int, int> vals;
    for (int i = 1, cur = an; i \le n; ++i) {
        if (!vals.count(cur))
            vals[cur] = i;
        cur = ((long long)cur * an) % m;
    for (int i = 0, cur = b; i <= n; ++i) {
        if (vals.count(cur)) {
            int res = (long long)vals[cur] * n - i;
            if (res < m)
                return res;
        cur = ((long long)cur * a) % m;
    return -1;
// usage example
int main() {
    // 2^x = 3 \pmod{5}, x = 3
    cout << discrete_log(2, 3, 5) << endl;</pre>
```

2 Geometry

2.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// fdefined.
//
// Running time: O(n log n)
//
// INPUT: a vector of input points, unordered.
// OUTPUT: a vector of points in the convex hull, counterclockwise, starting
with bottommost/leftmost point
finclude <br/>vits/stdc++.h>

using namespace std;
// Convex hull construction in O(n*log(n)): https://cp-algorithms.com/geometry/grahams-scan-convex-hull.html

struct point {
    int x, y;
}.
```

```
bool isNotRightTurn(const point &a, const point &b, const point &c) {
    \textbf{long long cross = (long long)} \; (\texttt{a.x - b.x}) \; \star \; (\texttt{c.y - b.y}) \; - \; (\texttt{long long)} \; (\texttt{a.y - b.y}) \; \star \; (\texttt{c.x - b.x});
    long long dot = (long long)(a.x - b.x) * (c.x - b.x) + (long long)(a.y - b.y) * (c.y - b.y);
    return cross < 0 || (cross == 0 && dot <= 0);
vector<point> convex_hull(vector<point> points) {
    sort (points.begin (), points.end (), [] (\textbf{auto a, auto b}) \{ \textbf{return a.} x < b.x || (a.x == b.x \&\& a.y < b.x \} \}
    int n = points.size();
    vector<point> hull;
    for (int i = 0; i < 2 * n - 1; i++) {
         while (hull.size() >= 2 && isNotRightTurn(hull.end()[-2], hull.end()[-1], points[j]))
             hull.pop_back();
         hull.push_back(points[j]);
    hull.pop_back();
    return hull;
// usage example
int main() {
    vector<point> hull1 = convex_hull({{0, 0}, {3, 0}, {0, 3}, {1, 1}});
    cout << (3 == hull1.size()) << endl:
    vector<point> hull2 = convex_hull({{0, 0}, {0, 0}});
    cout << (1 == hull2.size()) << endl;
```

2.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std:
double INF = 1e100;
double EPS = 1e-12;
  double x, y;
  PT() {}
  \texttt{PT}\,(\texttt{double}\ x,\ \texttt{double}\ y)\ :\ x\,(x)\,,\ y\,(y)\ \{\,\}
  PT(const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y);
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y);
  PT operator * (double c)
                                     const { return PT(x*c, y*c );
  PT operator / (double c)
                                    const { return PT(x/c, y/c ); }
1:
double dot(PT p, PT q)
double dist2(PT p, PT q)
double cross(PT p, PT q)

{    return p.x*q.x*p.y*q.y; }

{    return dot(p-q,p-q); }

double cross(PT p, PT q)

{    return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
    return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x);
PT RotateCW90(PT p)
                           { return PT(p.y,-p.x); ]
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
   r = dot(c-a, b-a)/r;
  if (r < 0) return a:
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
```

```
double DistancePointSegment(PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                           double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line seament from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
   if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
      return false;
    return true:
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
 // compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
  b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
   tests for checking point on polygon boundary
bool PointInPolygon (const vector <PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
      p[j].y \le q.y && q.y < p[i].y) &&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
      c = !c;
  return c:
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true;
    return false;
// compute intersection of line through points a and b with
// circle centered at c with radius r > |
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret:
  b = b-a:
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;
```

```
ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R | | d+min(r, R) < max(r, R)) return ret;
double x = (d*d-R*R+r*r)/(2*d);</pre>
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if(y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
\ensuremath{//} This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as // the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0:
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
 return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {
  int j = (i+1) % p.size();</pre>
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
 return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
int l = (k+1) % p.size();
      if (i == 1 || j == k) continue;
if (SegmentsIntersect(p[i], p[j], p[k], p[1]))
        return false:
 return true:
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5.2)
  cerr << RotateCCW(PT(2,5),M PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  << ProjectPointSegment (PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;
  // expected: 1 0 1
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
       << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
       << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
```

```
<< LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 1 1 1 0
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "</pre>
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
// expected: (1,2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
vector<PT> v:
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "
      << PointInPolygon(v, PT(2,0)) << " "
       << PointInPolygon(v, PT(0,2)) << " "
       << PointInPolygon(v, PT(5,2)) << " "
       << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "
      << PointOnPolygon(v, PT(2,0)) << " "
       << PointOnPolygon(v, PT(0,2)) << " "
       << PointOnPolygon(v, PT(5,2)) << " "
       << PointOnPolygon(v, PT(2,5)) << endl;
                  (5,4) (4,5)
                 blank line
                  (4,5) (5,4)
                 blank line
                 (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << end];
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << end];</pre>
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
 u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0); \\  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl; 
// area should be 5.0
// centroid should be (1.166666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;</pre>
return 0:
```

2.3 Slow Delaunay triangulation

```
// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
// Running time: O(n^4)
// INPUT: x[] = x-coordinates
// y[] = y-coordinates
// // OUTPUT: triples = a vector containing m triples of indices
corresponding to triangle vertices

#include<vector>
using namespace std;

typedef double T;

struct triple {
   int i, j, k;
   triple() {}
   triple(int i, int j, int k) : i(i), j(j), k(k) {}
}
```

```
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
         int n = x.size();
         vector<T> z(n);
          vector<triple> ret;
          for (int i = 0; i < n; i++)
              z[i] = x[i] * x[i] + y[i] * y[i];
         for (int i = 0; i < n-2; i++)
              for (int j = i+1; j < n; j++) {
    for (int k = i+1; k < n; k++) {
        if (j == k) continue;
    }</pre>
                        double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);

double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                        double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i])*(y[j]-y[i]);
                        bool flag = zn < 0;
                        for (int m = 0; flag && m < n; m++)</pre>
                             flag = flag && ((x[m]-x[i])*xn +
                                                 (z[m]-z[i])*zn <= 0);
                        if (flag) ret.push_back(triple(i, j, k));
         return ret;
int main()
    T xs[]={0, 0, 1, 0.9};
T ys[]={0, 1, 0, 0.9};
     vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
     vector<triple> tri = delaunayTriangulation(x, y);
     //expected: 0 1 3
     for(i = 0; i < tri.size(); i++)</pre>
         printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
     return 0:
```

2.4 Point in Polygon

