

Lab Session 11

Separation Logic and Concurrency

Software Verification

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For verification tasks that involve the concurrency control mechanisms you will have to install the spec `java.util.concurrent.locks.javaspec` (available online in CLIP) for the package `java.util.concurrent.locks` by copying it to the `bin/rt/` folder of your local verifast installation and append the `rt.jarspec` file with the name of the spec file.

1 Exercises

Exercise 1. Consider the `Stack` implementation from Appendix A. Implement a concurrent version of `Stack` using monitors, such that the `pop` method will *wait* for the stack to be non-empty. □

Exercise 2. Consider the `Queue` implementation from Appendix B. Implement a concurrent version of `Queue` using monitors, such that the `take` method will *wait* for the queue to be non-empty. □

Exercise 3. Implement a concurrent version of the `Queue` using monitors, such that the `take` method will *wait* for the queue to be non-empty and the `add` method will *wait* for the queue to be not full. □

Exercise 4. Implement (in Java) and specify using Verifast the class `PosQueue` which represents a queue of positive numbers. The API of the queue is:

```
//Queue based on a circular buffer.
class Queue {

    //creates a new Queue with capacity max.
    Queue(int max);

    //places the int v at the end of this Queue
    void enqueue(int v);

    //retrieves the element at the start of this Queue.
    int dequeue();

    //returns true if this Queue has reached its capacity.
    boolean isFull();

    //returns true if this Queue has no elements.
    boolean isEmpty();
}
```

You will likely want to define a predicate of the form: `predicate QueueInv(Queue q; int n, int m)`, where `n` is the number of elements in the queue and `m` is the queue capacity. Note that your invariant will have to characterize the array using `array_slice` and `array_slice_deep` predicates in terms of the head index, the tail index and the beginning (index 0) and the length of the array. □

A Stack Implementation

```
import java.util.concurrent.locks.*;

/*@ predicate Node(Node n; Node nxt, int v) = n.next |-> nxt &*& n.val |-> v;

public class Node {
    Node next;
    int val;

    public Node()
    //@ requires true;
    //@ ensures Node(this, null, 0);
    {
        next = null;
        val = 0;
    }

    public void setnext(Node n)
    //@ requires Node(this, _, ?vv);
    //@ ensures Node(this, n, vv);
    {
        next = n;
    }

    public void setval(int v)
    //@ requires Node(this, ?nn, _);
    //@ ensures Node(this, nn, v);
    {
        val = v;
    }

    public Node getnext()
    //@ requires Node(this, ?nn, ?vv);
    //@ ensures Node(this, nn, vv) &*& result == nn;
    {
        return next;
    }

    public int getval()
    //@ requires Node(this, ?nn, ?vv);
    //@ ensures Node(this, nn, vv) &*& result == vv;
    {
        return val;
    }
}

/*@
predicate List(Node n; list<int> elems) =
    n == null ? (emp &*& elems == nil) :
```

```

        Node(n,?nn,?v) &*& List(nn,?tail) &*& elems == cons(v,tail);

predicate StackInv(Stack s; list<int> l, int max) =
    s.head |-> ?h &*& List(h, l) &*& s.max |-> max; @*/

public class Stack {
    Node head;
    int max;

    public Stack(int max)
        //@ requires max > 0;
        //@ ensures StackInv(this, nil, max);
    {
        head = null;
        this.max = max;
    }

    public void push(int v)
        //@ requires StackInv(this, ?l, ?m);
        //@ ensures StackInv(this, cons(v, l), m);
    {
        Node n = new Node();
        n.setval(v);
        n.setnext(head);
        head = n;
    }

    public int pop()
        //@ requires StackInv(this, cons(?x, ?vs), ?m);
        //@ ensures StackInv(this, vs, m) &*& result == x;
    {
        //@ open StackInv(this, _, _);
        //@ open List(head, _);
        int v = head.getval();
        head = head.getnext();
        return v;
    }

    public boolean isEmpty()
        //@ requires StackInv(this, ?l, ?m);
        //@ ensures StackInv(this, l, m) &*& result == (l == nil);
    {
        //@ open StackInv(this, _, _);
        //@ open List(head, _);
        return head == null;
    }

    public void clear ()
        //@ requires StackInv(this, ?l, ?m);

```

```
    //@ ensures StackInv(this, nil, m);  
    {  
        head = null;  
    }  
}
```

