

PART 4

Imperative Programming with Memory Updates

Mutable Update in Scala

Mutable Variables

➤ Mutable Variables

- Use “var” instead of “val” and “def”
- We can update the value stored in a variable.

```
class Main(i: Int) {  
  var a = i  
}
```

```
val m = new Main(10)  
m.a    // 10  
m.a = 20  
m.a    // 20  
m.a += 5    // m.a = m.a + 5  
m.a    // 25
```

While loop

➤ While loop

- Syntax: **while** (*cond*) *body*
Executes *body* while *cond* holds.
- It is equivalent to:

```
def mywhile(cond: =>Boolean)(body: =>Unit) : Unit =  
  if (cond) { body; mywhile(cond)(body) } else ()
```

➤ Example

```
var i = 0  
var sum = 0  
while (i <= 100) { // mywhile (i <= 100) {  
  sum += i  
  i += 2  
}  
sum // 2550
```

For loop

➤ For loop

- Syntax: **for** (i <- *collection*) *body*
Executes *body* for each i in *collection*.
- It is equivalent to:

```
def myfor[A](xs: Traversable[A])(f: A => Unit) : Unit =  
  xs.foreach(f)
```

➤ Example

```
var sum = 0  
for (i <- 0 to 100 by 2) { // myfor (0 to 100 by 2) { i =>  
  sum += i  
}  
sum // 2550
```

Immutability, Mutability & Ownership

Immutability for Guarantee & Mutability for Efficiency

➤ Immutable array

```
arr = Array.new([1;2;3])
```

```
... using arr ...
```

```
arr' = Array.set(arr, 0, 42)
```

```
f(arr')
```

```
... using arr' ...
```

```
g()
```

```
... using arr' ...
```

➤ Mutable array

```
arr = Array.new([1;2;3])
```

```
... using arr ...
```

```
Array.set(arr, 0, 42);
```

```
f(arr)
```

```
... using arr ...
```

```
g()
```

```
... using arr ...
```

Can we not have both? Mutability with Ownership!

➤ Immutable array

```
arr = Array.new([1;2;3])
```

```
... using arr ...
```

```
arr' = Array.set(arr, 0, 42)
```

```
f(arr')
```

```
... using arr' ...
```

```
g()
```

```
... using arr' ...
```

➤ Mutation with Ownership

```
// own
```

```
arr = Array.new([1;2;3])
```

```
... using arr ...
```

```
// mutable borrow
```

```
Array.set(&mut arr, 0, 42)
```

```
// immutable borrow
```

```
f(&arr)
```

```
... using arr ...
```

```
g()
```

```
... using arr ...
```


Types with Ownership

➤ Ownership Types

- expresses and guarantees immutability
 - making code behavior more predictable
- automatically deallocates memory when its ownership has gone
 - guaranteeing absence of use-after-free, double-free, memory-leak
- disallows mutating the same value at the same time
 - guaranteeing data-race freedom in concurrent code

Programming with Mutation in a Principled Way!

