

# **Evaluate reduced versions of BLAKE, Grøstl and Keccak, by near collisions**

by

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# Dedication

Dedicated to my parents.

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# Abstract

## Evaluate reduced versions of BLAKE, Grøstl and Keccak, by near collisions

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Digital signatures used in computer security, are created and verified using cryptographic hash functions [12]. Standard Hashing Algorithm(SHA) is the nomenclature for hashing functions that have been standardized by National Institute of Standards and Technology(NIST).

Due to advances in cryptanalysis of SHA-2, NIST announced a competition in November, 2007 to choose SHA-3. All the 64 submissions to NIST were open to public, and were cryptanalysed by the cryptographic community. In October, 2012 Keccak was selected to be the basis for SHA-3. Keccak was chosen for its robustness and high security margin.

Of the 64 submissions to SHA-3 competition, 56 were selected for the first round, and 5 made it final round. In this project I have compared the resistance to near collisions, on reduced versions of Keccak and two other finalist BLAKE, and Grøstl. A  $\epsilon/n$  bit near collision for hash function  $h$  and two messages  $M_1$  and  $M_2$ , where  $M_1 \neq M_2$  is defined as

$$HW(h(M_1, CV) \oplus h(M_2, CV)) = n - \epsilon$$

where HW is the Hamming weight, and CV is the chaining value, and  $n$  is the hash size in bits. I have also compared the near collision resistance amongst different versions of

Keccak based on reduced internal state size.

Hill climbing, simulated annealing, tabu search and random selection are applied to find a neighbourhood of chaining value that is appended to 2 messages that dissimilar by a bit to find near collisions in the reduced versions of mentioned hashing algorithms.

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# Chapter 1

## Introduction

### 1.1 Cryptographic Hash Functions

A cryptographic hash function, is a function that accepts an arbitrarily long input string, and outputs a fixed length string that is uniquely related to the input string. The output string is either called message digest, or hash value of the given input string. Hash functions should be one way, that is given the message digest, it should not be possible to find the input string. Two input strings even differing in a single bit should output two different hash values, that are not close or display any co-relation to each other.

### 1.2 The need for cryptographic hash function

Digital signatures are method of authenticating electronic documents, and hash functions are used to ease the process of digitally signing a document. Digital signatures based on asymmetric algorithm like RSA, have a input size limitation of around 128 to 324 bytes [19]. However most documents in practice are longer than 324 bytes.

One approach would be to divide the message into blocks of size acceptable by that of the signing algorithm, and sign each block separately. However, there are disadvantages to this are -

1. *Computationally intensive:* Modular exponentiation of large integers used in asymmetric algorithms are resource intensive. For signing, multiple blocks of message,

the resource utilization is pronounced. Additionally, not only the sender but the receiver will also have to do the same resource intensive operations.

2. *Overheads:* The signature is of the same length as the message. This increases the overheads in storage and transmission.
3. *Security concerns:* An attacker could remove, or reorder, or reconstruct new message and signatures from the previous message and signature pairs. Though attacker, cannot manipulate the individual blocks, but safety of the entire message is compromised.

Thus to eliminate the overheads, and security limitations; a method is required to uniquely generate fixed size finger print of arbitrarily large message blocks. Hash functions, fill this void of signing large messages.

## 1.3 Standards and NIST Competition

### 1.3.1 Secure Hashing Algorithm SHA-0 and SHA-1

In 1993 National Security Agency(NSA) proposed SHA-0 as standard hashing algorithm, which was later standardised by NIST. However by 1995, Florent Chabaud and Antoine Joux, found collisions in SHA-0 with complexity of  $2^{61}$ . In 2004, Eli Biham and Chen found near collisions for SHA-0, about 142 out of 160 bits to be equal. Full collisions were also found, when the number of rounds for the algorithm were reduced from 80 to 62.

In 1995 SHA-0 was replaced by SHA-1 designed by NSA [10, 13]. It has block size of 512 bits and output of 160 bits, which is similar to that of SHA-0. SHA-1 has an additional circular shift operation, that rectifies the weakness in SHA-0.

In 2005 a team from Shandong University in China consisting of Xiaoyun Wang, Yiqun Lisa Yin, and Hongbo Yu, found a way to find collisions on full version of SHA-1 requiring  $2^{69}$  operations. This is less than the number of operations required by brute force search



for finding collisions, which is  $2^{80}$  [24]. An ideal hash function requires the number of operations to find a collision to be equal to a brute force search, to be considered secure.

In October 2012, Jesse Walker from Skein team estimated the computational cost for finding collisions in SHA-1 to be \$ 2.77 million, based on HashClash developed by Marc Stevens [25].

Algorithm	Message Size	Block Size	Word Size	Hash Value Size
SHA-1	$<2^{64}$ bits	512 bits	32 bits	160 bits
SHA-224	$<2^{64}$ bits	512 bits	32 bits	224 bits
SHA-256	$<2^{64}$ bits	512 bits	32 bits	256 bits
SHA-384	$<2^{128}$ bits	1024 bits	64 bits	384 bits
SHA-512	$<2^{128}$ bits	1024 bits	64 bits	512 bits

Table 1.1: Secure Hash Algorithms as specified in FIPS 180-2

### 1.3.2 SHA-2

SHA-2 was designed by NSA, and released in 2001 by NIST. It is a family of hash functions consisting of SHA-224, SHA-256, SHA-384, SHA-512. Table 1.1 gives a brief overview of specifications of SHA-1 and family of SHA-2 hash functions. The number suffix after the SHA acronym, indicates the bit length, of the output of that hash function. Although SHA-2 family of algorithms were influenced by SHA-1 design, but the attacks on SHA-1 have not been successfully extended completely to SHA-2.

Collisions were found on SHA-256 and SHA-512, that were reduced to 22 permutation rounds. Computational operations to find collisions in 23 and 24 reduced rounds of SHA-256 were  $2^{11.5}$  and  $2^{28.5}$  corresponding calls to their respectively reduced functions. For SHA-512 reduced versions of 23 and 24 rounds, the corresponding values for were  $2^{16.5}$  and  $2^{32.5}$  calls [23]. Since, SHA-2 family relies on the Merkle-Damgård construction, the whole process of creation of hash can be considered as repeated application of certain operations generally called as compression function, on the input cumulatively. The rounds here refer to the number of times the permutations applied to the input during its compression.

Preimage attacks have been a success, on 41 reduced rounds of SHA-256 and 46 reduced rounds of SHA-512. As per specifications, SHA-256 contains 64 rounds, while SHA-512 consists of 80 rounds [2]. It can be said that SHA-2 functions are partially susceptible to preimage attacks.

### 1.3.3 NIST competition and SHA-3

In November 2, 2007 NIST through a Federal Register Notice announced a public competition for selection of new SHA-3 algorithm. This was done in response to advances made in cryptanalysis of SHA-2. Submission requirements stated to provide a cover sheet, algorithm specifications and supporting documentation, optimized implementations as per specifications of NIST, and intellectual property statements.

Submissions for the competition were accepted till October 31, 2008, and 51 candidates from 64 submissions for first round of competition were announced on December 9, 2008. On October 2, 2012 NIST announced the winner of the competition to be Keccak, amongst the other four finalist, which were BLAKE, Grøstl, JH and Skein. Keccak was chosen for its' large security margin, efficient hardware implementation, and flexibility.

## 1.4 Our work

In this project we compare the reduced versions of three hashing algorithms, that made it to final round in SHA-3. They are BLAKE, Grøstl, and Keccak. The three algorithms are first compared on the basis of reduced number of permutation rounds, that are kept equal for all three of them. After that we compare variants of Keccak hashing algorithm, that have their internal state reduced by different amounts. For this part also we reduce the number of permutation rounds.

