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# REPORT: HENSEL LIFTING AND NEWTON ITERATION IN VALUTATION RINGS

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# **Chapter 1**

## **Introduction**

# Chapter 2

## Hensel Lifting

The hensel method described here will lift an approximate factorization of a polynomial over a Hensel Ring  $R$  with valuation  $v$  where the factors are relatively prime. We will show a linear convergence and a quadratic convergence behavior for the liftings.

### 2.1 Hensel Ring

#### Definition 2.1.1: Hensel Ring

A ring with valuation  $v : R \rightarrow \mathbb{R}_{\geq 0}$  is called a Hensel Ring if:

- (i)  $\forall a \in R, v(a) \leq 1$
- (ii)  $\forall a, b \in R, \forall \epsilon > 0, \exists c \in R$  such that  $(v(a) \leq v(b) \implies v(a - bc) \leq \epsilon)$

In other words  $R$  is Hensel iff it is contained and dense in the valuation ring of its quotient field (with respect to the unique extension of  $v$ ). We sometimes call such  $v$  a Hensel Valuation.

In condition (ii) we assume we can compute the  $c$  efficiently.

#### Theorem 2.1.1

Condition (i) of Hensel Ring  $\implies v$  is Non-Archimedean.

*Proof.* Let  $a, b \in R$ . Now

$$\begin{aligned}
 v(a + b)^k &= v((a + b)^k) \\
 &= v\left(\sum_{i=0}^k \binom{k}{i} a^{n-i} b^i\right) \leq \sum_{i=0}^k v\left(\binom{k}{i}\right) v(a)^{n-i} v(b)^i \\
 &\leq \sum_{i=0}^k v(a)^{n-i} v(b)^i \leq \sum_{i=0}^k M^{n-i} M^i \quad [m = \max v(a), v(b)] \\
 &= \sum_{i=0}^k M^k = M^k (k + 1)
 \end{aligned}$$

Hence

$$\left(\frac{v(a + b)}{M}\right)^k \leq (k + 1) \iff \frac{v(a + b)}{M} \leq (1 + k)^{\frac{1}{k}}$$

As  $k \rightarrow \infty$  the RHS approaches 1 so  $v(a + b) \leq M$ . ■

### Example 2.1 ( $p$ -adic Valuations)

- $\mathbb{Z}$  with  $p$ -adic valuation  $v_p$  where  $p \in \mathbb{N}$  is prime is a Hensel ring. Here  $v_p(a) = p^{-n}$  where  $n = \max\{k \geq 0 \mid p^k \mid a\}$
- $\mathbb{F}[y]$  with  $p$ -adic valuation  $v_p$  where  $p \in \mathbb{F}[y]$  is an irreducible polynomial is a Hensel Ring. Here  $v_p(f) = 2^{-n \deg p}$  where  $n = \max\{k \geq 0 \mid p^k \mid f\}$

#### Note:-

From the valuation  $v$  over  $R$  we naturally get a valuation  $v$  over the polynomial ring  $R[x]$  by defining

$$\forall f \in R[x], \text{ let } f = \sum_{i=0}^n f_i x^i, \text{ then } v\left(\sum_{i=0}^n f_i x^i\right) = \max_i \{v(f_i)\}$$

## 2.2 Conditions related to Hensel's Lemma

We will define 5 conditions. First suppose we have:

- (1)  $f \in R[x]$
- (2)  $f_0, \dots, f_m \in R[x] \quad \mathcal{F} = \{f_i : 0 \leq i \leq m\}$
- (3)  $f_0^*, \dots, f_m^* \in R[x] \quad \mathcal{F}^* = \{f_i^* : 0 \leq i \leq m\}$
- (4)  $s_0, \dots, s_m \in R[x] \quad \mathcal{S} = \{s_i : 0 \leq i \leq m\}$
- (5)  $s_0^*, \dots, s_m^* \in R[x] \quad \mathcal{S}^* = \{s_i^* : 0 \leq i \leq m\}$
- (6)  $z \in R$
- (7)  $\alpha, \delta, \epsilon \in \mathbb{R}$
- (8)  $\delta^* \in \mathbb{R}$
- (9)  $\gamma = \max\{\delta, \alpha\epsilon\}$

As you can see the set  $\mathcal{F}^*$  basically represents the lift of  $\mathcal{F}$  but here since we are saying the conditions in more generality we are not assuming any relations among them and we define some conditions involving them.

