



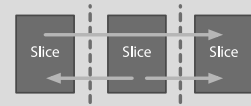
Software Engineering in der industriellen Praxis (SEIP)

Dr. Ralf S. Engelschall

Slicing Principle

Vertically split code or data into two or more logically, optionally also spatially, clearly distinct, named, and unranked slices.

The particular slicing should minimize the total amount of individual relationships between the resulting slices. Per type of relationship, there should be no cycle in the transitive relationships.



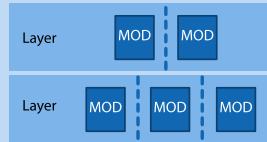
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MOD Concerned Module

Split related code or data (usually across a single Layer) into two or more logically distinct domain- or technology-induced Modules.

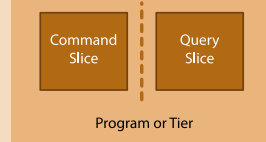
Rationale: Separation of Concern, Single Responsibility Principle, Mastering Complexity.



CQRS Command-Query Responsibility Segregation

Split code and data of a Program (across all Layers) or a Tier into exactly two slices to segregate operations that read data (queries) from the operations that update data (commands).

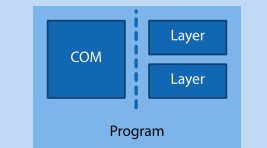
Rationale: Separated Scalability, Separated Data Access Patterns, Event Sourcing Approach.



COM Common Slice

Factor out common or cross concern code or data of a Program (across all Layers) into a single spatially distinct, separate slice.

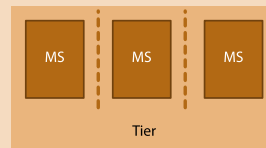
Rationale: Lack of Redundancy, Single Point of Truth, Reusability.



MS Microservice

Split code and data of a Tier (across all Layers) into two or more distinct, loosely-coupled, domain-enclosed, functional services, each forming a stand-alone Program.

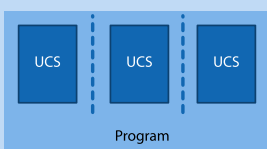
Rationale: Heterogeneity, Resilience, (Scalability), (Easy Deployment), (Organizational Alignment), (Composability), (Reusability), Replaceability.



UCS Use-Case Slice

Split the code and data of a Program (across all Layers) into two or more purely logical slices, one for each distinct, domain-specific Use-Case.

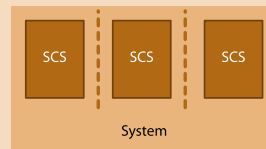
Rationale: Comprehensibility, Domain Alignment, Mastering Complexity.



SCS Self-Contained System

Split code and data of a System (across all Layers and Tiers) into two or more distinct, loosely-coupled, domain-enclosed, functional systems, each forming a stand-alone sub-System.

Rationale: Mastering Complexity, Heterogeneity, Resilience, Scalability, Easy Deployment, Organizational Alignment, Reusability, Replaceability.



When **Slicing**, code or data are split into two or more **logically** — if necessary also “physically” (**spatially**) — **Slices**. These slices are **clearly distinct**, **isolated** from each other, and **named**. Slices are always drawn **vertically**.

Slices in the same Layer should be as independent of each other as possible. In the case of relationships, at least no cycle should exist. There are different special variants of slices, each of which has its own name.

Concerned Modules are Slices of a Layer that realize a specific domain-specific or technical concern. **Common Package** is a Slice of a Tier, where commonalities of other Layers were moved to. **Use-Case Slices** are Slices of a Tier that are dedicated to certain domain-specific use cases.

With the **Command-Query Responsibility Segregation** architecture, a Tier is split into two Slices for Commands/Writes and Queries/Reads. A **Microservice** is a Slice of a Tier, which is executed as a separate Program and which is concerned with a closed domain-specific functionality. A **Self-Contained System** is a Slice of a whole System that is executed as a separate Sub-System.

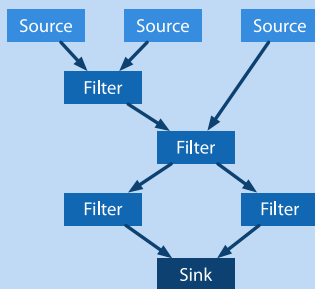
Questions

- ❓ What do one call the resulting units when code or data is split **vertically**?
- ❓ What do one call the Slices of a Tier, which are executed as separate Programs and which are concerned with closed domain-specific functionalities?