The comparison is done on the basis of near collisions found in the reduced version of the algorithms. A  $\epsilon/n$  bit near collision for hash function  $h$  and two messages  $M_1$  and  $M_2$ , where  $M_1 \neq M_2$  is defined as

$$HW(h(M_1, CV) \oplus h(M_2, CV)) = n - \epsilon$$

where HW is the Hamming weight, and CV is the chaining value, and n is the hash size in bits. 60 different message pairs are made from a seed string "A quick brown fox jumps over the lazy dog", and 60 other strings that differ from seed string in one bit. The output of the message pairs is compared for a near collision, and improvements on it are made by iterating it over search algorithms that are: hill climbing, simulated annealing, tabu search and random selection.

#### HYPOTHESIS

- Reduced state Keccak, has better resistance to near collisions than BLAKE and Grøstl. For the attack algorithms hill climbing, simulated annealing, tabu search and random selection.
- Simulated annealing and tabu search, are better at finding near collisions compared to hill climbing and random selection.
- State size has no effect on efficiency of Keccak permutation rounds, for reduced number of rounds.

Through our experiments as formulated through our hypothesis, we aim to determine a search algorithm, that would be optimum in finding near collisions in reduced versions of hashing algorithm. In our comparison of the three SHA-3 finalist algorithm, we try to evaluate the winner Keccak and its competitors efficiency, when the number of permutation rounds are reduced. Lastly we investigate if reduced internal state size of Keccak, in any way reduces effectiveness of permutation rounds, if the permutation rounds are reduced.

# Chapter 2

## Background

### 2.1 Hashing

A cryptographic hash function, is a function that intakes a string of arbitrary, and outputs a string of fixed length, that is unique to the input string given. The aforementioned is description of a single fixed hash function. But, hash functions can be tweaked with an extra key parameter. This gives rise multiple hash functions or *hash family* as defined below [26]

A *hash family* is a four-tuple  $(\mathcal{X}, \mathcal{Y}, \mathcal{K}, \mathcal{H})$ , satisfying the following conditions.

- $\mathcal{X}$  is a set of possible messages
- $\mathcal{Y}$  is a finite set of hash function output
- $\mathcal{K}$ , the *keyspace*, is a finite set of possible keys
- For each  $K \in \mathcal{K}$ , there is a hash function  $h_K \in \mathcal{H}$ . Each  $h_K : \mathcal{X} \rightarrow \mathcal{Y}$

In the above definition,  $\mathcal{X}$  could be finite or infinite set, but  $\mathcal{Y}$  is always a finite set, since the length of bit string or hash function output, that defines  $\mathcal{Y}$  is finite. A pair  $(x, y) \in \mathcal{X} \times \mathcal{Y}$  is a *valid pair* under key  $K$ , if  $h_K(x) = y$ .

A hash family  $\mathcal{F} \subseteq \mathcal{F}^{\mathcal{X}\mathcal{Y}}$  is called as  $(N, M)$  - hash family, where  $\mathcal{F}^{\mathcal{X}\mathcal{Y}}$  denotes set of all functions that map from domain  $\mathcal{X}$  to co-domain  $\mathcal{Y}$ ; if  $|\mathcal{X}| = N$  and  $|\mathcal{Y}| = M$ , and  $|\mathcal{F}^{\mathcal{X}\mathcal{Y}}| = M^N$ .

An *unkeyed hash function* is a function  $h_k : \mathcal{X} \rightarrow \mathcal{Y}$ , where  $\mathcal{X}$  and  $\mathcal{Y}$  are as defined previously, and where  $|\mathcal{K}| = 1$ . In our work we have worked only with unkeyed hash family of fixed hash function only. Although the algorithms we deal with can have their keys varied, but we keep them as constant.

The output of a hash function is generally called as a message digest. Since, it can viewed as a unique snapshot of the message, that cannot be replicated if the bits in message are tampered with.

### 2.1.1 Properties of an ideal hash function

An ideal hash function should be easy to evaluate in practice. However, it should satisfy the following three properties primarily, for a hash function to be considered *secure* [26].

#### 1. Preimage resistance

##### PREIMAGE

**Given:** A hash function  $h : \mathcal{X} \rightarrow \mathcal{Y}$  and an element  $y \in \mathcal{Y}$ .

**Find:**  $x \in \mathcal{X}$  such that  $h(x) = y$ .

The problem preimage suggests that can we find an input  $x \in \mathcal{X}$ , given we have the hash output  $y$ , such that  $h(x) = y$ . If the preimage problem for a hash function cannot be efficiently solved, then it is preimage resistant. That is the hash function is one way, or rather it is difficult to find the input, given the output alone.

#### 2. Second preimage resistance

##### SECOND PREIMAGE

**Given:** A hash function  $h : \mathcal{X} \rightarrow \mathcal{Y}$  and an element  $x \in \mathcal{X}$ .

**Find:**  $x' \in \mathcal{X}$  such that  $x' \neq x$  and  $h(x) = h(x')$ .

Second preimage problem suggests that given an input  $x$ , can another input  $x'$  be found, such that  $x \neq x'$  and hash output of both the inputs are same, that is  $h(x) = h(x')$ . A hash function for which a different input given another input, that compute to same hash cannot be found easily, is called as having second preimage resistance.

### 3. Collision resistance

#### COLLISION

**Given:** A hash function  $h : \mathcal{X} \rightarrow \mathcal{Y}$

**Find:**  $x, x' \in \mathcal{X}$  such that  $x' \neq x$  and  $h(x') = h(x)$ .

Collision problem states that, can two different input strings be found, such that they hash to the same value given the same hash function. If the collision problem for the hash function, is computationally complex, then the hash function is said to be collision resistant.

Basically, the above properties make sure that hash function has one to one mapping from input to output, and is one way. That is if a two different input strings with even minute differences should map to two different hash values. And it should be practically infeasible, to find a input given a hash value.

## 2.2 Security Model

A generic model of security fulfillment, for any hash function can be set based on previously mentioned properties for it to be considered as secure. Two such ideas, are described as follows

### 2.2.1 Random Oracle

Hash functions being built on mathematical operations, cannot be truly random, but are efficient approximations of fixed random output mapping to an input. An ideal hash function can be abstracted as a random oracle, to show that algorithm is secure modulo the way it creates random outputs. Proofs can be formalized on this basis [7].

Random oracle model, proposed by Bellare and Rogaway, is a mathematical model of ideal hash function. It can be interpreted as, the only way to know the hash value for an input  $x$  would be to ask the Oracle or rather compute the hash of the input itself. There is no way of formulating or guessing the hash value for input, even if you are provided with

substantial number of input and output pairs. It is analogous to looking up for corresponding value of the key in a large table. To know the value for an input, you look into the table. A well designed hash function, mimics the behaviour of random oracle as closely as possible.

### 2.2.2 Birthday Paradox

If we randomly choose 23 people, then the probability that two people from the group will have identical birthday is around 0.507. This is because, the first person can be paired with rest of 22 people in group, to form 22 pairs. The next person in group can be paired with remaining 21 people to get 21 pairs. Thus we end up with  $22 + 21 + 20 + \dots + 1 = 253$  pairs. Thus the probability is ratio of pairs 253 to the sample space 365 days in a year (ignoring the leap year).

Two people with same birthday can be seen analogous to two inputs hashing to the same value, that is collision. Say the sample space of hash as  $M$ , and denote the number of samples to be taken as  $N$ . Then by birthday problem described above, the minimum number of people required ( $N$ ) to have the same birthday within a year ( $M = 365$ ) with probability 0.5, would be  $N = 23$ .

It can be formally proved for any sample size  $M$ , to find two values that are identical with probability 0.5 can be given by the equation  $N \approx 1.17\sqrt{M}$ . This can be interpreted as hashing over  $\sqrt{M}$  values roughly will give us two entries, with probability of collision being 0.5.

The above theorem can be generalised as, that if we were to brute force our way to finding collisions in a hash function that has a message digest length of  $2^{128}$  bits. Then at minimum we would need to calculate  $2^{64}$  instances of hash, to find a collision with a probability of 0.5. Any good hash function in practice should be resistant to attacks, that require operations less than estimated by the birthday paradox on message digest size.

## 2.3 Applications

Applications of cryptographic hash functions, can be broadly classified in areas of verification, data integrity and pseudo random generator functions.

