An Introduction to Database Systems

chapter 4. An Introduction to SQL

4.1 Introduction

- \square SQL2: "SQL/92" \rightarrow SQL3(2000)
 - International Standard Database Language SQL(1992)

□ Remarks

- SQL was originally intended to be a data sublanguage. With the incorporation of the persistent stored Modules feature(PSM) (late 1996), SQL became computationally complete.
 - CALL, RETURN, SET, CASE, IF, LOOP, LEAVE, WHILE, REPEAT, as well as several related features such as variables and exception handlers
- SQL was the term table in place of both of the terms relation and relvar.
- SQL is an enormous language
 - Standard document > 600 pages, SQL3 > 1200 pages
- SQL is very far from being the perfect relational language

□ Definitional operations

Fig.4.1 The suppliers and parts database (SQL definition)

```
CREATE TABLE S
  (S#
                  CHAR(5),
                  CHAR(20),
   SNAME
   STATUS
                  NUMERIC(5),
   CITY
                  CHAR(15),
  PRIMARY KEY
                  (S#));
CREATE TABLE P
  (P#
                   CHAR(6),
                   CHAR(20),
   PNAME
   COLOR
                   CHAR(6),
                   NUMERIC(5, 1),
    WEIGHT
   CITY
                   CHAR(15),
  PRIMARY KEY
                   (P#));
```

```
CREATE TABLE SP

( S# CHAR(5),

P# CHAR(6),

QTY NUMERIC(9),

PRIMARY KEY (S#, P#)),

FOREIGN KEY (S#) REFERENCES S,

FOREIGN KEY (P#) REFERENCES P;
```

□ CREATE TABLE statement

• SQL does not permit users to define their own types. Columns must be defined in terms of built-in (system-defined) types only:

CHARACTER	[VARYING]	(n)	INTEGER	DATA
BIT	[VARYING]	(n)	SMALLINT	TIME
NUMERIC	(p, g)		FLOAT (p)	TIMESTAMP
DECIMAL	(p, g)			INTERVAL

lengths or precisions → integrity constraints

□ Manipulate operations

• SELECT, INSERT, UPDATE, DELETE

Restrict :				
	Result:	S#	P#	QTY
SELECT S#, P#, QTY		S1	P5	100
FROM SP		S 1	P6	100
WHERE QTY <150 ;				
Project :	Result:	S#	CITY	
		S1	London	
SELECT S#, CITY		S2	Paris	
FROM S;		S 3	Paris	
		S4	London	
		S 5	Athens	

□ Manipulate operations

SELECT, INSERT, UPDATE, DELETE

Join: SELECT S.S#, SNAME, STATUS, CITY, P#, QTY FROM S, SP WHERE S.S# = P.P#; **STATUS** Result: S# SNAME CITY P# **QTY S1 Smith** 20 London **P1** 300 **S1 Smith** London 200 20 **P2 S1** Smith 20 London **P3** 400 ... **S4 P5** Clark 20 London 400

□ Manipulate operations

- Qualified name (S.S#, SP.S#)
- A shorthand form

SELECT * or SELECT S.*
FROM S FROM S

INSERT, UPDATE, DELETE – set level operations

INSERT
INTO TEMP (P#, WEIGHT)
SELECT P#, WEIGHT
FROM P
WHERE COLOR='Red';

```
UPDATE S
SET STATUS = STATUS * 2
WHERE CITY = 'Paris";
```

```
DELETE
FROM SP
WHRE P# = 'P2';
```

- SQL does not include a direct analog of the relational assignment operation
 - INSERT... SELECT...

