

# Principles of Programming

(4190.210)

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# Syllabus

## ➤ Lecture

- Tue & Thu: 14:00 ~ 15:15 (@302-107)
- <https://github.com/snu-sf-class/pp202302>
- Bring your laptop to lectures

## ➤ Instructor

- Chung-Kil Hur
- <http://sf.snu.ac.kr/gil.hur/>

## ➤ Teaching Assistant

- Hoyoung Joo & Jaehyung Lee
- Contact TAs at [pp@sf.snu.ac.kr](mailto:pp@sf.snu.ac.kr)

## ➤ Grading (tentative)

- Attendance: 5%
- Assignments: 25%
- Midterm exam: 30%
- Final exam: 40%

# Who am I?

➤ Prof. Chung-Kil (Gil) Hur [허충길]

- Education: KAIST (B.S.), Univ of Cambridge (Ph.D.)
- (High school) Bronze medal in IMO 1994 @Hongkong.
- Software Foundations Lab  
<http://sf.snu.ac.kr>
- Research Topics
  - Software Verification
  - Low-level Language Semantics (C/C++/LLVM/Rust)
  - Relaxed-Memory Concurrency
- Our collaborators
  - Many in USA, UK, Germany, France, Israel.
- Publications in Programming Languages (PL) area
  - POPL and PLDI are the two most prestigious conferences in PL area (see <https://csrankings.org>)
  - Published 12 PLDI and 5 POPL papers for the last 10 years (6<sup>th</sup> in the world)
  - Published either POPL or PLDI paper for 13 years in a row (2<sup>nd</sup> in the world)
  - Received **distinguished paper awards** from PLDI 2017, 2021 & POPL 2020

# Introduction

# Overview

- Part 1  
Functional Programming (with Scala)
- Part 2  
Object-Oriented Programming (with Scala)
- Part 3  
Type Class (Type-Oriented) Programming (with Scala)
- Part 4  
Imperative Programming (with Scala & Rust)

# Imperative vs. Functional Programming

## ➤ Imperative Programming

- Computation by memory reads/writes
- Sequence of read/write operations
- Repetition by loop
- More procedural (ie, describe how to do)
- Easier to write efficient code

```
sum = 0;
i = n;
while (i > 0) {
    sum = sum + i;
    i = i - 1;
}
```

## ➤ Functional Programming

- Computation by function application
- Composition of function applications
- Repetition by recursion
- More declarative (ie, describe what to do)
- Easier to write safe code

```
def sum(n: Int): Int
  if (n <= 0)
    0
  else
    n + sum(n-1)
```

# Both Imperative & Functional Style Supported

- Many languages support both imperative & functional style
  - More imperative: Java, Javascript, C++, Python, Rust, ...
  - More functional: OCaml, SML, Lisp, Scheme, ...
  - Middle: Scala
  - Purely functional: Haskell (monadic programming)

# Object-Oriented vs. Type-Oriented Programming

## ➤ Goal

- To support module systems for writing large software
- Specifically, to support code abstraction & reuse

## ➤ Object-Oriented Programming

- Class Instance (ie, Object): data + methods (ie, functions)
- Code abstraction & reuse together via inheritance

## ➤ Type Class (Type-Oriented) Programming

- Type Class Instance: type + methods (ie, functions)
- Abstraction and Reuse are separated
  - Code abstraction via type class instantiation
  - Code reuse via composition



# Why Scala & Rust?

## ➤ Why Scala?

- Equally well support both imperative & functional style
- Many advanced features (both OOP & Type class supported)
- Compatible with Java
- Memory management with Garbage Collection (automatic, dynamic)

## ➤ Why Rust?

- Can write (low-level) efficient code like C and C++ but safely
- Using ownership types
- Memory management with ownership types (automatic, static) + lightweight Garbage Collection

# PART 1

## Functional Programming with Function Applications

# Values, Names, Functions and Evaluations

# Values, Expressions, Names

## ➤ Types and Values

- A type is a set of values
- `Int`:  $\{-2147483648, \dots, -1, 0, 1, \dots, 2147483647\}$  //32-bit integers
- `Double`: 64-bit floating point numbers // real numbers in practice
- `Boolean`:  $\{\text{true}, \text{false}\}$
- ...

## ➤ Expressions

- Composition of  
values, names, primitive operations, ...

## ➤ Name Binding

- Binding expressions to names

## ➤ Examples

```
def a = 1 + (2 + 3)
def b = 3 + a * 4
```

# Evaluation

## ➤ Evaluation

- Reducing an expression into a value
- Strategy
  1. Take a name or an operator (outer to inner)
  2. (name) Replace the name with its associated expression
  3. (name) Evaluate the expression
  4. (operator) Evaluate its operands (left to right)
  5. (operator) Apply the operator to its operands

## ➤ Examples

$5+b \sim 5+(3+(a*4)) \sim 5+(3+(1+(2+3))*4) \sim \dots \sim 32$

# Functions and Substitution

## ➤ Functions

- Expressions with Parameters
- Binding functions to names

```
def f(x: Int): Int = x + a
```

## ➤ Evaluation by substitution

- ...
- (function) Evaluate its operands (left to right)
- (function)  
Replace the function application by the expression of the function  
Replace its parameters with the operands

$$\begin{aligned} 5+f(f(3)+1) &\sim 5+f((3+a)+1) \sim \dots \sim 5+f(10) \sim 5+(10+a) \\ &\sim \dots \sim 21 \end{aligned}$$

# Simple Recursion

## ➤ Recursion

- Use X in the definition of X
- Powerful mechanism for repetition
- Nothing special but just rewriting

```
def sum(n: Int) : Int =  
  if (n <= 0)  
    0  
  else  
    n + sum(n-1)
```

```
sum(2) ~ if (2<=0) 0 else (2+sum(2-1)) ~  
2+sum(1) ~ 2+(if (1<=0) 0 else (1+sum(1-1))) ~  
2+(1+sum(0)) ~ 2+(1+(if (0<=0) 0 else (0+sum(0-1))))  
~ 2+(1+0) ~ 3
```

# Termination/Divergence

Evaluation may not terminate

## ➤ Termination

- An expression may reduce to a value

## ➤ Divergence

- An expression may reduce forever

```
def loop: Int = loop
```

```
loop ~ loop ~ loop ~ ...
```



# Evaluation strategy: Call-by-value, Call-by-name

$f(e1, e2)$

## ➤ Call-by-value

- Evaluate the arguments first, then apply the function to them

## ➤ Call-by-name

- Just apply the function to its arguments, without evaluating them.

```
def square (x: Int) = x * x
```

```
[cbv]square(1+1) ~ square(2) ~ 2*2 ~ 4
```

```
[cbn]square(1+1) ~ (1+1)*(1+1) ~ 2*(1+1) ~ 2*2 ~ 4
```

# CBV, CBN: Differences

## ➤ Call-by-value

- Evaluates arguments once

## ➤ Call-by-name

- Do not evaluate unused arguments

## ➤ Question

- Do both always result in the same value?

# Scala's evaluation strategy

## ➤ Call-by-value

- By default

## ➤ Call-by-name

- Use “=>”

```
def one(x: Int, y: =>Int) = 1
```

```
one(1+2, loop)
```

```
one(loop, 1+2)
```

# Scala's name binding strategy

## ➤ Call-by-value

- Use “val” (also called “field”) e.g. `val x = e`
- Evaluate the expression first, then bind the name to it

## ➤ Call-by-name

- Use “def” (also called “method”) e.g. `def x = e`
- Just bind the name to the expression, without evaluating it
- Mostly used to define functions

```
def a = 1 + 2 + 3
```

```
val a = 1 + 2 + 3 // 6
```

```
def b = loop
```

```
val b = loop
```

```
def f(a: Int, b: Int): Int = a*b - 2
```

# Conditional Expressions

## ➤ If-else

- `if (b) e1 else e2`
- `b` : Boolean expression
- `e1, e2`: expressions of the same type

## ➤ Rewrite rules:

- `if (true) e1 else e2 → e1`
- `if (false) e1 else e2 → e2`

```
def abs(x: Int) = if (x >= 0) x else -x
```

# Boolean Expressions

## ➤ Boolean expression

- true, false
- !b
- b && b
- b || b
- e <= e, e >= e, e < e, e > e, e == e, e != e

## ➤ Rewrite rules:

- !true → false
- !false → true
- true && b → b
- false && b → false
- true || b → true
- false || b → b

true && (loop == 1) ~ loop == 1 ~ loop == 1

## Exercise: and, or

➤ Write two functions

- `and(x,y) == x && y`
- `or(x,y) == x || y`
- Do not use `&&`, `||`

```
and(false,loop==1)
```

```
~ if (false) loop==1 else false
```

```
~ false
```

```
and(true,loop==1)
```

```
~ if (true) loop==1 else false
```

```
~ loop==1 ~ loop==1 ...
```

# Solution

```
def and(x: Boolean, y: =>Boolean) =  
  if (x) y else false
```

```
def or(x: Boolean, y: =>Boolean) =  
  if (x) true else y
```



## Exercise: square root calculation

➤ Calculate square roots with Newton's method

```
def isGoodEnough(guess: Double, x: Double) =  
    ??? // guess*guess is 99.9% close to x
```

```
def improve(guess: Double, x: Double) =  
    (guess + x/guess) / 2
```

```
def sqrtIter(guess: Double, x: Double): Double =  
    ??? // repeat improving guess until it is good enough
```

```
def sqrt(x: Double) =  
    sqrtIter(1, x)
```

```
sqrt(2)
```

# Solution

➤ Calculate square roots with Newton's method

```
def isGoodEnough(guess: Double, x: Double) =  
    guess*guess/x > 0.999 && guess*guess/x < 1.001
```

```
def improve(guess: Double, x: Double) =  
    (guess + x/guess) / 2
```

```
def sqrtIter(guess: Double, x: Double): Double =  
    if (isGoodEnough(guess, x)) guess  
    else sqrtIter(improve(guess, x), x)
```

```
def sqrt(x: Double) =  
    sqrtIter(1, x)
```

```
sqrt(2)
```

# Blocks and Name Scoping

# Blocks in Scala

## ➤ Block

- { **val** x1 = e1  
    **def** x2 = e2  
    e  
}
- Is an expression
- Allow nested name binding
- Allow arbitrary order of “**def**”s, but not “**val**”s (think about why)

# Scope of names

## ➤ Block

```
{ val t = 0
  def f(x: Int) = t + g(x+1)
  def g(y: Int) = y * y
  val x = f(5)
  val r = {
    val t = 10
    val s = f(5)
    s - t
  }
  t + r }
```

- A definition inside a block is only accessible within the block
- A definition inside a block shadows definitions of the same name outside the block
- A definition is accessible inside a nested block unless it is shadowed
- The same name cannot be defined twice in the same block
- A function is evaluated under the environment where it is defined, not the environment where it is invoked.

# Problem

// should be allowed

{

def f(x: Int) = g(x)

def g(x: Int) = 10

val x = f(10)

x

}

// should not be allowed

{

def f(x: Int) = g(x)

val x = f(10)

def g(x: Int) = 10

x

}

# Safety Checking

Safety checking for ruling out accessing undefined values is simple.

- For “val x = e”,  
all names in “e” should be defined before this definition.
- For “def x = e”,  
all names in “e” should be defined before the next “val” definition.

/\* The following code passes the safety checking \*/

```
{ def f(x: Int) = g(x)
  def g(x: Int) = 10
  val a = 10
  f(10) }
```

/\* The following code fails at the safety checking \*/

```
{ def f(x: Int) = g(x)
  val a = 10
  def g(x: Int) = 10
  f(10) }
```

# Rewriting for Blocks

```
1: { val t = 0
2:   def f(x: Int) = t + g(x+1)
3:   def g(y: Int) = y*y
4:   val x = f(5)
5:   val r = {
6:     val t = 10
7:     val s = f(5)
8:     s - t
9:   t + r }
```

## ➤ Evaluation with Environment

$[E0 \mid t = \_, f = (x) t + g(x), g = (y) y * y, x = \_, r = \_], 1 \sim \emptyset$

$[E0 \mid t = 0, f = (x) t + g(x), g = (y) y * y, x = \_, r = \_], 4 \sim f(5) \sim 36$

$f(5): E0 :: [E1 \mid x = 5], t + g(x+1) \sim 0 + g(6) \sim 36$

$g(6): E0 :: [E2 \mid y = 6], y * y \sim 6 * 6 \sim 36$

$[E0 \mid t = 0, f = (x) t + g(x), g = (y) y * y, x = 36, r = \_], 5 \sim 26$

$5: E0 :: [E3 \mid t = \_, s = \_], 6 \sim 10$

$E0 :: [E3 \mid t = 10, s = \_], 7 \sim f(5) \sim 36$

$f(5): E0 :: [E4 \mid x = 5], t + g(x+1) \sim 0 + g(6) \sim 36$

$g(6): E0 :: [E5 \mid y = 6], y * y \sim 6 * 6 \sim 36$

$E0 :: [E3 \mid t = 10, s = 36], 8 \sim s - t \sim 26$

$[E0 \mid t = 0, f = (x) t + g(x), g = (y) y * y, x = 36, r = 26], 9 \sim t + r \sim 26$



## Example: def with no arguments

```
1: { val t = 0
2:   def x = t+t //basically treated same as def x() = t+t
3:   val r = {
4:     val t = 10
5:     x+t }
6:   x+t }
```

### ➤ Evaluation with Environment

$[E0 \mid t=_, x=t+t, r=_], 1 \sim 0$

$[E0 \mid t=0, x=t+t, r=_], 3 \sim 10$

3:  $E0 :: [E1 \mid t=_, 4 \sim 10$

$E0 :: [E1 \mid t=10], 5 \sim x+t \sim 0+10 \sim 10$

x:  $E0, t+t \sim 0+0 \sim 0$

$[E0 \mid t=0, x=t+t, r=10], 6 \sim x+t \sim 0+0 \sim 0$

x:  $E0, t+t \sim 0+0 \sim 0$

# Semi-colons and Parenthesis

## ➤Block

- Can write two definitions/expressions in a single line using ;
- Can write one definition/expression in two lines using (), but can omit () when clear

*// ok*

```
val r = {  
    val t = 10; val s = square(5); t +  
    s }
```

*// Not ok*

```
val r = {  
    val t = 10; val s = square(5); t  
    + s }
```

*// ok*

```
val r = {  
    val t = 10; val s = square(5); (t  
    + s) }
```

# Exercise: Writing Better Code using Blocks

➤ Make the following code better

```
def isGoodEnough(guess: Double, x: Double) =  
    guess*guess/x > 0.999 && guess*guess/x < 1.001
```

```
def improve(guess: Double, x: Double) =  
    (guess + x/guess) / 2
```

```
def sqrtIter(guess: Double, x: Double): Double = {  
    if (isGoodEnough(guess,x)) guess  
    else sqrtIter(improve(guess,x),x)  
}
```

```
def sqrt(x: Double) =  
    sqrtIter(1, x)
```

```
sqrt(2)
```