```
#include <bits/stdc++.h>
using namespace std:
using 11 = long long:
 \textbf{int} \  \, \texttt{pointInPolygon}(\textbf{int} \  \, \texttt{qx, int} \  \, \texttt{qy, const} \  \, \texttt{vector} \\ < \textbf{int} > \  \, \texttt{\&x, const} \  \, \texttt{vector} \\ < \textbf{int} > \  \, \texttt{\&y)} \  \, \{ \textbf{a}, \textbf{const}, \textbf{c
                    int n = x.size();
                    int cnt = 0;
                    for (int i = 0, j = n - 1; i < n; j = i++) {
                                       return 0; // boundary
                                        if ((y[i] > qy) != (y[j] > qy)) {
                                                               ll det = ((11) \times [i] - qx) * ((11)y[j] - qy) - ((11)x[j] - qx) * ((11)y[i] - qy);
                                                             if (det == 0)
                                                                                  return 0; // boundary
                                                             if ((det > 0) != (y[j] > y[i]))
                                                                                  ++cnt;
                    return cnt % 2 == 0 ? -1 /* exterior */ : 1 /* interior */;
  // usage example
int main() {
                    vector<int> x{0, 0, 2, 2};
                     vector<int> y{0, 2, 2, 0};
                   cout << (1 == pointInPolygon(1, 1, x, y)) << endl;
cout << (0 == pointInPolygon(0, 0, x, y)) << endl;</pre>
                    cout << (-1 == pointInPolygon(0, 3, x, y)) << endl;
```

3 Numerical algorithms

3.1 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
               a[][] = an nxn matrix
               b[][] = an nxm matrix
// OUTPUT: X
                      = an nxm matrix (stored in b[][])
               A^{-1} = an nxn matrix (stored in a[][])
               returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std:
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a. VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T det = 1:
   for (int i = 0; i < n; i++) {
    for (int 1 = 0; 1 < n, 1+++, 1
int pj = -1, pk = -1;
for (int j = 0; j < n; j++) if (!ipiv[j])
for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
         if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
det *= a[pk][pk];
a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] \star = c; for (int p = 0; p < m; p++) b[pk][p] \star = c; for (int p = 0; p < n; p++) if (p != pk) {
      c = a[p][pk];
       for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;</pre>
       for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
   for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
  return det:
int main() {
  const int n = 4;
   const int m = 2;
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  for (int i = 0; i < n; i++) {
    a[i] = VT(A[i], A[i] + n),
    b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
   // expected: 60
```

```
cout << "Determinant: " << det << endl;</pre>
// expected: -0.233333 0.166667 0.133333 0.0666667
               0.166667 0.166667 0.333333 -0.333333
               0.233333 0.833333 -0.133333 -0.0666667
               0.05 -0.75 -0.1 0.2
cout << "Inverse: " << endl;</pre>
for (int i = 0; i < n; i++) {</pre>
 for (int j = 0; j < n; j++)
  cout << a[i][j] << ' ';</pre>
  cout << endl;
// expected: 1.63333 1.3
              -0.166667 0.5
               2.36667 1.7
               -1.85 -1.35
cout << "Solution: " << endl;</pre>
for (int i = 0; i < n; i++) {
 for (int j = 0; j < m; j++)
  cout << b[i][j] << ' ';</pre>
  cout << endl;</pre>
```

3.2 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT: a[][] = an nxm matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
              returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std:
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a)
 int n = a.size();
int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m && r < n; c++) {
    int j = r;
    for (int i = r + 1; i < n; i++)
    if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
for (int i = 0; i < n; i++) if (i != r) {</pre>
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];</pre>
    r++;
  return r:
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
   {16, 2, 3, 13},
    { 5, 11, 10, 8},
    { 9, 7, 6, 12},
    { 4, 14, 15, 1},
    {13, 21, 21, 13}};
  VVT a(n);
for (int i = 0; i < n; i++)</pre>
    a[i] = VT(A[i], A[i] + m);
  int rank = rref(a);
  // expected: 3
```

```
cout << "Rank: " << rank << endl;

// expected: 1 0 0 1
// 0 1 0 3
// 0 0 1 -3
// 0 0 0 3.10862e-15
// 0 0 0 2.22045e-15
cout << "rref: " << endl;
for (int i = 0; i < 5; i++) {
   for (int j = 0; j < 4; j++)
      cout << a[i][j] << ' ';
   cout << endl;
}
}</pre>
```

3.3 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear programs of the form
         maximize
         subject to Ax <= b
                         x >= 0
// INPUT: A -- an m x n matrix
            b -- an m-dimensional vector
             c -- an n-dimensional vector
             x \ensuremath{\text{--}} a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
              above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std:
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
   VI B. N:
   VVD D;
   LPSolver(const VVD &A, const VD &b, const VD &c) :
     m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
     for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] = A[i][j];

for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + 1] = b[i]; }

for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
     N[n] = -1; D[m + 1][n] = 1;
   void Pivot(int r, int s) {
     double inv = 1.0 / D[r][s];
    double inv = 1.0 / D[r][s];
for (int i = 0; i < m + 2; i++) if (i != r)
  for (int j = 0; j < n + 2; j++) if (j != s)
    D[i][j] = D[r][j] * D[i][s] * inv;
for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;</pre>
     D[r][s] = inv;
     swap(B[r], N[s]);
   bool Simplex(int phase) {
     int x = phase == 1 ? m + 1 : m;
     while (true) {
       int s = -1;
        for (int j = 0; j <= n; j++) {</pre>
          if (phase == 2 \& \& N[j] == -1) continue;
if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] \& \& N[j] < N[s]) s = j;
        if (D[x][s] > -EPS) return true;
       int r = -1;
for (int i = 0; i < m; i++) {</pre>
          if (D[i][s] < EPS) continue;
if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||</pre>
             (D[i][n+1] / D[i][s]) == (D[r][n+1] / D[r][s]) && B[i] < B[r]) r = i;
```

```
if (r == -1) return false;
      Pivot(r, s);
  DOUBLE Solve(VD &x) {
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -numeric_limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1;
        for (int j = 0; j <= n; j++)

if (s == -1 \mid \mid D[i][j] < D[i][s] \mid \mid D[i][j] == D[i][s] && N[j] < N[s]) <math>s = j;
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return D[m][n + 1];
};
int main() {
  const int m = 4:
  const int n = 3;
  DOUBLE A[m][n] = {
   { 6, -1, 0 },
    { 1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE _{b[m]} = \{ 10, -4, 5, -5 \};
  DOUBLE _{c[n]} = \{ 1, -1, 0 \};
  VVD A(m);
  VD b(_b, _b + m);
  VD c(_c, _c + n);
for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);
  LPSolver solver(A, b, c);
  DOUBLE value = solver.Solve(x);
  cerr << "VALUE: " << value << endl; // VALUE: 1.29032
  cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1</pre>
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
  cerr << endl:
  return 0:
```