### 2.3.1 Verification and data integrity

1. Digital Forensics: When digital data is seized and to be used as evidence, a hash of the original digital media is taken. A copy of the digital evidence is made under the regulations, and the hash of the copied digital media is made, before it can be examined. After the evidence has been examined, then another hash value of the copy of the evidence that was used in examination is made. This ensures, that evidence has not been tampered [22].
2. Password verification: Hash value of salted passwords are stored in database. When a password is to be verified, a hash value of the given salted password is evaluated. This hash value is compared with the value stored in the database, and password is authenticated, if it matches the one in database.
3. Integrity of files: Hash values can be used to check, that data files have not been modified over time. Hash value of a file or files is taken regularly, and compared to previous stored hash value, when it is again examined. If the values do not match, it means the file or files have been compromised.

### 2.3.2 Pseudo random generator function

Cryptographic hash functions can be used as pseudo random bit generators. The hash function is initialised with a random seed, and then hash function is queried iteratively to get a sequence of bits, which look random. Since, the cryptographic hash algorithm is a mathematical function, so the sequence of two pseudo random bits would be similar if they come from same hash function with the same key.



## Chapter 3

# BLAKE, Grøstl, Keccak and SHA-3

For our project we chose to study BLAKE, Grøstl and Keccak, out of the 5 finalists. All the five finalist were found to be secure, devoid of any weakness that can be practically exploited. Keccak was chosen, since it was the winner of SHA-3 competition, and is the basis for the current SHA-3 in draft, thus required for project to be relevant. Keccak has sponge as design basis, BLAKE is based on HAIFA, and Grøstl is based on Merkle-Damgård design. The selection of these three, gives an opportunity to observe the effects of different design for reduced hash functions. Thus I chose to go with BLAKE and Grøstl for comparison with Keccak in this project. Rebound attacks in JH have been studied by Naveen Kandakumar [14], while Darryl Clyde Eychner did a statistical analysis on BLAKE and Skein [9], which were the other SHA-3 finalist.

### 3.1 BLAKE

BLAKE [3] hash function is built on HAIFA (HAsH Iterative FrAmework) structure [6] which is an improved version of Merkle-Damgård design. It provides resistance to long-message, second pre-image attack and provides a salting option, that BLAKE uses [16]. The design is local wide-pipe which avoids internal collisions. The compression function in BLAKE is tweaked version of ChaCha, a stream cipher.

BLAKE has 4 variations of the algorithm shown in table 3.1. Figure 3.1 shows how the individual message blocks are consumed. The construction takes in 4 inputs, a message; salt, that makes hash function parameter specific; a counter, which is count of all the bits

Algorithm	Word	Message	Block	Digest	Salt
BLAKE-224	32	$< 2^{64}$	512	224	128
BLAKE-256	32	$< 2^{64}$	512	256	128
BLAKE-384	64	$< 2^{128}$	1024	384	256
BLAKE-512	64	$< 2^{128}$	1024	512	256

Table 3.1: Specification of available input, output, block and salt sizes for various BLAKE hash functions [3].

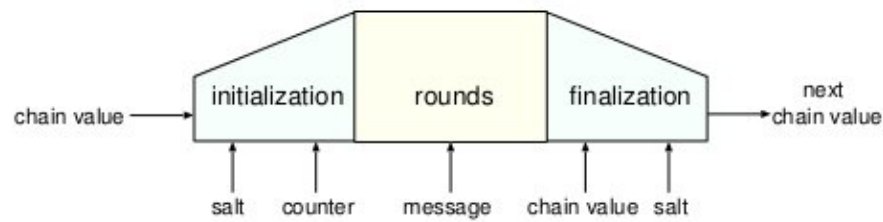


Figure 3.1: Local wide construction of BLAKE's compression function [3]

hashed till then; and lastly a chaining value which is input of the previous operation or initial value in case of hash initiation. The compression function is composed of a  $4 \times 4$  matrix of words, where a word is equal to 32 bits for BLAKE-256 variant, while 64 bit for variant BLAKE-512.

Symbol	Meaning
$\leftarrow$	variable assignment
$+$	addition modulo $2^{32}$ or (modulo $2^{64}$ )
$\ggg k$	rotate $k$ bits to least significant bits
$\lll k$	rotate $k$ bits to most significant bits
$\langle l \rangle_k$	encoding of integer $l$ over $k$ bits

Table 3.2: Convention of symbols used in BLAKE algorithm

### 3.1.1 BLAKE-256

The compression function takes following as input

- a chaining value of  $h = h_0, \dots, h_7$
- a message block  $m = m_0, \dots, m_{15}$
- a salt  $s = s_0, \dots, s_3$
- a counter  $t = t_0, t_1$

These four inputs of 30 words or 120 bytes, are processed as  $h' = \text{compress}(h, m, s, t)$  to provide a new chain value of 8 words.

#### Compression function

- **Constants**

$c_0 = 243F6A88$	$c_1 = 85A308D3$	$c_2 = 13198A2E$	$c_3 = 03707344$
$c_4 = A4093822$	$c_5 = 299F31D0$	$c_6 = 082EFA98$	$c_7 = EC4E6C89$
$c_8 = 452821E6$	$c_9 = 38D01377$	$c_{10} = BE5466CF$	$c_{11} = 34E90C6C$
$c_{12} = C0AC29B7$	$c_{13} = C97C50DD$	$c_{14} = B5470917$	$c_{15} = 3F84D5B5$

Table 3.3: 16 constants used for BLAKE-256 [3]

- **Initialization:** The constants mentioned are used with the salts, and counter along with initial value used as chaining input, to create a initial matrix of  $4 \times 4$ , 16 word state.

$$\begin{pmatrix} v_0 & v_1 & v_2 & v_3 \\ v_4 & v_5 & v_6 & v_7 \\ v_8 & v_9 & v_{10} & v_{11} \\ v_{12} & v_{13} & v_{14} & v_{15} \end{pmatrix} \leftarrow \begin{pmatrix} h_0 & h_1 & h_2 & h_3 \\ h_4 & h_5 & h_6 & h_7 \\ s_0 \oplus c_0 & s_1 \oplus c_1 & s_2 \oplus c_2 & s_3 \oplus c_3 \\ t_0 \oplus c_4 & t_0 \oplus c_5 & t_1 \oplus c_6 & t_1 \oplus c_7 \end{pmatrix}$$

$\sigma_0$	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
$\sigma_1$	14	10	4	8	9	15	13	6	1	12	0	2	11	7	5	3
$\sigma_2$	11	8	12	0	5	2	15	13	10	14	3	6	7	1	9	4
$\sigma_3$	7	9	3	1	13	12	11	14	2	6	5	10	4	0	15	8
$\sigma_4$	9	0	5	7	2	4	10	15	14	1	11	12	6	8	3	13
$\sigma_5$	2	12	6	10	0	11	8	3	4	13	7	5	15	14	1	9
$\sigma_6$	12	5	1	15	14	13	4	10	0	7	6	3	9	2	8	11
$\sigma_7$	13	11	7	14	12	1	3	9	5	0	15	4	8	6	2	10
$\sigma_8$	6	15	14	9	11	3	0	8	12	2	13	7	1	4	10	5
$\sigma_9$	10	2	8	4	7	6	1	5	15	11	9	14	3	12	13	0

Table 3.4: Round permutations to be used [3]

$$\begin{aligned}
IV_0 &= 6A09E667 & IV_1 &= BB67AE85 & IV_2 &= 3C6EF372 & IV_3 &= A54FF53A \\
IV_4 &= 510E527F & IV_5 &= 9B05688C & IV_6 &= 1F83D9AB & IV_7 &= 5BE0CD19
\end{aligned}$$

Table 3.5: Initial values which become the chaining value for the first message block [3]

- **Round function:** After initialisation, the state is subjected to column and diagonal operations, 14 times. A round operation G acts as per following

$$\begin{aligned}
&G_0(v_0, v_8, v_{12}) & G_1(v_1, v_5, v_9, v_{13}) & G_2(v_2, v_6, v_{10}, v_{14}) & G_3(v_3, v_7, v_{11}, v_{15}) \\
&G_4(v_0, v_5, v_{10}, v_{15}) & G_5(v_1, v_6, v_{11}, v_{12}) & G_6(v_2, v_7, v_8, v_{13}) & G_7(v_3, v_4, v_9, v_{14})
\end{aligned}$$

where the round function  $G_i(a, b, c, d)$  sets

$$a \leftarrow a + b + (m_{\sigma_r(2i)} \oplus c_{\sigma_r(2i+1)})$$

$$d \leftarrow (d \oplus a) \ggg 16$$

$$c \leftarrow c + d$$

$$b \leftarrow (b \oplus c) \ggg 12$$

$$a \leftarrow a + b + (m_{\sigma_r(2i+1)} \oplus c_{\sigma_r(2i)})$$

$$d \leftarrow (d \oplus a) \ggg 8$$

$$c \leftarrow c + d$$

$$b \leftarrow (b \oplus c) \ggg 7$$

The implementation of the G function is shown below.