# Solution

```
def sqrt(x: Double) = {  
  def sqrtIter(guess: Double, x: Double): Double = {  
    if (isGoodEnough(guess, x)) guess  
    else sqrtIter(improve(guess, x), x)  
  }  
  def isGoodEnough(guess: Double, x: Double) = {  
    val ratio = guess * guess / x  
    ratio > 0.999 && ratio < 1.001  
  }  
  def improve(guess: Double, x: Double) =  
    (guess + x/guess) / 2  
  
  sqrtIter(1, x)  
}
```

sqrt(2)

# Lazy Call-By-Value

# Lazy call-by-value

## ➤ Lazy call-by-value

- Use “lazy val” e.g. `lazy val x = e`
- Evaluate the expression **first time it is used**, then bind the name to it

```
def f(c: Boolean, i: =>Int): Int = {  
  lazy val iv = i  
  if (c) 0  
  else iv * iv * iv  
}
```

```
f(true, {println("ok"); 100+100+100+100})  
f(false, {println("ok"); 100+100+100+100})
```

# Tail Recursion & Tail Call

# Recursion needs care

## ➤ Summation function

- Write a summation function `sum` such that  
 $\text{sum}(n) = 1+2+\dots+n$
- Test  
`sum(10), sum(100), sum(1000), sum(10000),`  
`sum(100000), sum(1000000)`
- What's wrong? (Think about evaluation)



# Recursion: Try

```
def sum(n: Int): Int =  
  if (n <= 0) 0 else (n+sum(n-1))
```

# Recursion: Tail Recursion

```
import scala.annotation.tailrec
```

```
def sum(n: Int): Int = {  
    @tailrec def sumItr(res: Int, m: Int): Int =  
        if (m <= 0) res else sumItr(m+res,m-1)  
    sumItr(0,n)  
}
```

# Mutual Recursion: Try

```
{  
  def sum(acc: Int, n: Int): Int =  
    if (n <= 0) acc else sum2(n + acc, n-1)  
  
  def sum2(acc: Int, n: Int): Int =  
    if (n <= 0) acc else sum(2*n + acc, n-1)  
  
  sum(0, 20000) // stack overflow  
}
```

# Mutual Recursion: Tail Call Optimization

```
import scala.util.control.TailCalls._

{
  def sum(acc: Int, n: Int): TailRec[Int] =
    if (n <= 0) done(acc) else tailcall(sum2(n + acc, n-1))

  def sum2(acc: Int, n: Int): TailRec[Int] =
    if (n <= 0) done(acc) else tailcall(sum(2*n + acc, n-1))

  sum(0, 20000).result
}
```

# Higher-Order Functions

# Functions as Values

## ➤ Functions

- Functions are normal values of function types  $(A_1, \dots, A_n \Rightarrow B)$ .
- They can be copied, passed and returned.
- Functions that take functions as arguments or return functions are called higher-order functions.
- Higher-order functions increase code reusability.

# Examples

```
def sumLinear(n: Int): Int =  
  if (n <= 0) 0 else n + sumLinear(n-1)
```

```
def sumSquare(n: Int): Int =  
  if (n <= 0) 0 else n*n + sumSquare(n-1)
```

```
def sumCubes(n: Int): Int =  
  if (n <= 0) 0 else n*n*n + sumCubes(n-1)
```

Q: How to write reusable code?

# Examples

```
def sum(f: Int=>Int, n: Int): Int =  
  if (n <= 0) 0 else f(n) + sum(f, n-1)
```

```
def linear(n: Int) = n  
def square(n: Int) = n * n  
def cube(n: Int) = n * n * n
```

```
def sumLinear(n: Int) = sum(linear, n)  
def sumSquare(n: Int) = sum(square, n)  
def sumCubes(n: Int) = sum(cube, n)
```



# Anonymous Functions

## ➤ Anonymous Functions

- Syntax

$(x_1: T_1, \dots, x_n: T_n) \Rightarrow e$

or

$(x_1, \dots, x_n) \Rightarrow e$

```
def sumLinear(n: Int) = sum((x: Int) => x, n)
def sumSquare(n: Int) = sum((x: Int) => x * x, n)
def sumCubes(n: Int) = sum((x: Int) => x * x * x, n)
```

Or simply

```
def sumLinear(n: Int) = sum((x) => x, n)
def sumSquare(n: Int) = sum((x) => x * x, n)
def sumCubes(n: Int) = sum((x) => x * x * x, n)
```

# Exercise

```
def sum(f: Int=>Int, a: Int, b: Int): Int =  
  if (a <= b) f(a) + sum(f, a+1, b) else 0
```

```
def product(f: Int=>Int, a: Int, b: Int): Int =  
  if (a <= b) f(a) * product(f, a+1, b) else 1
```

DRY (Do not Repeat Yourself) using a higher-order function, called “mapReduce”.

# Exercise

```
def mapReduce(reduce: (Int,Int)=>Int, inival: Int,  
              f: Int=>Int, a: Int, b: Int): Int = {  
  if (a <= b) reduce(f(a),mapReduce(reduce,inival,f,a+1,b))  
  else inival  
}
```

```
def sum(f: Int=>Int, a: Int, b: Int): Int =  
  mapReduce((x,y)=>x+y,0,f,a,b)
```

```
def product(f: Int=>Int, a: Int, b: Int): Int =  
  mapReduce((x,y)=>x*y,1,f,a,b)
```

# Evaluation for functional values

```
1: { val t = 0
2:   val f : Int=>Int = {    // Note: What if using “def f” ?
3:     val t = 10
4:     def g(x: Int) : Int = x + t
5:     g _ }
6:   f(20) }
```

\* Try: Evaluation without Closures

$[E0 \mid t=_, f=_], 1 \sim \emptyset$

$[E0 \mid t=0, f=_], 2 \sim (x)x+t$

$E0 :: [E1 \mid t=_, g=(x)x+t], 3 \sim 10$

$E0 :: [E1 \mid t=10, g=(x)x+t], 5 \sim (x)x+t$

$[E0 \mid t=0, f=(x)x+t], 6 \sim f(20) \sim 20$

$f(20): E0 :: [E2 \mid x=20], x+t \sim 20+0 \sim 20$

# Closures for functional values

```
1: { val t = 0
2:   val f : Int=>Int = {    // Note: What if using “def f” ?
3:     val t = 10
4:     def g(x: Int) : Int = x + t
5:     g _ }
6:   f(20) }
```

## \* Evaluation with Closures

$[E0 \mid t=_, f=_], 1 \sim \theta$

$[E0 \mid t=0, f=_], 2 \sim (E1, (x)x+t)$

$E0 :: [E1 \mid t=_, g=(x)x+t], 3 \sim 10$

$E0 :: [E1 \mid t=10, g=(x)x+t], 5 \sim (E1, (x)x+t)$

$[E0 \mid t=0, f=(E1, (x)x+t)], 6 \sim f(20) \sim 30$

$f(20): E1 :: [E2 \mid x=20], x+t \sim 20+10 \sim 30$

# (Parameterized) expression vs. (closure) value

- Functions defined using “def” are not values but parameterized expressions.
- Parameterized expression  $f$  can be converted to a value  $(f \ \_)$ .
- The compiler often infers and inserts missing conversions automatically.
- Anonymous functions are values.
- Anonymous functions can be seen as syntactic sugar:  
     $(x:T) \Rightarrow e$   
is equivalent to  
     $\{ \text{def } \text{noname}(x:T) = e; (\text{noname } \_) \}$   
    where  $\text{noname}$  is not used in  $e$ .
- One can even write a recursive anonymous function in this way.
- Q: what's the difference between param. exps and function values?  
  A: function values are “closures” (ie, param. exp. + env.)
- Q: how to implement call-by-name?  
  A: The argument expression is converted to a closure.

# Example: call by name with closures

```
1: { val t = 0
2:   def f(x: =>Int) = t + x
3:   val r = {
4:     val t = 10
5:     f(t*t) }           // t*t is treated as ()=>t*t
6:   r }
```

## ➤ Evaluation with Closures

$[E0 \mid t = \_, f = (x :=> \text{Int}) t + x, r = \_], 1 \sim 0$

$[E0 \mid t = 0, f = (x :=> \text{Int}) t + x, r = \_], 3 \sim 100$

$E0 :: [E1 \mid t = \_], 4 \sim 10$

$E0 :: [E1 \mid t = 10], 5 \sim f(t*t) \sim 100$

$f(t*t): E0 :: [E2 \mid x = (E1, t*t)], t+x \sim 0+x \sim 0+100 \sim 100$

$x: E1, t*t \sim 10*10 \sim 100$

$[E0 \mid t = 0, f = (x) t + x, r = 100], 6 \sim r \sim 100$

# Currying



# Motivation

```
def sum(f: Int=>Int, a: Int, b: Int): Int =  
  if (a <= b) f(a) + sum(f, a+1, b) else 0  
def sumLinear(a: Int, b: Int) = sum((n)=>n, a, b)  
def sumSquare(a: Int, b: Int) = sum((n)=>n*n, a, b)  
def sumCubes(a: Int, b: Int) = sum((n)=>n*n*n, a, b)
```

We want the following. How?

```
def sumLinear = sum(linear)  
def sumSquare = sum(square)  
def sumCubes = sum(cube)
```

# Solution

```
def sum(f: Int=>Int): (Int,Int)=>Int = {  
  def sumF(a: Int, b: Int): Int =  
    if (a <= b) f(a) + sumF(a+1, b) else 0  
  sumF // sumF _  
}
```

```
def sumLinear = sum((n)=>n)  
def sumSquare = sum((n)=>n*n)  
def sumCubes = sum((n)=>n*n*n)
```

# Benefits

```
def sumLinear = sum((n)=>n)  
def sumSquare = sum((n)=>n*n)  
def sumCubes = sum((n)=>n*n*n)
```

`sumSquare(3, 10) + sumCubes(5, 20)`

We don't need to define the wrapper functions.

`sum((n)=>n*n)(3, 10) + sum((n)=>n*n*n)(5, 20)`

# Multiple Parameter List

```
def sum(f: Int=>Int): (Int,Int)=>Int = {  
  def sumF(a: Int, b: Int): Int =  
    if (a <= b) f(a) + sumF(a+1, b) else 0  
  sumF _  
}
```

Or simply:

```
def sum(f: Int=>Int)(a: Int, b: Int): Int =  
  if (a <= b) f(a) + sum(f)(a+1, b) else 0
```

Note that `sum(f)` is just a parameterized expression.  
`(sum(f) _)` creates a closure like `(sumF _)`

The following code is slightly inefficient. Think about why.

```
def sum(f: Int=>Int): (Int,Int)=>Int =  
  (a,b) => if (a <= b) f(a) + sum(f)(a+1, b) else 0
```

# Comparison between Param. Exp vs. Closure

```
{
  def sum(f: Int=>Int)(a: Int, b: Int): Int =
    if (a <= b) f(a) + sum(f)(a+1, b) else 0
  def sumLinear = sum((n)=>n) _ // sum((n)=>n) incorrect

  sumLinear(0,100)
}
{
  def sum(f: Int => Int): (Int, Int) => Int = {
    def sumF(a: Int, b: Int): Int =
      if (a <= b) f(a) + sumF(a + 1, b) else 0
    sumF _
  }
  def sumLinear = sum((n)=>n) // sum((n)=>n) _ incorrect

  sumLinear(0,100)
}
```

# Currying and Uncurrying

- A function of type

$$(T_1, T_2, \dots, T_n) \Rightarrow T$$

can be turned into that of type

$$T_1 \Rightarrow (T_2 \Rightarrow (\dots \Rightarrow (T_n \Rightarrow T) \dots))$$

- This is called “currying” named after Haskell Brooks Curry.
- The opposite direction is called “uncurrying”.

# Currying using Anonymous Functions

```
def foo(x: Int, y: Int, z: Int)(a: Int, b: Int) =  
  x + y + z + a + b
```

```
val f1 = (x: Int, z: Int, b: Int) => foo(x, 1, z)(2, b)
```

```
val f2 = foo(_: Int, 1, _: Int)(2, _: Int)
```

```
val f3 = (x: Int, z: Int) => ((b: Int) => foo(x, 1, z)(2, b))
```

```
f1(1, 2, 3)
```

```
f2(1, 2, 3)
```

```
f3(1, 2)(3)
```

# Exercise

Curry the mapReduce function.

```
def mapReduce(reduce: (Int, Int) => Int, inival: Int,  
              f: Int => Int, a: Int, b: Int): Int = {  
  if (a <= b) reduce(f(a), mapReduce(reduce, inival, f, a+1, b))  
  else inival  
}
```

```
def sum(f: Int => Int, a: Int, b: Int): Int =  
  mapReduce((x, y) => x+y, 0, f, a, b)
```

```
def product(f: Int => Int, a: Int, b: Int): Int =  
  mapReduce((x, y) => x*y, 1, f, a, b)
```



# Solution

```
def mapReduce(reduce:(Int,Int)=>Int,inival: Int)
    (f: Int=>Int) (a: Int, b: Int): Int = {
    if (a <= b) reduce(f(a),mapReduce(reduce,inival)(f)(a+1,b))
    else inival
}
```

// need to make a closure since mapReduce is param. code.

```
def sum = mapReduce((x,y)=>x+y,0) _
```

// val is better than def. Think about why.

```
val product = mapReduce((x,y)=>x*y,1) _
```

# Exceptions

# Exception & Handling

```
class factRangeException(val arg: Int) extends Exception
```

```
def fact(n : Int): Int =  
  if (n < 0) throw new factRangeException(n)  
  else if (n == 0) 1  
  else n * fact(n-1)
```

```
def foo(n: Int) = fact(n + 10)
```

```
try {  
  println (fact(3))  
  println (foo(-100))  
} catch {  
  case e : factRangeException => {  
    println("fact range error: " + e.arg)  
  }  
}
```

# Datatypes

# Types so far

Types have introduction operations and elimination ones.

- Introduction: how to construct elements of the type
- Elimination: how to use elements of the type

## ➤ Primitive types

- Int, Boolean, Double, String, ...
- Intro for Int: ..., -2, -1, 0, 1, 2,
- Elim for Int: +, -, \*, /, <, <=, ...

