

TDT4136 - INTRODUCTION TO ARTIFICIAL INTELLIGENCE  
ASSIGNMENT 5

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# Propositional and First Order Logic

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## 1. Models and entailment in propositional logic

### 1.1. Validity and Soundness

- (a) Generate the vocabulary of the following argument.
- (b) Translate the argument into propositional logic statements.
- (c) Add a premise (P4) to make the conclusion of the argument valid.

P1 to P3 are the premises, C is the conclusion:

- (P1) If Peter's argument is valid and all the premises of Peter's argument are true, then Peter's argument is sound.
- (P2) If the premises of Peter's argument entail the conclusion of Peter's argument, then Peter's argument is valid.
- (P3) The premises of Peter's argument entail the conclusion of Peter's argument.
- (C) Peter's argument is sound.

### 1.2. Modelling

A sentence is satisfiable if it is true in or satisfied by some model.

- (a)  $(p \Rightarrow q) \Rightarrow (p \Rightarrow r) \Rightarrow (q \Rightarrow r)$

p	q	r	$p \Rightarrow q$	$p \Rightarrow r$	$q \Rightarrow r$	$(p \Rightarrow r) \Rightarrow (q \Rightarrow r)$	Total
0	0	0	1	1	1	1	1
0	0	1	1	1	1	1	1
0	1	0	1	1	0	0	0
0	1	1	1	1	1	1	1
1	0	0	0	0	1	1	1
1	0	1	0	1	1	1	1
1	1	0	1	0	0	1	0
1	1	1	1	1	1	1	1

This complete model is not a tautology, as not all instances of the model is true.

(b)  $(p \vee (\neg q \Rightarrow r)) \Rightarrow (q \vee (\neg p \Rightarrow r))$

p	q	r	$\neg q \Rightarrow r$	$p \vee (\neg q \Rightarrow r)$	$\neg p \Rightarrow r$	$q \vee (\neg p \Rightarrow r)$	Total
0	0	0	0	0	0	0	1
0	0	1	1	1	1	1	1
0	1	0	1	1	0	1	1
0	1	1	1	1	1	1	1
1	0	0	0	1	1	1	1
1	0	1	1	1	1	1	1
1	1	0	1	1	1	1	1
1	1	1	1	1	1	1	1

This complete model is a tautology, as all instances of the model is true.

(c)  $\neg(p \vee (q \Rightarrow \neg r)) \Rightarrow ((p \Rightarrow q) \wedge (p \Rightarrow r))$

p	q	r	$q \Rightarrow \neg r$	$\neg(p \vee (q \Rightarrow \neg r))$	$p \Rightarrow q$	$p \Rightarrow r$	Total
0	0	0	1	0	1	1	1
0	0	1	0	1	1	1	1
0	1	0	1	0	1	1	1
0	1	1	1	0	1	1	1
1	0	0	1	0	0	0	1
1	0	1	1	0	0	1	1
1	1	0	0	0	1	0	1
1	1	1	1	0	1	1	1

This complete model is a tautology.

(d)  $\neg(\neg p \vee (q \wedge r)) \Rightarrow (\neg(p \vee q) \wedge r)$

p	q	r	$q \wedge r$	$\neg(\neg p \vee (q \wedge r))$	$\neg(p \vee q)$	$\neg(p \vee q) \wedge r$	Total
0	0	0	0	0	1	0	1
0	0	1	0	0	1	1	1
0	1	0	0	0	0	0	1
0	1	1	1	0	0	0	1
1	0	0	0	1	0	0	0
1	0	1	0	1	0	0	0
1	1	0	0	1	0	0	0
1	1	1	1	0	0	0	1

This complete model is not a tautology.

### 1.3. Modelling 2

p	q	r	$\phi$	$(r \Rightarrow q)$
0	0	0	1	1
0	0	1	0	0
0	1	0	1	1
0	1	1	1	1
1	0	0	1	1
1	0	1	0	0
1	1	0	1	1
1	1	1	1	1

The expression is independent of  $p$ , meaning it doesn't affect the value of  $\phi$ .  $\phi$  is equivalent to  $(r \Rightarrow q)$ . If we want to include  $p$ , we can write  $(r \Rightarrow q) \wedge (p \vee \neg p)$ .

## **2. Resolution in propositional logic**

### **2.1. Conjunctive Normal Form**

### **2.2. Inference in propositional logic**

### **3. Representation in First-Order Logic (FOL)**

#### **3.1. a) Argument A**

#### **3.2. b) Argument B**



## 4. Resolution in FOL