
ALGEBRAIC GEOMETRIC CODES

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Contents

1	Mathematics	2
1.1	Divisors	2
1.2	Reimann-Roch Spaces	2
1.3	Differentials	2
1.4	Reimann-Roch Theorem	2
1.5	Index of speciality	3
2	Codes from Algebraic Curves	4
2.1	Preliminaries	4
2.2	Geometric Reed Solomon Codes	5
3	Bibliography	6

1.1 Divisors

1.2 Reimann-Roch Spaces

Definition 1.2.1 (Reimann-Roch Spaces). *For any divisor $\mathcal{D} \in \tilde{\mathfrak{D}}$*

$$\mathcal{L}(\mathcal{D}) = \{f \in \mathbb{F}(\mathcal{X})^* \mid (f) + \mathcal{D} \succcurlyeq 0\} \cup \{0\}$$

The dimension of $\mathcal{L}(\mathcal{D})$ over \mathbb{F} is denoted by $l(\mathcal{D})$

Theorem 1.2.1. (i) *If $\deg(\mathcal{D}) < 0$ then $l(\mathcal{D}) = 0$*

(ii) $l(\mathcal{D}) \leq 1 + \deg(\mathcal{D})$

Theorem 1.2.2. $\mathcal{L}(0) = \mathbb{F}$. Hence $l(0) = 1$

1.3 Differentials

1.4 Reimann-Roch Theorem

Theorem 1.4.1 (Reimann-Roch Theorem). *\mathcal{D} is a divisor on a smooth projective curve with genus g . Then for any canonical divisor W*

$$l(\mathcal{D}) - l(W - \mathcal{D}) = \deg(\mathcal{D}) - (g - 1)$$

Corollary 1.4.2. *For any canonical divisor W , $\deg(W) = 2g - 2$*

Proof: Take $\mathcal{D} = W$. Then $l(W - \mathcal{D}) = l(0) = 1$ by [Theorem 1.2.2](#). So we have

$$l(W) - 1 = \deg(W) - (g - 1)$$

. By definition $l(W) = g$. Hence we have $g - 1 = \deg(W) - (g - 1) \iff \deg(W) = 2g - 2$. ■

With the help of this corollary we can finally focus on the divisors which we will actually use to define codes. The following corollary gives the dimension of the Reimann-Roch Spaces of divisors with degree more than $2g - 2$.

Corollary 1.4.3. *Let \mathcal{D} be a divisor on a smooth projective curve of genus g and let $\deg(\mathcal{D}) > 2g - 2$. Then*

$$l(\mathcal{D}) = \deg(\mathcal{D}) - (g - 1)$$

Proof: We have $\deg(W - \mathcal{D}) = \deg(W) - \deg(\mathcal{D})$. Now by [Corollary 1.4.2](#) $\deg(W - \mathcal{D}) < 0$. So $l(W - \mathcal{D}) = 0$ by [Theorem 1.2.1](#) part (ii). So We have $l(\mathcal{D}) = \deg(\mathcal{D}) - (g - 1)$. ■

1.5 Index of speciality

Definition 1.5.1 (Index of speciality). *Let \mathcal{D} be a divisor on a curve \mathcal{X} . We define*

$$\Omega(\mathcal{D}) = \{\omega \in \Omega(\mathcal{X}) \mid (\omega) - \mathcal{D} \succcurlyeq 0\}$$

and we denote the dimension of $\Omega(\mathcal{D})$ over \mathbb{F} by $\delta(\mathcal{D})$ called the index of speciality of \mathcal{D} .

Theorem 1.5.1. $\delta(\mathcal{D}) = l(W - \mathcal{D})$

Proof: If $W = (\omega)$. Define the linear map $\varphi : \mathcal{L}(W - \mathcal{D}) \rightarrow \Omega(\mathcal{D})$ by $\varphi(f) = f\omega$.

$$f \in \mathcal{L}(W - \mathcal{D}) \implies (f) + W - \mathcal{D} \succcurlyeq 0 \iff (f) + (\omega) - \mathcal{D} \succcurlyeq 0 \iff (f\omega) - \mathcal{D} \succcurlyeq 0 \iff f \in \Omega(\mathcal{D})$$

Hence φ is an isomorphism. Therefore $\delta(\mathcal{D}) = l(W - \mathcal{D})$ ■

Codes from Algebraic Curves

We have now come to define the Algebraic Geometric Codes.