## ➤ Function types

- $\text{Int} \Rightarrow \text{Int}$ ,  $(\text{Int} \Rightarrow \text{Int}) \Rightarrow (\text{Int} \Rightarrow \text{Int})$ , ...
- Intro:  $(x:T) \Rightarrow e$
- Elim:  $f(v)$

# Tuples

## ➤ Tuples

Intro:

- $(1, 2, 3) : (\text{Int}, \text{Int}, \text{Int})$
- $(1, \text{"a"}) : (\text{Int}, \text{String})$

Elim:

- $(1, \text{"a"}, 10)._1 = 1$
- $(1, \text{"a"}, 10)._2 = \text{"a"}$
- $(1, \text{"a"}, 10)._3 = 10$

Only up to length 22

# Structural Types (a.k.a. Record Types): Examples

```
import reflect.Selectable.reflectiveSelectable
def bar (x: Int) = x+1
object foo
{ val a = 1+2
  def b = a + 1
  def f(x: Int) = b + x
  val g : Int => Int = bar _
}
foo.b
foo.f(3)
val ff : Int=>Int = foo.f _
def g(x: {val a: Int; def b: Int;
         def f(x: Int): Int; val g: Int => Int}) =
  x.f(3)
g(foo)
```

# Type Alias

```
import reflect.Selectable.reflectiveSelectable
```

```
val gn = 0
```

```
object foo {
```

```
  val a = 3
```

```
  def b = a + 1
```

```
  def f(x: Int) = b + x + gn
```

```
}
```

```
foo.f(3)
```

```
type Foo = {val a: Int; def b: Int; def f(x: Int): Int}
```

```
def g(x: Foo) = {
```

```
  val gn = 10
```

```
  x.f(3)
```

```
}
```

```
g(foo)
```



# Structural Types: Evaluation

```
1: def bar (x: Int) = x+1
2: val foo = new //or, object foo
3: { val a = 2 + 1
4:   def f(x: Int) = a + x
5:   val g : Int => Int = bar _
6: }
7: val b = foo.f(1)
8: foo.g(b)
```

## ➤ Evaluation with Closures

```
E1[], 1 ~ E1[bar=(x)x+1], 2 ~ E1[...]:E2[], 3 ~
E1[...]:E2[a=3], 4 ~
E1[...]:E2[a=3, f=(x)a+x], 5 ~
E1[...]:E2[a=3, f=(x)a+x, g=((x)x+1, E1)], 6 ~
E1[bar=(x)x+1, foo=(E2)], 7 ~
E1[bar=(x)x+1, foo=(E2), b=4], 8 ~ 5
7: E1:E2:E3[x=1], a+x ~ 3+1 ~ 4
8: E1:E4[x=4], x+1 ~ 4+1 ~ 5
```

# Algebraic Datatypes

## ➤ Ideas

- $T = C \text{ of } T * \dots * T$   
|  $C \text{ of } T * \dots * T$   
|  $\dots$   
|  $C \text{ of } T * \dots * T$

- E.g.

Attr = Name of String  
| Age of Int  
| DOB of Int \* Int \* Int  
| Height of Double

Intro:

Name(“Chulsoo Kim”), Name(“Younghee Lee”), Age(16),  
DOB(2000,3,10), Height(171.5), ...

# Algebraic Datatypes: Recursion

## ➤ Recursive ADT

- E.g.

$\text{IList} = \text{INil}$

|  $\text{ICons of Int} * \text{IList}$

Intro:

$\text{INil}()$ ,

$\text{ICons}(3, \text{INil}())$ ,

$\text{ICons}(2, \text{ICons}(3, \text{INil}()))$ ,

$\text{ICons}(1, \text{ICons}(2, \text{ICons}(3, \text{INil}())))$ ,

...

# Algebraic Datatypes In Scala

## ➤ Attr

```
sealed abstract class Attr
case class Name(name: String) extends Attr
case class Age(age: Int) extends Attr
case class DOB(year: Int, month: Int, day: Int) extends Attr
case class Height(height: Double) extends Attr
```

```
val a : Attr = Name("Chulsoo Kim")
val b : Attr = DOB(2000, 3, 10)
```

## ➤ IList

```
sealed abstract class IList
case class INil() extends IList
case class ICons(hd: Int, tl: IList) extends IList
```

```
val x : IList = ICons(2, ICons(1, INil()))
def gen(n: Int) : IList =
  if (n <= 0) INil()
  else ICons(n, gen(n-1))
```

# Exercise

IOption = INone  
          | ISome of Int

BTree = Leaf  
        | Node of Int \* BTree \* BTree

```
sealed abstract class IList  
case class INil() extends IList  
case class ICons(hd: Int, tl: IList) extends IList  
  
def x : IList = ICons(2, ICons(1, INil()))
```

# Solution

```
sealed abstract class IOption  
case class INone() extends IOption  
case class ISome(some: Int) extends IOption
```

```
sealed abstract class BTree  
case class Leaf() extends BTree  
case class Node(value: Int, left: BTree, right: BTree)  
      extends BTree
```

# Pattern Matching

## ➤ Pattern Matching

- A way to use algebraic datatypes

```
e match {  
  case C1(...) => e1 : T  
  ...  
  case Cn(...) => en : T  
} : T
```

# Pattern Matching: An Example

```
def length(xs: IList) : Int =  
  xs match {  
    case /Ni/() => 0  
    case /Cons(x, tl) => 1 + length(tl)  
  }
```

length(x)



# Advanced Pattern Matching

## ➤ Advanced Pattern Matching

```
e match {  
  case P1 => e1  
  
  ...  
  
  case Pn => en  
}
```

- One can combine constructors and use `_` and `|` in a pattern.  
(E.g) `case ICons(x, INil()) | ICons(x, ICons(_, INil())) => ...`
- The given value `e` is matched against the first pattern `P1`.  
If succeeds, evaluate `e1`.  
If fails, `e` is matched against `P2`.  
If succeeds, evaluate `e2`.  
If fails, ...
- The compiler checks exhaustiveness (ie, whether there is a missing case).

# Advanced Pattern Matching: An Example

```
def secondElmt(xs: IList) : IOption =  
  xs match {  
    case /Nil() | /Cons(_, /Nil()) => /None()  
    case /Cons(_, /Cons(x, _)) => /Some(x)  
  }
```

Vs.

```
def secondElmt2(xs: IList) : IOption =  
  xs match {  
    case /Nil() | /Cons(_, /Nil()) => /None()  
    case /Cons(_, /Cons(x, /Nil())) => /Some(x)  
    case _ => /None()  
  }
```

Vs.

```
def secondElmt2(xs: IList) : IOption =  
  xs match {  
    case /Cons(_, /Cons(x, /Nil())) => /Some(x)  
    case _ => /None()  }
```

# Pattern Matching on Int

```
def factorial(n: Int) : Int =  
  n match {  
    case 0 => 1  
    case _ => n * factorial(n-1)  
  }
```

```
def fib(n: Int) : Int =  
  n match {  
    case 0 | 1 => 1  
    case _ => fib(n-1) + fib(n-2)  
  }
```

# Pattern Matching with If

```
def f(n: Int) : Int =  
  n match {  
    case 0 | 1 => 1  
    case _ if (n <= 5) => 2  
    case _ => 3  
  }
```

```
def f(t: BTree) : Int =  
  t match {  
    case Leaf() => 0  
    case Node(n,_,_) if (n <= 10) => 1  
    case Node(_,_,_) => 2  
  }
```

# Exercise

Write a function “find(t: BTree, x: Int) : Boolean” that checks whether x is in t.

```
sealed abstract class BTree
case class Leaf() extends BTree
case class Node(value: Int, left: BTree, right: BTree)
  extends BTree
```

# Solution

```
def find(t: BTree, i: Int) : Boolean =  
  t match {  
    case Leaf() => false  
    case Node(n, lt, rt) =>  
      if (i == n) true  
      else find(lt, i) || find(rt, i)  
  }
```

```
def find(t: BTree, i: Int) : Boolean = {  
  t match {  
    case Leaf() => false  
    case Node(v, _, _) if v == i => true  
    case Node(_, lt, rt) => find(lt, i) || find(rt, i)  
  }  
}
```

```
def t: BTree = Node(5, Node(4, Node(2, Leaf(), Leaf()), Leaf()),  
  Node(7, Node(6, Leaf(), Leaf()), Leaf()))  
find(t, 7), find(t, 1)
```

# Type Checking & Inference (Concept)

# What Are Types For?

## ➤ Typed Programming

```
def id1(x: Int): Int = x  
def id2(x: Double): Double = x
```

- At run time, type information is erased (ie, `id1 = id2`)

## ➤ Untyped Programming

```
def id(x) = x
```

- Do not care about types at compile time.
- But, many such languages check types at run time paying cost.
- Without run-time type check, errors can be badly propagated.

## ➤ What is compile-time type checking for?

- Can detect type errors at compile time.
- Increase Readability (Give a good abstraction).
- Soundness: Well-typed programs raise no type errors at run time.



# Type Checking and Inference

## ➤ Type Checking

$$x_1:T_1, x_2:T_2, \dots, x_n:T_n \vdash e : T$$

- `def f(x: Boolean): Boolean = x > 3`  
=> Type error
- `def f(x: Int): Boolean = x > 3`  
=> OK. `f: (x: Int)Boolean`

## ➤ Type Inference

$$x_1:T_1, x_2:T_2, \dots, x_n:T_n \vdash e : ?$$

- `def f(x: Int) = x > 3`  
=> OK by type inference. `f: (x: Int)Boolean`
- Too much type inference is not good. Why?

You can learn how type checking & inference work in  
**4190.310 Programming Languages**

# Parametric Polymorphism

# Parametric Polymorphism: Functions

## ➤ Problem

```
def id1(x: Int): Int = x
def id2(x: Double): Double = x
```

- Can we avoid copy and paste?
- Polymorphism to the rescue!

## ➤ Parametric Polymorphism (a.k.a. For-all Types)

```
def id[A](x: A) : A = x
```

- The type of `id` is `[A](val x:A)=>A`
- `id` is a parametric expression.
- `id[T] _` is a value of type `T=>T` for any type `T`.
- Function types do not support polymorphism.  
(E.g.) `[A](A=>A)` is not a valid function value type.

[We will learn other kinds of polymorphism later.]

# Examples

```
def id[A](x:A) = x
id(3)
id("abc")
```

```
def applyn[A](f: A => A, n: Int, x: A): A =
  n match {
    case 0 => x
    case _ => f(applyn(f, n - 1, x))
  }
```

```
applyn((x:Int)=>x+1, 100, 3)
applyn((x:String)=>x+"!", 10, "gil")
applyn(id[String], 10, "hur")
```

```
def foo[A,B](f: A=>A, x: (A,B)) : (A,B) =
  (applyn[A](f, 10, x._1), x._2)
```

```
foo[String, Int]((x:String)=>x+"!", ("abc", 10))
```

# Parametric Polymorphism: Datatypes

```
sealed abstract class MyOption[A]  
case class MyNone[A]() extends MyOption[A]  
case class MySome[A](some: A) extends MyOption[A]
```

```
sealed abstract class MyList[A]  
case class MyNil[A]() extends MyList[A]  
case class MyCons[A](hd: A, tl: MyList[A]) extends MyList[A]
```

```
sealed abstract class BTree[A]  
case class Leaf[A]() extends BTree[A]  
case class Node[A](value: A, left: BTree[A], right: BTree[A])  
extends BTree[A]
```

```
def x: MyList[Int] = MyCons(3, MyNil())  
def y: MyList[String] = MyCons("abc", MyNil())
```

# Revisit: Solution with Tail Recursion

```
def find[A](t: BTree[A], x: A) : Boolean = {  
  def findIter(ts: MyList[BTree[A]]) : Boolean =  
    ts match {  
      case MyNil() => false  
      case MyCons(Leaf(), tl) => findIter(tl)  
      case MyCons(Node(v, _, _), _) if v == x => true  
      case MyCons(Node(_, l, r), tl) =>  
        findIter(MyCons(l, MyCons(r, tl)))    }  
  findIter(MyCons(t, MyNil()))  
}
```

```
def genTree(v: Int, n: Int) : BTree[Int] = {  
  def genTreeIter(t: BTree[Int], m : Int) : BTree[Int] =  
    if (m == 0) t  
    else genTreeIter(Node(v, t, Leaf()), m-1)  
  genTreeIter(Leaf(), n)  
}
```

```
find(genTree(0, 10000), 1)
```

# Exercise

BSTree[A] = Leaf

| Node of Int \* A \* BSTree[A] \* BSTree[A]

```
def lookup[A](t: BSTree[A], k: Int) : MyOption[A] =  
  ???
```

```
def t : BSTree[String] =  
  Node(5, "My5",  
    Node(4, "My4", Node(2, "My2", Leaf(), Leaf()), Leaf()),  
    Node(7, "My7", Node(6, "My6", Leaf(), Leaf()), Leaf()))
```

lookup(t, 7)

lookup(t, 3)

# Solution

```
sealed abstract class BSTree[A]
case class Leaf[A]() extends BSTree[A]
case class Node[A](key: Int, value: A, left: BSTree[A], right:
BSTree[A]) extends BSTree[A]
def lookup[A](t: BSTree[A], key: Int) : MyOption[A] =
  t match {
    case Leaf() => MyNone()
    case Node(k,v,l t,r t) =>
      k match {
        case _ if key == k => MySome(v)
        case _ if key < k => lookup(l t,key)
        case _ => lookup(r t, key)
      }
  }
def t : BSTree[String] =
  Node(5, "My5",
    Node(4, "My4", Node(2, "My2", Leaf(), Leaf()), Leaf()),
    Node(7, "My7", Node(6, "My6", Leaf(), Leaf()), Leaf()))
lookup(t, 7)
lookup(t, 3)
```



# A Better Way

```
sealed abstract class BTree[A]  
case class Leaf[A]() extends BTree[A]  
case class Node[A](value: A, left: BTree[A], right: BTree[A])  
  extends BTree[A]
```

```
type BSTree[A] = BTree[(Int, A)]
```

```
def lookup[A](t: BSTree[A], k: Int) : MyOption[A] =  
  ???
```

```
def t : BSTree[String] =  
  Node((5, "My5"),  
    Node((4, "My4"), Node((2, "My2"), Leaf(), Leaf()), Leaf()),  
    Node((7, "My7"), Node((6, "My6"), Leaf(), Leaf()), Leaf()))
```

```
lookup(t, 7)
```

# Solution

```
type BSTree[A] = BTree[(Int,A)]
```

```
def lookup[A](t: BTree[(Int,A)], key: Int) : MyOption[A] =  
  t match {  
    case Leaf() => MyNone()  
    case Node((k,v), lt, rt) =>  
      k match {  
        case _ if key == k => MySome(v)  
        case _ if key < k => lookup(lt, key)  
        case _ => lookup(rt, key)  
      }  
  }
```

```
def t : BTree[String] =  
  Node((5, "My5"),  
    Node((4, "My4"), Node((2, "My2"), Leaf(), Leaf()), Leaf()),  
    Node((7, "My7"), Node((6, "My6"), Leaf(), Leaf()), Leaf()))
```

```
lookup(t, 7)  
lookup(t, 3)
```

# Polymorphic Option (Library)

## ➤ Option[T]

Intro:

- None
- Some(x)
- Library functions

Elim:

- Pattern matching
- Library functions

Some(3) : Option[Int]

Some("abc"): Option[String]

None: Option[Int]

None: Option[String]

# Polymorphic List (Library)

## ➤ List[T]

Intro:

- Nil
- $x :: L$
- Library functions

Elim:

- Pattern matching
- Library functions

`“abc”::Nil : List[String]`

`List(1,3,4,2,5) = 1::3::4::2::5::Nil : List[Int]`

# Summary: Parametric Polymorphism

## ➤ Parametric Polymorphism

- Program against unknown datatypes
- How is it possible?