Pipes & Filters

Pass data through a directed graph of **Components** and connecting **Pipes**. The components can be **Sources**, where data is produced, **Filters**, where data is processed, or **Sinks**, where data is captured. Source and Filter components can have one or more output Pipes. Filter and Sink components can have one or more input Pipes. Components are independent processing units and operate fully asynchronously.

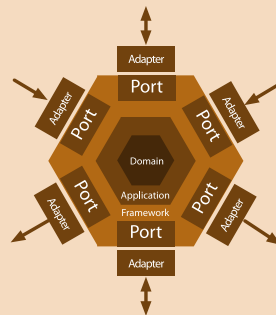
Examples: Unix commands with stdin/stdout/stderr and the Unix shell connecting them with pipes; Apache Spark or Apache Camel data stream processing pipelines.



Ports & Adapters (Hexagonal)

Perform communication in a Hub & Spoke fashion by structuring a solution into the three "Layers" **Domain**, **Application** and **Framework** and use the Framework layer to connect with the outside through **Ports** (general Interfaces) and **Adapters** (particular Implementations). Often some Ports & Adapters are user-facing sources and some are data-facing sinks, although the motivation for the Ports & Adapters architecture is to remove this distinction between user and data sides of a solution.

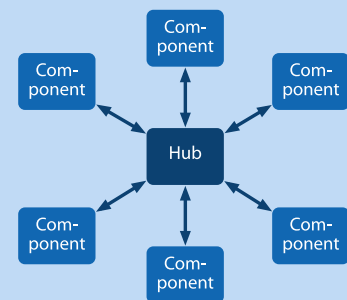
Examples: Message Queue, Enterprise Service Bus or Media Streaming Service internal realization.



Hub & Spoke

Perform communication (the **Spoke**) between multiple Components through a central **Hub** Component. Instead of having to communicate with $N \times (N-1) / 2$ bi-directional interconnects between N Components, use the intermediate Hub to communicate with just N interconnects only. Sometimes one distinguishes between K ($0 < K < N$) source and $N - K$ target Components and then $K \times (N - K)$ uni-directional interconnects are reduced to just N interconnects, too.

Examples: Message Queue, Enterprise Service Bus, Module Group Facade, GNU Compiler Collection, ImageMagick, etc.



The **Flow Architectures** are concerned with the primary data flow or the primary communication of an application. Here the following three classical architectural approaches exist.

With **Pipes & Filters** a directed **Graph** is built. The nodes of the graph are the **Components**, which are either of type **Source**, **Filter** or **Sink**. The edges of the graph are the **Pipes**: the data transmission links between the **Components**.

With the special **Ports & Adapters** (aka **Hexagonal Architecture**) a "Hub & Spoke" structure is set up. The "Hub" are the Components of the application core. The "Spokes" each consist of a **Component**, which is composed of the **Port** (the interface) and the **Adapter** (the implementation).

With **Hub & Spoke** in general, a central **Hub** Component acts as the communication center between **Spoke** components which are star-shaped around the **Hub**. The crux is that the maximum $N \times N$ communication paths between the Spoke components, thanks to the Hub, can be reduced to $N+M$ communication paths only.

Questions

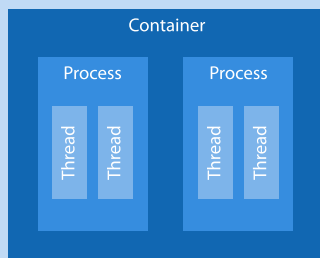
- ? With the help of which **Flow Architecture** can N components be connected with each other in a way that instead of $N \times N$ communication paths only $N+N$ are created?

Container, Process, Thread

The Operating System manages and orchestrates the run-time execution of applications in **Containers**, programs in **Processes** and control flows in **Threads**.

Containers are the ultimate enclosures, separating and controlling both the computing resources processor, memory, storage and network. Processes are the primary enclosures, still separating and controlling at least the computing resources processor and memory. Threads are the light-weight enclosures, just separating and controlling the computing resource processor. Containers can contain one or more Processes, and Processes can contain one or more Threads.