4 Graph algorithms

4.1 Fast Dijkstra's algorithm

```
// use priority queue in which top element has the "smallest" priority
         priority_queue<PII, vector<PII>, greater<PII> > Q;
         vector<int> dist(N, INF), dad(N, -1);
         Q.push(make_pair(0, s));
         dist[s] = 0;
         while (!Q.empty()) {
                  PII p = Q.top();
                  Q.pop();
                  int here = p.second;
                  if (here == t) break;
                  if (dist[here] != p.first) continue;
                  for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++) {
    if (dist(here] + it->first < dist[it->second]) {
        dist[it->second] = dist[here] + it->first;
    }
                                     dad[it->second] = here;
                                     Q.push(make_pair(dist[it->second], it->second));
         printf("%d\n", dist[t]);
         if (dist[t] < INF)</pre>
                  for (int i = t; i != -1; i = dad[i])
                            printf("%d%c", i, (i == s ? '\n' : ' '));
         return 0:
Sample input:
5 0 4
3 1 4 3 3 4 1
2 1 5 2 1
Expected:
4 2 3 0
```

4.2 Bellman Ford's shortest path for negative cycle detection

```
// This function runs the Bellman-Ford algorithm for single source
// shortest paths with negative edge weights. The function returns
\ensuremath{//} false if a negative weight cycle is detected. Otherwise, the
// function returns true and dist[i] is the length of the shortest
// path from start to i.
// Running time: O(|V|^3)
     INPUT: start, w[i][j] = cost of edge from i to j
OUTPUT: dist[i] = min weight path from start to i
               prev[i] = previous node on the best path from the
                           start node
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool BellmanFord (const VVT &w, VT &dist, VI &prev, int start) {
  int n = w.size();
  prev = VI(n, -1);
  dist = VT(n, 100000000);
  dist[start] = 0;
  for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
    if (dist[j] > dist[i] + w[i][j]) {
           if (k == n-1) return false;
           dist[j] = dist[i] + w[i][j];
prev[j] = i;
```

```
}
}
return true;
```

4.3 Strongly connected components

```
#include <vector>
#include <algorithm>
using namespace std;
vector<bool> visited; // keeps track of which vertices are already visited
// runs depth first search starting at vertex \mathbf{v}.
// each visited vertex is appended to the output vector when dfs leaves it.
void dfs(int v, vector<vector<int> > const& adj, vector<int> &output) {
    visited[v] = true;
    for (auto u : adj[v])
       if (!visited[u])
           dfs(u, adj, output);
    output.push_back(v); // This is used to record the t_out of each vertices
// input: adj -- adjacency list of G
// output: components -- the strongy connected components in G
// output: adj_cond -- adjacency list of G^SCC (by root vertices)
void strongly_connected_components(vector<vector<int> > const& adj,
                                  vector<vector<int> > &components,
                                  vector<vector<int> > &adi cond) {
    int n = adj.size();
    components.clear(), adj_cond.clear();
    vector<int> order; // will be a sorted list of G's vertices by exit time
    visited.assign(n, false);
    // first series of depth first searches
    for (int i = 0; i < n; i++)
        if (!visited[i])
            dfs(i, adj, order);
    // create adjacency list of G^T
    vector<vector<int> > adj_rev(n);
    for (int v = 0; v < n; v++)
        for (int u : adj[v])
            adj_rev[u].push_back(v);
    visited.assign(n, false);
    reverse(order.begin(), order.end());
    vector<int> roots(n, 0); // gives the root vertex of a vertex's SCC
    // second series of depth first searches
    for (auto v : order)
        if (!visited[v]) {
            std::vector<int> component;
            dfs(v, adj_rev, component);
            components.push back(component);
            int root = *min element(begin(component), end(component));
            // actually, we can choose any element in the component!!!.
            for (auto u : component)
                roots[u] = root;
    // add edges to condensation graph
    adj_cond.assign(n, {});
    for (int v = 0; v < n; v++)
        for (auto u : adj[v])
            if (roots[v] != roots[u])
                adj_cond[roots[v]].push_back(roots[u]);
```

4.4 Bridges And Articulation Points

```
// Official version
#include <bits/stdc++.h>
using namespace std;
const int maxN = 10010;
```

```
int n, m;
bool joint[maxN];
int timeDfs = 0, bridge = 0;
int low[maxN], num[maxN];
vector <int> g[maxN];
void dfs(int u, int pre) {
    int child = 0;
    num[u] = low[u] = ++timeDfs;
    for (int v : g[u]) {
         if (v == pre) continue;
         if (!num[v]) {
             dfs(v, u);
low[u] = min(low[u], low[v]);
if (low[v] == num[v]) bridge++;
             if (pre != -1 && low[v] >= num[u]) joint[u] = true;
         else low[u] = min(low[u], num[v]);
    if (pre == -1) {
         if (child > 1) joint[u] = true;
int main() {
    cin >> n >> m:
    for (int i = 1; i <= m; i++) {
        int u, v;
         cin >> u >> v;
         g[u].push_back(v);
         g[v].push_back(u);
    for (int i = 1; i <= n; i++)
         if (!num[i]) dfs(i, -1);
    int cntJoint = 0;
    for (int i = 1; i <= n; i++) cntJoint += joint[i];</pre>
    cout << cntJoint << ' ' << bridge;</pre>
```

4.5 Eulerian path

```
struct Edge;
typedef list<Edge>::iterator iter;
struct Edge
        int next vertex:
        iter reverse_edge;
        Edge(int next_vertex)
                :next_vertex(next_vertex)
};
const int max_vertices = ;
int num vertices:
list<Edge> adj[max_vertices];
                                          // adjacency list
vector<int> path:
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                 adj[v].pop_front();
                find_path(vn);
        path.push_back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
```

5 Combinatorial optimization

5.1 Dinic max-flow for sparse graph

```
#include <bits/stdc++.h>
using namespace std;
const int N = 5001;
struct TEdge
    int v, rit; //rit: reverse edge
    long long cap,flow;
map<pair<int, int> , long long> ww;
int n,m;
void enter()
    cin >> n >> m;
    for (int i=0,a,b,c; i<m; i++)</pre>
        cin >> a >> b >> c;
        ww[{a,b}] += c;
        ww[{b,a}] += c;
vector< TEdge> g[N];
void init()
    int ru, rv;
    for (pair<pair<int, int>, int> p: ww)
        if (p.first.first < p.first.second)</pre>
            ru = g[p.first.first].size();
            rv = g[p.first.second].size();
            g[p.first.first].push_back({p.first.second, rv, p.second, 0});
            g[p.first.second].push_back({p.first.first, ru, p.second, 0});
int MF = 1;
int tt;
int d[N], ni[N];
bool bfs()
    for (int i=1; i<=n; i++)</pre>
       d[i] = 0;
    d[1] = 1;
    queue<int> qu;
    int u;
    qu.push(1);
    while (!qu.empty())
        u = qu.front();
        qu.pop();
        for (auto v : g[u])
            if (!d[v.v])
                if (v.flow + MF <= v.cap)</pre>
                    d[v.v] = d[u] + 1;
                    qu.push(v.v);
    return d[n];
long long dfs (int u, long long ff)
    if (u == n)
        return ff;
    for (;ni[u] < g[u].size(); ++ni[u])</pre>
        if (d[g[u][ni[u]].v] == d[u] + 1)
            int fff = dfs(g[u][ni[u]].v, min(g[u][ni[u]].cap - g[u][ni[u]].flow, ff));
            if (fff >= MF)
                g[u][ni[u]].flow += fff;
```

```
g[g[u][ni[u]].v][g[u][ni[u]].rit].flow -= fff;
                return fff;
    return 0;
long long max_flow()
    long long res = 0,d;
    MF = 1 << 30;
    while (MF)
        while (bfs())
             for (int i=1; i<=n; i++)</pre>
                ni[i] = 0;
                d = dfs(1, 1 << 30);
            } while (d);
        MF >>= 1:
    return res;
int main()
    ios_base::sync_with_stdio(0);
    init();
    cout << max_flow();</pre>
    return 0;
```