# **PART 2**

## **Object-Oriented Programming**

# Sub Type Polymorphism (Concept)

# Motivation

We want:

```
object tom {  
  val name = "Tom"  
  val home = "02-880-1234"  
}
```

```
object bob {  
  val name = "Bob"  
  val mobile = "010-1111-2222"  
}
```

```
def greeting(r: ???) = "Hi " + r.name + ", How are you?"  
greeting(tom)  
greeting(bob)
```

Note that we have

```
tom: {val name: String; val home: String}  
bob: {val name: String; val mobile: String}
```



# Sub Types to the Rescue!

```
import reflect.Selectable.reflectiveSelectable
```

```
type NameHome = { val name: String; val home: String }  
type NameMobile = { val name: String; val mobile: String }  
type Name = { val name: String }
```

```
NameHome <: Name    (NameHome is a sub type of Name)
```

```
NameMobile <: Name  (NameMobile is a sub type of Name)
```

```
def greeting(r: Name) = "Hi " + r.name + ", How are you?"  
greeting(tom)  
greeting(bob)
```

# Sub Types

- The sub type relation is kind of the subset relation.
- But they are **NOT** the same.
- $T <: S$   
Every element of T **can be used as** that of S.
- *Cf.* T is a subset of S.  
Every element of T **is** that of S.
- Why polymorphism?  
A function of type  $S \Rightarrow R$  can be used as  $T \Rightarrow R$  for many sub types T of S.  
Note that  $S \Rightarrow R <: T \Rightarrow R$  when  $T <: S$ .

# Summary: Subtype Polymorphism

## ➤ Subtype Polymorphism

- Program against known datatypes with common structures
- How is it possible?

# Two Kinds of Sub Types

## ➤ Structural Sub Types (a.k.a. Duck Typing)

- The system implicitly determines the sub type relation by the structures of data types.
- Structurally equivalent types are treated the same.

## ➤ Nominal Sub Types (a.k.a. Ad hoc Polymorphism)

- The user explicitly specify the sub type relation using the names of data types.
- Structurally equivalent types with different names may be treated differently.

# Structural Sub Types

# General Sub Type Rules

- Reflexivity:

For any type T, we have:

$$T <: T$$

- Transitivity:

For any types T, S, R, we have:

$$T <: R \quad R <: S$$

=====

$$T <: S$$

# Sub Types for Special Types

- Nothing: The empty set
- Any: The set of all values

- For any type T, we have:

$\text{Nothing} <: T <: \text{Any}$

- Example

```
val a : Int = 3
val b : Any = a
def f(a: Nothing) : Int = a
```

# Sub Types for Records

- Permutation

$$\begin{array}{c} \text{=====} \\ \{...\; x: T1; y: T2; ...\} <: \{...\; y: T2, x: T1; ...\} \end{array}$$

- Width

$$\begin{array}{c} \text{=====} \\ \{...\; x: T; ...\} <: \{...\; ...\} \end{array}$$

- Depth

$$\begin{array}{c} T <: S \\ \text{=====} \\ \{...\; x: T ; ...\} <: \{...\; x: S; ...\} \end{array}$$



# Sub Types for Records

- Example

{val x: { val y: Int; val z: String}, val w: Int}

<: (by permutation)

{val w: Int; val x: { val y: Int; val z: String} }

<: (by depth & width)

{val w: Int; val x: {val z: String} }

# Sub Types for Tuples

- Depth

$$T <: S$$

=====

$$(\dots, T, \dots) <: (\dots, S, \dots)$$

# Sub Types for Functions

- Function Sub Type

$$T <: T' \quad S <: S'$$

=====

$$(T' \Rightarrow S) <: (T \Rightarrow S')$$

- Example

```
import reflect.Selectable.reflectiveSelectable
```

```
def foo(s: {val a: Int; val b: Int}) : {val x: Int; val y: Int} = {  
  object tmp {  
    val x = s.b  
    val y = s.a  
  }  
  tmp  
}  
val gee: {val a: Int; val b: Int; val c: Int} => {val x: Int} =  
  foo _
```

# Classes

# Class: Parameterized Record

```
import reflect.Selectable.reflectiveSelectable
```

```
type gee_type = {val name:String; val age: Int; def getPP(): String}  
def gee_fun(_name: String, _age: Int) : gee_type = {  
  if (!(_age >= 0 && _age < 200)) throw new Exception("Out of range")  
  object tmp {  
    val name : String = _name  
    val age : Int = _age  
    def getPP() : String = name + " of age " + age.toString() }  
  tmp }  
val gee : gee_type = gee_fun("David Jones",25)  
  
gee.getPP()
```

# Class: Parameterized Record

```
class foo_type(_name: String, _age: Int) {  
  if (!(_age >= 0 && _age < 200)) throw new Exception("Out of range")  
  val name: String = _name  
  val age: Int = _age  
  def getPP(): String = name + " of age " + age.toString() }  
val foo: foo_type = new foo_type("David Jones", 25)  
foo.getPP()
```

use: foo.name foo.age foo.getPP

- foo is a value of foo\_type
- gee is a value of gee\_type

# Class: No Structural Sub Typing

## ➤ Records: Structural sub-typing

`foo_type <: gee_type`

## ➤ Classes: Nominal sub-typing

`gee_type <: foo_type`

```
val v1 : gee_type = foo
```

```
val v2 : foo_type = gee // type error
```

```
def greeting(r:{val name:String}) =  
  "Hi " + r.name + ", How are you?"  
greeting(foo)
```

# Structural Types vs. Nominal Types

## ➤ Structural Types

- Includes arbitrary values with the required structures as elements
- Allows arbitrary types with the required structures as sub types
- Cannot assume any properties on their elements

## ➤ Nominal Types

- Includes only specific values as elements
- Allows only specific types as sub types
- Can assume specific properties on their elements



# Class: Can be Recursive!

```
class MyList[A](v: A, nxt: Option[MyList[A]]) {  
  val value : A = v  
  val next : Option[MyList[A]] = nxt  
}
```

```
type YourList[A] = Option[MyList[A]]
```

```
val t : YourList[Int] =  
  Some(new MyList(3, Some (new MyList(4, None))))
```

# Note on Null value

- `null`: The special element of every class & structural type
- `null` is often used to represent None instead of using an Option type (Efficient but Not Safe)
- It is discouraged to use `null` in Scala although Scala supports `null` for compatibility with Java.

# Simplification using Argument Members

```
class MyList[A](v: A, nxt: Option[MyList[A]]) {  
  val value : A = v  
  val next : Option[MyList[A]] = nxt  
}
```

```
class MyList[A](val value:A, val next:Option[MyList[A]]) {  
}
```

```
class MyList[A](val value:A, val next:Option[MyList[A]])
```

# Simplification using Companion Object

```
class MyList[A](val value:A, val next:Option[MyList[A]])  
object MyList  
{ def apply[A](v: A, nxt: Option[MyList[A]]) =  
    new MyList(v,nxt)  
}
```

```
type YourList[A] = Option[MyList[A]]  
object YourList  
{ def apply[A](v: A, nxt: Option[MyList[A]]) =  
    Some(new MyList(v,nxt))  
}
```

```
val t0 = None
```

```
val t1 = Some(new MyList(3,Some(MyList(4,None))))
```

```
val t2 = YourList(3,(YourList(4,None)))
```

# Exercise

Define a class “MyTree[A]” for binary trees:

```
MyTree[A] =  
  (value: A) *  
  (left: Option[MyTree[A]]) *  
  (right: Option[MyTree[A]])
```

# Solution

```
class MyTree[A](v: A,  
                lt: Option[MyTree[A]],  
                rt: Option[MyTree[A]]) {  
    val value = v  
    val left = lt  
    val right = rt  
}
```

```
type YourTree[A] = Option[MyTree[A]]
```

```
val t0 : YourTree[Int] = None  
val t1 : YourTree[Int] = Some(new MyTree(3, None, None))  
val t2 : YourTree[Int] =  
    Some(new MyTree(3, Some (new MyTree(4, None, None)), None))
```

# Simplified Solution

```
class MyTree[A](val value : A,  
                val left  : Option[MyTree[A]],  
                val right : Option[MyTree[A]])
```

```
type YourTree[A] = Option[MyTree[A]]
```

```
object YourTree  
{ def apply[A](v:A, lt:Option[MyTree[A]], rt:Option[MyTree[A]]) =  
    Some(new MyTree(v,lt,rt))  
}
```

```
val t0: YourTree[Int] = None  
val t1: YourTree[Int] = YourTree(3,None,None)  
val t2: YourTree[Int] = YourTree(3,YourTree(4,None,None),None)
```

# Nominal Sub Typing for Classes



# Nominal Sub Typing, a.k.a. Inheritance

```
class foo_type(x: Int, y: Int) {  
  val a : Int = x  
  def b : Int = a + y  
  def f(z: Int) : Int = b + y + z  
}
```

```
class gee_type(x: Int) extends foo_type(x+1,x+2) {  
  val c : Int = f(x) + b  
}
```

```
gee_type <: foo_type
```

```
(new gee_type(30)).c  
def test(f: foo_type) = f.a + f.b  
test(new foo_type(10,20))  
test(new gee_type(30))
```

# Overriding

```
class foo_type(x: Int, y: Int) {  
  val a : Int = x  
  def b : Int = 0  
  def f(z: Int) : Int = b * z  
}  
  
class gee_type(x: Int) extends foo_type(x+1,x+2) {  
  override def b = 10  
  // or, override def b = super.b + 10  
  val c : Int = f(x) + b  
}
```

```
(new gee_type(30)).c  
def test(v: foo_type) =  
  println(v.f(42))  
test(new foo_type(1,2))  
test(new gee_type(0))
```

# Overriding vs. Overloading

```
class foo_type(x: Int, y: Int) {  
  val a : Int = x  
  def b : Int = 0  
  def f(z: Int) : Int = b * z  
}  
class gee_type(x: Int) extends foo_type(x+1,x+2) {  
  override def b = 10  
}
```

**Q:** Can we override with a different type?

```
override def f(z: String): Int = 77    //No, arg: diff type  
def f(z: String): Int = 77             // Overloading, arg: diff type  
override def f(z: Int): Int = 77      //Yes, arg: same type
```

## Example: MyList

```
class MyList[A](v: A, nxt: Option[MyList[A]]) {  
  val value: A = v  
  val next: Option[MyList[A]] = nxt  
}  
type YourList[A] = Option[MyList[A]]  
val t: YourList[Int] =  
  Some(new MyList(3, Some(new MyList(4, None))))
```

```
class MyList[A]()  
class MyNil[A]() extends MyList[A]  
class MyCons[A](val hd: A, val tl: MyList[A])  
  extends MyList[A]  
val t: MyList[Int] =  
  new MyCons(3, new MyCons(4, new MyNil()))
```

# Simplification: MyList

```
class MyList[A]
```

```
class MyNil[A]() extends MyList[A]
```

```
object MyNil { def apply[A]() = new MyNil[A]() }
```

```
class MyCons[A](val hd: A, val tl: MyList[A])  
  extends MyList[A]
```

```
object MyCons {  
  def apply[A](hd:A, tl:MyList[A]) = new MyCons[A](hd, tl)}
```

```
val t: MyList[Int] = MyCons(3, MyNil())
```

```
def length(x: MyList[Int]) = ???
```

# Case Class

```
sealed abstract class MyList[A] { ... }  
case class MyNil[A]() extends MyList[A] { ... }  
object MyNil { def apply[A]() = new MyNil[A]() }  
case class MyCons[A](val hd: A, val tl: MyList[A])  
  extends MyList[A] { ... }  
object MyCons {  
  def apply[A](hd:A, tl:MyList[A]) = new MyCons[A](hd, tl) }  
val t: MyList[Int] = MyCons(3, MyNil())
```

## Allow Pattern Matching

```
def length(x: MyList[Int]): Int =  
  x match {  
    case MyNil() => 0  
    case MyCons(hd, tl) => 1 + length(tl)  
  }
```

*Cf.* sealed abstract class MyList[A]

# Exercise

Define “MyTree[A]” using sub class.

```
class MyTree[A](v: A,  
                lt: Option[MyTree[A]],  
                rt: Option[MyTree[A]]) {  
  val value = v  
  val left = lt  
  val right = rt  
}
```

```
type YourTree[A] = Option[MyTree[A]]
```

# Solution

```
sealed abstract class MyTree[A]  
case class Empty[A]() extends MyTree[A]  
case class Node[A](value:A, left:MyTree[A], right:MyTree[A])  
    extends MyTree[A]
```

```
val t : MyTree[Int] =  
    Node(3, Node(4, Empty(), Empty()), Empty())
```

```
t match {  
    case Empty() => 0  
    case Node(v, l, r) => v  
}
```



# Solution with Monotonicity

```
// sealed abstract class MyTree[A]  
// case class Empty[A]() extends MyTree[A]  
// case class Node[A](value:A, left:MyTree[A], right:MyTree[A])  
//   extends MyTree[A]
```

```
// MyTree[+A]:  $A <: B \Rightarrow \text{MyTree}[A] <: \text{MyTree}[B]$ 
```

```
// MyTree[-A]:  $A <: B \Rightarrow \text{MyTree}[B] <: \text{MyTree}[A]$ 
```

```
sealed abstract class MyTree[+A]  
case object Empty extends MyTree[Nothing]  
case class Node[A](value:A, left:MyTree[A], right:MyTree[A])  
  extends MyTree[A]
```

```
val t : MyTree[Int] = Node(3, Node(4, Empty, Empty), Empty)  
t match {  
  case Empty => 0  
  case Node(v, l, r) => v  
}
```

# Solution with enum

```
// sealed abstract class MyTree[+A]  
// case object Empty extends MyTree[Nothing]  
// case class Node[A](value:A, left:MyTree[A], right:MyTree[A])  
//   extends MyTree[A]
```

```
enum MyTree[+A]:  
  case Empty //: MyTree[Nothing]  
  case Node(value: A, left: MyTree[A], right: MyTree[A])  
import t MyTree._
```

```
val t : MyTree[Int] = Node(3, Node(4, Empty, Empty), Empty)  
t match {  
  case Empty => 0  
  case Node(v, l, r) => v  
}
```

# Encoding ADT using classes: Option

```
sealed abstract class MyOption[+A] {  
  def matches[R](e1 : =>R, e2: A=>R) : R  
}  
object MyNone extends MyOption[Nothing] {  
  def matches[R](e1 : =>R, e2: Nothing=>R) = e1  
}  
class MySome[A](v: A) extends MyOption[A] {  
  def matches[R](e1 : =>R, e2: A=>R) = e2(v)  
}  
object MySome { def apply[A](v: A) = new MySome(v) }  
  
val x: MyOption[Int] = MySome(42)  
val y: MyOption[Int] = MyNone  
x.matches(-1,(n)=>n+1)  
y.matches(-1,(n)=>n+1)
```

# Encoding ADT using classes: List

```
sealed abstract class MyList[+A] {  
  def matches[R](nilE: =>R, consE: (A,MyList[A])=>R) : R  
}  
object MyNil extends MyList[Nothing] {  
  def matches[R](nilE: =>R, consE: (Nothing,MyList[Nothing])=>R) = nilE  
}  
class MyCons[A](val hd: A, val tl: MyList[A]) extends MyList[A] {  
  def matches[R](nilE: =>R, consE: (A,MyList[A])=>R) = consE(hd,tl)  
}  
object MyCons{ def apply[A](hd:A, tl:MyList[A]) = new MyCons[A](hd, tl) }  
  
def length[A](l: MyList[A]) : Int =  
  l.matches(  
    0,  
    (_,tl) => 1 + length(tl) )  
  
length(MyCons(3, MyCons(2, MyNil)))
```

# Encoding ADT using classes: Monotonicity

```
sealed abstract class MyList[+A] {  
  def matches[R](nilE: =>R, consE: (A,MyList[A])=>R) : R  
  def append[B>:A](l: MyList[B]) : MyList[B]  
}  
  
object MyNil extends MyList[Nothing] {  
  def matches[R](nilE: =>R, consE: (Nothing,MyList[Nothing])=>R) = nilE  
  def append[B](l: MyList[B]) = l  
}  
  
class MyCons[A](val hd: A, val tl: MyList[A]) extends MyList[A] {  
  def matches[R](nilE: =>R, consE: (A,MyList[A])=>R) = consE(hd,tl)  
  def append[B>:A](l: MyList[B]) = new MyCons[B](hd, tl.append(l))  
}  
  
object MyCons { def apply[A](hd:A, tl:MyList[A]) = new MyCons[A](hd, tl) }  
def length[A](l: MyList[A]) : Int =  
  l.matches(  
    0,  
    (_,tl) => 1 + length(tl) )  
length(MyCons(3, MyCons(2, MyNil)).append(MyCons(1,MyNil)))
```

# **Abstract Classes for Interface**

# Abstract Class: Interface

## ➤ Abstract Classes

- Can be used to abstract away the implementation details.