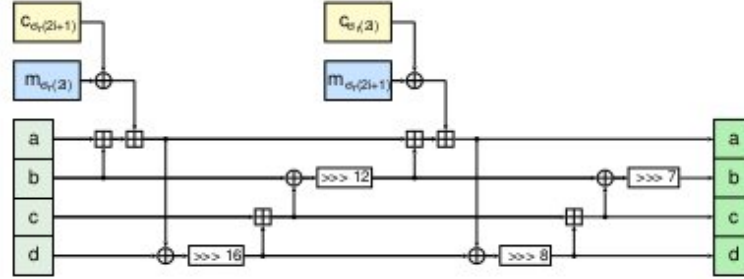


Figure 3.2: The  $G_i$  function in BLAKE [3]

- **Finalization:** The chaining values for the next stage are obtained by XOR of the words from the state matrix, the salt and the initial value.

$$h'_0 \leftarrow h_0 \oplus s_0 \oplus v_0 \oplus v_8$$

$$h'_1 \leftarrow h_1 \oplus s_1 \oplus v_1 \oplus v_9$$

$$h'_2 \leftarrow h_2 \oplus s_2 \oplus v_2 \oplus v_{10}$$

$$h'_3 \leftarrow h_3 \oplus s_3 \oplus v_3 \oplus v_{11}$$

$$h'_4 \leftarrow h_4 \oplus s_0 \oplus v_4 \oplus v_{12}$$

$$h'_5 \leftarrow h_5 \oplus s_1 \oplus v_5 \oplus v_{13}$$

$$h'_6 \leftarrow h_6 \oplus s_2 \oplus v_6 \oplus v_{14}$$

$$h'_7 \leftarrow h_7 \oplus s_3 \oplus v_7 \oplus v_{15}$$

### Hashing the message

A given input message is padded with a bit '1' followed followed by at most 511 bits of zeros, so that the message size is equal to 447 modulo 512. This padding is followed by a bit '1' and a 64-bit unsigned big-endian representation of block length  $l$ . The padding to a message, can be represented as  $m \leftarrow m \parallel 1000 \dots 0001 \langle l \rangle_{64}$

---

**Algorithm 3.1** BLAKE Compression procedure [3]

---

```

1:  $h^0 \leftarrow IV$ 
2: for  $i = 0, \dots, N - 1$  do
3:    $h^{i+1} \leftarrow \text{compress}(h^i, m^i, s, l^i)$ 
4: end for
5: return  $h^N$ 

```

---

$IV_0 = 6A09E667F3BCC908$     $IV_1 = BB67AE8584CAA73B$     $IV_2 = 3C6EF372FE94F82B$   
 $IV_3 = A54FF53A5F1D36F1$     $IV_4 = 510E527FADE682D1$     $IV_5 = 9B05688C2B3E6C1F$   
 $IV_6 = 1F83D9ABFB41BD6B$     $IV_7 = 5BE0CD19137E2179$

Table 3.6: Initial values used for BLAKE-512 [3]

As shown in algorithm 3.1, the BLAKE compression function ingests the padded message block by block, in a loop starting from the initial value, and then sends the last chained value obtained from the finalization to the  $\Omega$  truncation function, to obtain the hash value.

### 3.1.2 BLAKE-512

BLAKE-512 operates on 64-bit words and returns a 64-byte hash value. The chaining value is 512 bit long, message blocks are 1024 bits, salt is 256 bits, and counter size is 128 bits. The difference from BLAKE-256 are in initial values and constants (tables 3.6 and 3.7) used. Compression function and padding is also tweaked to better suit 64 bit word.

Compression function in BLAKE-512 gets 16 iterations instead of 14 as in BLAKE-256, the rotations are updated and word size increased from 32 bits to 64 bits.  $G_i(a, b, c, d)$  is given as

```

 $a \leftarrow a + b + (m_{\sigma_r(2i)} \oplus c_{\sigma_r(2i+1)})$ 
 $d \leftarrow (d \oplus a) \ggg 32$ 
 $c \leftarrow c + d$ 
 $b \leftarrow (b \oplus c) \ggg 25$ 
 $a \leftarrow a + b + (m_{\sigma_r(2i+1)} \oplus c_{\sigma_r(2i)})$ 

```

$c_0 = 243F6A8885A308D3$	$c_1 = 13198A2E03707344$	$c_2 = A4093822299F31D0$
$c_3 = 082EFA98EC4E6C89$	$c_4 = 452821E638D01377$	$c_5 = BE5466CF34E90C6C$
$c_6 = C0AC29B7C97C50DD$	$c_7 = 3F84D5B5B5470917$	$c_8 = 9216D5D98979FB1B$
$c_9 = D1310BA698DFB5AC$	$c_{10} = 2FFD72DBD01ADFB7$	$c_{11} = B8E1AFED6A267E96$
$c_{12} = BA7C9045F12C7F99$	$c_{13} = 24A19947B3916CF7$	$c_{14} = 0801F2E2858EFC16$
$c_{15} = 636920D871574E69$		

Table 3.7: 16 constants used for BLAKE-512 [3]

$$d \leftarrow (d \oplus a) \ggg 16$$

$$c \leftarrow c + d$$

$$b \leftarrow (b \oplus c) \ggg 11$$

Once more than 9 rounds are done, then permutations to be applied from table are selected by modulo 10 of round. For example if round  $r > 9$  then permutation used is  $\sigma_{r \bmod 10}$ , say  $r = 15$  then permutation would be  $\sigma_{15 \bmod 10} = \sigma_5$ .

For the padding, the message is first padded with bit 1 and then as many zeros required to make the bit length equivalent to 895 modulo 1024. After that another bit of value 1 is appended followed by 128-bits unsigned big-endian representation of message length as  $m \leftarrow m \parallel 100 \dots 001 \langle l \rangle_{128}$ .

### 3.1.3 BLAKE-224 and BLAKE-384

#### BLAKE-224

BLAKE-224 is similar to BLAKE-256, but differs slightly. It has different initial values, different padding and the output bits are truncated to first 224 bits. The padding differs

$$\begin{array}{llll}
 IV_0 = C1059ED8 & IV_1 = 367CD507 & IV_2 = 3070DD17 & IV_3 = F70E5939 \\
 IV_4 = FFC00B31 & IV_5 = 68581511 & IV_6 = 64F98FA7 & IV_7 = BEFA4FA4
 \end{array}$$

Table 3.8: Initial values for BLAKE-224 which are taken from SHA-224 [3]

from BLAKE-256 in way that the bit preceding the message length is replaced by a 0 bit. Which is represented as  $m \leftarrow m \parallel 100 \dots 000 \langle l \rangle_{64}$ .

