

Information-Theoretic Security

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What is information?

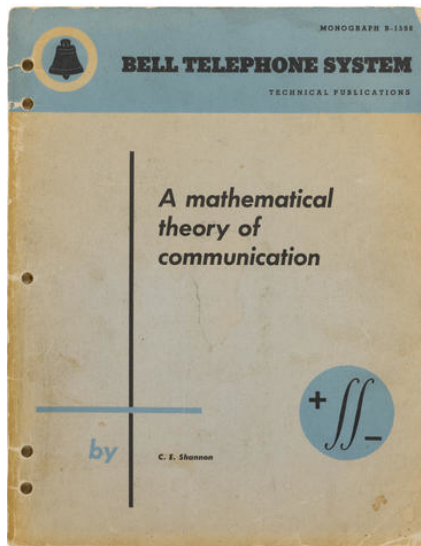
Claude E. Shannon

Henri Cartier-Bresson / Magnum Photos, 1962



A Mathematical Theory of Communication

C. E. Shannon, "A mathematical theory of communication," *Bell System Technical Journal*, vol. 27, pp. 379-423, Jul. 1948.



Information Source

- ▶ Alphabet: \mathcal{X}
- ▶ Probability distribution: P_X

Entropy

Rate at which information is produced

$$H(X) = - \sum_{x \in \mathcal{X}} P_X(x) \log_2 P_X(x)$$

Mutual Information

Information Shared by Two Random Variables

$$I(X; Y) = \sum_{x,y \in \mathcal{X} \times \mathcal{Y}} P_{XY}(x, y) \log_2 \frac{P_{XY}(x, y)}{P_X(x)P_Y(y)}$$

Motivation for Information Theoretic Security

- ▶ Cryptographic approach
 - ▶ Complexity: Limited computational power
 - ▶ Inversion of mathematical functions: Mathematical progress

Motivation for Information Theoretic Security

I am a new slide, really!

- ▶ Cryptographic approach
 - ▶ Complexity: Limited computational power
 - ▶ Inversion of mathematical functions: Mathematical progress
- ▶ Alternative approach: **Information-theoretic secrecy**
 - ▶ Security guarantees independent of adversary's computing power
 - ▶ Exploit intrinsic uncertainty in channel/source

Information Leakage in Written English

Information Leakage in Written English

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Information Leakage in Written English

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Shannon's Perfect Secrecy

One Time Pad:

- ▶ Key \mathbf{k} is shared by transmitter and legitimate receiver
- ▶ Encoding: $\mathbf{y} = \mathbf{x} \oplus \mathbf{k}$

Perfect Secrecy?

Best Strategy for Eavesdropper

Guessing under complete uncertainty

$$H(\mathbf{k}) \geq H(\mathbf{x}) \iff I(\mathbf{x}; \mathbf{y}) = 0$$

Shortcoming of One Time Pad

- ▶ Key must be exchanged prior to communication
- ▶ Key can only be used once

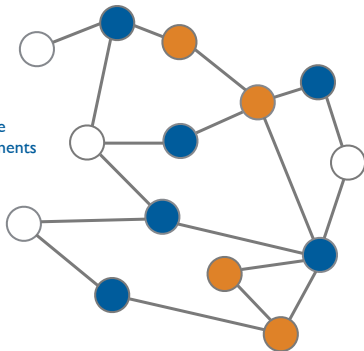
Secrecy in a Network

Operator

Eavesdropper

Legitimate
Measurements

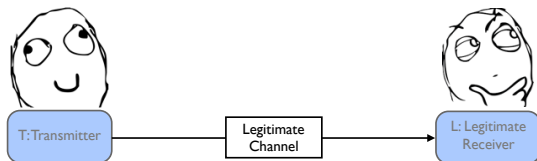
Compromised
Measurements



What are the conditions for secrecy?