Abstract classes for Interface

Concrete sub-classes for Implementation

# Abstract Class: Interface

## ➤ Example Interface

*// Written by Alice*

*// if getValue(i) returns None, you should not use i.getNext()*

```
abstract class Iter[A] {  
  def getValue: Option[A]  
  def getNext: Iter[A]  
}
```

```
def sumElements[A](f: A=>Int)(xs: Iter[A]) : Int =  
  xs.getValue match {  
    case None => 0  
    case Some(n) => f(n) + sumElements(f)(xs.getNext)  
  }
```

```
def sumElementsId(xs: Iter[Int]) =  
  sumElements((x: Int)=>x)(xs)
```



# Concrete Class: Implementation

*// Written by Bob*

```
sealed abstract class MyList[A] extends Iter[A]
case class MyNil[A]() extends MyList[A] {
  def getValue = None
  def getNext = throw new Exception("...")
}
case class MyCons[A](hd: A, tl: MyList[A])
  extends MyList[A]
{
  def getValue = Some(hd)
  def getNext = tl
}
```

```
val t1 = MyCons(3, MyCons(5, MyCons(7, MyNil())))
```

```
sumElementsId(t1)
```

# Exercise

Define `IntCounter(n)` that implements the interface `Iter[A]`.

*// Written by Catherine*

```
class IntCounter(n: Int) extends Iter[Int] {  
  def getValue = ???  
  def getNext = ???  
}
```

```
sumElementsId(new IntCounter(100))
```

# Solution

Define `IntCounter(n)` that implements the interface `Iter[A]`.

*// Written by Catherine*

```
class IntCounter(n: Int) extends Iter[Int] {  
  def getValue = if (n >= 0) Some(n) else None  
  def getNext = new IntCounter(n-1)  
}
```

```
sumElementsId(new IntCounter(100))
```

# A Better Interface

```
abstract class Iter[A] {  
  def get: Option[(A, Iter[A])]  
}  
  
def sumElements[A](f: A=>Int)(xs: Iter[A]) : Int =  
  xs.get match {  
    case None => 0  
    case Some(n,nxt) => f(n) + sumElements(f)(nxt)  
  }  
  
def sumElementsId(xs:Iter[Int]) = sumElements((x:Int)=>x)(xs)  
sealed abstract class MyList[A] extends Iter[A]  
case class MyNil[A]() extends MyList[A] {  
  def get = None }  
case class MyCons[A](hd: A, tl: MyList[A]) extends MyList[A] {  
  def get = Some(hd,tl) }  
class IntCounter(n: Int) extends Iter[Int] {  
  def get = if (n >= 0) Some(n, new IntCounter(n-1)) else None }
```

## **More on Abstract Classes**

# Problem: Iter for MyTree

```
abstract class Iter[A] {  
  def getValue: Option[A]  
  def getNext: Iter[A]  
}
```

*// Written by David*

```
sealed abstract class MyTree[A]  
case class Empty[A]() extends MyTree[A]  
case class Node[A](value: A,  
                   left: MyTree[A],  
                   right: MyTree[A]) extends MyTree[A]
```

Q: Can MyTree[A] implement Iter[A]?

Try it, but it is not easy.

# Possible Solution

*// Written by David*

```
sealed abstract class MyTree[A] extends Iter[A]
case class Empty[A]() extends MyTree[A] {
  def getValue = None
  def getNext = this }
case class Node[A](value: A, left: MyTree[A], right: MyTree[A])
  extends MyTree[A] {
  def getValue = Some(value)
  def getNext: MyTree[A] = {
    def merge_right(l : MyTree[A]): MyTree[A] = l match {
      case Empty() => right
      case Node(v, lt, rt) => Node(v, lt, merge_right(rt)) }
    merge_right(left) } }
val t1 = Node(3, Node(7, Node(2, Empty(), Empty()), Empty()),
              Node(8, Empty(), Empty()))
sumElements[Int]((x)=>x*x)(t1)
```

# Solution: Better Interface

```
abstract class Iter[A] {  
  def getValue: Option[A]  
  def getNext: Iter[A]  
}  
abstract class Iterable[A] {  
  def iter : Iter[A]  
}
```

```
def sumElements[A](f: A=>Int)(xs: Iter[A]) : Int =  
  xs.getValue match {  
    case None => 0  
    case Some(n) => f(n) + sumElements(f)(xs.getNext)  
  }
```

```
def sumElementsGen[A](f: A=>Int)(xs: Iterable[A]) : Int =  
  sumElements(f)(xs.iter)
```



# Let's Use MyList

```
sealed abstract class MyList[A] extends Iter[A]
case class MyNil[A]() extends MyList[A] {
  def getValue = None
  def getNext = throw new Exception("...")
}
case class MyCons[A](val hd: A, val tl: MyList[A])
  extends MyList[A] {
  def getValue = Some(hd)
  def getNext = tl
}
```

# MyTree <: Iterable (Try)

```
sealed abstract class MyTree[A] extends Iterable[A]
```

```
case class Empty[A]() extends MyTree[A] {  
  val iter = MyNil()  
}
```

```
case class Node[A](value: A,  
                  left: MyTree[A],  
                  right: MyTree[A]) extends MyTree[A] {  
  // "val iter" is more specific than "def iter",  
  // so it can be used in a sub type.  
  // In this example, "val iter" is also  
  // more efficient than "def iter".  
  val iter = MyCons(value, ???(left.iter, right.iter))  
}
```

# Extend MyList with append

```
sealed abstract class MyList[A] extends Iter[A] {  
  def append(lst: MyList[A]) : MyList[A]  
}  
  
case class MyNil[A]() extends MyList[A] {  
  def getValue = None  
  def getNext = throw new Exception("...")  
  def append(lst: MyList[A]) = lst  
}  
  
case class MyCons[A](val hd: A, val tl: MyList[A])  
  extends MyList[A]  
{  
  def getValue = Some(hd)  
  def getNext = tl  
  def append(lst: MyList[A]) = MyCons(hd, tl.append(lst))  
}
```

# MyTree <: Iterable

```
sealed abstract class MyTree[A] extends Iterable[A] {  
  def iter : MyList[A]  
  // Note:  
  // def iter : Int // Type Error because not (Int <: Iter[A])  
}  
case class Empty[A]() extends MyTree[A] {  
  val iter = MyNil()  
}  
case class Node[A](value: A,  
                   left: MyTree[A],  
                   right: MyTree[A]) extends MyTree[A] {  
  def iter = MyCons(value, left.iter.append(right.iter))  
  // def iter = left.iter.append(MyCons(value, right.iter))  
  // def iter = left.iter.append(right.iter.append(  
  //   MyCons(value, MyNil())))  
}
```

# Test

```
def generateTree(n: Int) : MyTree[Int] = {  
  def gen(lo: Int, hi: Int) : MyTree[Int] =  
    if (lo > hi) Empty()  
    else {  
      val mid = (lo+hi)/2  
      Node(mid, gen(lo,mid-1), gen(mid+1,hi))  
    }  
  gen(1,n)  
}
```

```
sumElementsGen((x: Int) => x)(generateTree(100))
```

# Iter <: Iterable

```
abstract class Iterable[A] {  
  def iter : Iter[A]  
}
```

```
abstract class Iter[A] extends Iterable[A] {  
  def getValue: Option[A]  
  def getNext: Iter[A]  
  def iter = this  
}
```

```
val lst : MyList[Int] =  
  MyCons(3, MyCons(4, MyCons(2, MyNil())))
```

```
sumElementsGen ((x:Int)=>x)(lst)
```

## Note: tail-recursive “append”

```
sealed abstract class MyList[A] extends Iter[A] {  
  def append(lst: MyList[A]) : MyList[A] =  
    MyList.revAppend(MyList.revAppend(this, MyNil()), lst)  
}  
  
object MyList { // Mutual references are allowed between class T and object T  
  // Tail-recursive functions should be written in “object”, or as final methods  
  def revAppend[A](lst1: MyList[A], lst2: MyList[A]): MyList[A] =  
    lst1 match {  
      case MyNil() => lst2  
      case MyCons(hd, tl) => revAppend(tl, MyCons(hd, lst2))  
    }  
}  
  
case class MyNil[A]() extends MyList[A] {  
  def getValue = None  
  def getNext = throw new Exception("...")  
}  
  
case class MyCons[A](val hd:A, val tl:MyList[A]) extends MyList[A] {  
  def getValue = Some(hd)  
  def getNext = tl  
}
```

# Lazy List



# Problem: Inefficiency

```
def time[R](block: => R): R = {  
  val t0 = System.nanoTime()  
  val result = block    // call-by-name  
  val t1 = System.nanoTime()  
  println("Elapsed time: " + ((t1 - t0)/1000000) + "ms"); result  
}  
  
def sumN[A](f: A=>Int)(n: Int, xs: Iterable[A]) : Int = {  
  def sumIter(res : Int, n: Int, xs: Iter[A]) : Int =  
    if (n <= 0) res  
    else xs.getValue match {  
      case None => res  
      case Some(v) => sumIter(f(v) + res, n-1, xs.getNext)  
    }  
  sumIter(0, n, xs.iter)  
}  
  
// Problem: takes a few seconds to get a single value  
{ val t: MyTree[Int] = generateTree(200000)  
  time (sumN((x: Int) => x)(1, t)) }
```

# Solution 1: Using Lists of Trees

```
class MyTreeIter[A](val lst: MyList[MyTree[A]]) extends Iter[A] {  
  val getValue = lst match {  
    case MyCons(Node(v,_,_), _) => Some(v)  
    case _ => None  
  }  
  def getNext = {  
    val remainingTrees : MyList[MyTree[A]] = lst match {  
      case MyNil() => throw new Exception("...")  
      case MyCons(hd,tl) => hd match {  
        case Empty() => throw new Exception("...")  
        case Node(_,Empty(),Empty()) => tl  
        case Node(_,lt,Empty()) => MyCons(lt,tl)  
        case Node(_,Empty(),rt) => MyCons(rt,tl)  
        case Node(_,lt,rt) => MyCons(lt,MyCons(rt,tl))  
      }  
    }  
    new MyTreeIter(remainingTrees)  
  }  
}
```

# Lazy Iteration using Lists of Trees

```
sealed abstract class MyTree[A] extends Iterable[A]
case class Empty[A]() extends MyTree[A] {
  val iter = new MyTreeIter(MyNil())
}
case class Node[A](value: A,
                  left: MyTree[A],
                  right: MyTree[A]) extends MyTree[A]
{
  val iter = new MyTreeIter(MyCons(this, MyNil()))
}

{ val t: MyTree[Int] = generateTree(200000)
  time (sumN((x: Int) => x)(100, t))
  time (sumN((x: Int) => x)(100000, t))
}
```

## Solution 2: Lazy List

```
sealed abstract class LazyList[A] extends Iter[A] {  
  def append(lst: LazyList[A]) : LazyList[A]  
}
```

```
case class LNil[A]() extends LazyList[A] {  
  def getValue = None  
  def getNext = throw new Exception("")  
  def append(lst: LazyList[A]) = lst  
}
```

```
class LCons[A](hd: A, _tl: =>LazyList[A]) extends LazyList[A] {  
  lazy val tl = _tl  
  def getValue = Some(hd)  
  def getNext = tl  
  def append(lst: LazyList[A]) = LCons(hd, tl.append(lst))  
object LCons {  
  def apply[A](hd: A, tl: =>LazyList[A]) = new LCons(hd, tl)  
}
```

Note: “append” is not recursive!!!

# Lazy Iteration using LazyList

```
sealed abstract class MyTree[A] extends Iterable[A] {
  def iter : LazyList[A]
}
case class Empty[A]() extends MyTree[A] {
  val iter = LNil()
}
case class Node[A](value: A,
                   left: MyTree[A],
                   right: MyTree[A]) extends MyTree[A] {
  lazy val iter = LCons(value, left.iter.append(right.iter))
  // lazy val iter = left.iter.append(LCons(value, right.iter))
  // lazy val iter = left.iter.append(right.iter.append(
  //   LCons(value, LNil())))
}
{ val t: MyTree[Int] = generateTree(200000)
  time (sumN((x: Int) => x)(100, t))
  time (sumN((x: Int) => x)(100000, t))
}
```

Note: “i t e r” is not recursive!!!

