# Part 20 – Tasklet | Static Method

[ <https://embetronicx.com/tutorials/linux/device-drivers/tasklet-static-method/> ]

This is the Linux Device Driver Tutorial Part 20 – Tasklet Static Method Tutorial.

# Bottom Half

When Interrupt triggers, Interrupt Handler should be execute very quickly and it should not run for more time (it should not perform time-consuming tasks). If we have the interrupt handler which is doing more tasks then we need to divide into two halves.

1. Top Half
2. Bottom Half

Top Half is nothing but our interrupt handler. If our interrupt handler is doing less task, then top half is more than enough. No need of bottom half in that situation. But if our we have more work when interrupt hits, then we need bottom half. The bottom half runs in the future, at a more convenient time, with all interrupts enabled. So, The job of bottom halves is to perform any interrupt-related work not performed by the interrupt handler.

There are 4 bottom half mechanisms are available in Linux:

1. Work-queue
2. Threaded IRQs
3. Softirqs
4. **Tasklets**

# Tasklets in Linux Kernel

Tasklets are used to queue up work to be done at a later time. Tasklets can be run in parallel, but the same tasklet cannot be run on multiple CPUs at the same time. Also each tasklet will run only on the CPU that schedules it, to optimize cache usage. Since the thread that queued up the tasklet must complete before it can run the tasklet, race conditions are naturally avoided. However, this arrangement can be suboptimal, as other potentially idle CPUs cannot be used to run the tasklet. Therefore workqueues can, and should be used instead, and workqueues were already discussed [here](https://www.embetronicx.com/tutorials/linux/device-drivers/workqueue-in-linux-kernel/).

In short, a **tasklet** is something like a very small thread that has neither stack, not context of its own. Such “threads” work quickly and completely.

## Points To Remember

Before using Tasklets, you should consider these below points.

* Tasklets are atomic, so we cannot use **sleep()** and such synchronization primitives as [mutexes](https://en.wikipedia.org/wiki/Mutual_exclusion), [semaphores](https://en.wikipedia.org/wiki/Semaphore_(programming)), etc. from them. But we can use [spinlock.](https://en.wikipedia.org/wiki/Spinlock)
* A tasklet only runs on the same core (CPU) that schedules it.
* Different tasklets can be running in parallel. But at the same time, a tasklet cannot be called concurrently with itself, as it runs on one CPU only.
* Tasklets are executed by the principle of non-preemptive scheduling, one by one, in turn. We can schedule them with two different priorities: **normal**and **high**.

We can create tasklet in Two ways.

1. **Static Method**
2. **Dynamic Method**

In this tutorial we will see static method.

# Tasklet Structure

This is the important data structure for the tasklet.

struct tasklet\_struct

{

struct tasklet\_struct \*next;

unsigned long state;

atomic\_t count;

void (\*func)(unsigned long);

unsigned long data;

};

Here,

next – The next tasklet in line for scheduling.

state – This state denotes Tasklet’s State. TASKLET\_STATE\_SCHED (**Scheduled**) or TASKLET\_STATE\_RUN (**Running**).

count – It holds a nonzero value if the tasklet is disabled and 0 if it is enabled.

func –  This is the main function of the tasklet. Pointer to the function that needs to scheduled for execution at a later time.

data –  Data to be passed to the function “func”.

# Create Tasklet

The below macros used to create a tasklet.

## DECLARE\_TASKLET

|  |
| --- |
| This macro used to create the tasklet structure and assigns the parameters to that structure.  If we are using this macro then tasklet will be in enabled state.  **DECLARE\_TASKLET(name, func, data);**  name – name of the structure to be create.  func – This is the main function of the tasklet. Pointer to the function that needs to scheduled for execution at a later time.  data – Data to be passed to the function “func”. |

### Example

DECLARE\_TASKLET(tasklet,tasklet\_fn, 1);

Now we will see how the macro is working. When I call the macro like above, first it creates tasklet structure with the name of tasklet. Then it assigns the parameter to that structure. It will be looks like below.

struct tasklet\_struct tasklet = { NULL, 0, 0, tasklet\_fn, 1 };