The Rust Programming Language Book

<https://doc.rust-lang.org/book/>

Types, Values & Memory

Values and Types

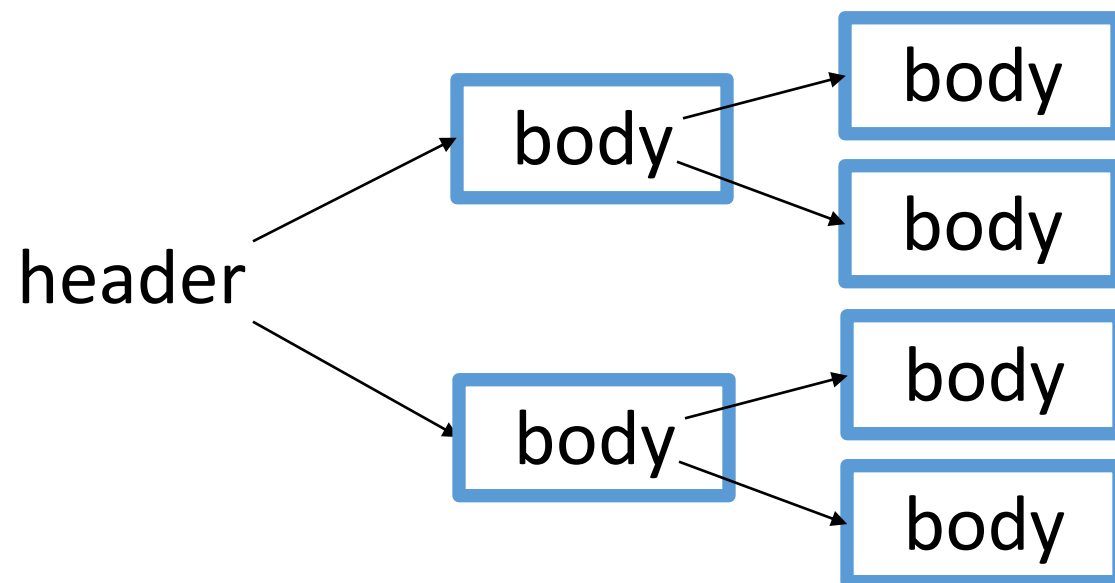
A type defines a set of values.

- For functionality, we only need to know the denotation of values.
- For efficiency, we also need to know the shape of values.

A value has a tree structure and consists of a header and bodies.

- A header is data that may contain scalar values and pointers to bodies.
- A body is data that must be stored in memory and may contain scalar values and pointers to other bodies.
- When passing or storing a value, only its header is passed or stored.

Types in Rust: `i64`, `u64`, `(T1,T2)`, `[T; 5]`, `String`, `Vec<T>`, ...



Stack memory with static size and static lifetime

Stack memory consists of variables of **sized types** (ie, static header sizes) with **scopes** (ie, static lifetime).

A stack variable of a type T stores a header of the type T.

```
// variable arg of type i64 (8-byte header) with Scope 1
fn foo(arg: i64) -> i64 {
    let mut x = 10; // variable x of type i64 with Scope 1
    {
        let y = arg + x; // variable y of type i64 with Scope 2
        x = x + y;
    } // Scope 2
    let x2 = x; // variable x2 of type i64 with Scope 1
    let z = arg + x2; // variable z of type i64 with Scope 1
    return z;
} // Scope 1
```

Heap memory with dynamic size and dynamic lifetime

Heap memory consists of blocks of **all types (ie, dynamic sizes)** with **no scopes (ie, dynamic lifetime)**.

A heap block can be a body of a value, which can also store headers of other values.

```
// variable sz of type usize (8-byte header) with Scope 1
fn gee(sz: usize) -> Vec<i64> {
    // variable v of type Vec<i64> (24-byte header) with Scope 1
    // The header points to a heap block of (sz * 8) bytes, filled with zeros
    let v : Vec<i64> = vec![0; sz];
    // The header goes out of Scope 1 and the heap block is still alive
    return v;
} // Scope 1
```

Principles of Ownership & Lifetime

Issues with mutation and deallocation for memory

<Problematic mutations>

Multiple parties modify a value logically simultaneously.

- Problem

The value may become inconsistent despite valid individual updates.

A value is read logically simultaneously while being modified.

- Problem

An inconsistent intermediate state may be read.

<Problematic deallocation>

A value's body is deallocated while its header remains accessible.

- Problem

The deallocated body may be accessed via the header (use-after-free bug).

Principles of Ownership

<Ownership rules>

- A header must be stored in exactly one location, which owns the header.
- Every heap block must be owned by exactly one header that directly points to it.
- Headers can be moved between locations.
- Headers are dropped when their owner is freed or overwritten.
- When dropped, a header's owned heap blocks are also freed.
- A stack variable is freed when it goes out of scope.

<Borrowing rules> (known as, “mutation-xor-sharing” checking)

- Mutable: single borrower has exclusive read/write access.
- Immutable: multiple borrowers and owner can read.

Principles of Lifetime

<Lifetime definition>

- Stack variables have a static lifetime bound by their scope.
- Heap blocks have a dynamic lifetime that ends when their owner is dropped.
- Headers have a dynamic lifetime that ends when their owner is freed or overwritten.