# Wrapper for Inheritance

# Using a Wrapper Class

```
abstract class Iter[A] {  
  def getValue: Option[A]  
  def getNext: Iter[A]  
}
```

```
class ListIter[A](val list: List[A]) extends Iter[A] {  
  def getValue = list.headOption  
  def getNext = new ListIter(list.tail)  
}
```

```
sumElements((x: Int) => x)(new ListIter(List(1, 2, 3, 4)))
```

# MyTree Using ListIter

```
abstract class Iterable[A] {  
  def iter : Iter[A]  
}  
sealed abstract class MyTree[A] extends Iterable[A] {  
  override def iter : ListIter[A]  
}  
case class Empty[A]() extends MyTree[A] {  
  val iter : ListIter[A] = new ListIter(Ni)  
}  
case class Node[A](value: A,  
                   left: MyTree[A],  
                   right: MyTree[A])  
  extends MyTree[A] {  
  val iter : ListIter[A] = new ListIter(  
    value::(left.iter.list ++ right.iter.list))  
}
```



# Test

```
val t : MyTree[Int] =  
  Node(3, Node(4, Node(2, Empty(), Empty()),  
    Node(3, Empty(), Empty()))),  
  Node(5, Empty(), Empty()))  
  
sumElementsGen((x: Int) => x)(t)
```

# **Abstract Class With Associate Types**

# Using an Associate Type

```
abstract class Iterable[A] {  
  type iter_t  
  def iter: iter_t  
  def getValue(i: iter_t) : Option[A]  
  def getNext(i: iter_t) : iter_t  
}  
  
def sumElements[A](f:A=>Int)(xs: Iterable[A]) : Int = {  
  def sumElementsIter(i: xs.iter_t) : Int =  
    xs.getValue(i) match {  
      case None => 0  
      case Some(n) => f(n) + sumElementsIter(xs.getNext(i))  
    }  
  sumElementsIter(xs.iter)  
}
```

# MyTree Using List

```
sealed abstract class MyTree[A] extends Iterable[A] {
  type iter_t = List[A]
  def getValue(i: List[A]): Option[A] = i.headOption
  def getNext(i: List[A]): List[A] = i.tail
}

case class Empty[A]() extends MyTree[A] {
  val iter: List[A] = Nil
}

case class Node[A](value: A,
                   left: MyTree[A], right: MyTree[A])
  extends MyTree[A] {
  val iter = value :: (left.iter ++ right.iter) //Pre-order
  //val iter = left.iter ++ (value :: right.iter) // In-order
  //val iter = left.iter ++ (right.iter ++ List(value))
  //Post-order
}
```

# Test

```
val t : MyTree[Int] =  
  Node(3, Node(4, Node(2, Empty(), Empty()),  
    Node(3, Empty(), Empty()))),  
  Node(5, Empty(), Empty()))  
  
sumElements((x: Int) => x)(t)
```

# Abstract Class with Arguments

# Abstract Class with Arguments

```
abstract class IterableHE[A](eq: (A,A) => Boolean)
  extends Iterable[A]
{
  def hasElement(a: A) : Boolean = {
    def hasElementIter(i: iter_t) : Boolean =
      getValue(i) match {
        case None => false
        case Some(n) =>
          if (eq(a,n)) true
          else hasElementIter(getNext(i))
      }
    hasElementIter(iter)
  }
}
```

# MyTree

```
sealed abstract class MyTree[A](eq: (A, A) => Boolean)
  extends Iterable[A](eq) {
    type iter_t = List[A]
    def getValue(i: List[A]) : Option[A] = i.headOption
    def getNext(i: List[A]) : List[A] = i.tail
  }

case class Empty[A](eq: (A, A) => Boolean)
  extends MyTree[A](eq) {
    val iter : List[A] = Nil
  }

case class Node[A](eq: (A, A) => Boolean,
                  value: A, left: MyTree[A], right: MyTree[A])
  extends MyTree[A](eq) {
    val iter : List[A] = value :: (left.iter ++ right.iter)
  }
```



# Test

```
val leq = (x: Int, y: Int) => x == y
```

```
val lEmpty = Empty(leq)
```

```
def lNode(n: Int, t1: MyTree[Int], t2: MyTree[Int]) =  
  Node(leq, n, t1, t2)
```

```
val t : MyTree[Int] =  
  lNode(3, lNode(4, lNode(2, lEmpty, lEmpty),  
                    lNode(3, lEmpty, lEmpty)),  
        lNode(5, lEmpty, lEmpty))
```

```
sumElements((x: Int) => x)(t)
```

```
t.hasElement(5)
```

```
t.hasElement(10)
```

# Alternatively, Argument Elimination

```
abstract class IterableHE[A]  
  extends Iterable[A]  
{  
  def eq(a:A, b:A) : Boolean  
  def hasElement(a: A) : Boolean = {  
    def hasElementIter(i: iter_t) : Boolean =  
      getValue(i) match {  
        case None => false  
        case Some(n) =>  
          if (eq(a,n)) true  
          else hasElementIter(getNext(i))  
      }  
    hasElementIter(iter)  
  }  
}
```

# MyTree

```
sealed abstract class MyTree[A] extends Iterable[A] {  
  type iter_t = List[A]  
  def getValue(i : List[A]) : Option[A] = i.headOption  
  def getNext(i: List[A]) : List[A] = i.tail  
}  
  
case class Empty[A](_eq: (A,A)=>Boolean) extends MyTree[A] {  
  def eq(a:A, b:A) = _eq(a,b)  
  val iter : List[A] = Nil  
}  
  
case class Node[A](_eq: (A,A)=>Boolean,  
                  value: A, left: MyTree[A], right: MyTree[A])  
  extends MyTree[A] {  
  def eq(a:A, b:A) = _eq(a,b)  
  val iter : List[A] = value :: (left.iter ++ right.iter)  
}
```

# Test

```
val leq = (x: Int, y: Int) => x == y
```

```
val lEmpty = Empty(leq)
```

```
def lNode(n: Int, t1: MyTree[Int], t2: MyTree[Int]) =  
  Node(leq, n, t1, t2)
```

```
val t : MyTree[Int] =  
  lNode(3, lNode(4, lNode(2, lEmpty, lEmpty),  
                    lNode(3, lEmpty, lEmpty)),  
        lNode(5, lEmpty, lEmpty))
```

```
sumElements((x: Int) => x)(t)
```

```
t.hasElement(5)
```

```
t.hasElement(10)
```

## **More on Classes**

# Motivating Example

```
class Primes(val prime: Int, val primes: List[Int]) {  
  def getNext: Primes = {  
    val p = computeNextPrime(prime + 2)  
    new Primes(p, primes ++ (p :: Nil))  
  }  
  def computeNextPrime(n: Int) : Int =  
    if (primes.forall((p: Int) => n%p != 0)) n  
    else computeNextPrime(n+2)  
}
```

```
def nthPrime(n: Int): Int = {  
  def go(primes: Primes, k: Int): Int =  
    if (k <= 1) primes.prime  
    else go(primes.getNext, k - 1)  
  if (n <= 0) 2 else go(new Primes(3, List(3)), n)  
}  
nthPrime(10000)
```

# Multiple Constructors

```
class Primes(val prime: Int, val primes: List[Int]) {  
  def this() = this(3, List(3))  
  def getNext: Primes = {  
    val p = computeNextPrime(prime + 2)  
    new Primes(p, primes ++ (p :: Nil))  
  }  
  def computeNextPrime(n: Int) : Int =  
    if (primes.forall((p: Int) => n%p != 0)) n  
    else computeNextPrime(n+2)  
}
```

```
def nthPrime(n: Int): Int = {  
  def go(primes: Primes, k: Int): Int =  
    if (k <= 1) primes.prime  
    else go(primes.getNext, k - 1)  
  if (n == 0) 2 else go(new Primes, n)  
}  
nthPrime(10000)
```

# Access Modifiers

## ➤ Access Modifiers

- Private: Only the class can access the member.
- Protected: Only the class and its sub classes can access the member.



# Using Access Modifiers

```
class Primes private (val prime: Int, protected val primes: List[Int])  
{  
  def this() = this(3, List(3))  
  def getNext: Primes = {  
    val p = computeNextPrime(prime + 2)  
    new Primes(p, primes ++ (p :: Nil))  
  }  
  private def computeNextPrime(n: Int) : Int =  
    if (primes.forall((p: Int) => n%p != 0)) n  
    else computeNextPrime(n+2)  
}
```

```
def nthPrime(n: Int): Int = {  
  def go(primes: Primes, k: Int): Int =  
    if (k <= 1) primes.prime  
    else go(primes.getNext, k - 1)  
  if (n == 0) 2 else go(new Primes, n)  
}  
nthPrime(10000)
```

# Traits for Multiple Inheritance

# Multiple Inheritance Problem

## ➤ Multiple Inheritance

- The famous “diamond problem”

```
class A(val a: Int)
class B extends A(10)
class C extends A(20)
class D extends B, C.
```

Problem 1: What is the value of (new D).a ?

Problem 2: The constructor of A must be executed once because A may contain side effects such as sending messages over the network.

# Java's Solution: Interface

## ➤ Interface

- An interface cannot contain any implementation but only types of its methods.
- A class can inherit implementations from only one parent class but implement multiple interfaces.

# Scala's Solution: Trait

## ➤ Traits

- A trait can implement any of its methods, but should have only one constructor with no arguments.
- An **[abstract] class** (resp. **trait**)  $X$  can “extends” one trait or **[abstract] class** with **any** (resp. **no**) arguments “with” multiple traits  $T_1, \dots, T_n$  such that, for each  $i$ , the least superclass of  $T_i$ , if exists, should be a superclass of  $X$  where  
 $C$  is a superclass of  $T$  if  $C$  is an (abstract) class and  $T$  transitively “extends”  $C$ .
- No cyclic inheritance is allowed.

## ➤ Property

- Among the ancestors of a class, for the same class,
  - A constructor with arguments can appear at most once
  - A constructor with no argument can appear multiple times

# Example

```
class A(val a : Int) {  
  def this () = this(0)  
}  
trait B {  
  def f(x: Int): Int = x  
}  
trait C extends A with B {  
  def g(x: Int): Int = x + a  
}  
trait D extends B {  
  def h(x: Int): Int = f(x + 50)  
}  
class E extends A(10) with C with D {  
  override def f(x: Int) = x * a  
}  
  
val e = new E
```

# Algorithm for Multiple Inheritance

## ➤ Algorithm

- Give a linear order among all ancestors by “post-order” traversing without revisiting the same node.
- Invoke the constructors once in that order.

Note. Post-order traversal of a class C means

- Recursively post-order traverse C’s first parent; ...;
- Recursively post-order traverse C’s last parent; and
- Visit C.

By post-order traversing from “E” in the previous example, we have the order: A(10) → B → C → D → E

```
val e = new E
```

```
e.a // 10
```

```
e.f(100) // 100*10
```

```
e.g(100) // 100 + 10
```

```
e.h(100) // (100 + 50) * 10
```

- A constructor with arguments is always visited before the same constructor with no arguments.
- Compile error if the same field is implemented by multiple classes

# A Simple Example With Traits



# Motivation

```
abstract class Iter[A] {  
  def getValue: Option[A]  
  def getNext: Iter[A]  
}
```

```
class ListIter[A](val list: List[A]) extends Iter[A] {  
  def getValue = list.headOption  
  def getNext = new ListIter(list.tail)  
}
```

```
abstract class Dict[K,V] {  
  def add(k: K, v: V): Dict[K,V]  
  def find(k: K): Option[V]  
}
```

Q: How can we extend ListIter and implement Dict?

# Interface using Traits

```
// abstract class Dict[K,V] {  
//   def add(k: K, v: V): Dict[K,V]  
//   def find(k: K): Option[V] }
```

```
trait Dict[K,V] {  
  def add(k: K, v: V): Dict[K,V]  
  def find(k: K): Option[V]  
}
```

# Implementing Traits

```
class ListIterDict[K,V]  
  (eq: (K,K)=>Boolean, list: List[(K,V)])  
  extends ListIter[(K,V)](list)  
    with Dict[K,V]  
{  
  def add(k:K,v:V): ListIterDict[K,V] =  
    new ListIterDict(eq,(k,v)::list)  
  def find(k: K) : Option[V] = {  
    def go(l: List[(K, V)]): Option[V] = l match {  
      case Nil => None  
      case (k1, v1) :: tl =>  
        if (eq(k, k1)) Some(v1) else go(tl) }  
    go(list) }  
}
```

# Test

```
def sumElements[A](f: A=>Int)(xs: Iter[A]) : Int =  
  xs.getValue match {  
    case None => 0  
    case Some(n) => f(n) + sumElements(f)(xs.getNext)  
  }
```

```
def find3(d: Dict[Int,String]) = {  
  d.find(3)  
}
```

```
val d0 = new ListIterDict[Int,String]((x,y)=>x==y,Nil)  
val d = d0.add(4,"four").add(3,"three")
```

```
sumElements[(Int,String)](x=>x._1)(d)  
find3(d)
```

# **Mixin with Traits**

# Motivation: Mixin Functionality

```
abstract class Iter[A] {  
  def getValue: Option[A]  
  def getNext: Iter[A]  
}
```

```
class ListIter[A](val list: List[A]) extends Iter[A]  
{  
  def getValue = list.headOption  
  def getNext: ListIter[A] = new ListIter(list.tail)  
}
```

```
trait MRIter[A] extends Iter[A] {  
  def mapReduce[B,C](combine: (B,C)=>C, ival: C, f: A=>B): C = ???  
}
```

# Mixin Composition

```
trait MRIter[A] extends Iter[A] {  
  override def getNext: MRIter[A]  
  def mapReduce[B,C](combine: (B,C)=>C, ival: C, f: A=>B): C =  
    getValue match {  
      case None => ival  
      case Some(v) =>  
        combine(f(v), getNext.mapReduce(combine, ival, f))  
    }  
}
```

```
class MRListIter[A](list: List[A])  
  extends ListIter (list) with MRIter[A]  
{  
  override def getNext = new MRListIter(super.getNext.list)  
    // new MRListIter(list.tail)  
}
```

```
val mr = new MRListIter[Int](List(3,4,5))  
mr.mapReduce[Int,Int]((b,c)=>b+c,0,(a)=>a*a)
```

# Mixin Composition: A Better Way

```
trait MRIter[A] extends Iter[A] {  
  def mapReduce[B,C](combine: (B,C)=>C, ival: C, f: A=>B): C = {  
    def loop(c: Iter[A]): C = c.getValue match {  
      case None => ival  
      case Some(v) => combine(f(v), loop(c.getNext))  
    }  
    loop(this)  
  }  
}
```

```
class MRListIter[A](list: List[A])  
  extends ListIter(list) with MRIter[A]
```

```
val mr = new MRListIter[Int](List(3,4,5))
```

```
// or, val mr = new ListIter(List(3,4,5)) with MRIter[Int]
```

```
mr.mapReduce[Int,Int]((b,c)=>b+c,0,(a)=>a*a)
```



# Syntactic Sugar: new A with B with C { ... }

```
new A(...) with B1 ... with Bm {  
    code  
}
```

is equivalent to

```
{  
    class _tmp_(args) extends A(args) with B1 ... with Bm {  
        code  
    }  
    new _tmp_(...)  
}
```

# Intersection Types

# Intersection Types

## ➤ Typing Rule

$$\frac{t : T1 \quad t : T2}{t : T1 \text{ with } T2}$$

## ➤ Example

```
trait A { val a: Int = 0 }  
trait B { val b: Int = 0 }  
class C extends A with B {  
  override val a = 10  
  override val b = 20  
  val c = 30  
}
```

```
val x = new C  
val y: A with B = x
```

```
y.a // 10
```

```
y.b // 20
```

```
y.c // type error
```

# Subtype Relation for “with”

The subtype relation for “with” is structural.