5.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
                          O(|V|^3) augmentations
      max flow:
       min cost max flow: O(|V|^4 * MAX_EDGE_COST) augmentations
       - graph, constructed using AddEdge()
      - source
      - sink
// OUTPUT:
       - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow {
  int N;
  VVL cap, flow, cost;
  VI found;
  VL dist, pi, width;
  VPII dad;
  MinCostMaxFlow(int N) :
   N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)), found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
    this->cap[from][to] = cap;
this->cost[from][to] = cost;
```

void Relax(int s, int k, L cap, L cost, int dir) { L val = dist[s] + pi[s] - pi[k] + cost;if (cap && val < dist[k]) {</pre> dist[k] = val; dad[k] = make_pair(s, dir); width[k] = min(cap, width[s]); L Dijkstra(int s, int t) { fill(found.begin(), found.end(), false); fill(dist.begin(), dist.end(), INF); fill(width.begin(), width.end(), 0); dist[s] = 0;width[s] = INF; **while** (s != -1) { int best = -1; found[s] = true; for (int k = 0; k < N; k++) { if (found[k]) continue; Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);Relax(s, k, flow[k][s], -cost[k][s], -1); if (best == -1 || dist[k] < dist[best]) best = k; s = best: for (int k = 0; k < N; k++) pi[k] = min(pi[k] + dist[k], INF);return width[t]; pair<L, L> GetMaxFlow(int s, int t) { L totflow = 0, totcost = 0; while (L amt = Dijkstra(s, t)) { totflow += amt; for (int x = t; x != s; x = dad[x].first) {
 if (dad[x].second == 1) { flow[dad[x].first][x] += amt; totcost += amt * cost[dad[x].first][x]; } else { flow[x][dad[x].first] -= amt; totcost -= amt * cost[x][dad[x].first]; return make_pair(totflow, totcost); }; // REGIN CUT // The following code solves UVA problem #10594: Data Flow int main() { int N. M: while (scanf("%d%d", &N, &M) == 2) { VVL v(M, VL(3)); for (int i = 0; i < M; i++)</pre> scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]); scanf("%Ld%Ld", &D, &K); MinCostMaxFlow mcmf(N+1); for (int i = 0; i < M; i++) {
 mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i][2]);</pre> mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i][2]); mcmf.AddEdge(0, 1, D, 0); pair<L, L> res = mcmf.GetMaxFlow(0, N); if (res.first == D) { printf("%Ld\n", res.second); printf("Impossible.\n"); return 0: // END CUT

5.3 Min Cost Matching

```
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
     cost[i][j] = cost for pairing left node i with right node j
    Lmate[i] = index of right node that left node i pairs with
    Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cstdio>
#include <cmath>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
  int n = int(cost.size()):
  // construct dual feasible solution
  VD u(n);
  VD v(n);
  for (int i = 0; i < n; i++) {</pre>
    u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
  for (int j = 0; j < n; j++)
v[j] = cost[0][j] - u[0];</pre>
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
  // construct primal solution satisfying complementary slackness
  Lmate = VI(n, -1);
  Rmate = VI(n, -1);
  int mated = 0;
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
   if (Rmate[j] != -1) continue;</pre>
      if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
        Lmate[i] = j;
        Rmate[j] = i;
        mated++;
        break:
  VD dist(n);
  VI dad(n);
  VI seen(n);
  // repeat until primal solution is feasible
  while (mated < n) {
    // find an unmatched left node
    while (Lmate[s] != -1) s++;
    // initialize Dijkstra
    fill(dad.begin(), dad.end(), -1);
    fill(seen.begin(), seen.end(), 0);
    for (int k = 0; k < n; k++)
     dist[k] = cost[s][k] - u[s] - v[k];
    int j = 0;
    while (true) {
      // find closest
      for (int k = 0; k < n; k++) {
        if (seen[k]) continue;
        if (j == -1 || dist[k] < dist[j]) j = k;</pre>
      seen[j] = 1;
      // termination condition
      if (Rmate[j] == -1) break;
```

```
// relax neighbors
    const int i = Rmate[j];
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
      if (dist[k] > new_dist) {
        dist[k] = new_dist;
        dad[k] = j;
  // update dual variables
 for (int k = 0; k < n; k++) {
   if (k == j || !seen[k]) continue;
   const int i = Rmate[k];</pre>
    v[k] += dist[k] - dist[j];
    u[i] -= dist[k] - dist[j];
  u[s] += dist[j];
  // augment along path
  while (dad[j] >= 0) {
    const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
    j = d;
  Rmate[j] = s;
  Lmate[s] = j;
  mated++;
double value = 0;
for (int i = 0; i < n; i++)
  value += cost[i][Lmate[i]];
return value;
```

5.4 Push-relabel max-flow

```
// Adjacency list implementation of FIFO push relabel maximum flow
// with the gap relabeling heuristic. This implementation is
// significantly faster than straight Ford-Fulkerson. It solves
// random problems with 10000 vertices and 1000000 edges in a few
// seconds, though it is possible to construct test cases that
// achieve the worst-case.
// Running time:
// INPUT:
       - graph, constructed using AddEdge()
      - source
      - sink
// OUTPUT:
      - maximum flow value
       - To obtain the actual flow values, look at all edges with
        capacity > 0 (zero capacity edges are residual edges).
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std;
typedef long long LL;
struct Edge {
  int from, to, cap, flow, index;
  Edge(int from, int to, int cap, int flow, int index) :
    from(from), to(to), cap(cap), flow(flow), index(index) {}
struct PushRelabel {
 int N;
  vector<vector<Edge> > G:
  vector<LL> excess;
  vector<int> dist, active, count;
  queue<int> 0:
  PushRelabel(int N): N(N), G(N), excess(N), dist(N), active(N), count(2*N) {}
  void AddEdge(int from, int to, int cap) {
```