<Lifetime rules for borrowing> (known as, “borrow lifetime checking”)

- A location's lifetime must be longer than any of its borrowers' lifetimes.

An example with ownership and borrowing

An example showing how ownership and borrowing work.

```
// arg owns a string, say ARG
```

```
fn gee(arg: String) -> String {
```

```
    let mut x = String::from("abc"); // x owns the string "abc"
```

```
{
```

```
    let y = x + &arg; // y owns the string "abc"+ARG, x not accessible
```

```
    x = y + &arg; // x owns the string "abc"+ARG+ARG, y not accessible
```

```
} // y is deallocated
```

```
let z = x; // z owns the string "abc"+ARG+ARG, x not accessible
```

```
return z; // the string "abc"+ARG+ARG is returned, z not accessible
```

```
} // arg, x, z are deallocated, the string ARG in arg is deallocated
```

Rust's approach

- References to a location of type T:
 - &T: immutable reference
 - &mut T: mutable reference
- Function signatures specify ownership/lifetime conditions
- Compiler checks ownership/lifetime in Safe Rust at two levels:
 - Stack variables (assuming used functions' signatures)
 - A function's argument and return value (verifying against its signature)
- Heap Memory Management:
 - Experts implement primitive heap types in Unsafe Rust with ownership/lifetime guarantees specified in function signatures
 - Users compose these primitives in Safe Rust to build complex data structures

Datatypes with Ownership & Lifetime

Scalar Types

```
fn main() {  
    let mut x = [1,2,3,4];  
    println!("{}", x[2]);  
    x[2] = 42;  
    println!("{}", x[2]);  
    let y = x;  
    println!("{}", x[0]);  
    println!("{}", y[0]);  
}
```

String and str types

```
fn main() {  
    let mut s : &str = "abc";  
    {  
        let mut st : String = String::from(s);  
        st.push_str("def"); // String::push_str(&mut st, "def");  
        println!("{}", st);  
        let s2 : &str = &st[1..4];  
        // st.push_str("ghi");  
        println!("{}", s2);  
        // s = s2;  
    }  
    // println!("{}", s);  
}
```

Borrow and re-borrow

// We track borrow & re-borrow relations

```
fn main() {  
    let mut st : String = String::from("abc");  
    {  
        let stp1 : &mut String = &mut st; // borrow  
        {  
            let stp2 : &mut String = stp1; // re-borrow  
            // stp1.push_str("def");  
            stp2.push_str("def");  
        }  
        stp1.push_str("ghi");  
    }  
    println!("{}", st);  
}
```


Function Types with lifetime annotation

```
fn foo<'a> (s: &'a mut String) -> &'a str { return &s[0..2]; }  
fn gee(f: for <'a> fn(&'a mut String) -> &'a str) {  
    let s;  
    {  
        let mut a = String::from("test");  
        let s1 = f(&mut a); // &mut a is assumed to be expired within f  
        let s2 = s1; // let s2 = &a  
        println!("{}", s1, s2); // s1, s2 are dropped here since no longer used  
        a.push_str("123"); println!("{}", a);  
        s = f(&mut a); println!("{}", s);  
    } // s is dropped here because a is freed here  
    // println!("{}", s);  
}  
fn main() { gee(foo); }
```

Lifetime Elision Rules

Elision rules are as follows:

- Each elided lifetime in input position becomes a distinct lifetime parameter.
- If there is exactly one input lifetime position (elided or not), that lifetime is assigned to *all* elided output lifetimes.
- If there are multiple input lifetime positions, but one of them is `&self` or `&mut self`, the lifetime of `self` is assigned to *all* elided output lifetimes.

```
fn foo(s: &mut String) -> &str { return &s[0..2]; }
```

```
fn gee(f: fn(&mut String) -> &str) {
```

```
    ...
```

```
}
```

```
fn main() { gee(foo); }
```