(or)

struct tasklet\_struct tasklet;

tasklet.next = NULL;

taklet.state = TASKLET\_STATE\_SCHED; //Tasklet state is scheduled

tasklet.count = 0; //taskelet enabled

tasklet.func = tasklet\_fn; //function

tasklet.data = 1; //data arg

## DECLARE\_TASKLET\_DISABLED

|  |
| --- |
| The tasklet can be declared and set at disabled state, which means that tasklet can be scheduled, but will not run until the tasklet is specifically enabled. You need to use **tasklet\_enable**to enable.  **DECLARE\_TASKLET\_DISABLED(name, func, data);**  name – name of the structure to be create.  func – This is the main function of the tasklet. Pointer to the function that needs to schedule for execution at a later time.  data – Data to be passed to the function “func”. |

# Enable and Disable Tasklet

## tasklet\_enable

|  |
| --- |
| This used to enable the tasklet.  **void tasklet\_enable(struct);**  t – pointer to the tasklet struct |

## tasklet\_disable

|  |
| --- |
| This used to disable the tasklet , however it wait(s) for the completion of tasklet’s operation.  **void tasklet\_disable(struct tasklet\_struct \*t);**  t – pointer to the tasklet struct |

## tasklet\_disable\_nosync

|  |
| --- |
| This is used to disable tasklet immediately.  **void tasklet\_disable\_nosync(struct tasklet\_struct \*t);**  t – pointer to the tasklet struct |

**NOTE : If the tasklet has been disabled, we can still add it to the queue for scheduling, but it will not be executed on the CPU until it is enabled again. Moreover, if the tasklet has been disabled several times, it should be enabled exactly the same number of times, there is the count field in the structure for this purpose.**

# Schedule Tasklet

When we schedule the tasklet, then that tasklet is placed into one queue out of two, depending on the priority. Queues are organized as singly-linked lists. At that, each CPU has its own queues.

There are two priorities.

1. Normal Priority
2. High Priority

## tasklet\_schedule

|  |
| --- |
| Schedule a tasklet with normal priority. If a tasklet has previously been scheduled (but not yet run), the new schedule will be silently discarded.  **void tasklet\_schedule (struct tasklet\_struct \*t);**  t – pointer to the tasklet struct |

### Example

/\*Scheduling Task to Tasklet\*/

tasklet\_schedule(&tasklet);

## tasklet\_hi\_schedule

|  |
| --- |
| Schedule a tasklet with high priority. If a tasklet has previously been scheduled (but not yet run), the new schedule will be silently discarded.  **void tasklet\_hi\_schedule (struct tasklet\_struct \*t);**  t – pointer to the tasklet struct |

## tasklet\_hi\_schedule\_first

|  |
| --- |
| This version avoids touching any other tasklets. Needed for kmemcheck in order not to take any page faults while enqueueing this tasklet. Consider VERY carefully whether you really need this or tasklet\_hi\_schedule().  **void tasklet\_hi\_schedule\_first(struct tasklet\_struct \*t);**  t – pointer to the tasklet struct |

# Kill Tasklet

Finally, after a tasklet has been created, it’s possible to delete a tasklet through these below functions.

## tasklet\_kill

|  |
| --- |
| This will wait for its completion, and then kill it.  **void tasklet\_kill( struct tasklet\_struct \*t );**  t – pointer to the tasklet struct |

### Example

/\*Kill the Tasklet \*/

tasklet\_kill(&tasklet);

## tasklet\_kill\_immediate

|  |
| --- |
| This is used only when a given CPU is in the dead state.  **void tasklet\_kill\_immediate( struct tasklet\_struct \*t, unsigned int cpu );**  t – pointer to the tasklet struct  cpu – cpu num |

# Programming

## Driver Source Code

In that source code, When we read the /dev/etx\_device interrupt will hit (To understand interrupts in Linux go to [this tutorial](https://www.embetronicx.com/tutorials/linux/device-drivers/linux-device-driver-tutorial-part-13-interrupt-example-program-in-linux-kernel/)). Whenever interrupt hits, I’m scheduling the task to the tasklet. I’m not going to do any job in both interrupt handler and tasklet function,  since it is a tutorial post. But in real tasklet, this function can be used to carry out any operations that need to be scheduled.

NOTE: In this source code many unwanted functions will be there (which is not related to the Tasklet). Because I’m just maintaining the source code throughout these Device driver series.