Examples: Docker Container, Unix Processes, POSIX Threads.

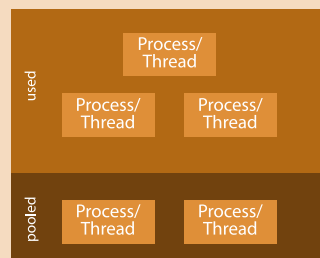


Process/Thread Pool

Instead of creating a Process/Thread for handling each incoming I/O request, pick a pre-created Process/Thread out of a resource **Pool** in order to increase performance and decouple I/O traffic (leading to threads of execution) from the actual computing resource usage and utilization.

The Process/Thread Pool usually has a lower and upper bound of processes/threads. The lower bound keeps the system "hot" between I/O requests. The upper bound limits the computing resource usage and avoids over-utilization.

Examples: Apache HTTP Daemon

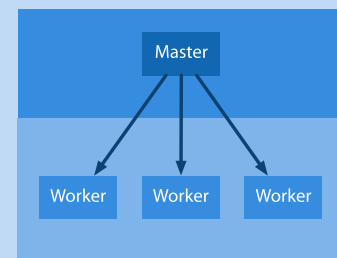


Master-Worker

The system has a single permanent **Master** container/process/thread and a Pool of many ephemeral **Worker** containers/processes/threads. The Master starts, restarts, pauses, resumes and stops the Workers and usually also delegates incoming I/O requests to them. The Workers process the I/O requests and deliver the responses.

Starting the Master usually implicitly starts an initial set of Workers (the initial Pool), stopping the Master implicitly stops all still pending Workers.

Examples: Unix init(8) daemon, Apache HTTP Daemon, SupervisorD, Node.js Cluster module



The **Process Architectures** are all about the interaction between different **Containers**, **Processes** or **Threads**. All three concepts encapsulate code and data.

Containers are the strongest capsule, which encapsulates both CPU, RAM, hard disk, and network (e.g. Docker Container). A **Process** encapsulates CPU and RAM (e.g., Unix process). In the case of a **Thread**, the weakest capsule, only the CPU is encapsulated (e.g., Unix thread).

Containers are the strongest capsule, which encapsulates both CPU, RAM, hard disk, and network (e.g. Docker Container). A **Process** encapsulates CPU and RAM (e.g., Unix process). In the case of a **Thread**, the weakest capsule, only the CPU is encapsulated (e.g., Unix thread).

In order to be able to answer several requests at the same time, server applications use multiple processes/threads per request. Since the constant creation of such processes/threads noteworthy reduces the runtime performance and the hardware load typically should be limited and not linearly be coupled to the incoming requests, a so-called **Pool** of one-time created worker processes/threads is used (e.g., Apache HTTPd or NGINX).

Classically, such a pool is split into a single **Master** Process/Thread and multiple **Worker** Processes/Threads. The permanently running Master generates, controls, and stops the Workers. Usually, the Workers are also permanently existent, but in the event of errors, the Master will actively stop them, or in case of a crash, automatically restart them (e.g., Node.js cluster module).

Questions

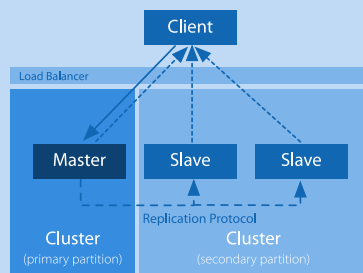
- With which **Process Architecture** is in practice a **Process/Thread Pool** usually managed?

Master-Slave (Static Replication)

Cluster of a single **Master** and multiple **Slave** nodes, where data is continuously copied from the Master to the Slave nodes in order to support high-availability (where a Slave will take over the Master role) in case of a Master outage and increased read performance (where regular read requests are also served by the Slaves).

In this static replication scenario the Master is usually assigned statically and in case of outages has to be reassigned usually semi-manually. Especially, the full reestablishment of the original Master assignment after a Master recovery usually is a manual process.

Examples: OpenLDAP Replication, PostgreSQL WAL Replication.