```
G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
    G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
    if (!active[v] && excess[v] > 0) { active[v] = true; Q.push(v); }
   void Push (Edge &e) {
    int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
    if (dist[e.from] <= dist[e.to] || amt == 0) return;</pre>
    e.flow += amt;
    G[e.to][e.index].flow -= amt;
    excess[e.to] += amt;
excess[e.from] -= amt;
    Enqueue (e.to);
   void Gap(int k) {
    for (int v = 0; v < N; v++) {
      if (dist[v] < k) continue;</pre>
       count[dist[v]]--;
      dist[v] = max(dist[v], N+1);
       count[dist[v]]++;
      Enqueue (v);
  void Relabel(int v) {
    count[dist[v]]--;
    dist[v] = 2*N;
    for (int i = 0; i < G[v].size(); i++)</pre>
      if (G[v][i].cap - G[v][i].flow > 0)
         dist[v] = min(dist[v], dist[G[v][i].to] + 1);
    count[dist[v]]++;
    Enqueue (v);
  void Discharge(int v) {
   for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[v][i]);</pre>
    if (excess[v] > 0) {
   if (count[dist[v]] == 1)
         Gap(dist[v]);
       else
         Relabel(v);
   LL GetMaxFlow(int s, int t) {
    count[0] = N-1;
    count[N] = 1;
    dist[s] = N;
    active[s] = active[t] = true;
for (int i = 0; i < G[s].size(); i++) {
  excess[s] += G[s][i].cap;</pre>
      Push(G[s][i]);
    while (!Q.empty()) {
      int v = Q.front();
      Q.pop();
       active[v] = false;
      Discharge(v);
    LL totflow = 0;
    for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;
    return totflow;
};
// The following code solves SPOJ problem 4110: Fast Maximum Flow (FASTFLOW)
/* Input
The first line contains the two integers N and M. The next M lines each contain
three integers A, B, and C, denoting that there is an edge of capacity C (1 <= C <= 109) between nodes
        A and B (1 \le A, B \le N). Note that it is possible for there to be duplicate edges, as well as
       an edge from a node to itself.
Print a single integer (which may not fit into a 32-bit integer) denoting the
maximum flow / minimum cut between 1 and N.
Example
Input:
```

```
Output:
5
*/
int main() {
   int n, m;
   scanf("%d%d", &n, &m);

PushRelabel pr(n);
   for (int i = 0; i < m; i++) {
    int a, b, c,
        scanf("%d%d%d", &a, &b, &c);
    if (a == b) continue;
    pr.AddEdge(b-1, b-1, c);
   printf("%Ld\n", pr.GetMaxFlow(0, n-1));
   return 0;</pre>
```

5.5 Unweighted Bipartite Matching

```
// Max matching for unweighted bipartie graph
// Kuhn's algorithm O(n^2)
Given a **bipartite graph** $G = (X \setminus y, E)$. The vertices of $X$ are denoted $x_1, x_2, \cdot d.
      x_m, and the vertices of $Y$ are denoted $y_1, y_2, \ldots, y_n$.
A **matching** on G is a set of edges E \subseteq E such that no two edges in E share a common
**Requirement:** Find a **maximum matching** (having the most edges) on $G$.
* **Line 1:** Contains two integers, $m$ and $n$ ($1 \le m, n \le 100$).
* **Subsequent lines: ** Each line contains two positive integers, $i$ and $j$, representing an edge $(
      x_i, y_j) \in E$.
## Output
* **Line 1:** The number of edges in the maximum matching found ($K$).
* **\$K$ subsequent lines:** Each line contains two numbers, \$u$ and \$v$, representing the edge \$(x_u,
      y_v)$ chosen for the maximum matching.
#include <bits/stdc++.h>
using namespace std;
const int N = 102;
int n, m, Assigned[N];
int Visited[N], t = 0;
vector<int> a[N];
bool visit(int u) {
    if (Visited[u] != t)
        Visited[u] = t;
        return false;
    for (int i = 0; i < a[u].size(); i++) {</pre>
        int v = a[u][i];
        if (!Assigned[v] || visit(Assigned[v])) {
            Assigned[v] = u;
            return true:
    return false:
int main() {
    scanf("%d%d", &m, &n);
    int x, y;
    while (scanf("%d%d", &x, &y) > 0)
        a[x].push_back(y);
    for (int i = 1; i \le m; i++) {
        Count += visit(i);
    printf("%d\n", Count);
    for (int i = 1; i <= n; i++)
        if (int j = Assigned[i])
            printf("%d %d\n", j, i);
```

5.6 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// Running time:
       0(1V1^3)
      - graph, constructed using AddEdge()
      - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI:
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
  int N = weights.size();
  VI used(N), cut, best_cut;
  int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
    VI w = weights[0];
    VI added = used;
    int prev. last = 0:
    for (int i = 0; i < phase; i++) {</pre>
      prev = last;
       last = -1;
      for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];
for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];</pre>
        used[last] = true;
        cut.push_back(last);
        if (best_weight == -1 || w[last] < best_weight) {</pre>
          best cut = cut;
          best_weight = w[last];
      l else (
        for (int j = 0; j < N; j++)
          w[j] += weights[last][j];
        added[last] = true;
  return make_pair(best_weight, best_cut);
// The following code solves UVA problem #10989: Bomb, Divide and Conquer
int main() {
  cin >> N:
  for (int i = 0; i < N; i++) {
    int n. m:
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
      int a, b, c;
      cin >> a >> b >> c;
      weights[a-1][b-1] = weights[b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
    cout << "Case #" << i+1 << ": " << res.first << endl;
// END CUT
```