B Queue Implementation

```
/*@

predicate Cell (Cell c; Cell nxt, int v) =
  c.next |-> nxt && c.content |-> v;

predicate lseg (Cell from, Cell to; list<int> l) =
  from == to ?
  l == nil :
  from != null && from.content |-> ?v &&
  from.next |-> ?next && lseg(next, to, ?nvs) &&
  l == cons (v, nvs);

predicate QueueInv (Queue q; list<int> vs) =
  q.first |-> ?first && q.last |-> ?last && q.length |-> ?ll &&
  last == null ?
  first == null && vs == nil && ll == 0 :
  lseg(last, null, cons(?v, nil)) && lseg(first, last, ?vs1) &&
  vs == append(vs1, cons(v, nil)) && ll == length(vs);

lemma void cell_null_lseg (Cell c)
  requires Cell(c, null, ?v); //??? && c != null
  ensures lseg(c, null, cons(v, nil));
{
  open Cell(c, null, v);
}

lemma void lseg_merge (Cell x, Cell y, Cell w)
  requires lseg(x, y, ?vs1) && lseg(y, w, ?vs2) && lseg(w, null, ?vs3);
  ensures lseg(x, w, append(vs1, vs2)) && lseg(w, null, vs3);
{
  open lseg(w, null, vs3);
  open lseg(x, y, vs1);
  if (x != y) {
    lseg_merge(x.next, y, w);
  }
  else { }
}

lemma_auto void length_one (list<int> vs, int x)
  requires true;
  ensures length(vs) + 1 == length(append(vs, cons(x, nil)));
{
  length_append(vs, cons(x, nil));
}

@*/
```

```

class Cell
{
    Cell next;
    int content;

    public Cell (int v, Cell next)
        //@ requires true;
        //@ ensures Cell(this, next, v);
    {
        this.next = next;
        content = v;
    }

    public Cell clone()
        //@ requires Cell(this, ?n, ?v);
        //@ ensures Cell(this, n, v) &*& Cell(result, n, v);
    {
        Cell c = new Cell(this.content, this.next);
        return c;
    }
}

public class Queue
{
    Cell first;
    Cell last;
    int length;

    public Queue ()
        //@ requires true;
        //@ ensures QueueInv(this, nil);
    {
        first = null;
        last = null;
        length = 0;

        //@ close QueueInv(this, nil);
    }

    public void clear ()
        //@ requires QueueInv(this, ?l);
        //@ ensures QueueInv(this, nil);
    {
        //@ open QueueInv(this, l);
        first = null;
        last = null;
        length = 0;
        //@close QueueInv(this, nil);
    }
}

```

```

public void add (int x)
/*@ requires QueueInv(this, ?l);
  @ ensures QueueInv(this, append(l, cons(x, nil)));
  {
    Cell cell = new Cell(x, null);

    @ open QueueInv(this, l);

    if (last == null) {
      last = cell;
      first = cell;
      length = 1;

      @ cell_null_lseg(cell);
      @ close QueueInv(this, cons(x, nil));
    }
    else {
      if (last == first) {
        @ open lseg(first, last, ?vs1);

        @ Cell old_last = last;
        last.next = cell;
        last = cell;
        @ open Cell(last, null, x);

        length = 2;

        @ close lseg(cell, null, cons(x, nil));
        @ close QueueInv(this, append(l, cons(x, nil)));
      }
      else {
        @ Cell old_last = last;
        last.next = cell;
        last = cell;

        @ close Cell(last, null, x);

        length = length + 1;

        @ cell_null_lseg(last);
        @ lseg_merge(first, old_last, last);
        @ close QueueInv(this, _);
      }
    }
  }

public boolean isEmpty ()
/*@ requires QueueInv(this, ?l);

```



```

    //@ ensures QueueInv(this, l) &*& result == (l == nil);
    {
        return length == 0;
    }

    public int take ()
    //@ requires QueueInv(this, cons(?v, ?vs));
    //@ ensures QueueInv(this, vs) &*& result == v;
    {
        //@ open QueueInv(this, cons(v, vs));
        //@ open lseg(first, last, ?vv);

        if (first == last) {
            int c = first.content;
            this.clear();

            return c;
        }
        else {
            int c = first.content;
            length = length - 1;
            first = first.next;

            return c;
        }
    }
}

```