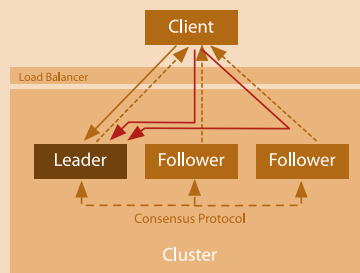


Leader-Follower (Dynamic Replication)

Cluster of a single **Leader** and multiple **Follower** nodes, where data is written on the current Leader node and data read on both the current Leader and all Follower nodes. For writing data to the cluster, the Leader node performs a consensus protocol (e.g. RAFT, Paxos or at least Two-Phase-Commit) with the Followers and this way automatically and consistently replicates the data to the Followers.

In this dynamic replication scenario the Leader is usually automatically assigned by the cluster nodes through an election protocol and in case of outages is automatically re-assigned. There is usually no re-establishment of the original Leader assignment.

Examples: Apache Zookeeper, Consul, EtcD, CockroachDB, InfluxDB.



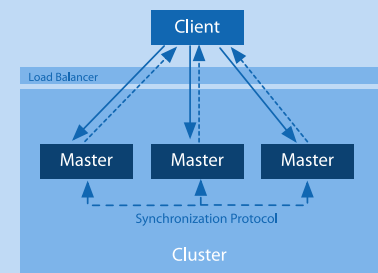
→ Write Operation
--- Read Operation

Master-Master (Synchronization)

Cluster of multiple **Master** nodes, where data is read and written on any Master node concurrently. The Master nodes either use Strict Consistency through writing to a mutual-exclusion-locked shared storage concurrently or use Eventual Consistency in a Shared Nothing storage scenario where they continuously synchronize their local data state to all other nodes with the help of a synchronization protocol.

The synchronization protocol usually is based on either Conflict-Free Replicated Data Types (CRDT) or at least Operational Transformation (OT). In any scenario, data update conflicts are explicitly avoided.

Examples: ORACLE RAC, MySQL/MariaDB Galera Cluster, Riak, Automerge/Hypermerge.



In **Cluster Architectures**, the merger of compute nodes to a cluster is addressed.

The **Master-Slave** architecture is a static replication of data from a Master server to one or more Slave servers. The Clients can send read requests to all Servers, but write requests must be run exclusively via the Master. This is usually used to increase the Read Performance.

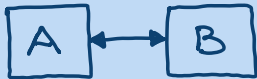
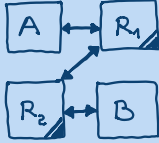
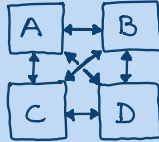
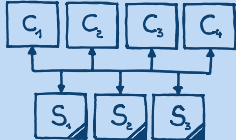
The **Leader-Follower** architecture is a kind of dynamic replication of data from a Leader server to multiple Follower servers. The Clients can send read and write requests to all servers. Since only the Leader server can handle write requests, the Follower servers internally and transparently forward these to the Leader server.

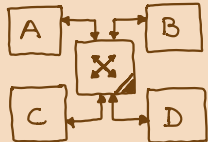
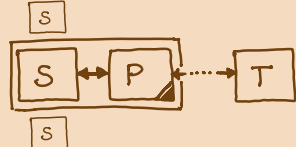
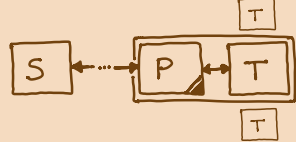
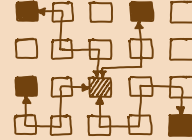
This is also the difference to Master-Slave: the Leader is selected automatically and dynamically between all servers via a Leader Election Protocol (in the event of a failure of the current Leader server). The advantage is that Leader-Follower to Clients feels like Master-Master, but the cluster does not require any complex conflict resolution strategy as is the case with Master-Master.

The **Master-Master** Architecture is a genuine synchronization of data between two or more equal Master servers. The Clients can send both read and write requests to any Master server. However, the Master servers internally must implement an elaborate conflict resolution strategy in order to resolve simultaneous changes to the same data.

Questions

- Which simple **Cluster Architecture** can be used if the read performance of a server application should be increased?