- Permutation

=====

$$\dots \text{ with } T1 \text{ with } T2 \dots <: \dots \text{ with } T2 \text{ with } T1 \dots$$

- Width

=====

$$\dots \text{ with } T \dots <: \dots \dots$$

- Depth

=====

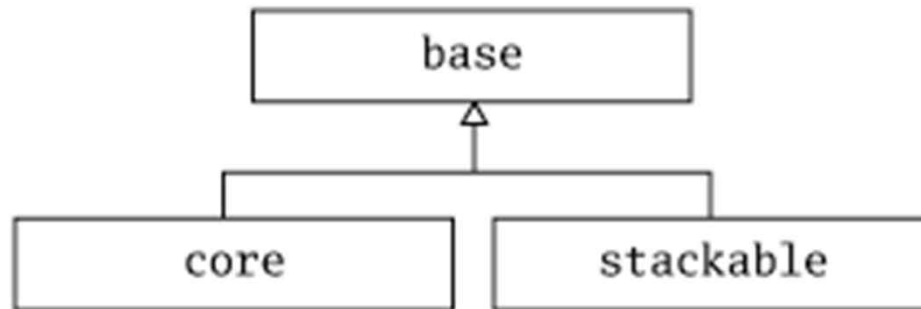
$$T <: S$$

=====

$$\dots \text{ with } T \dots <: \dots \text{ with } S \dots$$

# Stacking with Traits

# Typical Hierarchy in Scala



- **BASE**  
Interface (trait or abstract class)
- **CORE**  
Functionality (trait or concrete class)
- **CUSTOM**  
Modifications (each in a separate, composable trait)

# IntStack: Base

## ➤ BASE

```
trait Stack[A] {  
  def get(): (A, Stack[A])  
  def put(x: Int): Stack[A]  
}
```

# IntStack: Core

## ➤CORE

```
class BasicIntStack protected (xs: List[Int]) extends Stack[Int]
{
  override val toString = "Stack:" + xs.toString
  def this() = this(Nil)

  def get():(Int,Stack[Int]) = (xs.head,new BasicIntStack(xs.tail))
  def put(x:Int): Stack[Int] = new BasicIntStack(x :: xs)
}
```

```
val s0 = new BasicIntStack
val s1 = s0.put(3)
val s2 = s1.put(-2)
val s3 = s2.put(4)
val (v1,s4) = s3.get()
val (v2,s5) = s4.get()
```



# IntStack: Custom Modifications

## ➤CUSOM

```
trait Doubling extends Stack[Int] {  
  abstract override def put(x: Int): Stack[Int] = super.put(2 * x)  
}
```

```
trait Incrementing extends Stack[Int] {  
  abstract override def put(x: Int): Stack[Int] = super.put(x + 1)  
}
```

```
trait Filtering extends Stack[Int] {  
  abstract override def put(x: Int): Stack[Int] =  
    if (x >= 0) super.put(x) else this  
}
```

# IntStack: Stacking

## ➤ Stacking

```
class DIFIntStack protected (xs: List[Int])  
  extends BasicIntStack(xs)  
  with Doubling with Incrementing with Filtering  
{  
  def this() = this(Nil)  
}
```

```
val s0 = new DIFIntStack  
val s1 = s0.put(3)  
val s2 = s1.put(-2)  
val s3 = s2.put(4)  
val (v1,s4) = s3.get()  
val (v2,s5) = s4.get()  
val (v2,s6) = s5.get()
```

# IntStack: Core (Correct)

## ➤ CORE

```
class BasicIntStack protected (xs: List[Int]) extends Stack[Int]
{
  override val toString = "Stack:" + xs.toString
  def this() = this(Nil)
  protected def mkStack(xs: List[Int]): Stack[Int] =
    new BasicIntStack(xs)
  def get(): (Int, Stack[Int]) = (xs.head, mkStack(xs.tail))
  def put(x: Int): Stack[Int] = mkStack(x :: xs)
}
```

```
val s0 = new BasicIntStack
val s1 = s0.put(3)
val s2 = s1.put(-2)
val s3 = s2.put(4)
val (v1,s4) = s3.get()
val (v2,s5) = s4.get()
```

# IntStack: Stacking (Correct)

## ➤ Stacking

```
class DIFIntStack protected (xs: List[Int])  
  extends BasicIntStack(xs)  
  with Doubling with Incrementing with Filtering  
{  
  def this() = this(Nil)  
  override def mkStack(xs: List[Int]): Stack[Int] =  
    new DIFIntStack(xs)  
}
```

```
val s0 = new DIFIntStack  
val s1 = s0.put(3)  
val s2 = s1.put(-2)  
val s3 = s2.put(4)  
val (v1,s4) = s3.get()  
val (v2,s5) = s4.get()
```

# Additional Resources

## ➤ Traits

- <http://www.scala-lang.org/old/node/126>

## ➤ Mixin Composition

- <http://www.scala-lang.org/old/node/117>

## ➤ Stackable Trait Pattern

- [http://www.artima.com/scalazine/articles/stackable\\_trait\\_pattern.html](http://www.artima.com/scalazine/articles/stackable_trait_pattern.html)

## ➤ Multiple Inheritance via Traits

- <https://www.safaribooksonline.com/blog/2013/05/30/traits-how-scala-tames-multiple-inheritance/>

## ➤ UCSD CSE 130

- <http://cseweb.ucsd.edu/classes/wi14/cse130-a/lectures/scala/02-classes.html>

# **PART 3**

## **Type Classes for Interfaces**

# Problems with OOP

# Subtype Polymorphism

```
trait Ord {  
  // this cmp that < 0 iff this < that  
  // this cmp that > 0 iff this > that  
  // this cmp that == 0 iff this == that  
  def cmp(that: Ord): Int  
  
  def ==(that: Ord): Boolean = (this.cmp(that)) == 0  
  def < (that: Ord): Boolean = (this.cmp(that) < 0  
  def > (that: Ord): Boolean = (this.cmp(that) > 0  
  def <= (that: Ord): Boolean = (this.cmp(that) <= 0  
  def >= (that: Ord): Boolean = (this.cmp(that) >= 0  
}  
  
def max3(a: Ord, b: Ord, c: Ord) : Ord =  
  if (a <= b) { if (b <= c) c else b }  
  else      { if (a <= c) c else a }
```

\* Problem: hard (almost impossible) to implement Ord (e.g., using Int)



# Interface over Parameter Types

```
trait Ord[A] {  
  def cmp(that: A): Int  
  
  def ==(that: A): Boolean = (this.cmp(that)) == 0  
  def < (that: A): Boolean = (this.cmp that) < 0  
  def > (that: A): Boolean = (this.cmp that) > 0  
  def <= (that: A): Boolean = (this.cmp that) <= 0  
  def >= (that: A): Boolean = (this.cmp that) >= 0  
}  
  
def max3[A <: Ord[A]](a: A, b: A, c: A) : A =  
  if (a <= b) { if (b <= c) c else b }  
  else      { if (a <= c) c else a }  
  
class OInt(val value : Int) extends Ord[OInt] {  
  def cmp(that: OInt) = value - that.value  
}  
  
max3(new OInt(3), new OInt(2), new OInt(10)).value
```

## Further example: Ordered Bag

```
class Bag[U <: Ord[U]] protected (val toList: List[U]) {  
  def this() = this(Nil)  
  def add(x: U) : Bag[U] = {  
    def go(elmts: List[U]): List[U] =  
      elmts match {  
        case Nil => x :: Nil  
        case e :: _ if (x < e) => x :: elmts  
        case e :: _ if (x == e) => elmts  
        case e :: rest => e :: go(rest)  
      }  
    new Bag(go(toList))  
  }  
}  
  
val emp = new Bag[0Int]()  
val b = emp.add(new 0Int(3)).add(new 0Int(2)).  
  add(new 0Int(10)).add(new 0Int(2))  
b.toList.map((x)=>x.value)
```

# Problems with OOP

1. Needs “subtyping” like “`OInt <: Ord[OInt]`”, which is quite complex as we have seen (and moreover, involves more complex concepts like variance).
2. Needs a wrapper class like “`OInt`” in order to add a new interface to an existing type like “`Int`”.
3. Interface only contains only “elimination” functions, not “introduction” functions.
4. No canonical operator
5. ...

# Type Classes

# Completely Separating Ord from Int

```
trait Ord[A] {  
  def cmp(me: A, you: A): Int  
  
  def ==(me: A, you: A): Boolean = cmp(me, you) == 0  
  def < (me: A, you: A): Boolean = cmp(me, you) < 0  
  def > (me: A, you: A): Boolean = cmp(me, you) > 0  
  def <= (me: A, you: A): Boolean = cmp(me, you) <= 0  
  def >= (me: A, you: A): Boolean = cmp(me, you) >= 0  
}  
  
def max3[A](a: A, b: A, c: A)(implicit ord: Ord[A]) : A =  
  if (ord.<=(a, b)) {if (ord.<=(b, c)) c else b }  
  else {if (ord.<=(a, c)) c else a }  
  
// behaves like Int <: Ord in OOP  
implicit val intOrd : Ord[Int] = new Ord[Int] {  
  def cmp(me: Int, you: Int) = me - you }  
max3(3, 2, 10) // 10
```

# Implicit

## ➤ Implicit

- An argument is given “implicitly”

```
def foo(s: String)(implicit t: String) = s + t
```

```
implicit val exclamation : String = "!!!!!!"
```

```
foo("Hi")
```

```
foo("Hi")("???) // can give it explicitly
```

# Syntax for type class: syntactic sugar

```
trait Ord[A]:  
  extension (a:A)  
    def cmp(b:A): Int  
    def ===(b:A) = a.cmp(b) == 0  
    def < (b:A) = a.cmp(b) < 0  
    def > (b:A) = a.cmp(b) > 0  
    def <= (b:A) = a.cmp(b) <= 0  
    def >= (b:A) = a.cmp(b) >= 0  
  
def max3[A: Ord](a: A, b: A, c: A) : A =  
  if (a<=b) {if (b<=c) c else b }  
  else    {if (a<=c) c else a }  
  
given intOrd : Ord[Int] with  
  extension (a: Int)  
    def cmp(b: Int) = a - b
```

# Bag Example using type class

```
class Bag[A: Ord] protected (val toList: List[A])
{ def this() = this(Nil)
  def add(x: A) : Bag[A] = {
    def loop(elmts: List[A]) : List[A] =
      elmts match {
        case Nil => x :: Nil
        case e :: _ if (x < e) => x :: elmts
        case e :: _ if (x == e) => elmts
        case e :: rest => e :: loop(rest)
      }
    new Bag(loop(toList))
  }
}
```

```
(new Bag[Int]()).add(3).add(2).add(3).add(10).toList
```



# Bootstrapping Implicits

*// lexicographic order*

**given** tupOrd[A, B](using Ord[A], Ord[B]): Ord[(A,B)] **with**

extension (a: (A,B))

**def** cmp(b: (A, B)) : Int = {

**val** c1 = a.\_1.cmp(b.\_1)

**if** (c1 != 0) c1

**else** { a.\_2.cmp(b.\_2) }

}

**val** b = **new** Bag[(Int,(Int,Int))]

b.add((3,(3,4))).add((3,(2,7))).add((4,(0,0))).toList

# With Different Orders

```
// given intOrd : Ord[Int] with
```

```
//  extension (a: Int)
```

```
//    def cmp(b: Int) = a - b
```

```
def intOrdRev : Ord[Int] = new {
```

```
  extension (a: Int)
```

```
    def cmp(b: Int) = b - a
```

```
}
```

```
(new Bag[Int]()).add(3).add(2).add(10).toList
```

```
(new Bag[Int]()(intOrdRev)).add(3).add(2).add(10).toList
```

# Type Classes: Abstraction

# Interfaces I: elimination

```
trait Iter[I,A]:  
  extension (i: I)  
    def getValue: Option[A]  
    def getNext: I
```

```
trait Iterable[I,A]:  
  type Itr  
  given Itrl: Iter[Itr,A]  
  extension (i: I)  
    def iter: Itr
```

// behaves like `Iter[A] <: Iterable[A]` in OOP

```
given iter2iterable[I,A](using iterI: Iter[I,A]): Iterable[I,A] with  
  type Itr = I  
  def Itrl = iterI  
  extension (i:I)  
    def iter = i
```

# Programs for Testing: use Iter, Iterable

```
def sumElements[I](xs: I)(implicit ll: Iterable[I, Int]) = {  
  def loop(i: ll.Itr): Int =  
    i.getValue match {  
      case None => 0  
      case Some(n) => n + loop(i.getNext)  
    }  
  loop(xs.iter)  
}
```

```
def printElements[I, A](xs: I)(implicit ll: Iterable[I, A]) = {  
  def loop(i: ll.Itr): Unit =  
    i.getValue match {  
      case None =>  
      case Some(a) => {println(a); loop(i.getNext)}  
    }  
  loop(xs.iter)  
}
```

# Interfaces II: introduction + elimination

```
trait Listlike[L,A]:  
  extension(u:Unit)  
    def unary_! : L  
  extension(elem:A)  
    def ::(l: L): L  
  extension(l: L)  
    def head: Option[A]  
    def tail: L  
    def ++(l2: L): L  
  
trait Treelike[T,A]:  
  extension(u:Unit)  
    def unary_! : T  
  extension(a:A)  
    def node(lt: T, rt: T): T  
  extension(t: T)  
    def root : Option[A]  
    def left : T  
    def right : T
```

# Programs for Testing: use All

```
def testList[L](implicit ListI: Listlike[L,Int], IterI: Iterable[L,Int]) = {  
  val l = (3 :: !()) ++ (1 :: 2 :: !())  
  println(sumElements(l)) //sumElements(l)(listIter[Int])  
  printElements(l) //printElements(l)(listIter[Int])  
}
```

```
def testTree[T](implicit TreeI: Treelike[T,Int], IterI: Iterable[T,Int]) = {  
  val t = 3.node(4.node(!(), !()), 2.node(!(),!()))  
  println(sumElements(t))  
  printElements(t)  
}
```

# Implement Iter and Listlike for List

*// behaves like List[A] <: Iter[A] in OOP*

```
given listIter[A]: Iter[List[A],A] with
  extension (l: List[A])
    def getValue = l.headOption
    def getNext = l.tail
```

*// behaves like List[A] <: Listlike[A] in OOP*

```
given listListlike[A]: Listlike[List[A],A] with
  extension (u: Unit)
    def unary_! = Nil
  extension (a: A)
    def ::(l: List[A]) = a::l
  extension (l: List[A])
    def head = l.headOption
    def tail = l.tail
    def ++(l2: List[A]) = l ::: l2
```



# Implement Iterable for MyTree using Listlike, Iter

```
enum MyTree[+A]:  
  case Leaf  
  case Node(value: A, left: MyTree[A], right: MyTree[A])  
import MyTree._  
  
// behaves like MyTree[A] <: Iterable[A], but clumsy in OOP  
given treeIterable[L, A](using listl: Listlike[L, A], iterl: Iter[L, A])  
  : Iterable[MyTree[A], A] with  
  type Itr = L  
  def ltrl = iterl  
  extension (t: MyTree[A])  
    def iter: L = t match {  
      case Leaf => !()  
      case Node(v, lt, rt) => v :: (ltrl.iter ++ rtl.iter)  
    }  
}
```

# Implement Treelike for MyTree

```
// behaves like MyTree[A] <: Treelike[A] in OOP
given mytreeTreelike[A] : Treelike[MyTree[A],A] with
  extension (u: Unit)
    def unary_! = Leaf
  extension (a: A)
    def node(l: MyTree[A], r: MyTree[A]) = Node(a,l,r)
  extension (t: MyTree[A])
    def root = t match {
      case Leaf => None
      case Node(v,_,_) => Some(v)
    }
    def left = t match {
      case Leaf => t
      case Node(_,lt,_) => lt
    }
    def right = t match {
      case Leaf => t
      case Node(_,_,rt) => rt }
```

