

**SPCS CRYPTOGRAPHY**  
**HOMEWORK 9 (OPTIONAL)**

*Spend time on your project first - just do the ones in this p-set that interest you. You are strongly encouraged to work in groups, but you have to write up the solution on your own.*

**Reference for today's lecture:** Chapter 12.2-12.4; For  $p-1$ -method, see [https://math.berkeley.edu/~sagrawal/su14\\_math55/notes\\_pollard.pdf](https://math.berkeley.edu/~sagrawal/su14_math55/notes_pollard.pdf)

1. Compute the number of  $B$ -smooth number from 2 to  $X$  (inclusive):
  - (a)  $X = 25, B = 3$ .
  - (b)  $X = 35, B = 5$ .
2. If  $n$  is  $B$ -power smooth, prove that it is also  $B$ -smooth. Is the converse true?
3. Use  $p-1$  method to factorize
  - (a) 1739.
  - (b) 220459.

# 4. For each of the following numbers  $n$ , compute the values of

$$n + 1^2, n + 2^2, n + 3^2, \dots,$$

as we did in class, until you find a value  $n + b^2$  that is a perfect square  $a^2$ . Then use the values of  $a$  and  $b$  to factor  $n$ .

- (a)  $n = 53357$ .
- (b)  $n = 34571$ .

5. Use the quadratic sieve/Dixon's method to factor the following numbers.
  - (a)  $n = 61063$ . Hint:

|                                    |     |                                 |
|------------------------------------|-----|---------------------------------|
| $1882^2 \equiv 270 \pmod{61063}$   | and | $270 = 2 \cdot 3^3 \cdot 5$     |
| $1898^2 \equiv 60750 \pmod{61063}$ | and | $60750 = 2 \cdot 3^5 \cdot 5^3$ |

- (b)  $n = 52907$ . Hint:

|                                   |     |                             |
|-----------------------------------|-----|-----------------------------|
| $399^2 \equiv 480 \pmod{52907}$   | and | $480 = 2^5 \cdot 3 \cdot 5$ |
| $763^2 \equiv 192 \pmod{52907}$   | and | $192 = 2^6 \cdot 3$         |
| $773^2 \equiv 15552 \pmod{52907}$ | and | $15552 = 2^6 \cdot 3^5$     |
| $976^2 \equiv 250 \pmod{52907}$   | and | $250 = 2 \cdot 5^3$         |

# 6. Prove that  $n$  is  $B$ -power smooth if and only if  $n$  divides the least common multiple of  $1, 2, \dots, B$ .

7. Suppose Bob leaks his private key  $d$  to Eve. Instead of choosing a new  $n$ , he decides to choose a new public key  $e$  thus a new private key  $d$ , but with the same value of  $n$ . Is this safe? (That is, can Eve decrypt messages now?)

8. Alice and Bob are such good friends that they choose to use RSA with the same  $n$ , but their public keys  $e$  and  $f$  are different, and indeed, they are relatively prime. Charles wants to send the same message  $m$  to Alice and Bob. If Eve intercepts both of his messages, how can she recover the plaintext message  $m$ ?

- ★ 9. A *multiplicative function*  $f : \mathbb{N} \rightarrow \mathbb{R}$  is a function such that whenever  $\gcd(a, b) = 1$ , we have  $f(a)f(b) = f(ab)$ . We said in class that Euler's  $\phi$  function is one such example.
- (a) Prove that Euler's  $\phi$  function is multiplicative.
  - (b) The *divisor function*,  $d(n)$ , counts the number of divisors of  $n$ . For example, 3 has two divisors 1 and 3, so  $d(3) = 2$ . 12 has 6 divisors: 1, 2, 3, 4, 6, 12, so  $d(12) = 6$ . Prove that  $d(n)$  is multiplicative.
  - (c) Find a formula for  $d(p^e)$  for a prime  $p$  and positive integer  $e$ . Deduce a formula for  $d(n)$ , in terms of the prime factorization of  $n = p_1^{e_1} \cdots p_k^{e_k}$ .
  - (d) The *Mobius function*,  $\mu(n)$  is defined as follows.

$$\mu(n) = \begin{cases} 0 & \text{if } n \text{ is divisible by square of a prime} \\ (-1)^r & \text{if } n \text{ is the product of } r \text{ distinct primes} \end{cases}$$

Prove that  $\mu(n)$  is multiplicative.

- ★ 10. Here is another example of a public key cryptosystem. Bob chooses two large primes  $p$  and  $q$  and he publishes  $n = pq$ . It is assumed that  $n$  is hard to factor. Bob also chooses three random numbers  $g$ ,  $r_1$ , and  $r_2$  modulo  $n$  and computes

$$g_1 \equiv g^{r_1(p-1)} \pmod{n} \text{ and } g_2 \equiv g^{r_2(q-1)} \pmod{n}.$$

His public key is the triple  $(n, g_1, g_2)$  and his private key is the pair of primes  $(p, q)$ .

Alice wants to send the message  $m$  to Bob, where  $m$  is a number modulo  $n$ . She chooses two random integers  $s_1$  and  $s_2$  modulo  $n$  and computes

$$c_1 \equiv mg_1^{s_1} \pmod{n} \text{ and } c_2 \equiv mg_2^{s_2} \pmod{n}.$$

Alice then sends the ciphertext  $(c_1, c_2)$  to Bob.

Decryption is extremely fast and easy. Bob uses the Chinese remainder theorem to solve the pair of congruences

$$\begin{cases} x \equiv c_1 \pmod{p} \\ x \equiv c_2 \pmod{q} \end{cases}$$

- (a) Prove that Bob's solution  $x$  equals Alice's plaintext  $m$ .
- (b) Explain why this cryptosystem is not secure.