- $H_1(m, f, \mathcal{F}, \mathcal{S}, \epsilon) := v\left(f - \prod_{i=0}^m f_i\right) \leq \epsilon < 1$
- $H_2(m, f, \mathcal{F}, \mathcal{S}, z, \delta) := v\left(\sum_{i=0}^m s_i \prod_{j \neq i} f_j - z\right) < \log \delta < 1$
- $H_3(m, f, \mathcal{F}, \mathcal{S}, z, \alpha, \delta, \epsilon) :=$ 
  - (1)  $f_1, \dots, f_m$  are monic
  - (2)  $\deg\left(\prod_{i=0}^m f_i\right) \leq \deg f$
  - (3)  $\deg s_i \leq \deg f_i \forall i \in [m]$
  - (4)  $\alpha\delta \leq 1, \alpha\epsilon^2 \leq 1$
  - (5)  $1 \leq \alpha v(z)$
- $H_4(m, \mathcal{F}, \mathcal{F}^*, \mathcal{S}, \mathcal{S}^*, \alpha, \delta, \epsilon) :=$ 
  - (1)  $v(f_i^* - f_i) \leq \alpha\epsilon \quad \forall 0 \leq i \leq m$
  - (2)  $v(s_i^* - s_i) \leq \alpha\epsilon \quad \forall 0 \leq i \leq m$
  - (3)  $\deg f_i^* = \deg f_i \quad \forall i \in [m]$
  - (4)  $\deg s_i < \deg f_i \implies \deg s_i^* < \deg f_i^* \quad \forall i \in [m]$
- $H_5(m, f, \mathcal{F}, \mathcal{F}^*, \mathcal{S}, \mathcal{S}^*, \alpha, \delta, \epsilon, \delta^*) :=$  Let  $p \in [m]$ . Then suppose
  - $\mathcal{I}_p = \{I_0, I_1, \dots, I_p\}$  be a partition of  $\{0, \dots, m\}$  with  $0 \in I_0$ .
  - $\overline{\mathcal{F}}_p^m = \{\bar{f}_i : i \in [p]\} \subseteq R[x]$  be a set of monic polynomials

Then define:

$$F_i = \prod_{j \in I_i} f_j, \quad F_i^* = \prod_{j \in I_i} f_j^*, \quad \mathfrak{s}_i^* = \sum_{j \in I_i} s_j \frac{F_i^*}{f_j^*}$$

So now we denote:

$$\mathcal{F} = \{F_i : 0 \leq i \leq p\}, \quad \mathcal{F}^* = \{F_i^* : 0 \leq i \leq p\}, \quad \mathcal{S} = \{\mathfrak{s}_i^* : 0 \leq i \leq p\}$$

Assume:

1.  $v(\bar{f}_i - F_i) \leq \alpha\epsilon \forall i \in [p]$
2.  $\alpha v(s_i) \leq 1 \forall 0 \leq i \leq m$
3.  $\alpha\delta < 1, \alpha^2\delta \leq 1$
4.  $\alpha^2\epsilon < 1, \alpha^3\epsilon \leq 1$

Then the following are equivalent:

- (i)  $\exists \bar{f}_0, \bar{s}_0, \dots, \bar{s}_p \in R[x]$  denote

$$\bar{\mathcal{F}} = \{\bar{f}_i : 0 \leq i \leq p\}, \quad \bar{\mathcal{S}} = \{\bar{s}_i : 0 \leq i \leq p\}$$

then the following conditions are true:

- (a)  $H_1(p, f, \bar{\mathcal{F}}, \bar{\mathcal{S}}, \epsilon^*)$
- (b)  $H_2(p, f, \bar{\mathcal{F}}, \bar{\mathcal{S}}, z, \delta^*)$
- (c)  $H_3(p, f, \bar{\mathcal{F}}, \bar{\mathcal{S}}, z, \alpha^*, \delta^*, \epsilon^*)$
- (d)  $H_4(p, f, \bar{\mathcal{F}}, \bar{\mathcal{S}}, z, \alpha^*, \delta^*, \epsilon^*)$

where  $\alpha^* = \alpha, \epsilon^* = \alpha\epsilon\gamma$

- (ii)  $\exists \bar{f}_0 \in R[x]$  such that  $H_1(p, f, \bar{\mathcal{F}}, \bar{\mathcal{S}}, \epsilon^*)$  is true

- (iii)  $\forall i \in [p]$  we have  $v(\bar{f}_i - F_i^*) \leq \epsilon^*$ .

The first 3 conditions here together imply that: From  $H_1$  we get that  $f_0 \cdots f_m$  is a good approximation of factorization of  $f$  with  $\epsilon$ -precision,  $H_2 \implies z$  plays a similar role to the gcd of  $f_0, \dots, f_m$  and it shows the generalized bezout's identity for gcd for multiple elements. In the usual treatment of Hensel's Lemma  $f_0, \dots, f_m$  are relatively prime (more precisely their images in the residue class field or  $R$  modulo the maximal ideal  $\langle a \in R \mid v(a) < 1 \rangle$  satisfy the assumption then one can find  $s_0, \dots, s_m, \delta$  satisfying  $H_2$  with  $z = 1$ . One can set  $\alpha = 1$  or in general one can choose  $\alpha = \frac{1}{v(z)}$ . Thus  $H_2$  states that  $f_0, \dots, f_m$  are approximately pairwise relatively prime.