General lifetime annotation

```
fn foo1<'a>(s1: &'a mut String, s2: &'a String) -> &'a str {  
    ...  
}  
fn foo2<'a,'b>(s1: &'a String, s2: &'b String) -> &'a str {  
    ...  
}  
fn foo3<'a,'b,'c>(s1: &'a String, s2: &'b String) -> &'c str {  
    ...  
}  
fn foo4<'a,'b,'c, 'd, 'e>(s1: &'a String, s2: &'b String, s3: &'c String)  
    -> (&'d str, &'e str) where 'a: 'd, 'b: 'd, 'b: 'e, 'c: 'e {  
    ...  
}
```

General lifetime annotation

```
fn foo1<'a>(s1: &'a String, s2: &'a String) -> &'a str {  
    return if rand::random() { &s1[0..1] } else { &s2[0..2] };  
}  
  
fn foo2<'a,'b>(s1: &'a String, s2: &'b String) -> &'a str {  
    return &s1[0..1];  
}  
  
fn foo3<'a,'b,'c>(s1: &'a String, s2: &'b String) -> &'c str {  
    return "abc";  
}  
  
fn foo4<'a,'b,'c, 'd, 'e>(s1: &'a String, s2: &'b String, s3: &'c String)  
    -> (&'d str, &'e str) where 'a: 'd, 'b: 'd, 'b: 'e, 'c: 'e {  
    let r1 = if rand::random() { &s1[0..1] } else { &s2[0..1] };  
    let r2 = if rand::random() { &s2[0..1] } else { &s3[0..1] };  
    (r1,r2)  
}
```

Struct

```
struct User {  
    active: bool,  
    name: String,  
}  
  
fn main() {  
    println!("Size: {} bytes", std::mem::size_of::<User>());  
    let mut user_name = String::from("gil");  
    let mut u = User { active: true, name: user_name };  
    u.name.push_str(" hur");  
    println!("{}", {}, u.active, u.name);  
    user_name = u.name;  
    println!("{}", {}, u.active, user_name);  
    // println!("{}", {}, u.active, u.name);  
}
```

Struct with lifetime annotation

// Key Idea: Track the lifetime of each field of a struct separately.

```
struct User<'a,'b> { name: &'a mut String, email: &'b String }

fn main() {
    println!("Size: {} bytes", std::mem::size_of::<User>());
    let mut user_name = String::from("gil");
    let temp : &String;
    {
        let mut user_email = String::from("gil.hur@sf.snu.ac.kr");
        let mut u = User { name: &mut user_name, email: &user_email };
        u.name.push_str(" hur");
        println!("{}", {}, u.name, u.email);
        temp = u.name; // temp = u.email;
        // u.name.push_str("xxx");
    }
    println!("{}", temp);
}
```

Struct with lifetime annotation (Variations)

// Key Idea: Track the lifetime of each field of a struct separately.

// Easy to understand if you inline the definition of a struct.

```
struct User1<'a> {  
    name: &'a String,  
    email: &'a String }
```

```
// User1 : for <'a> fn(&'a String, &'a String) : User1<'a>
```

```
// User1::name : for <'a> fn(User1<'a>) : &'a String
```

```
// User1::email : for <'a> fn(User1<'a>) : &'a String
```

```
struct User2<'a,'b> {  
    user: User1<'a>,  
    addr: &'b String }
```

```
// User2 : for <'a,'b> fn(User1<'a>, &'b String) : User2<'a,'b>
```

```
// User2::user : for <'a,'b> fn(User2<'a,'b>) : User1<'a>
```

```
// User2::addr : for <'a,'b> fn(User2<'a,'b>) : &'b String
```

Struct with lifetime annotation (Variations)

```
struct Foo<'a,'b> {  
    x: &'a i32,  
    y: &'b i32,  
}
```

```
fn main() {  
    let x = 1;  
    let v;  
    {  
        let y = 2;  
        let f = Foo { x: &x, y: &y };  
        v = f.x;  
    }  
    println!("{}", *v);  
}
```