2.1 Preliminaries

First we will define the system where we will define the codes.

- Our alphabet will be \mathbb{F}_q
- We will consider the functions $f \in \mathbb{F}_q[X_1, \dots, X_n]$. Sometimes we will write \overline{X} to denote (X_1, \dots, X_n) . n depends on the context
- If the affine curve \mathcal{X} over \mathbb{F}_q is defined by a prime ideal I in $\mathbb{F}_q[\overline{X}]$ then its coordinate ring $\mathbb{F}_q[\mathcal{X}] = \mathbb{F}_q[\overline{X}]/I$ and its function field $\mathbb{F}_q(\mathcal{X})$ is the quotient field of $\mathbb{F}_q[\mathcal{X}]$.
- It is always assumed that the curve is *absolutely irreducible*, i.e. the defining ideal is also prime in $\mathbb{F}[\overline{X}]$ where $\mathbb{F} := \overline{\mathbb{F}_q}$ i.e. \mathbb{F} is the algebraic closure of \mathbb{F}_q .

Similar adaptations are made for projective curves.

Observation. For any $F \in \mathbb{F}_q[\overline{X}]$, $F(x_1, \dots, x_n)^q = F(x_1^q, \dots, x_n^q)$. So if (x_1, \dots, x_n) is a zero of F and F is defined over \mathbb{F}_q then (x_1^q, \dots, x_n^q) is also a zero of F .

We can extend the *Frobenius Map*, $Fr : x \mapsto x^q$ coordinate-wise to points in affine and projective space by $Fr(x_1, \dots, x_n) = (x_1^q, \dots, x_n^q)$. If \mathcal{X} is a curve defined over \mathbb{F}_q and P is a point of \mathcal{X} , then $Fr(P)$ is also a point of \mathcal{X} .

Definition 2.1.1 (Rational Divisor). A divisor \mathcal{D} on \mathcal{X} is called *rational* if the coefficients of P and $Fr(P)$ in \mathcal{D} are the same for any point P of \mathcal{X} .

Remark: Now on the space $\mathcal{L}(\mathcal{D})$ will only be considered for rational divisors and as before but with the restriction of the rational functions to $\mathbb{F}_q(\mathcal{X})$

Let \mathcal{W} be an absolutely irreducible nonsingular projective curve over \mathbb{F}_q . We will define two kinds of algebraic geometry codes from \mathcal{X} , *Geometric Reed Solomon Codes* and *Geometric Goppa Codes*. Let P_1, \dots, P_n are rational

points on \mathcal{X} and \mathcal{D} be the divisor $\mathcal{D} = P_1 + \cdots + P_n$. Furthermore \mathcal{G} is some other divisor that has support disjoint from \mathcal{D} . **Remark:** We will make more restrictions on \mathcal{G} , $\deg(\mathcal{G}) > 2g - 2$

2.2 Geometric Reed Solomon Codes

With the setting as above we define

Definition 2.2.1 (Geometric Reed Solomon Codes). *The linear code $C(\mathcal{D}, \mathcal{G})$ of length n over \mathbb{F}_q is the image of the linear map $\alpha : \mathcal{L}(\mathcal{G}) \rightarrow \mathbb{F}_q^n$ defined by $\alpha(f) = (f(P_1), \dots, f(P_n))$*

Theorem 2.2.1. *The code $C(\mathcal{D}, \mathcal{G})$ has dimension $k = \deg(\mathcal{G}) - (g - 1)$ and distance $d \geq n - \deg(\mathcal{G})$*

Corollary 2.2.2. $k + d \geq n - (g - 1)$

Proof: $k + n \geq \deg(\mathcal{G}) - (g - 1) + n - \deg(\mathcal{G}) = n - (g - 1)$ ■

CHAPTER 3

Bibliography