4.3 The Catalog (CONT.)

- □ SQL standard : Information schema
- SQL catalog consists of the descriptors for an individual database
- SQL schema consists of the descriptors for that portion of that database that belongs to some individual user
- □ Database → any number of catalogs → into an number of schemas

INFORMATION_SCHEMA

4.3 The Catalog (CONT.)

- □ Information Schema
 - A set of SQL tables whose contents echo all of the definitions from all of the other schemas in the catalog
 - A set of views of a hypothetical "Definition Schema"
- Implementation is required
 - a. to support some kind of Definition Schema and
 - b. to support views of that Definition Schema
- □ Rationale
 - Existing products do support something akin to the Definition Schema. However, those Definition Schemas vary widely from one to another
 - 2. an (not the) Information Schema

4.3 The Catalog (CONT.)

□ List some of the more important Information Schema Views

SCHEMATA REFERENTIAL_CONSTRAINTS

DOMAINS CHECK_CONSTRAINTS

TABLES KEY_COLUMN_USAGE

VIEWS ASSERTIONS

COLUMNS VIEW_TABLE_USAGE

TABLE_PRIVILEGES VIEW_COLUMN_USAGE

COLUMN_PRIVILEGES CONSTRAINT_TABLE_USAGE

USAGE_PRIVILEGES CONSTRAINT_COLUMN_USAGE

DOMAIN_CONSTRAINTS CONSTRAINT_COLUMN_USAGE

TABLE_CONSTRAINTS CONSTRAINT_DOMAIN_USAGE

4.4 Views (CONT.)

□ an SQL view

```
CREATE VIEW GOOD_SUPPLIER

AS SELECT S#, STATUS, CITY

FROM S

WHERE STATUS > 15;
```

□ An SQL query against this view

```
SELECT S#, STATUS

FROM GOOD_SUPPLIER

WHERE CITY = 'London';
```

4.4 Views (CONT.)

```
□ Substituting the view definition to the view name
    SELECT GOOD_SUPPLIER.S#, GOOD_SUPPLIER.STATUS
    FROM
           (SELECT S#, STATUS, CITY
            FROM S
            WHERE STATUS > 15) AS GOOD_SUPPLIER
    WHERE GOOD_SUPPLIER.CITY = 'London';
□ Simplify it
        SELECT S#, STATUS
        FROM
               S
        WHERE STATUS > 15
        AND CITY = 'London';
```

4.4 Views (CONT.)

□ DELETE operation

```
DELETE

FROM GOOD_SUPPLIER

WHERE CITY = 'London';

DELETE
```

WHERE STATUS > 15

AND CITY = 'London';

4.5 Transactions

- □ SQL's analog of COMMIT and ROLLBACK
 - COMMIT [WORK], ROLLBACK[WORK]
- □ BEGIN TRANSACTION (X)
 - A transaction is begun implictly whenever the program executes a transaction-initiating statement

□ dual-mode principle

 any SQL statement that can be used interactively can also be used in an application program

□ Embedded SQL statement

- are prefixed EXEC SQL
- are terminated special terminator
- can include references to host variables

□ executable statement

- can appear wherever an executable host statement can appear
- but, not declarative statement (DECLARE CURSOR, BEGIN, END, DECLARE SECTION, etc..)

□ INTO clause

- data type compatibility
- □ all host variable will be referenced in SQL
 - must be defined within an embedded SQL declare section

□ SQLSTATE

- every SQL program must include a host variable, SQLSTATE
- 00000 :executed successfully
- 02000 : executed but no data was found to satisfy the request

□ WHENEVER

- EXEC SQL WHENEVER condition action terminator
- condition: SQLERROR(error occurred), NOT FOUND(0200)
- action : CONTINUE, GO TO

□ example

```
EXEC SQL BEGIN DECLARE SECTION;
    DCL SQLSTATE
                   CHAR(5);
    DCL P#
                   CHAR(6);
    DCL WEIGHT
                   FIXED DECIMAL(3);
EXEC SQL END DECLARE SECTION;
P# = 'P2';
                                     /* for example
EXEC SQL
           SELECT P.WEIGHT
                            :WEIGHT
                    INTO
                    FROM P
                    WHERE P.P# = :P#;
IF SQLSTATE = '00000'
                            /* WEIGHT = retrieved value */
THEN
ELSE
                            /* some exception occurred */
```