### BLAKE-384

In BLAKE-384 the output of BLAKE-512 is truncated to 384 bits. The padding differs from BLAKE-512, in way that bit preceding the length encoding is 0 and not 1. It can be shown as  $m \leftarrow m \parallel 100 \dots 000 \langle l \rangle_{128}$ . The initial chaining values are given in table 3.9.

$$\begin{aligned} IV_0 &= \text{CBBB9D5DC1059ED8} & IV_1 &= \text{629A292A367CD507} & IV_2 &= \text{9159015A3070DD17} \\ IV_3 &= \text{152FECDD8F70E5939} & IV_4 &= \text{67332667FFC00B31} & IV_5 &= \text{8EB44A8768581511} \\ IV_6 &= \text{DB0C2E0D64F98FA7} & IV_7 &= \text{47B5481DBEFA4FA4} \end{aligned}$$

Table 3.9: Initial values for BLAKE-384 [3]

## 3.2 Grøstl

Grøstl is collection of hash functions which produce digest size, ranging from 1 to 64 bytes. The variant of Grøstl that returns a message digest of size  $n$ , is called Grøstl- $n$ .

Grøstl is an iterated hash function, with two two compression functions named P and Q, based on wide trail design and having distinct permutations. Grøstl has a byte oriented SP network, and its diffusion layers and S-box are identical to AES. The design is a wide-pipe construction, where the internal state size is larger than output size, thus preventing most of the generic attacks. None of the permutations are indexed by a key, to prevent attacks from a weak key schedule [27].

### 3.2.1 The hash function construction

The input is padded and then split into  $l$ -bit message blocks  $m_1, \dots, m_t$ , and each message block is processed sequentially. The initial  $l$ -bit chaining value  $h_0 = iv$  is defined, and the



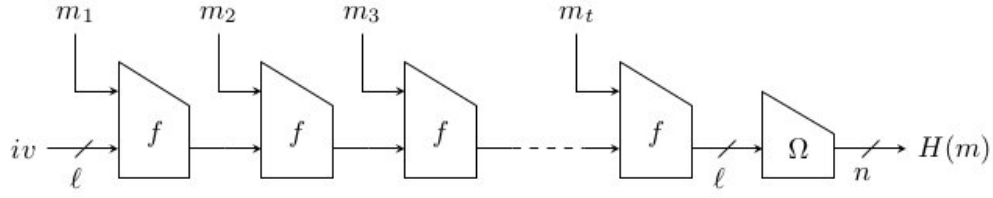


Figure 3.3: Grøstl hash function [27]

blocks  $m_i$  are processed as

$$h_i \leftarrow f(h_{i-1}, m_i) \text{ for } i = 1, \dots, t.$$

Thus  $f$  consumes two  $l$  – bits input, and maps to output of  $l$  – bits. For variants up to 256 bits output, size of  $l$  is 512 bits. And for digest sizes larger than 256 bits,  $l$  is 1024 bits.

After the last message block is processed, the last chaining value output is sent through a  $\Omega$  function, to get the hash output  $H(M)$ .

$$H(M) = \Omega(h_t),$$

Figure 3.3 displays the process. The permutation function  $f$ , is composed of two  $l$ -bit permutations called P and Q.

$$f(h, m) = P(h \oplus m) \oplus Q(m) \oplus h.$$

The  $\Omega$  function consists of a  $trunc_n(x)$  that outputs only the trailing  $n$  bits of input  $x$ .

$$\Omega(x) = trunc_n(P(x) \oplus x).$$

In order to fit the varying input length message to the block sizes of  $l$ , the message is padded to make it a multiple of block size  $l$ . First bit '1' is appended, then  $w = -N - 65 \bmod l$ , 0 bits are appended; where  $N$  is the length of the original message. Finally a 64 bit representation of  $(N + w + 65)/l$  is padded. Given the need for message length to be present in the padding, the maximum size of message digest in bits for Grøstl-512 version is  $2^{73} - 577$  bits, and that for 1024 version is  $2^{74} - 1089$ bits.

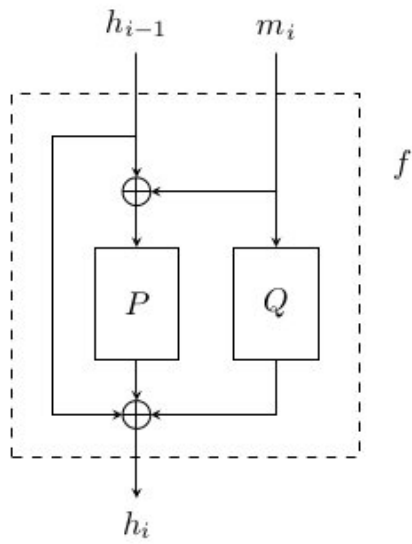


Figure 3.4: Compression functions, where  $P$  and  $Q$  are  $l - bit$  permutations [27]

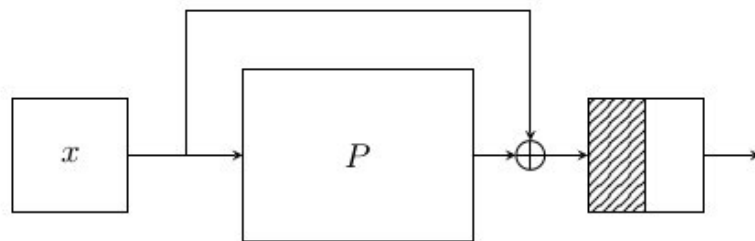


Figure 3.5: Omega truncation function [27]

Permutations	Digest size	Recommended value of r
$P_{512}$ and $Q_{512}$	8 - 256	10
$P_{1024}$ and $Q_{1024}$	264 - 512	14

Table 3.10: Recommended number of rounds [27]

n	$iv_n$
224	00 ... 00 00 e0
256	00 ... 00 01 00
384	00 ... 00 01 80
512	00 ... 00 02 00

Table 3.11: Initial values for Grøstl-n function. The numbers on left denote digest size in bits [27].

### 3.2.2 Design of P and Q permutations

There are two variations for P and Q permutations, one each for the digest size lower and higher than 256 bits. P and Q, make a round R, that has 4 transformations

$$R = MixBytes \cdot ShiftBytes \cdot SubBytes \cdot AddRoundConstant$$

The transformations *SubBytes* and *MixBytes* are same for all transformation while, *ShiftBytes* and *AddRoundConstant* differ for each of the transformations. The transformations operate on matrix of bytes, with the permutation of lower size digest having matrix of 8 rows and 8 columns, while that for larger variant is of 16 columns and 8 rows. The number of rounds for each R is given as recommendation in table 3.10 and the initial values are given in table 3.11

- **Mapping:** of a 64-byte sequence of 00 01 02 ... 3f to a  $8 \times 8$  matrix is shown in the following matrix. For a  $8 \times 16$  matrix, the mapping is extended the same way. Mapping the intermediate state values to byte sequence is reverse of this process.

$$\text{Input Mapping} = \begin{bmatrix} 00 & 08 & 10 & 18 & 20 & 28 & 30 & 38 \\ 01 & 09 & 11 & 19 & 21 & 29 & 31 & 39 \\ 02 & 0a & 12 & 1a & 22 & 2a & 32 & 3a \\ 03 & 0b & 13 & 1b & 23 & 2b & 33 & 3b \\ 04 & 0c & 14 & 1c & 24 & 2c & 34 & 3c \\ 05 & 0d & 15 & 1d & 25 & 2d & 35 & 3d \\ 06 & 0e & 16 & 1e & 26 & 2e & 36 & 3e \\ 07 & 0f & 17 & 1f & 27 & 2f & 37 & 3f \end{bmatrix}$$

- **AddRoundConstant:** XORs a round dependant constant to the state matrix say  $A$ . It is represented as  $A \leftarrow A \oplus C[i]$ , where  $C[i]$  is the round constant in round  $i$ . The constants for both P and Q for both variations are shown below.

$$P_{512} : C[i] = \begin{bmatrix} 00 \oplus i & 10 \oplus i & 20 \oplus i & 30 \oplus i & 40 \oplus i & 50 \oplus i & 60 \oplus i & 70 \oplus i \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \end{bmatrix}$$

and

$$Q_{512} : C[i] = \begin{bmatrix} ff & ff & ff & ff & ff & ff & ff & ff \\ ff & ff & ff & ff & ff & ff & ff & ff \\ ff & ff & ff & ff & ff & ff & ff & ff \\ ff & ff & ff & ff & ff & ff & ff & ff \\ ff & ff & ff & ff & ff & ff & ff & ff \\ ff & ff & ff & ff & ff & ff & ff & ff \\ ff & ff & ff & ff & ff & ff & ff & ff \\ ff \oplus i & ef \oplus i & df \oplus i & cf \oplus i & bf \oplus i & af \oplus i & 9f \oplus i & 8f \oplus i \end{bmatrix}$$