#include <linux/kernel.h>

#include <linux/init.h>

#include <linux/module.h>

#include <linux/kdev\_t.h>

#include <linux/fs.h>

#include <linux/cdev.h>

#include <linux/device.h>

#include<linux/slab.h> //kmalloc()

#include<linux/uaccess.h> //copy\_to/from\_user()

#include<linux/sysfs.h>

#include<linux/kobject.h>

#include <linux/interrupt.h>

#include <asm/io.h>

#include <asm/hw\_irq.h>

#define IRQ\_NO 11

void tasklet\_fn(unsigned long);

/\* Init the Tasklet by Static Method \*/

DECLARE\_TASKLET(tasklet,tasklet\_fn, 1);

/\*Tasklet Function\*/

void tasklet\_fn(unsigned long arg)

{

printk(KERN\_INFO "Executing Tasklet Function : arg = %ld\n", arg);

}

//Interrupt handler for IRQ 11.

static irqreturn\_t irq\_handler(int irq,void \*dev\_id) {

printk(KERN\_INFO "Shared IRQ: Interrupt Occurred");

/\*Scheduling Task to Tasklet\*/

tasklet\_schedule(&tasklet);

return IRQ\_HANDLED;

}

volatile int etx\_value = 0;

dev\_t dev = 0;

static struct class \*dev\_class;

static struct cdev etx\_cdev;

struct kobject \*kobj\_ref;

static int \_\_init etx\_driver\_init(void);

static void \_\_exit etx\_driver\_exit(void);

/\*\*\*\*\*\*\*\*\*\*\*\*\*\*\* Driver Fuctions \*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*/

static int etx\_open(struct inode \*inode, struct file \*file);

static int etx\_release(struct inode \*inode, struct file \*file);

static ssize\_t etx\_read(struct file \*filp,

char \_\_user \*buf, size\_t len,loff\_t \* off);

static ssize\_t etx\_write(struct file \*filp,

const char \*buf, size\_t len, loff\_t \* off);

/\*\*\*\*\*\*\*\*\*\*\*\*\*\*\* Sysfs Fuctions \*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*/

static ssize\_t sysfs\_show(struct kobject \*kobj,

struct kobj\_attribute \*attr, char \*buf);

static ssize\_t sysfs\_store(struct kobject \*kobj,

struct kobj\_attribute \*attr,const char \*buf, size\_t count);

struct kobj\_attribute etx\_attr = \_\_ATTR(etx\_value, 0660, sysfs\_show, sysfs\_store);

static struct file\_operations fops =

{

.owner = THIS\_MODULE,

.read = etx\_read,

.write = etx\_write,

.open = etx\_open,

.release = etx\_release,

};

static ssize\_t sysfs\_show(struct kobject \*kobj,

struct kobj\_attribute \*attr, char \*buf)

{

printk(KERN\_INFO "Sysfs - Read!!!\n");

return sprintf(buf, "%d", etx\_value);

}

static ssize\_t sysfs\_store(struct kobject \*kobj,

struct kobj\_attribute \*attr,const char \*buf, size\_t count)

{

printk(KERN\_INFO "Sysfs - Write!!!\n");

sscanf(buf,"%d",&etx\_value);

return count;

}

static int etx\_open(struct inode \*inode, struct file \*file)

{

printk(KERN\_INFO "Device File Opened...!!!\n");

return 0;

}

static int etx\_release(struct inode \*inode, struct file \*file)

{

printk(KERN\_INFO "Device File Closed...!!!\n");

return 0;

}

static ssize\_t etx\_read(struct file \*filp,

char \_\_user \*buf, size\_t len, loff\_t \*off)

{

struct irq\_desc \*desc;

printk(KERN\_INFO "Read function\n");

/\* New way of mapping irq lines on newer kernels \*/

desc = irq\_to\_desc(11);

if (!desc)

return -EINVAL;

\_\_this\_cpu\_write(vector\_irq[59], desc);

/\* Triggering Interrupt \*/

asm("int $0x3B"); // Corresponding to irq 11

return 0;

}

static ssize\_t etx\_write(struct file \*filp,

const char \_\_user \*buf, size\_t len, loff\_t \*off)

{

printk(KERN\_INFO "Write Function\n");

return 0;

}

static int \_\_init etx\_driver\_init(void)