<p>PTP Point-to-Point</p> <p>Communicate between two network nodes in a point-to-point fashion, usually through a direct link.</p> <p>Rationale: simple communication where both nodes know about each other and can directly reach each other.</p> 	
<p>RTG Routing</p> <p>Communicate between two network nodes in a point-to-point fashion, but by routing the network packets over intermediate forwarding nodes (routers).</p> <p>Rationale: simple communication where both nodes know about each other, but cannot directly reach each other.</p> 	
<p>P2P Peer-to-Peer</p> <p>Communicate between multiple network nodes (usually all in the client and server role at the same time) without involving a central hub node (in the role of a server) — except for the initial network entry discovery.</p> <p>Rationale: communication without central control (although a seed peer is required).</p> 	
<p>C/S Client/Server</p> <p>Communicate between multiple nodes in the client role (making requests, and usually with ephemeral addresses) and multiple nodes in the server role (serving responses, and usually with fixed addresses).</p> <p>Rationale: communication with central orchestration, control and data storage.</p> 	

<p>BUS Bus/Broker/Relay</p> <p>Communicate between multiple nodes with the help of a central packet forwarding hub node in a star network topology.</p> <p>Rationale: decouple communication nodes: instead of Point-to-Point (PTP) communications between all nodes, there are just PTP communications with the hub.</p> 	
<p>FPR (Forward) Proxy</p> <p>Communicate between two nodes by using an intermediate forwarding proxy node in front of the source node.</p> <p>Rationale: bridge network topology constraints (segmented networks); caching at source side; auditing of communication.</p> 	
<p>RPR Reverse Proxy</p> <p>Communicate between a source and a target node by using a masquerading proxy node directly in front of the target node.</p> <p>Rationale: load balancing for multiple target nodes; caching at target side; auditing of communication; security shielding of target nodes; protocol conversions.</p> 	
<p>VPN Virtual (Private) Network</p> <p>Communicate between nodes in a logical star network topology on top of an arbitrary physical routed network topology.</p> <p>Rationale: secure private network overlaying an unsecure public network; simplify network topology.</p> 	

In **Networking Architectures**, the network-topological communication between computer nodes is addressed. The simplest way is **Point-to-Point** communication via a direct connection of the nodes.

Usually, however, the communication today goes over a network of nodes, where the individual messages are exchanged with the help of **routing** via intermediate nodes.

If all nodes in both client and server roles communicate directly with each other, it is called a **Peer-to-Peer** architecture. If some nodes are only in the client role and others are only in the server role, it is called a **Client/Server** architecture.

In order to let several nodes communicate with each other, without these having to know and address each other, one usually uses a central **Bus/Broker**. and a star topology.

If between source and target intermediate nodes are active, which act as **Proxy** in the communication and not only forward the network packets like a **Router**, one speaks of either a **(Forward) Proxy** or a **Reverse Proxy** situation. The former, if the proxy acts on the side of the source node, the latter, if the proxy acts as a proxy of the destination node.

In addition, a so-called **Virtual Private Network** can be established, in which a logical secure “overlay network” is placed over a physical network.

Questions

- With which **Network Architecture** can several nodes communicate with each other without these nodes having to know each other exactly?
- What do you call a computer node that acts on behalf of a target node?

UCT Unicast (one-to-one)

Communicate messages from one source to exactly one destination node. The destination node is explicitly and individually addressed.

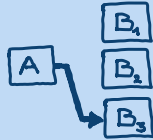
Rationale: private communication between exactly two nodes which both know each other beforehand.



ACT Anycast (one-to-any)

Communicate messages from one source to one of many destination nodes. The picked destination node usually is the network-topology-wise "nearest" or least utilized node in a group of nodes.

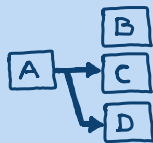
Rationale: Unicast, optimized for network failover scenarios, load balancing and CDNs.



MCT Multicast (one-to-many)

Communicate messages from one source to many destination nodes. The destination nodes usually form a group and are usually not individually addressed.

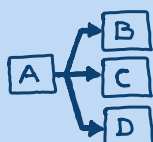
Rationale: node communication where destination nodes dynamically change or where total traffic should be reduced.



BCT Broadcast (one-to-all)

Communicate messages from one source to all available destination nodes. The destination nodes usually are implicitly defined by the extend of the local communication network segment.

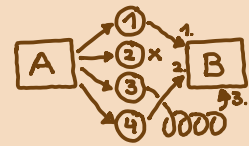
Rationale: spreading out messages to all available nodes for potential responses.