Enum

```
enum IpAddr {  
    V4(u8, u8, u8, u8),  
    V6(String),  
}  
  
use IpAddr::*;  
  
fn main() {  
    let ip1 : IpAddr = V4(147,36,52,255);  
    match ip1 { // match &ip1 // match &mut ip1  
        V4(a1,a2,a3,a4) => println!("{a1:?} {a2:?} {a3:?} {a4:?}"),  
        V6(addr) => println!("{addr:?}")    }  
  
    let ip2 = V6(String::from("123.123.123.123.123.123"));  
    match &ip2 {  
        V4(a1,a2,a3,a4) => println!("{a1:?} {a2:?} {a3:?} {a4:?}"),  
        V6(addr) => println!("{addr:?}")    }  
}
```

Recursive Enum with Box

```
enum List<T> { Nil, Cons(T, Box<List<T>>) }  
impl <T> List<T> {  
    fn new() -> List<T> { List::Nil }  
    fn cons(hd: T, tl: List<T>) -> List<T> { List::Cons(hd, Box::new(tl)) }  
}  
fn mylen<T>(l: &List<T>) -> u64 {  
    match l {  
        List::Nil => 0,  
        List::Cons(_, tl) => 1+mylen(tl) // 1+mylen(tl.as_ref())  
    }  
}  
fn main() {  
    let l = List::cons(0, List::cons(1, List::cons(2, List::new())));  
    println!("{}", mylen(&l));  
}
```

Update via a mutable reference

```
use std::mem;

fn test(x: &mut List<i32>) {
    // *x = List::cons(0, *x)
    let old = mem::replace(x, List::new());
    *x = List::cons(0, old)
}

fn main() {
    let mut l = List::cons(0, List::cons(1, List::cons(2, List::new())));
    println!("{}", mylen(&l));
    test(&mut l);
    println!("{}", mylen(&l));
}
```

Caveat: subtle immutability in let (seems a design bug)

- “let x : T” only guarantees its owned parts are unchanged
- “&x” guarantees that all data reachable from x are unchanged

```
struct Foo<'a> {  
    x: i64,  
    s: &'a mut String  
}  
  
fn main() {  
    let mut str = "abc".to_string();  
    let foo = Foo {x: 42, s: &mut str};  
    println!("{}", foo.x, foo.s);  
    foo.s.push_str("def");  
    // (&foo).s.push_str("def");  
    // foo.x = 37;  
    println!("{}", foo.x, foo.s);  
}
```

Hack: achieving proper immutability

```
struct Foo<'a> {  
    x: i64,  
    s: &'a mut String  
}  
  
fn main() {  
    let mut str = "abc".to_string();  
    let foo = Foo {x: 42, s: &mut str};  
    let _foo_immutable = &foo;  
    println!("{}", foo.x, foo.s);  
    // foo.s.push_str("def");  
    println!("{}", foo.x, foo.s);  
    let _foo_end = _foo_immutable;  
}
```

Smart Pointers

Box type

```
enum List<T> {  
  Nil,  
  Cons(T, Box<List<T>>)  
}
```

Rc type

Use Rc type to replace static lifetime checking with dynamic checking.

<Problem>

```
enum List<T> {  
    Cons(T, Rc<List<T>>),  
    Nil,  
}  
  
use crate::List::*;  
  
fn test() -> (List<i64>,List<i64>) {  
    let a = Cons(5, Box::new(Nil));  
    let b = Cons(3, Box::new(a));  
    let c = Cons(4, Box::new(a));  
    (b,c)  
}
```


Rc type

```
use std::rc::Rc;

enum List<T> { Cons(T, Rc<List<T>>), Nil, }

use crate::List::*;

impl <T> List<T> {
    fn unfold(&self) -> Option<(&T,&List<T>)> {
        if let Cons(hd,tl) = self { Some((hd, tl.as_ref())) } else { None }
    }
}

fn test() -> (List<i64>,List<i64>) {
    let a = Rc::new(Cons(5, Rc::new(Nil)));
    let b = Cons(3, Rc::clone(&a));
    let c = Cons(4, Rc::clone(&a));
    (b,c) }

fn main() {
    let (l1, l2) = test(); { let _t = l1; }
    println!("{}", l2.unfold().unwrap().1.unfold().unwrap().0) }
```

RefCell type

Use RefCell type to replace static mutation-xor-sharing checking with dynamic checking.