{

/\*Allocating Major number\*/

if((alloc\_chrdev\_region(&dev, 0, 1, "etx\_Dev")) <0){

printk(KERN\_INFO "Cannot allocate major number\n");

return -1;

}

printk(KERN\_INFO "Major = %d Minor = %d \n",MAJOR(dev), MINOR(dev));

/\*Creating cdev structure\*/

cdev\_init(&etx\_cdev,&fops);

/\*Adding character device to the system\*/

if((cdev\_add(&etx\_cdev,dev,1)) < 0){

printk(KERN\_INFO "Cannot add the device to the system\n");

goto r\_class;

}

/\*Creating struct class\*/

if((dev\_class = class\_create(THIS\_MODULE,"etx\_class")) == NULL){

printk(KERN\_INFO "Cannot create the struct class\n");

goto r\_class;

}

/\*Creating device\*/

if((device\_create(dev\_class,NULL,dev,NULL,"etx\_device")) == NULL){

printk(KERN\_INFO "Cannot create the Device 1\n");

goto r\_device;

}

/\*Creating a directory in /sys/kernel/ \*/

kobj\_ref = kobject\_create\_and\_add("etx\_sysfs",kernel\_kobj);

/\*Creating sysfs file for etx\_value\*/

if(sysfs\_create\_file(kobj\_ref,&etx\_attr.attr)){

printk(KERN\_INFO"Cannot create sysfs file......\n");

goto r\_sysfs;

}

if (request\_irq(IRQ\_NO, irq\_handler, IRQF\_SHARED, "etx\_device", (void \*)(irq\_handler))) {

printk(KERN\_INFO "my\_device: cannot register IRQ ");

goto irq;

}

printk(KERN\_INFO "Device Driver Insert...Done!!!\n");

return 0;

irq:

free\_irq(IRQ\_NO,(void \*)(irq\_handler));

r\_sysfs:

kobject\_put(kobj\_ref);

sysfs\_remove\_file(kernel\_kobj, &etx\_attr.attr);

r\_device:

class\_destroy(dev\_class);

r\_class:

unregister\_chrdev\_region(dev,1);

cdev\_del(&etx\_cdev);

return -1;

}

void \_\_exit etx\_driver\_exit(void)

{

/\*Kill the Tasklet \*/

tasklet\_kill(&tasklet);

free\_irq(IRQ\_NO,(void \*)(irq\_handler));

kobject\_put(kobj\_ref);

sysfs\_remove\_file(kernel\_kobj, &etx\_attr.attr);

device\_destroy(dev\_class,dev);

class\_destroy(dev\_class);

cdev\_del(&etx\_cdev);

unregister\_chrdev\_region(dev, 1);

printk(KERN\_INFO "Device Driver Remove...Done!!!\n");

}

module\_init(etx\_driver\_init);

module\_exit(etx\_driver\_exit);

MODULE\_LICENSE("GPL");

MODULE\_AUTHOR("EmbeTronicX <embetronicx@gmail.com>");

MODULE\_DESCRIPTION("A simple device driver - Tasklet part 1");

MODULE\_VERSION("1.15");

**MakeFile**

obj-m += driver.o

KDIR = /lib/modules/$(shell uname -r)/build

all:

make -C $(KDIR) M=$(shell pwd) modules

clean:

make -C $(KDIR) M=$(shell pwd) clean

# Building and Testing Driver

* Build the driver by using Makefile (sudo make)
* Load the driver using sudo insmod driver.ko
* To trigger interrupt read device file (sudo cat /dev/etx\_device)
* Now see the Dmesg (dmesg)

*$ dmesg*

[ 8592.698763 ] Major = 246 Minor = 0  
[ 8592.703380 ] Device Driver Insert…Done!!!  
[ 8601.716673 ] Device File Opened…!!!  
[ 8601.716697 ] Read function  
[ 8601.716727 ] Shared IRQ: Interrupt Occurred  
[ 8601.716732 ] Executing Tasklet Function : arg = 1  
[ 8601.716741 ] Device File Closed…!!!

[ 8603.916741 ] Device Driver Remove…Done!!!

* We can able to see the print “**Shared IRQ: Interrupt Occurred**“ and “**Executing Tasklet Function : arg = 1**“
* Unload the module using sudo rmmod driver