DGR Datagram (Single Packet)

Communicate messages as an unordered set of single packets, usually without any network congestion control, retries or other delivery guarantees.

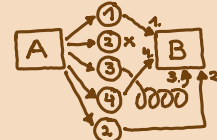
Rationale: simple low-overhead communication without prior communication establishment (handshake).



STR Stream (Sequence of Packets)

Communicate messages as an ordered sequence (stream) of packets, usually with network congestion control, retries and delivery guarantees (at-most-once, exactly-once, at-least-once).

Rationale: reliable communication between nodes.



PLL Pull (Request/Response, RPC)

Communicate by performing a request (from the client node) and pulling a corresponding response (from the server node).

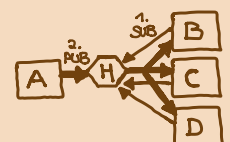
Rationale: Remote Procedure Call (RPC) like Unicast or Anycast communication.



PSH Push (Publish/Subscribe, Events)

Communicate by "subscribing" to "channels" of messages (on one or more receiver nodes or on an intermediate hub) once and then publishing events to those "channels" (on the sender node) multiple times.

Rationale: event-based Multicast or Broadcast communication.



The **Communication Architectures** address the kind of communication between components. One distinguishes primarily four different kinds of message transmission: with **Unicast**, a source node sends to exactly one directly addressed target node. With **Anycast**, a source node sends to a group of potential destination nodes, but the message is delivered to one destination node in the group only.

With **Multicast**, a source node also sends to a group of target nodes, but the message is delivered to all target nodes in the group. With **Broadcast**, a source node sends to all reachable destination nodes without these particular destination nodes being known to the source node.

With the kind of messages, one differentiates two variants: with **Datagram**, each message consists of exactly one network packet, and when sending, no guarantees are given whether and in which order the messages will arrive at the destination node. In contrast, with **Stream**, a message consists of a sequence of network packets and different guarantees are given:

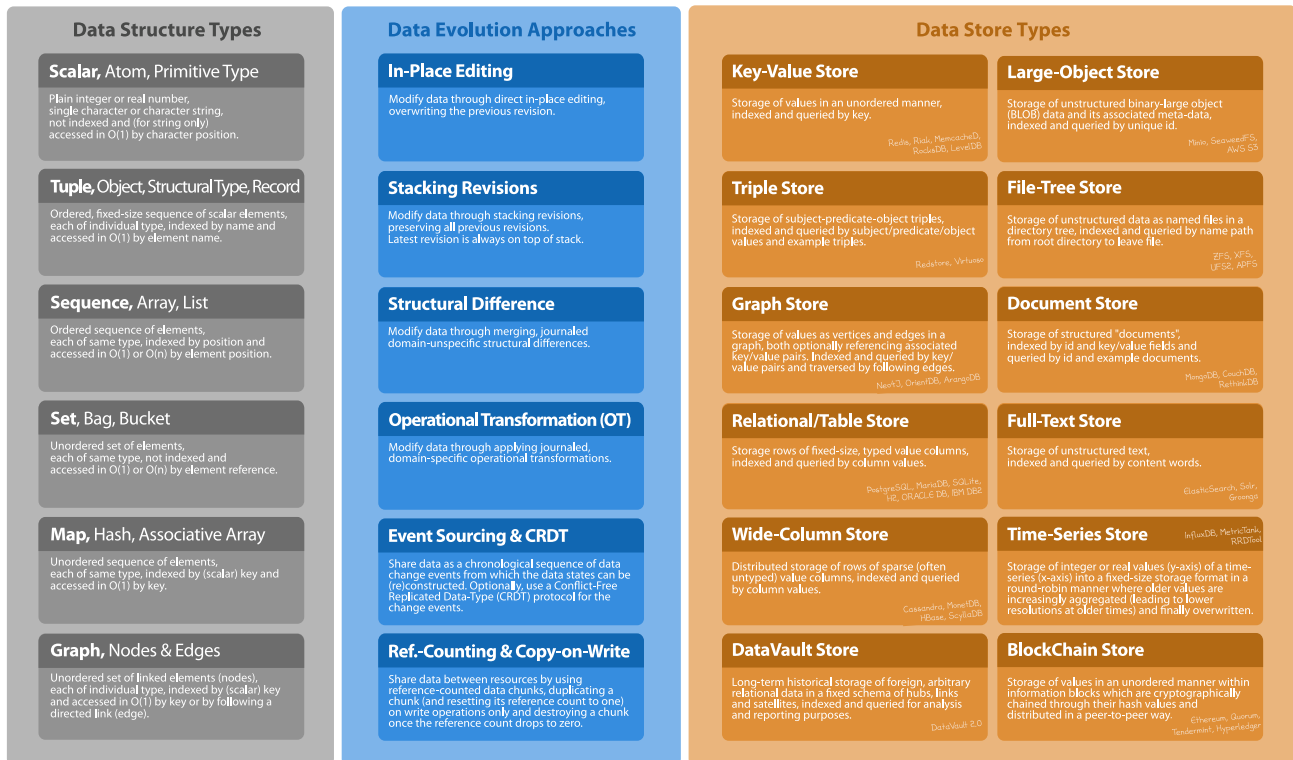
In case of packet congestion on intermediate nodes, the source of the **Stream** may be throttled. In case of packet loss, packets are resent. And one might get control over whether the packet will be delivered at most once, exactly once, or at least once at the destination node.