# Linking Modules

```
testList[List[Int]]
```

```
testTree[MyTree[Int]]
```

**Type class: Code Reuse**

# IntStack Spec

```
trait Stack[S,A]:  
  extension (u: Unit)  
    def empty : S  
  extension (s: S)  
    def get: (A,S)  
    def put(a: A): S
```

```
def testStack[S](implicit Stkl: Stack[S,Int]) = {  
  val s0 = ().empty  
  val s1 = s0.put(3)  
  val s2 = s1.put(-2)  
  val s3 = s2.put(4)  
  val (v1,s4) = s3.get  
  val (v2,s5) = s4.get  
  (v1,v2)  
}
```

# Implementation using List

```
given BasicStack[A] : Stack[List[A],A] with  
  extension (u: Unit)  
    def empty = List()  
  extension (s: List[A])  
    def get = (s.head, s.tail)  
    def put(a: A) = a :: s
```

# Modifying Traits

```
def StackOverridePut[S,A](newPut: (S,A)=>S)(implicit stkl: Stack[S,A]) =  
  new Stack[S,A] {  
    extension (u: Unit)  
      def empty = stkl.empty(u)  
    extension (s: S)  
      def get = stkl.get(s)  
      def put(a: A) = newPut(s,a)  
  }
```

```
def Doubling[S](implicit stkl: Stack[S,Int]) : Stack[S,Int] =  
  StackOverridePut((s,a) => s.put(2 * a))
```

```
def Incrementing[S](implicit stkl: Stack[S,Int]) : Stack[S,Int] =  
  StackOverridePut((s,a) => s.put(a + 1))
```

```
def Filtering[S](implicit stkl: Stack[S,Int]) : Stack[S,Int] =  
  StackOverridePut((s,a) => if (a >= 0) s.put(a) else s)
```

# Linking

```
// testStack(BasicStack)
```

```
testStack
```

```
// testStack(Filtering(Incrementing (Doubling(BasicStack))))
```

```
testStack(Filtering(Incrementing (Doubling)))
```

```
// testStack(Filtering(Incrementing(Incrementing(Doubling(BasicStack))))))
```

```
testStack(Filtering(Incrementing(Incrementing(Doubling))))
```



# Implementation: Sorted Stack

```
def SortedStack : Stack[List[Int],Int] = new {  
  extension (u: Unit)  
    def empty = List()  
  extension (s: List[Int])  
    def get = (s.head, s.tail)  
    def put(a: Int) : List[Int] = {  
      def loop(l: List[Int]) : List[Int] = l match {  
        case Nil => a :: Nil  
        case hd :: tl => if (a <= hd) a :: l else hd :: loop(tl)  
      }  
      loop(s)  
    }  
  }  
}  
  
testStack(Filtering(Incrementing(Doubling(SortedStack))))
```

# A trick for reducing boiler plate code

```
//def StackOverridePut[S,A](newPut: (S,A)=>S)(implicit stkl: Stack[S,A]) =  
// new Stack[S,A] {  
//   extension (u: Unit)  
//     def empty = stkl.empty(u)  
//   extension (s: S)  
//     def get = stkl.get(s)  
//     def put(a: A) = newPut(s,a) }  
class mkStack[S,A](stkl: Stack[S,A]) extends Stack[S,A]:  
  extension (u: Unit)  
    def empty = stkl.empty(u)  
  extension (s: S)  
    def get = stkl.get(s)  
    def put(a: A) = stkl.put(s)(a)  
def StackOverridePut[S,A](newPut: (S,A)=>S)(implicit stkl: Stack[S,A]) =  
  new mkStack(stkl) {  
    extension (s:S)  
      override def put(a: A) = newPut(s, a)  
  }
```

# Higher Type Classes

# Interfaces I

// eg. Iter[List]

```
trait Iter[I[_]]:  
  extension [A](i: I[A])  
    def getValue: Option[A]  
    def getNext: I[A]
```

```
// trait Iter[I,A]:  
//   extension (i: I)  
//     def getValue: Option[A]  
//     def getNext: I
```

// eg. Iterable[MyTree]

```
trait Iterable[I[_]]:  
  type Itr[_]  
  given Itrl: Iter[Itr]  
  extension [A](i: I[A])  
    def iter: Itr[A]
```

```
// trait Iterable[I,A]:  
//   type Itr  
//   given Itrl: Iter[Itr,A]  
//   extension (i: I)  
//     def iter: Itr
```

```
given iter2iterable[I[_]](using iterI: Iter[I]): Iterable[I] with  
  type Itr[A] = I[A]  
  def Itrl = iterI  
  extension [A](i: I[A])  
    def iter = i
```

# Programs for Testing: use Iter, Iterable

```
def sumElements[I[_]](xs: I[Int])(implicit ll: Iterable[I]) = {  
  def loop(i: ll.Itr[Int]): Int =  
    i.getValue match {  
      case None => 0  
      case Some(n) => n + loop(i.getNext)  
    }  
  loop(xs.iter)  
}
```

```
def printElements[I[_], A](xs: I[A])(implicit ll: Iterable[I]) = {  
  def loop(i: ll.Itr[A]): Unit =  
    i.getValue match {  
      case None =>  
      case Some(a) => {println(a); loop(i.getNext)}  
    }  
  loop(xs.iter)  
}
```

# Interfaces II

```
trait Listlike[L[_]]:  
  extension[A](u:Unit)  
    def unary_! : L[A]  
  extension[A](elem:A)  
    def ::(l: L[A]): L[A]  
  extension[A](l: L[A])  
    def head: Option[A]  
    def tail: L[A]  
    def ++(l2: L[A]): L[A]  
  
trait Treelike[T[_]]:  
  extension[A](u:Unit)  
    def unary_! : T[A]  
  extension[A](a:A)  
    def node(lt: T[A], rt: T[A]): T[A]  
  extension[A](t: T[A])  
    def root : Option[A]  
    def left : T[A]  
    def right : T[A]
```

# Programs for Testing: use All

```
def testList[L[_]](implicit ListI: Listlike[L], ltrl: Iterable[L]) = {  
  val l = (3 :: !()) ++ (1 :: 2 :: !())  
  println(sumElements(l))  
  printElements(l)  
}
```

```
def testTree[T[_]](implicit Treel: Treelike[T], ltrl: Iterable[T]) = {  
  val t = 3.node(4.node(!(), !()), 2.node(!(),!()))  
  println(sumElements(t))  
  printElements(t)  
}
```

# List: provide Iter, ListIF

*// behaves like List[A] <: Iter[A] in OOP*

```
given listIter: Iter[List] with
  extension [A](l: List[A])
    def getValue = l.headOption
    def getNext = l.tail
```

*// behaves like List[A] <: Listlike[A] in OOP*

```
given listListlike: Listlike[List] with
  extension [A](u: Unit)
    def unary_! = Nil
  extension [A](a: A)
    def ::(l: List[A]) = a::l
  extension [A](l: List[A])
    def head = l.headOption
    def tail = l.tail
    def ++(l2: List[A]) = l ::: l2
```



# MyTree: use Iter, ListIF, provide Iterable, TreeIF

```
enum MyTree[+A]:  
  case Leaf  
  case Node(value: A, left: MyTree[A], right: MyTree[A])  
import MyTree._
```

*// behaves like MyTree[A] <: Iterable[A], but clumsy in OOP*

```
given treeIterable[L[_]](using listl: Listlike[L], iterl: Iter[L]): Iterable[MyTree]  
with  
  type Itr[A] = L[A]  
  def Itrl = iterl  
  extension [A](t: MyTree[A])  
    def iter: L[A] = t match {  
      case Leaf => !()  
      case Node(v, lt, rt) => v :: (lt.iter ++ rt.iter)
```

# MyTree: use Iter, ListIF, provide Iterable, TreeIF

*// behaves like MyTree[A] <: Treelike[A] in OOP*

given mytreeTreelike: Treelike[MyTree] with

extension [A](u: Unit)

def unary\_! = Leaf

extension [A](a: A)

def node(l: MyTree[A], r: MyTree[A]) = Node(a,l,r)

extension [A](t: MyTree[A])

def root = t match {

case Leaf => None

case Node(v,\_,\_) => Some(v)

}

def left = t match {

case Leaf => t

case Node(\_,lt,\_) => lt

}

def right = t match {

case Leaf => t

case Node(\_,\_,rt) => rt }

# Linking Modules

testList[List]

testTree[MyTree]

# List with Map

```
trait MapListlike[L[_]]:  
  given LLI: Listlike[L]  
  extension[A](l: L[A])  
    def map[B](f: A => B): L[B]
```

```
def testMapList[L[_]](implicit LI: MapListlike[L], IT: Iter[L]) = {  
  val l1 = 3.3 :: 2.2 :: 1.5 :: !()  
  val l2 = l1.map((n:Double)=>n.toInt)  
  val l3 = l2.map((n:Int)=>n.toString)  
  printElements(l3)  
}
```

# List with Map

```
given listMapListlike(using LLI: Listlike[List]): MapListlike[List] with
  def ListlikeI = LLI
  extension [A](l: List[A])
    def map[B](f: A => B) = l.map(f)
```

```
testMapList[List]
```

# Turning Type Classes into OO Classes

# Interfaces

```
trait DataProcessor[D]:  
  extension (d: D)  
    def input(s: String) : D  
    def output : String
```

```
trait DPFactory:  
  extension (u: Unit)  
    def getTypes: List[String]  
    def makeDP(dptype: String) : ???
```

```
def run(implicit factory: DPFactory) : Unit
```

How to return data with associated functions like OOP?

# Turning Type Classes into OO Classes

```
import scala.language.implicitConversions
```

```
trait Box[S[_]:  
  type Data  
  val * : Data  
  given DI: S[Data]
```

```
object Box {  
  implicit // needed for implicit conversion of D into Box[S]  
  def apply[S[_],D](d: D)(implicit i: S[D]): Box[S] =  
    new Box[S] {  
      type Data = D  
      val * = d  
      val DI = i  
    }  
}
```



# Interfaces

```
trait DataProcessor[D]:  
  extension (d: D)  
    def input(s: String): D  
    def output: String
```

```
trait DPFactory:  
  extension (u: Unit)  
    def getTypes: List[String]  
    def makeDP(dptype: String): Box[DataProcessor]
```

```
trait UserInteraction:  
  def run(implicit factory: DPFactory): Unit
```

# User Interaction

```
val userInteraction: UserInteraction = new UserInteraction {  
  def run(implicit factory: DPFactory) = {  
    val dptype = scala.io.StdIn.readLine(  
      "Input a processor type " + ().getTypes.toString + ": ")  
    val d = ().makeDP(dptype)  
    val d_done = getInputs(d)  
    printOutputs(d_done.*)  
  }  
  def getInputs(d: Box[DataProcessor]): Box[DataProcessor] = {  
    val d2 = d.*.input(scala.io.StdIn.readLine("Input Data: "))  
    val done = scala.io.StdIn.readLine("More inputs? [Y/N]: ")  
    if (done.toLowerCase() == "n") d2  
    else getInputs(d2)  
  }  
  def printOutputs[D](d: D)(implicit DP: DataProcessor[D]) = {  
    println("The result of processing your inputs is:")  
    println(d.output)  
  }  
}
```

# Data Processor

```
implicit val dpfactory: DPFactory = new {  
  extension (u: Unit)  
    def getTypes = List("sum", "mult")  
    def makeDP(dptype: String) = {  
      if (dptype == "sum")  
        makeProc(0, (x, y) => x + y)  
      else  
        makeProc(1, (x, y) => x * y)  
    }  
  
    def makeProc(init: Int, op: (Int, Int) => Int): Box[DataProcessor] = {  
      given dp: DataProcessor[Int] with  
        extension (d: Int)  
          def input(s: String) = op(d, s.toInt)  
          def output = d.toString()  
      init // Box.apply[Int, DataProcessor](init)(dp)  
    }  
}
```

# Linking

```
userInteraction.run(dpfactory)
```

# Heterogeneous List of Iter

```
trait Iter[I,A]:  
  extension (i: I)  
    def getValue: Option[A]  
    def getNext: I  
  
def sumElements[I](xs: I)(implicit I: Iter[I,Int]) : Int = {  
  xs.getValue match {  
    case None => 0  
    case Some(n) => n + sumElements(xs.getNext)  
  }  
}  
  
def sumElementsList(xs: List[Box2[Iter,Int]]) : Int =  
  xs match {  
    case Nil => 0  
    case hd :: tl => sumElements(hd.*) + sumElementsList(tl)  
  }
```

# Turning Type Classes into OO Classes

```
import scala.language.implicitConversions
```

```
trait Box2[S[_],A]:  
  type Data  
  val * : Data  
  given DI: S[Data,A]
```

```
object Box2 {  
  implicit def apply[S[_], D,A](d: D)(implicit i: S[D,A]): Box2[S,A] =  
    new Box2[S,A] {  
      type Data = D  
      val * = d  
      val DI = i  
    }  
}
```

# Test

```
given listIter[A]: Iter[List[A],A] with  
  extension (l: List[A])  
    def getValue = l.headOption  
    def getNext = l.tail
```

```
given declter : Iter[Int,Int] with  
  extension (i: Int)  
    def getValue = if (i >= 0) Some(i) else None  
    def getNext = i - 1
```

```
sumElementsList(List(  
  100,  
  List(1,2,3),  
  10))
```

# **PART 4**

## **Imperative Programming with Memory Updates**



# Mutable Variables

## ➤ Mutable Variables

- Use “var” instead of “val” and “def”
- We can update the value stored in a variable.

```
class Main(i: Int) {  
    var a = i  
}
```

```
val m = new Main(10)  
m.a    // 10  
m.a = 20  
m.a    // 20  
m.a += 5    // m.a = m.a + 5  
m.a    // 25
```

# While loop

## ➤ While loop

- Syntax: **while** (*cond*) *body*  
Executes *body* while *cond* holds.
- It is equivalent to:

```
def mywhile(cond: =>Boolean)(body: =>Unit) : Unit =  
  if (cond) { body; mywhile(cond)(body) } else ()
```

## ➤ Example

```
var i = 0  
var sum = 0  
while (i <= 100) { // mywhile (i <= 100) {  
  sum += i  
  i += 2  
}  
sum // 2550
```

# For loop

## ➤ For loop

- Syntax: `for (i <- collection) body`  
Executes *body* for each *i* in *collection*.
- It is equivalent to:

```
def myfor[A](xs: Traversable[A])(f: A => Unit) : Unit =  
  xs.foreach(f)
```

## ➤ Example

```
var sum = 0  
for (i <- 0 to 100 by 2) { // myfor (0 to 100 by 2) { i =>  
  sum += i  
}  
sum // 2550
```

# Additional Resources

## ➤ UCSD CSE 130

- <http://cseweb.ucsd.edu/classes/wi14/cse130-a/lectures/scala/00-crash.html>
- <http://cseweb.ucsd.edu/classes/wi14/cse130-a/lectures/scala/01-iterators.html>

Thanks for your hard work!