□ operations not involving cursor

Singleton SELECT: Get status and city for the supplier whose supplier number is given by the host variable GIVENS#

EXEC SQL SELECT STATUS, CITY

INTO :RANK, :CITY

FROM S

S# = :GIVENS#;

INSERT: Insert a new part into table P

(part no., name, and weight given by host variable P#, PNAME,PWT; color and city unknown)

EXEC SQL INSERT

INTO P (P#, PNAME, WEIGHT)

VALUSE (:P#, :PNAME, :PWT)

UPDATE: increase the status of all London suppliers by the amount given by the host variable RAISE

EXEC SQL UPDATE S

SET STATUS = STATUS + :RAISE

WHERE CITY = 'London';

If no rows satisfy the WHERE, SQLSTATE will be set to 02000

DELETE: Delete all shipments for suppliers whose city is given by the host variable CITY

```
EXEC SQL DELETE
```

FROM SP

WHERE :CITY = (SELECT CITY

FROM S

WHERE S.S# = SP.S#);

- □ operations Involving Cursors
- □ declare

EXEC SQL DECLARE cursor CURSOR

FOR <table-expression>

[ORDER BY order-item-commalist]

- □ executable statements to operate on CURSOR
 - OPEN, FETCH, CLOSE

□ Multi-row retrieval example

```
EXEC SQL DECLARE X CURSOR FOR /* define the cursor */
            SELECT S.S#, S.SNAME, S.STATUS
            FROM
                      S
            WHERE S.CITY = :Y
            ORDER BY S# ASC;
EXEC SQL OPEN X;
                                      /* execute the query */
            DO for all S rows accessible via X;
                     EXEC SQL FETCH X INTO :S#, :SNAME, :STATUS;
                                              /* fetch next supplier */
            END;
EXEC SQL
                                      /* deactivate cursor X */
            CLOSE X;
```

dynamic SQL

 a set of embedded SQL facilities that are provided specifically to allow the construction of generalized, online, and possibly interactive applications.

online application

- 1. Accept a command from the terminal
- Analyze that command
- 3. Execute appropriate SQL statements on the database
- 4. Return a message and/or results to the terminal

□ dynamic statements

- PREPARE and EXECUTE
- example

```
DCL SQLSOURCE CHAR VARYING (65000);
SQLSOURCE = 'DELETE FROM SP WHERE QTY <300';
EXEC SQL PREPARE SLQPREPPED FROM :SQLSOURCE;
EXEC SQL EXECUTE SQLPREPPED;
```

□ SQLSOURCE

identifies a PL/I a varying length character string variable

□ SQLPREPPED

 identifies an SQL variable, not a PL/I variable, that will be used to hold the compiled form of the SQL statement whose source form is given in SQLSOURCE

□ assignment " SQLSOURCE = ..."

 assigns to SQLSOURCE the source form of an SQL DELETE statement

□ PREPARE

 "prepares"(i.e., compiles) it to produce an executable version, which it stores in SQLPREPPED

EXECUTE

executes that SQLPREPPED version and thus causes the actual DELETE to occur

- □ SQL Call-level interface (SQL/CLI, CLI)
 - based on MS ODBC
 - CLI permits an application written in one of the usual host languages to issue database requests by invoking certain vendorprovided CLI routines
 - address the same general problem as dynamic SQL does.
- SQL/CLI represent a better approach than dynamic SQL does
 - Dynamic SQL is a source code standard → require the srvices of some kind of SQL complier (PREPARE, EXECUTE).
 - CLI merely standardize the details of certain routine invocations. (no special complier) (shrink-wrapped object code form)
 - 2. Those applications can be DBMS-independent.

4.7 SQL in not Perfect

There is no product on the market today that supports every detail of the relational model