5.7 Graph cut inference

```
// Special-purpose {0,1} combinatorial optimization solver for // problems of the following by a reduction to graph cuts: // minimize sum_i psi_i(x[i]) // x[1]...x[n] in {0,1} + sum_[i < j] phi_(ij)(x[i], x[j]) // where
```

```
psi_i : {0, 1} --> R
    phi_{ij}: {0, 1} x {0, 1} --> R
    phi_{ij}(0,0) + phi_{ij}(1,1) \le phi_{ij}(0,1) + phi_{ij}(1,0)  (*)
// This can also be used to solve maximization problems where the
// direction of the inequality in (*) is reversed.
// INPUT: phi -- a matrix such that phi[i][j][u][v] = phi_{ij}(u, v)
          psi -- a matrix such that psi[i][u] = psi_i(u)
           x \ -- \ a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution
// To use this code, create a GraphCutInference object, and call the
// DoInference() method. To perform maximization instead of minimization,
// ensure that #define MAXIMIZATION is enabled.
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef vector<VVI> VVVI;
typedef vector<VVVI> VVVVI;
const int INF = 1000000000;
 // comment out following line for minimization
#define MAXIMIZATION
struct GraphCutInference {
  int N;
  VVI cap, flow;
  VI reached;
  int Augment(int s, int t, int a) {
    reached[s] = 1;
    if (s == t) return a;
    for (int k = 0; k < N; k++) {
   if (reached[k]) continue;</pre>
      if (int aa = min(a, cap[s][k] - flow[s][k])) {
   if (int b = Augment(k, t, aa)) {
          flow[s][k] += b;
           flow[k][s] -= b;
          return b;
      }
    return 0:
  int GetMaxFlow(int s. int t) {
    N = cap.size():
    flow = VVI(N, VI(N));
    reached = VI(N);
    int totflow = 0;
    while (int amt = Augment(s, t, INF)) {
      totflow += amt;
      fill(reached.begin(), reached.end(), 0);
    return totflow;
  int DoInference(const VVVVI &phi, const VVI &psi, VI &x) {
    int M = phi.size();
    cap = VVI (M+2, VI (M+2));
    VI b(M);
    int c = 0;
    for (int i = 0; i < M; i++) {
      b[i] += psi[i][1] - psi[i][0];
      c += psi[i][0];
        b[i] += phi[i][j][1][1] - phi[i][j][0][1];
      for (int j = i+1; j < M; j++) {
   cap[i][j] = phi[i][j][0][1] + phi[i][j][1][0] - phi[i][j][0][0] - phi[i][j][1][1];</pre>
        b[i] += phi[i][j][1][0] - phi[i][j][0][0];
c += phi[i][j][0][0];
#ifdef MAXIMIZATION
    for (int i = 0; i < M; i++) {
  for (int j = i+1; j < M; j++)
    cap[i][j] *= -1;</pre>
      b[i] *= -1;
```

```
c \star = -1;
#endif
    for (int i = 0; i < M; i++) {
        cap[M][i] = b[i];
      } else {
        cap[i][M+1] = -b[i];
        c += b[i];
    int score = GetMaxFlow(M, M+1);
    fill(reached.begin(), reached.end(), 0);
Augment(M, M+1, INF);
    x = VI(M);
    for (int i = 0; i < M; i++) x[i] = reached[i] ? 0 : 1;</pre>
#ifdef MAXIMIZATION
    score *= -1;
#endif
    return score:
};
int main() {
  // solver for "Cat vs. Dog" from NWERC 2008
  int numcases;
  for (int caseno = 0; caseno < numcases; caseno++) {</pre>
    cin >> c >> d >> v;
    VVVVI phi(c+d, VVVI(c+d, VVI(2, VI(2))));
    VVI psi(c+d, VI(2));
    for (int i = 0; i < v; i++) {
      char p, q;
      int u, v;
      cin >> p >> u >> q >> v;
      u--: v--:
      if (p == 'C') {
        phi[u][c+v][0][0]++;
        phi[c+v][u][0][0]++;
      } else {
        phi[v][c+u][1][1]++;
        phi[c+u][v][1][1]++;
    GraphCutInference graph;
    cout << graph.DoInference(phi, psi, x) << endl;</pre>
  return 0:
```

6 Data structures

6.1 Binary Indexed Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];
int N = (1<<LOGSZ);
// add v to value at x
void set(int x, int v) {
    while (x < N) {
        tree[x] += v;
        x += (x & -x);
    }
}
// get cumulative sum up to and including x
int get(int x) {
    int res = 0;
    while(x) {</pre>
```

```
res += tree[x];
    x -= (x & -x);
}
return res;
}
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-l and add l to result
int getind(int x) {
    int idx = 0, mask = N;
    while(mask && idx < N) {
     int t = idx + mask;
        if(x >= tree[t]) {
        idx = t;
        x -= tree[t];
    }
    mask >>= 1;
}
return idx;
}
```

6.2 DSU rollback

```
#include <bits/stdc++.h>
using namespace std;
using 11 = long long;
// BeginCodeSnip{DSU}
class DSU {
 private:
        vector<11> p, sz, sum;
// stores all history info related to merges
        vector<pair<11 &, 11>> history;
  public:
        DSU(int n) : p(n), sz(n, 1), sum(n) { iota(p.begin(), p.end(), 0); }
        void init sum(const vector<11> a) {
                for (int i = 0; i < (int)a.size(); i++) { sum[i] = a[i]; }</pre>
        int get(int x) \{ return (p[x] == x) ? x : get(p[x]); \}
        11 get_sum(int x) { return sum[get(x)]; }
        void unite(int a, int b) {
                a = get(a);
                b = get(b);
                if (a == b) { return; }
                if (sz[a] < sz[b]) { swap(a, b); }</pre>
                 // add to history
                 history.push_back({p[b], p[b]});
                 history.push_back({sz[a], sz[a]});
                history.push_back({sum[a], sum[a]});
                p[b] = a;
sz[a] += sz[b];
                 sum[a] += sum[b];
        void add(int x, 11 v) {
                 x = get(x);
                 history.push_back({sum[x], sum[x]});
        int snapshot() { return history.size(); }
        void rollback(int until) {
                while (snapshot() > until) {
                         history.back().first = history.back().second;
                         history.pop_back();
// EndCodeSnip
const int MAXN = 3e5;
DSU dsu (MAXN):
struct Ouerv {
        int t, u, v, x;
vector<Query> tree[MAXN * 4];
```

```
void update(Query &q, int v, int query_1, int query_r, int tree_1, int tree_r) {
    if (query_1 > tree_r || query_r < tree_1) { return; }
    if (query_1 <= tree_1 && query_r >= tree_r) {
                                   tree[v].push_back(q);
                                   return;
                 int m = (tree_1 + tree_r) / 2;
                  update(q, v * 2, query_1, query_r, tree_1, m);
                  update(q, v * 2 + 1, query_1, query_r, m + 1, tree_r);
\begin{tabular}{ll} \beg
                 int snapshot = dsu.snapshot();
                  // perform all available operations upon entering
                 for (Query &q : tree[v]) {
   if (q.t == 1) { dsu.unite(q.u, q.v); }
                                   if (q.t == 2) { dsu.add(q.v, q.x); }
                                    // answer type 3 query if we have one
                                   for (Query &q : tree[v]) {
                                                     if (q.t == 3) { ans[1] = dsu.get_sum(q.v); }
                 } else {
                                   // go deeper into the tree
int m = (1 + r) / 2;
                                  dfs(2 * v, 1, m, ans);
dfs(2 * v + 1, m + 1, r, ans);
                  // undo operations upon exiting
                 dsu.rollback(snapshot);
int main() {
                  int n, q;
                  cin >> n >> q;
                  vector<11> a(n);
                  for (int i = 0; i < n; i++) { cin >> a[i]; }
                 dsu.init_sum(a);
                  map<pair<int, int>, int> index_added;
                 for (int i = 0; i < q; i++) {
                                   int t:
                                    cin >> t;
                                   if (t == 0) {
                                                     int u, v;
                                                      cin >> u >> v;
                                                      if (u > v) swap(u, v);
                                                      // store index this edge is added, marks beginning of interval
                                                      index_added[{u, v}] = i;
                                   } else if (t == 1) {
                                                     int u, v;
                                                      cin >> u >> v;
                                                     if (u > v) swap(u, v);
                                                     Query cur_q = \{1, u, v\};
// add all edges that are deleted to interval [index added, i - 1]
                                                     update(cur_q, 1, index_added[{u, v}], i - 1, 0, q - 1); index_added[{u, v}] = -1;
                                   } else if (t == 2) {
                                                     int v, x;
                                                      cin >> v >> x;
                                                      Query cur_q = \{2, -1, v, x\};
                                                      // add all sum queries to interval [i, q - 1]
                                                      update(cur_q, 1, i, q - 1, 0, q - 1);
                                    } else if (t == 3) {
                                                     int v;
                                                      cin >> v;
                                                      Query cur_q = \{3, -1, v\};
                                                      // add all output queries to interval [i, i]
                                                     update(cur_q, 1, i, i, 0, q - 1);
                  // add all edges that are not deleted to interval [index added, q-1]
                  for (auto [edge, index] : index_added) {
                                   if (index != -1) {
                                                      Query cur_q = {1, edge.first, edge.second};
                                                      update(cur_q, 1, index, q - 1, 0, q - 1);
                  vector<ll> ans(q, -1);
                 dfs(1, 0, q - 1, ans);
for (int i = 0; i < q; i++) {
    if (ans[i] != -1) { cout << ans[i] << "\n"; }</pre>
```

7 String algorithms

7.1 Suffix array