Similarly, the P and Q for the wider variants are written.

$$P_{1024} : C[i] = \begin{bmatrix} 00 \oplus i & 10 \oplus i & 20 \oplus i & 30 \oplus i & 40 \oplus i & 50 \oplus i & 60 \oplus i & 70 \oplus i \dots f0 \oplus i \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \dots 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \dots 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \dots 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \dots 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \dots 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \dots 00 \\ 00 & 00 & 00 & 00 & 00 & 00 & 00 & 00 \dots 00 \end{bmatrix}$$

and

$$Q_{1024} : C[i] = \begin{bmatrix} ff & ff & ff & ff & ff & ff & ff & ff \dots ff \\ ff & ff & ff & ff & ff & ff & ff & ff \dots ff \\ ff & ff & ff & ff & ff & ff & ff & ff \dots ff \\ ff & ff & ff & ff & ff & ff & ff & ff \dots ff \\ ff & ff & ff & ff & ff & ff & ff & ff \dots ff \\ ff & ff & ff & ff & ff & ff & ff & ff \dots ff \\ ff & ff & ff & ff & ff & ff & ff & ff \dots ff \\ ff \oplus i & ef \oplus i & df \oplus i & cf \oplus i & bf \oplus i & af \oplus i & 9f \oplus i & 8f \oplus i \dots 0f \oplus i \end{bmatrix}$$

where  $i$  is the round number represented as 8 bits value, and all other numbers are represented as hexadecimals.

- **SubBytes:** substitutes each byte in state by value from S-box shown in table 3.12. Say  $a_{i,j}$  a element in row  $i$  and column  $j$  of the state matrix, then the transformation done is  $a_{i,j} \leftarrow S(a_{i,j}), 0 \leq i < 8, 0 \leq j < v$ .
- **ShiftBytes:** cyclically shifts the bytes in a row to left by that number. Let list vector of a number denote the shift, with the index of the element indicating the row. The vector representation for  $P_{512} = [0, 1, 2, 3, 4, 5, 6, 7]$  and  $Q_{512} = [1, 3, 5, 7, 0, 2, 4, 6]$ . The shift is shown in figure 3.4. Those for the larger permutation are  $P_{1024} = [0, 1, 2, 3, 4, 5, 6, 11]$  and  $Q_{1024} = [1, 3, 5, 11, 0, 2, 4, 6]$ . This shifting is illustrated in figure 3.6 and figure 3.7.

	00	01	02	03	04	05	06	07	08	09	0a	0b	0c	0d	0e	0f
00	63	7c	77	7b	f2	6b	6f	c5	30	01	67	2b	fe	d7	ab	76
10	ca	82	c9	7d	fa	59	47	f0	ad	d4	a2	af	9c	a4	72	c0
20	b7	fd	93	26	36	3f	f7	cc	34	a5	e5	f1	71	d8	31	15
30	04	c7	23	c3	18	96	05	9a	07	12	80	e2	eb	27	b2	75
40	09	83	2c	1a	1b	6e	5a	a0	52	3b	d6	b3	29	e3	2f	84
50	53	d1	00	ed	20	fc	b1	5b	6a	cb	be	39	4a	4c	58	cf
60	d0	ef	aa	fb	43	4d	33	85	45	f9	02	7f	50	3c	9f	a8
70	51	a3	40	8f	92	9d	38	f5	bc	b6	da	21	10	ff	f3	d2
80	cd	0c	13	ec	5f	97	44	17	c4	a7	7e	3d	64	5d	19	73
90	60	81	4f	dc	22	2a	90	88	46	ee	b8	14	de	5e	0b	db
a0	e0	32	3a	0a	49	06	24	5c	c2	d3	ac	62	91	95	e4	79
b0	e7	c8	37	6d	8d	d5	4e	a9	6c	56	f4	ea	65	7a	ae	08
c0	ba	78	25	2e	1c	a6	b4	c6	e8	dd	74	1f	4b	bd	8b	8a
d0	70	3e	b5	66	48	03	f6	0e	61	35	57	b9	86	c1	1d	9e
e0	e1	f8	98	11	69	d9	8e	94	9b	1e	87	e9	ce	55	28	df
f0	8c	a1	89	0d	bf	e6	42	68	41	99	2d	0f	b0	54	bb	16

Table 3.12: Grøstl S-box. For an input  $x$ , you do a logical AND of  $x$  with  $f0$  and with  $0f$ . The first value obtained is used for column location and second for row location. The row and column location is used to identify the cell that will be used for substitution [27].

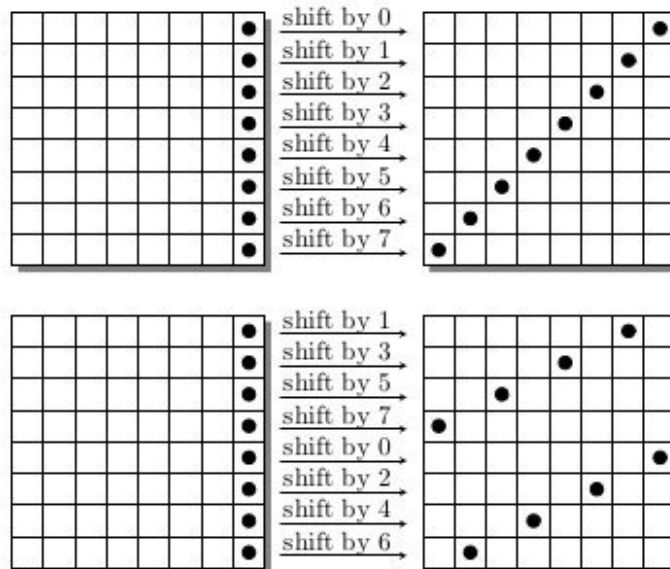


Figure 3.6: ShiftBytes transformation of permutation  $P_{512}$ (top) and  $Q_{512}$ (bottom) [27]

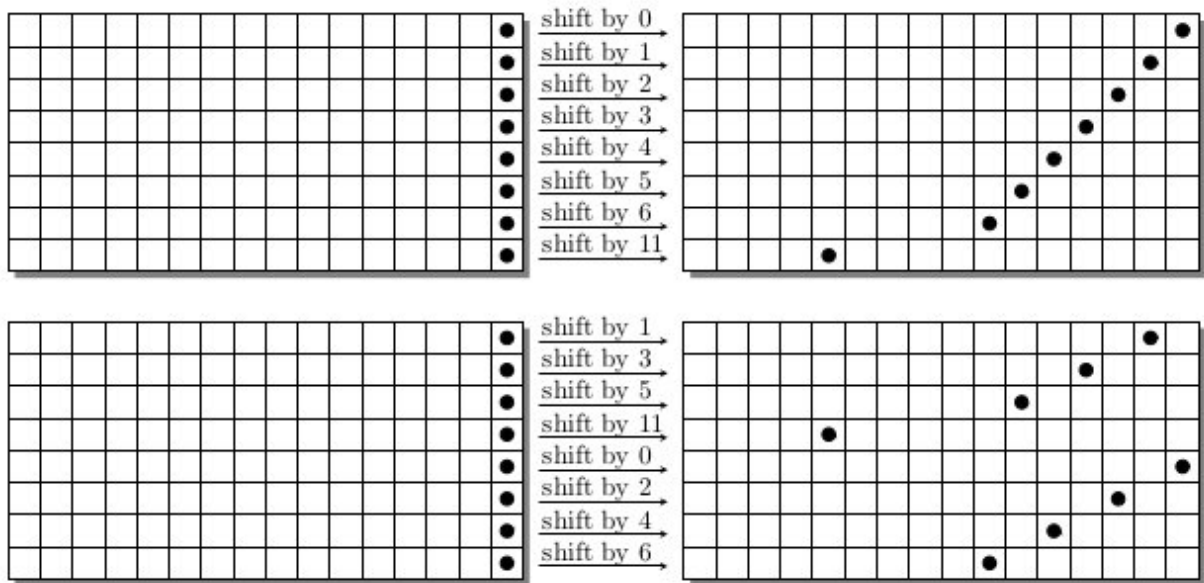


Figure 3.7: ShiftBytes transformation of permutation  $P_{1024}$ (top) and  $Q_{1024}$ (bottom) [27]