There are usually two modes of client/server communication: in **Pull** mode, the client sends a request, and the server sends a response. The server cannot proactively (without a prior request) send a message. In **Push** mode, the client sends a message in advance to the server to subscribe to certain types of messages. After that, the server can send a message to all subscribed clients at any time.

Usually, **Pull** is implemented via **Unicast/Anycast** and as a **Stream**, for example, in the HTTP protocol. On the other hand, **Push** is usually implemented via **Multicast/Broadcast** as a **Datagram**, for example, in the DHCP protocol.

Questions

- Which well known Web-protocol uses a communication based on **Unicast**, **Stream** and **Pull**?



The Software Architect distinguishes only 6 **Data Structure Types** for data elements: **Scalar** (e.g. Integer, String, etc), **Tuple** (ordered fixed-size sequence of Scalars), **Sequence** (ordered sequence of elements), **Set** (unordered set of elements), **Map** (unordered set of elements, each indexed by key) and **Graph** (unordered set of elements, each indexed by key or by following a link between elements). All complex specific data structures in practice, for the Software Architect, are only the combination of these 6 types.

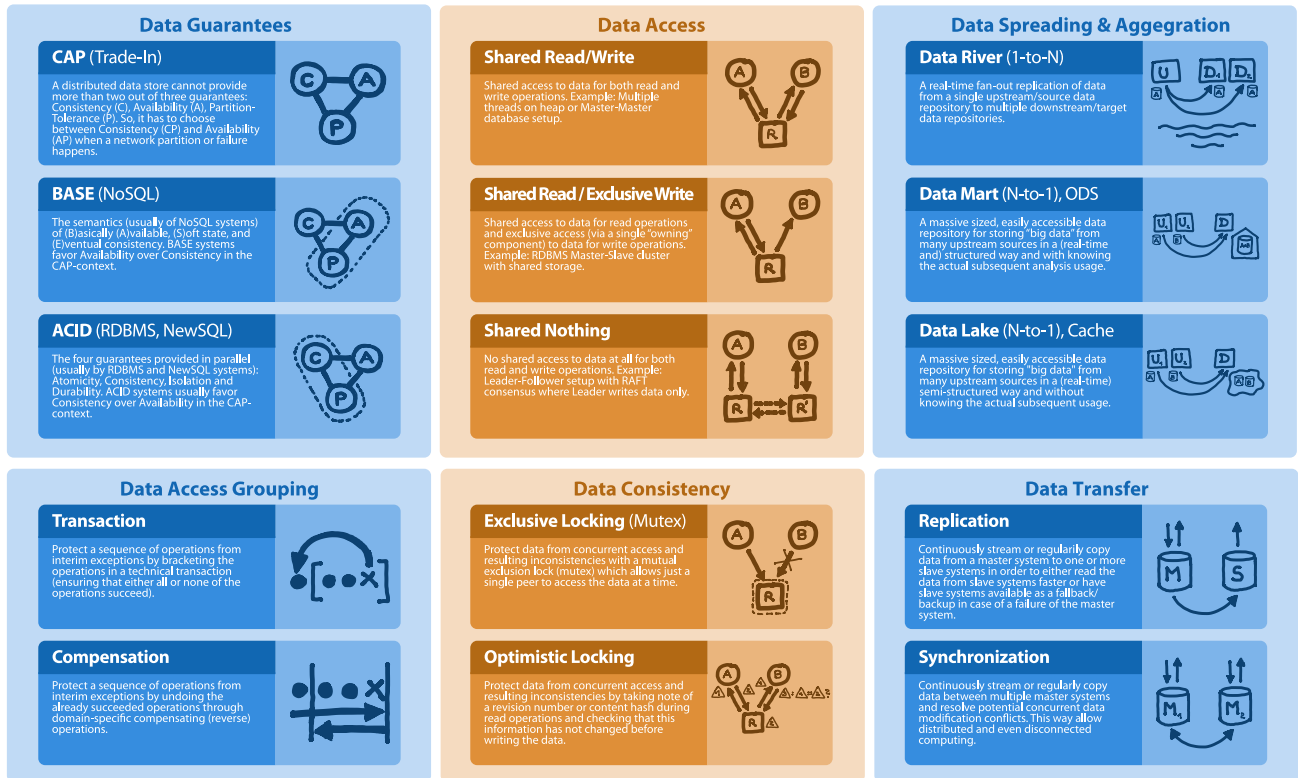
There are numerous **Data Evolution Approaches**, with which data can change over time: in the simplest case, **In-Place Editing**, data is simply changed directly. Access to previous states does not exist. If one wants to access previous states, one can use **Stacking Revisions**, in which the entire data set is copied before each change. So that the entire data record does not have to be copied, **Structural Difference** stores only a technical difference between the old and the new data records. Alternatively, with **Operational Transformation**, the technical change operations can be stored as a journal.