```
#include <bits/stdc++.h>
using namespace std;
struct SA {
    string s;
     vector<int> p;
    int n:
    SA (string s) : s(s) {
    s = s + "$";
         n = s.size():
        p.resize(n);
         for (int i=0; i<n; ++i)</pre>
             p[i] = i;
         sort (p.begin(), p.end(), [&] (int a, int b) {
             return s[a] < s[b];
         vector<int> rank(n, 0);
         for (int i=0; i<n; ++i) {
   rank[i] = lower_bound(p.begin(), p.end(), i, [&] (int a, int b) {</pre>
                  return s[a] < s[b];</pre>
             }) - p.begin();
         vector<int> rank_new(n), p_new(n), cnt(n);
         for (int k=1; k<n; k*=2) {
             for (int i = 0; i < n; i++)
                  p_new[i] = p[i] - k;
if (p_new[i] < 0) p_new[i] += n;</pre>
              cnt.assign(n, 0); rank_new.assign(n, 0);
             for (int i = 0; i < n; i++)
                  cnt[rank[p_new[i]]]++;
             for (int i = 1; i < n; i++)
                 cnt[i] += cnt[i-1];
             for (int i = n-1; i >= 0; i--)
             p[--cnt[rank[p_new[i]]]] = p_new[i];
rank_new[p[0]] = 0;
             int classes = 0;
             for (int i = 1; i < n; i++) {
                  pair<int, int> cur = {rank[p[i]], rank[(p[i] + k) % n]};
                  pair<int, int> prev = {rank[p[i-1]], rank[(p[i-1] + k) % n]};
                  if (cur != prev)
                      ++classes;
                  rank_new[p[i]] = classes;
             rank.swap(rank_new);
    }
// Input "ppppplppp" -> Output "9 5 8 4 7 3 6 2 1 0"
// "ababba" -> "6 5 0 2 4 1 3"
int main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);
    string s; cin >> s;
     for (auto x : sa.p) cout << x << " ";
```

7.2 Knuth-Morris-Pratt

```
/*
Finds all occurrences of the pattern string p within the text string t. Running time is O(n + m), where n and m are the lengths of p and t, respecitvely.

*/
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
```

```
pi = VI(p.length());
  int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
    while (k >= -1 &  p[k+1] != p[i])
      k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP (string& t, string& p)
  buildPi(p, pi);
  int k = -1;
for(int i = 0; i < t.length(); i++) {</pre>
    while (k >= -1 &  p[k+1] != t[i])
      k = (k == -1) ? -2 : pi[k];
    if(k == p.length() - 1) {
      // p matches t[i-m+1, ..., i]
      cout << "matched at index " << i-k << ": ";
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
  return 0:
int main()
  string a = "AABAACAADAABAABA", b = "AABA";
  KMP(a, b); // expected matches at: 0, 9, 12
  return 0;
```

8 Miscellaneous

8.1 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
     INPUT: a vector of integers
    OUTPUT: a vector containing the longest increasing subsequence
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
  VI dad(v.size(), -1);
  for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASING
   PII item = make_pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.end(), item);
    item.second = i:
    PII item = make_pair(v[i], i);
    VPII::iterator it = upper_bound(best.begin(), best.end(), item);
   if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().second);
      best.push_back(item);
   } else {
     dad[i] = it == best.begin() ? -1 : prev(it)->second;
      *it = item;
  for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
```

```
reverse(ret.begin(), ret.end());
return ret;
```

8.2 Dates

```
// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.
#include <iostream>
#include <string>
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
  return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
 int x, n, i, j;
  x = jd + 68569;
  n = 4 \times x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;

x = 1461 * i / 4 - 31;
  j = 80 * x / 2447;

d = x - 2447 * j / 80;
 x = j / 11;

m = j + 2 - 12 * x;
  y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
 return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
 int jd = dateToInt (3, 24, 2004);
  int m, d, y;
 intToDate (jd, m, d, y);
string day = intToDay (jd);
  // expected output:
  // 2453089
// 3/24/2004
// Wed
 << day << endl;
```

8.3 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
int main()
{
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);

    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout < 100.0 << endl;
    cout.unsetf(ios::showpoint);</pre>
```

```
// Output a '+' before positive values
cout.setf(ios::showpos);
cout << 100 << " " << -100 << endl;
cout.unsetf(ios::showpos);

// Output numerical values in hexadecimal
cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;</pre>
```

8.4 Latitude/longitude

```
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
#include <iostream>
#include <cmath>
using namespace std;
struct 11
  double r, lat, lon;
};
struct rect
  double x, y, z;
};
11 convert (rect& P)
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
  Q.lat = 180/M_PI*asin(P.z/Q.r);
  Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
  return Q;
rect convert(11& Q)
  P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.z = Q.r*sin(Q.lat*M_PI/180);
  return P;
int main()
 11 B:
 A.x = -1.0; A.y = 2.0; A.z = -3.0;
  B = convert(A);
  cout << B.r << " " << B.lat << " " << B.lon << endl;
  A = convert(B);
cout << A.x << " " << A.y << " " << A.z << endl;
```

8.5 Random STL stuff