$$B = \begin{bmatrix} 02 & 02 & 03 & 04 & 05 & 03 & 05 & 07 \\ 07 & 02 & 02 & 03 & 04 & 05 & 03 & 05 \\ 05 & 07 & 02 & 02 & 03 & 04 & 05 & 03 \\ 03 & 05 & 07 & 02 & 02 & 03 & 04 & 05 \\ 05 & 03 & 05 & 07 & 02 & 02 & 03 & 04 \\ 04 & 05 & 03 & 05 & 07 & 02 & 02 & 03 \\ 03 & 04 & 05 & 03 & 05 & 07 & 02 & 02 \\ 02 & 03 & 04 & 05 & 03 & 05 & 07 & 02 \end{bmatrix}$$

- **MixBytes:** multiplies each column of the state matrix  $A$ , by a constant  $8 \times 8$  matrix  $B$ . The transformation, can be shown as  $A \leftarrow B \times A$ . The matrix  $B$ , can be seen as a finite field over  $\mathbb{F}_{256}$ , which is shown above. The finite field is defined over  $\mathbb{F}_2$  by the irreducible polynomial  $x^8 \oplus x^4 \oplus x^3 \oplus x \oplus 1$ .

### 3.3 Keccak

Keccak hash function, is built on sponge construction, which can input and output arbitrary length strings. The sponge construction has two phases. First is absorb, where the input message is ingested in blocks of defined bit rate interleaved with the permutations. The second phase is squeeze phase, where the blocks of output are squeezed out as per bit rate. The Keccak, state is different in the sense that, the permutations work on a 3 dimensional block, cube structure rather than linear strings, or 2 dimensional arrays.

#### 3.3.1 Keccak state, sponge functions and padding

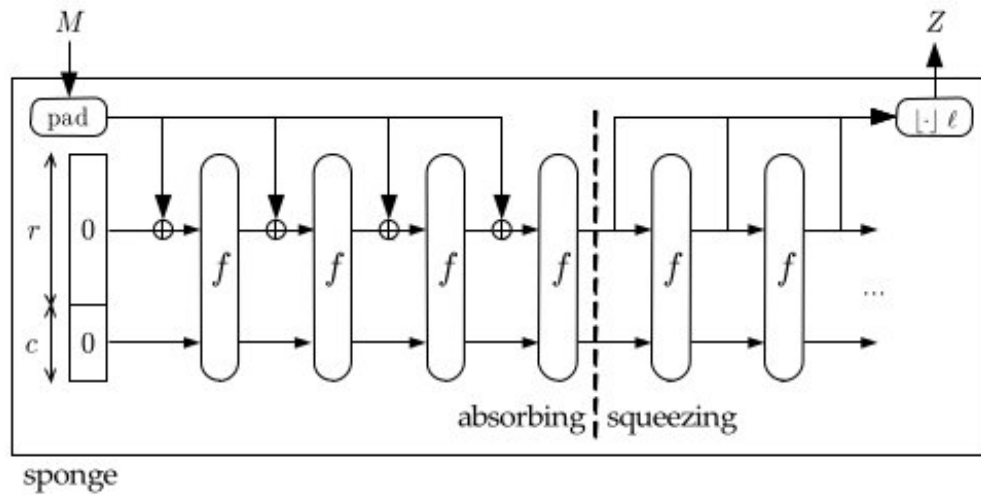


Figure 3.8: Sponge construction  $Z = \text{Sponge}[f, \text{pad}, r](M, l)$  [4]

The sponge construction is used to build function  $\text{SPONGE}[f, \text{pad}, r]$  which inputs and outputs variable length strings [4]. It uses fixed length permutation  $f$ , a padding "pad", and parameter bit rate 'r'. The permutations are operated on fixed number of bits, width  $b$ . The value  $c = b - r$  is the capacity of the sponge function. The width  $b$  in Keccak defines the state size which can be any of the following  $\{25, 50, 100, 200, 400, 800, 1600\}$  number of bits.



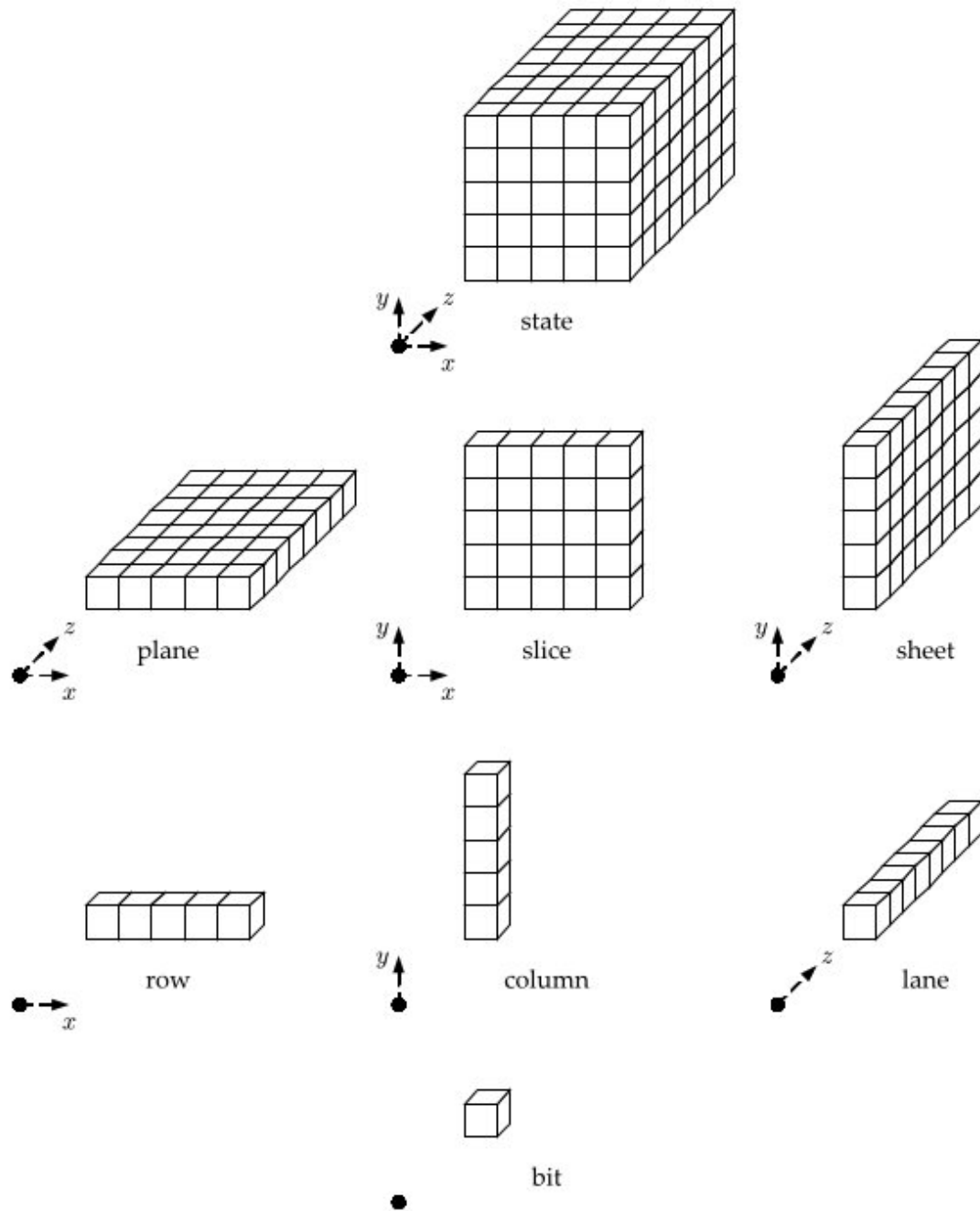


Figure 3.9: Sponge construction  $Z = \text{Sponge}[f, \text{pad}, r](M, l)$  [5]

The state in Keccak can be represented as a cube having bits, as shown in figure 3.9. The initial state to the sponge construction has value 'b' number of 0 bits (represented as  $0^b$ ), and called the root state. The root state has fixed value and should not be considered as initial value to sponge construction. The different number of state produces the Keccak family of hash function variations denoted by  $KECCAK - f[b]$ .