$H_4$  shows the connection between the lifts  $f_i^*, s_i^*$  and  $f_i, s_i$ .

$H_5$  basically states that the lifts are unique in the sense that one can group some of the  $f_i^*$ s to form  $F_0, \dots, F_p$  and change  $F_i$  to  $\bar{f}_i$  with precision  $\epsilon^*$  and still one will have the factorization of  $f$  with precision  $\epsilon^*$ .  $H_5$  is very important for the factorization algorithm in chapter 6.

Now we will state the Hensel's Lemma and will later give the algorithm to obtain the lifts.

## 2.3 Hensel's Lemma

First we will prove a helping lemma which will be very much usefull in the proof of Hensel's Lemma then we will state the actual theorem.

### Theorem 2.3.1

- (i) Let  $a, f, p, s \in R[x]$  such that  $f$  is monic and  $s = pf + a$  with  $\deg a < \deg f$ . Then we have  $v(p) \leq v(s)$  and  $v(a) \leq v(s)$
- (ii) Let  $h_0, \dots, h_m \in R[x]$  and  $h_0^*, \dots, h_m^* \in R[x]$  such that we have  $v(h_i^* - h_i) \leq \epsilon$  for all  $0 \leq i \leq m$ . Then we have  $v\left(\prod_{i=0}^m h_i - \prod_{i=0}^m h_i^*\right) \leq \epsilon$

**Theorem 2.3.2 Hensel's Lemma**

Assume that we have  $f \in R[x]$ ,  $\mathcal{F} = \{f_0, \dots, f_m\} \subseteq R[x]$ ,  $\mathcal{S} = \{s_0, \dots, s_m\} \subseteq R[x]$ ,  $z \in R$  and  $\alpha, \delta, \epsilon \in \mathbb{R}$  which satisfy:

1.  $H_1(m, f, \mathcal{F}, \mathcal{S}, \epsilon)$
2.  $H_2(m, f, \mathcal{F}, \mathcal{S}, z, \delta)$
3.  $H_3(m, f, \mathcal{F}, \mathcal{S}, z, \alpha, \delta, \epsilon)$

Then we can compute efficiently

$$\mathcal{F}^* = \{f_i^* : 0 \leq i \leq m\} \quad \text{and} \quad T = \{t_0, \dots, t_m\}$$

such that

- (i) **Linear Case:**  $\mathcal{S}^* = \mathcal{S}$  and  $\delta^* = \gamma$ ,  $\epsilon^* = \alpha\gamma\epsilon$ . Then we have the following conditions hold:

- (a)  $H_1(m, f, \mathcal{F}^*, \mathcal{S}^*, \epsilon^*)$
- (b)  $H_2(m, f, \mathcal{F}^*, \mathcal{S}^*, z, \delta^*)$
- (c)  $H_3(m, f, \mathcal{F}^*, \mathcal{S}^*, z, \alpha, \delta^*, \epsilon^*)$
- (d)  $H_4(m, f, \mathcal{F}, \mathcal{F}^*, \mathcal{S}, \mathcal{S}^*, \alpha, \delta, \epsilon)$
- (e)  $H_5(m, f, \mathcal{F}, \mathcal{F}^*, \mathcal{S}, \mathcal{S}^*, \alpha, \delta, \epsilon, \delta^*)$

- (ii) **Quadratic Case:**  $\mathcal{S}^* = T$  and  $\delta^* = \alpha\gamma^2$ ,  $\epsilon^* = \alpha\gamma\epsilon$ . Assume that  $\deg s_i > \deg f_i$  for  $0 \leq i \leq m$ . Then we have the following conditions hold:

- (a)  $H_1(m, f, \mathcal{F}^*, \mathcal{S}^*, \epsilon^*)$
- (b)  $H_2(m, f, \mathcal{F}^*, \mathcal{S}^*, z, \delta^*)$
- (c)  $H_3(m, f, \mathcal{F}^*, \mathcal{S}^*, z, \alpha, \delta^*, \epsilon^*)$
- (d)  $H_4(m, f, \mathcal{F}, \mathcal{F}^*, \mathcal{S}, \mathcal{S}^*, \alpha, \delta, \epsilon)$
- (e)  $H_5(m, f, \mathcal{F}, \mathcal{F}^*, \mathcal{S}, \mathcal{S}^*, \alpha, \delta, \epsilon, \delta^*)$

**2.4 Hensel's Algorithm****2.5 Proof of Hensel's Lemma**

# **Chapter 3**

## **Newton Iteration**



## **Chapter 4**

# **Solving Differential Equations**

## **Chapter 5**

### **Finding Short Vectors in Modules**

## **Chapter 6**

# **Factorization of Polynomials**