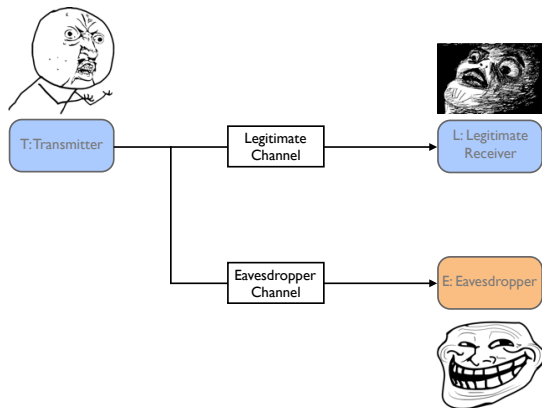
Wiretap channel

A. D. Wyner, "The wire-tap channel," *Bell System Technical Journal*, vol. 54, pp. 1355-1387, Oct. 1975



Wiretap channel

A. D. Wyner, "The wire-tap channel," *Bell System Technical Journal*, vol. 54, pp. 1355-1387, Oct. 1975



Secrecy types

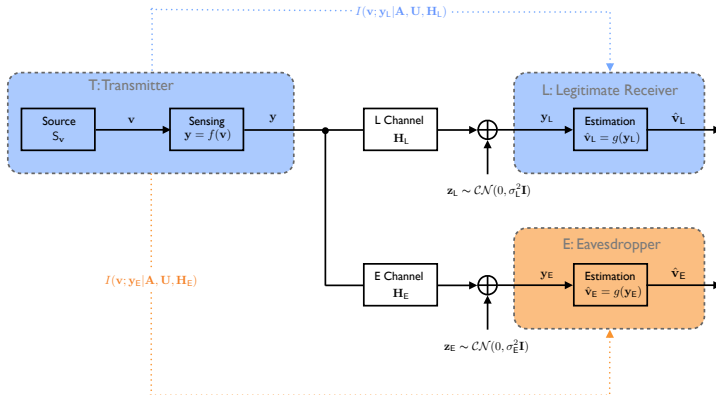
- ▶ Weak secrecy

$$\lim_{n \rightarrow \infty} \frac{1}{n} I(\mathbf{x}; \mathbf{y}_E) = 0$$

- ▶ Strong secrecy

$$\lim_{n \rightarrow \infty} I(\mathbf{x}; \mathbf{y}_E) = 0$$

MIMO Wiretap channel model



Let \mathbf{V} and \mathbf{Z} be independent random variables with $\mathbf{V} \sim \mathbf{F}_V$ and $\mathbf{Z} \sim \mathcal{CN}(0, 1)$ and $\mathbf{R}_L = \mathbf{A}^\dagger \mathbf{U}^\dagger \mathbf{H}_L^\dagger \mathbf{H}_L \mathbf{A} \mathbf{U}$ and $\mathbf{R}_E = \mathbf{A}^\dagger \mathbf{U}^\dagger \mathbf{H}_E^\dagger \mathbf{H}_E \mathbf{A} \mathbf{U}$, respectively. Then,

$$\begin{aligned} \mathcal{I}_S = & \left(I_s(\mathbf{V}, \eta_L, \eta_E) - \log e(\eta_L \chi_L - \eta_E \chi_E) \right. \\ & \left. + \log e \left(\int_0^{\chi_L} \mathcal{R}_{\mathbf{R}_L}(-u) du - \int_0^{\chi_E} \mathcal{R}_{\mathbf{R}_E}(-u) du \right) \right)^+ \end{aligned}$$

where

$$I_s(\mathbf{V}, \eta_L, \eta_E) \triangleq I \left(\mathbf{V}; \mathbf{V} + \frac{1}{\sqrt{\eta_L}} \mathbf{Z} \right) - I \left(\mathbf{V}; \mathbf{V} + \frac{\sigma_E}{\sqrt{\eta_E}} \mathbf{Z} \right)$$

and $\eta_L, \eta_E, \chi_L, \chi_E$ are the non-negative solutions to the system of fixed point equations given by

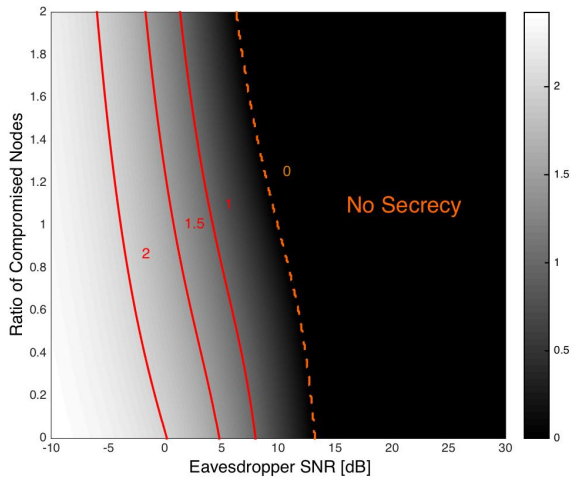
$$\begin{aligned} \eta_L &= \mathcal{R}_{\mathbf{R}_L}(-\chi_L) \\ \chi_L &= \text{mmse}(\eta_L) \\ \eta_E &= \mathcal{R}_{\mathbf{R}_E}(-\chi_E) \\ \chi_E &= \text{mmse}(\eta_E) \end{aligned}$$

where we define

$$\text{mmse}(\eta) \triangleq \mathbb{E} \left[|\mathbf{V} - \mathbb{E}(\mathbf{V} | \mathbf{V} + \eta^{-\frac{1}{2}} \mathbf{Z})|^2 \right]$$

Secrecy region

$\text{SNR}_L = 10 \text{ dB}$



Thanks!