If such a journal is used also to keep replicas of the data sets up-to-date, one refers to it as **Event Sourcing**. If the journal is used as the protocol of so-called **Conflict-Free Replicated Data-Types (CRDT)**, instead of (unidirectional) replication, a (bidirectional) synchronization can be achieved. If several processes/threads logically operate on copies, but physically on the same data sets, **Copy-on-Write** and **Reference Counting** can be used to achieve common access and the life cycle of the data sets can nevertheless be reasonably controlled.

For the storage of data in databases, there are numerous **Data Store Types**. These differ primarily in the type and flexibility of the data structure and the guarantees provided. The most common type is the **Relational/Table Store**. The most elegant type is the **Graph Store**. The most convenient is the **Document Store**.

Questions

- ? Name 3 **Data Evolution Approaches**, each of which allows to access the previous states of the data?



In the area of **Data Guarantees**, there are three main aspects: The **CAP** theorem addresses the so-called "trade-in": In practice, one usually has to choose between Consistency + Partition-Tolerance (CP) or Availability + Partition-Tolerance (AP). Both at the same time is not possible. With **BASE** systems, AP is usually favored. For a traditional RDBMS with **ACID** guarantees, one usually favors CP.

With **Data Access Grouping** one knows about **Transaction** and **Compensation**. The former is a "technical bracket" that allows you to revert to the previous state in case of an error. The latter is a "domain-specific bracket," where so-called compensation operations allow to "cancel" the earlier changes in order to regain a previous consistent state.

With the **Data Access** of two or more processes/threads on the same data one distinguishes between the approaches **Shared Read/Write** (all read and write the same data), **Shared Read/Exclusive Write** (all read and only one writes the same data) and **Shared Nothing** (all read and write to the equal synchronized data).

With the **Data Consistency** one knows about **Exclusive Locking** (per time unit only one writes) and **Optimistic Locking** (all try to write, but recognize and resolve a conflict).

With **Data Spreading & Aggregation** one differentiates three kinds: with the **Data River** the data are replicated from a master system to many slave systems to achieve, among other things, a higher read performance. With the **Data Mart** (structured data) and **Data Lake** (semi-structured data), data is replicated from one master system to many slave systems in order to centrally report or cache the data.

With the **Data Transfer** we finally distinguish between the unidirectional and conflict-free **Replication** and the bidirectional and conflict-rich **Synchronization**.

Questions

- ? What is the name of the approach in which data is replicated from a master system to many slave systems?