<Problem>

Want to mutate the shared list in the previous example.

RefCell type

```
use std::rc::Rc; use std::cell::{RefCell, Ref, RefMut};

enum List<T> {
    Cons(T, Rc<RefCell<List<T>>>),
    Nil,
}

use crate::List::{Cons, Nil};

impl <T> List<T> {
    fn unfold<'a>(&'a self) -> Option<(&'a T, Ref<'a, List<T>>)> {
        if let Cons(hd,tl) = self { Some((hd, tl.borrow())) } else { None }
        // tl.as_ref().borrow()
    }

    fn unfold_mut<'a>(&'a mut self) -> Option<(&'a mut T, RefMut<'a,
List<T>>)> {
        if let Cons(hd,tl) = self { Some((hd,tl.borrow_mut())) } else { None }
    }
}
```

RefCell type

```
fn test() -> (List<i64>,List<i64>) {
    let a = Rc::new(RefCell::new(Cons(5, Rc::new(RefCell::new(Nil))));
    let b = Cons(3, Rc::clone(&a));
    let c = Cons(4, Rc::clone(&a));
    (b,c)
}

fn main() {
    let (mut l1, l2) = test();
    {
        let mut r = l1.unfold_mut().unwrap().1;
        *r = Cons(42, Rc::new(RefCell::new(Nil)));
    }
    println!("{}", l2.unfold().unwrap().1.unfold().unwrap().0)
}
```

Type class and Closure

Type class

```
trait Describable {  
    fn describe(&self) -> String;  
}  
  
struct Book { title: String, author: String, }  
  
impl Describable for Book {  
    fn describe(&self) -> String {  
        format!("{}", self.title, self.author)  
    }  
}  
  
//fn print_desc(a: &impl Describable) -> ()  
fn print_desc<T>(a: &T) -> () where T : Describable  
{ println!("{}", a.describe()) }  
  
fn main() {  
    print_desc(&Book{title: "PP".to_string(), author: "Gil Hur".to_string()})  
}
```

Dyn: dynamic dispatch

```
impl Describable for i64 {  
    fn describe(&self) -> String {  
        format!("64-bit-signed-integer: {}", *self)  
    }  
}  
  
fn main() {  
    let book = Book{title: "PP".to_string(), author: "Gil Hur".to_string()};  
    let ds : Vec<Box<dyn Describable>> =  
        vec![Box::new(42), Box::new(book)];  
    for d in &ds {  
        println!("{}", d.describe())  
        // cannot use print_desc : why?  
    }  
}
```

Closure type: Fn

```
fn test1<F>(f: F) -> usize
where F: Fn(usize)->usize {
    f(10) + f(20)
}

fn main() {
    let mut s = "abc".to_string();
    let f1 = |n| n+s.len();
    println!("{}", test1(f1));
}
```


Closure type: FnMut

```
fn test2<F>(mut f: F)
where F: FnMut(&str)->() {
    f("gil");
    f("hur");
}

fn main() {
    let mut s = "abc".to_string();
    let f1 = |n| n+s.len();
    println!("{}", test1(f1));
    let f2 = |x:&str| s.push_str(x);
    test2(f2);
    println!("{}", s);
}
```

Closure type: FnOnce

```
fn test3<F>(f: F)->String
where F: FnOnce(&str)->String {
    f("test")
}

fn main() {
    let mut s = "abc".to_string();
    let f1 = |n| n+s.len();
    println!("{}", test1(f1));
    let f2 = |x:&str| s.push_str(x);
    test2(f2);
    println!("{}", s);
    let f3 = move |x:&str| {s.push_str(x); s}; // can omit "move"
    let s2 = test3(f3);
    println!("{}", s2);
}
```

Closure construction: move

```
fn test1<F>(f: F) -> usize
where F: Fn(usize)->usize {
    f(10) + f(20)
}

fn main() {
    let clo;
    {
        let s = "abc".to_string();
        clo = move |n| n+s.len();
    }
    println!("{}", test1(clo));
}
```

Thank you for your hard work!