```
// Example for using stringstreams and next_permutation
#include <algorithm>
#include <sostream>
#include <sstream>
#include <vector>
using namespace std;
int main(void) {
    vector<int> v;

    v.push_back(1);
    v.push_back(2);
    v.push_back(3);
    v.push_back(4);
```

```
// Expected output: 1 2 3 4
                      4 3 2 1
 ostringstream oss;
oss << v[0] << " " << v[1] << " " << v[2] << " " << v[3];
  // for input from a string s,
  // istringstream iss(s);
  // iss >> variable;
  cout << oss.str() << endl;</pre>
} while (next_permutation (v.begin(), v.end()));
v.push_back(1);
v.push_back(2);
v.push_back(1);
v.push_back(3);
// To use unique, first sort numbers. Then call
// unique to place all the unique elements at the beginning
// of the vector, and then use erase to remove the duplicate
// elements.
sort(v.begin(), v.end());
v.erase(unique(v.begin(), v.end()), v.end());
// Expected output: 1 2 3
for (size_t i = 0; i < v.size(); i++)
cout << v[i] << " ";</pre>
cout << endl;</pre>
```

8.6 Longest common subsequence

```
#include <bits/stdc++.h>
using namespace std;
int main() {
    int n; cin >> n; int m; cin >> m;
    vector<int> a(n); for (int &x:a) cin >> x;
    vector<int> b(m); for (int &x:b) cin >> x;
    vector< vector< int > > f(n+1, vector<int> (m+1, 0));
    for (int i=0; i<n; ++i) {</pre>
        for (int j=0; j<m; ++j) {
            if (a[i] == b[j]) {
                f[i+1][j+1] = f[i][j] + 1;
                f[i+1][j+1] = max(f[i][j+1], f[i+1][j]);
    cout << f[n][m] << endl;
    int x=n,y=m; vector<int> trace;
    while (x>0&&v>0) {
        if (a[x-1] == b[y-1]) {
            trace.push_back(a[x-1]);
        } else if (f[x][y] == f[x-1][y]) x--;
        else y--;
    for (int i=trace.size()-1; i>=0; --i) cout << trace[i] << " ";</pre>
```

8.7 Miller-Rabin Primality Test (C)

```
while(b)
                if(b&1) ret=(ret+c)%m;
                b>>=1; c=(c+c)%m;
LL ModularExponentiation(LL a, LL n, LL m)
        LL ret=1, c=a;
        while (n)
                if(n&1) ret=ModularMultiplication(ret, c, m);
                n>>=1; c=ModularMultiplication(c, c, m);
        return ret;
bool Witness(LL a, LL n)
        LL u=n-1;
  int t=0:
        while(!(u&1)){u>>=1; t++;}
        LL x0=ModularExponentiation(a, u, n), x1;
        for (int i=1; i <=t; i++)</pre>
                x1=ModularMultiplication(x0, x0, n);
                if(x1==1 && x0!=1 && x0!=n-1) return true;
                x0=x1:
        if(x0!=1) return true;
        return false;
LL Random(LL n)
  LL ret=rand(); ret*=32768;
        ret+=rand(); ret*=32768;
        ret+=rand(); ret *= 32768;
        ret+=rand();
  return ret%n;
bool IsPrimeFast (LL n, int TRIAL)
  while (TRIAL--)
    LL a=Random(n-2)+1;
   if(Witness(a, n)) return false;
  return true;
```

8.8 Super Duper Fast IO

```
#include <memory.h>
#include <cstdio>
const int BUF SIZE = 65536;
char input[BUF_SIZE];
struct scanner {
    char* curPos;
        fread(input, 1, sizeof(input), stdin);
        curPos = input;
    void ensureCapacity() {
   int size = input + BUF_SIZE - curPos;
        if (size < 100) {
            memcpy(input, curPos, size);
            fread(input + size, 1, BUF_SIZE - size, stdin);
            curPos = input;
    int nextInt() {
        ensureCapacity();
        while (*curPos <= ' ')</pre>
            ++curPos:
        bool sign = false;
        if (*curPos == '-') {
            sign = true;
            ++curPos;
        int res = 0;
        while (*curPos > ' ')
```

8.9 FFT

```
// Convolution using the fast Fourier transform (FFT).
// INPUT:
      a[1...n]
      b[1...m]
// OUTPUT:
      c[1...n+m-1] such that c[k] = sum_{i=0}^k a[i] b[k-i]
// Alternatively, you can use the DFT() routine directly, which will
// zero-pad your input to the next largest power of 2 and compute the
#include <iostream>
#include <vector>
#include <complex>
using namespace std:
typedef long double DOUBLE;
typedef complex<DOUBLE> COMPLEX;
typedef vector<DOUBLE> VD;
typedef vector<COMPLEX> VC;
struct FFT {
  VC A;
  int ReverseBits(int k) {
    int ret = 0;
    for (int i = 0; i < L; i++) {
     ret = (ret << 1) | (k & 1);
     k >>= 1:
    return ret;
  void BitReverseCopy(VC a) {
    for (n = 1, L = 0; n < a.size(); n <<= 1, L++);
    A.resize(n);
    for (int k = 0; k < n; k++)
     A[ReverseBits(k)] = a[k];
  VC DFT(VC a, bool inverse) {
   BitReverseCopy(a);
for (int s = 1; s <= L; s++) {</pre>
      int m = 1 << s;
      COMPLEX wm = \exp(\text{COMPLEX}(0, 2.0 * M_PI / m));
      if (inverse) wm = COMPLEX(1, 0) / wm;
      for (int k = 0; k < n; k += m) {
        COMPLEX w = 1;
        for (int j = 0; j < m/2; j++) {
         COMPLEX t = w * A[k + j + m/2];
         COMPLEX u = A[k + j];
          A[k + j] = u + t;
         A[k + j + m/2] = u - t;
          w = w * wm;
    if (inverse) for (int i = 0; i < n; i++) A[i] /= n;</pre>
    return A;
  // c[k] = sum_{i=0}^k a[i] b[k-i]
```

```
VD Convolution(VD a, VD b) {
   int L = 1;
   while ((1 << L) < a.size()) L++;
   while ((1 << L) < b.size()) L++;
   int n = 1 << (L+1);

VC aa, bb;
   for (size_t i = 0; i < n; i++) aa.push_back(i < a.size() ? COMPLEX(a[i], 0) : 0);
   for (size_t i = 0; i < n; i++) bb.push_back(i < b.size() ? COMPLEX(b[i], 0) : 0);

VC AA = DFT(aa, false);
   VC BB = DFT(bb, false);
   VC CC;
   for (size_t i = 0; i < AA.size(); i++) CC.push_back(AA[i] * BB[i]);
   VC cc = DFT(CC, true);

VD c;
   for (int i = 0; i < a.size() + b.size() - 1; i++) c.push_back(cc[i].real());
   return c;
}</pre>
```

```
};
int main() {
    double a[] = {1, 3, 4, 5, 7};
    double b[] = {2, 4, 6};

FFT fft;
VD c = fft.Convolution(VD(a, a + 5), VD(b, b + 3));

// expected output: 2 10 26 44 58 58 42
for (int i = 0; i < c.size(); i++) cerr << c[i] << " ";
    cerr << endl;
    return 0;
}</pre>
```