The varying number of states can be visualized as state having varying number or  $l$  number of slices. The width  $b$  is defined as  $b = 25 \times 2^l$ , where  $l$  takes values from 0 to 6.

---

**Algorithm 3.2** The sponge construction  $SPONGE[f, pad, r]$  [4]

---

**Require:**  $r < b$

```

1: Interface:  $Z = \text{sponge}(M, l)$  with  $M \in \mathbb{Z}_2^*$ , integer  $l > 0$  and  $Z \in \mathbb{Z}_2^l$ 
2:  $P = M \parallel pad[r](|M|)$ 
3:  $s = 0^b$ 
4:
5: for  $i = 0$  to  $|P|_r - 1$  do
6:    $s = s \oplus (P_i \parallel 0^{b-r})$ 
7:    $s = f(s)$ 
8:
9: end for
10:  $Z = \lfloor s \rfloor_r$ 
11:
12: while do  $|Z|_r r < l$ 
13:    $s = f(s)$ 
14:    $Z = Z \parallel \lfloor s \rfloor_r$ 
15:
16: end while
17: return  $\lfloor Z \rfloor_l$ 

```

---

Algorithm 3.2 shows how the sponge construction applied to  $KECCAK - f[r + c]$ , with multi-rate padding, length of a string  $M$  is denoted by  $|M|$ . The string  $M$  can also be considered as having blocks of size say  $x$ , and those number of blocks are shown as  $|M|_x$ . The  $\lfloor M \rfloor_l$  denotes the string  $M$  truncated to its first  $l$  bits.

The multi-rate padding in Keccak is denoted as  $pad\ 10^*1$ , where a bit '1' is appended to message, followed by minimum number of zeros, and a single bit 1, so that resultant block is multiple of block length  $b$ .

$$KECCAK[r, c] \doteq SPONGE[KECCAK - f[r + c], pad10^*1, r]$$

Where  $r > 0$  and  $r + c$  is the width. The default value of  $r$  is  $1600 - c$ , and the default value of  $c$  is 576.

$$KECCAK[c] \doteq KECCAK[r = 1600 - c, c],$$

$$KECCAK[] \doteq KECCAK[c = 576]$$

### 3.3.2 Permutations

The  $KECCAK - f[b]$  permutations are operated on state represented as  $a[5][5][w]$ , with  $w = 2^l$ , where  $l$  can be any value from 0 to 6. The position in this 3 dimensional state is given by  $a[x][y][z]$  where  $x, y \in \mathbb{Z}_5$  and  $z \in \mathbb{Z}_w$ . The mapping of the bits from the input message 's' to state 'a' is like this  $s[w(5y + x) + z] = a[x][y][z]$ . The  $x, y$  coordinates are taken modulo 5, while the  $z$  coordinate is taken as modulo  $w$  [5].

There are five steps, for a permutation round R.

$$R = \iota \circ \chi \circ \pi \circ \rho \circ \theta$$

The permutations are repeated for  $12 + 2l$  times, with  $l$  dependent on the variant chosen.

$$\theta : a[x][y][z] \leftarrow a[x][y][z] + \sum_{y'=0}^4 a[x-1][y'][z] + \sum_{y'=0}^4 a[x+1][y'][z-1],$$

$$\rho : a[x][y][z] \leftarrow a[x][y][z - (t+1)(t+2)/2],$$

$$t \text{ satisfying } 0 \leq t < 24 \text{ and } \begin{pmatrix} 0 & 1 \\ 2 & 3 \end{pmatrix}^t \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} x \\ y \end{pmatrix} \text{ in } GF(5)^{2 \times 2},$$

$$\text{or } t = -1 \text{ if } x = y = 0,$$

$$\pi : a[x][y] \leftarrow a[x'][y'], \text{ with } \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 2 & 3 \end{pmatrix} \begin{pmatrix} x' \\ y' \end{pmatrix},$$

$$\chi : a[x] \leftarrow a[x] + (a[x+1] + 1) a[x+2],$$

$$\iota : a \leftarrow a + RC[i_r].$$

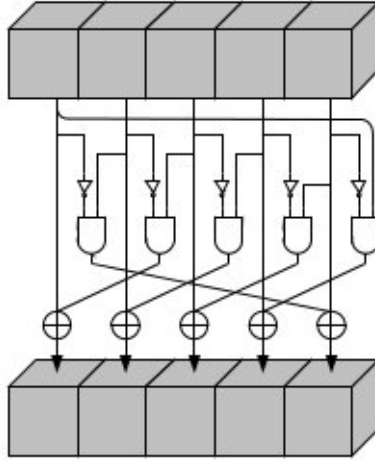


Figure 3.10:  $\chi$  applied to a single row [5].

The addition and the multiplications are in Galois field  $GF(2)$ , except for the round constants  $RC[i_r]$ . The round constants are given by

$$RC[i_r][0][0][2^j - 1] = rc[j + 7i_r] \text{ for all } 0 \leq j \leq l,$$

and the rest are zeros. The value of  $rc[t] \in GF(2)$  is output of linear feedback shift register given as

$$rc[t] = (x^t \bmod x^8 + x^6 + x^5 + x^4 + 1) \bmod x \text{ in } GF(2)[x].$$

---

**Algorithm 3.3**  $\chi$  transformation KECCAK [5].

---

```

1:
2: for  $y = 0$  to 4 do
3:
4:   for  $x = 0$  to 4 do  $A[x, y] = a[x, y] \oplus ((NOT\ a[x + 1, y])\ AND\ a[x + 2, y])$ 
5:
6:   end for
7:
8: end for

```

---

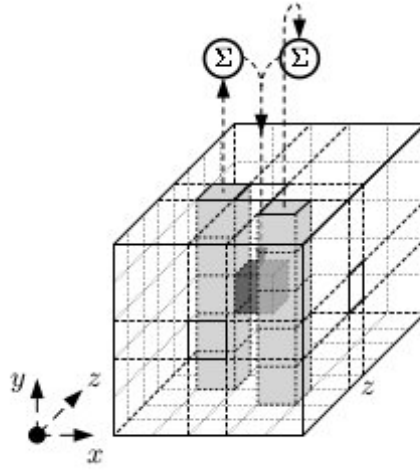


Figure 3.11:  $\theta$  applied to a single bit [5].

---

**Algorithm 3.4**  $\theta$  transformation KECCAK [5].

---

```

1:
2: for  $x = 0$  to 4 do
3:    $C[x] = a[x, 0]$ 
4:
5:   for  $y = 1$  to 4 do  $C[x] = C[x] \oplus a[x, y]$ 
6:
7:   end for
8:
9: end for
10:
11: for  $x = 0$  to 4 do
12:    $D[x] = C[x - 1] \oplus ROT(C[x + 1], 1)$ 
13:
14:   for  $y = 0$  to 4 do
15:      $A[x, y] = a[x, y] \oplus D[x]$ 
16:
17:   end for
18:
19: end for

```

---

---

**Algorithm 3.5**  $\pi$  transformation KECCAK [5].

---

```

1:
2: for  $x = 0$  to 4 do
3:
4:   for  $y = 1$  to 4 do
5:      $\begin{pmatrix} X \\ Y \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 2 & 3 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}$ 
6:      $A[X, Y] = a[x, y]$ 
7:
8:   end for
9:
10: end for
```

---

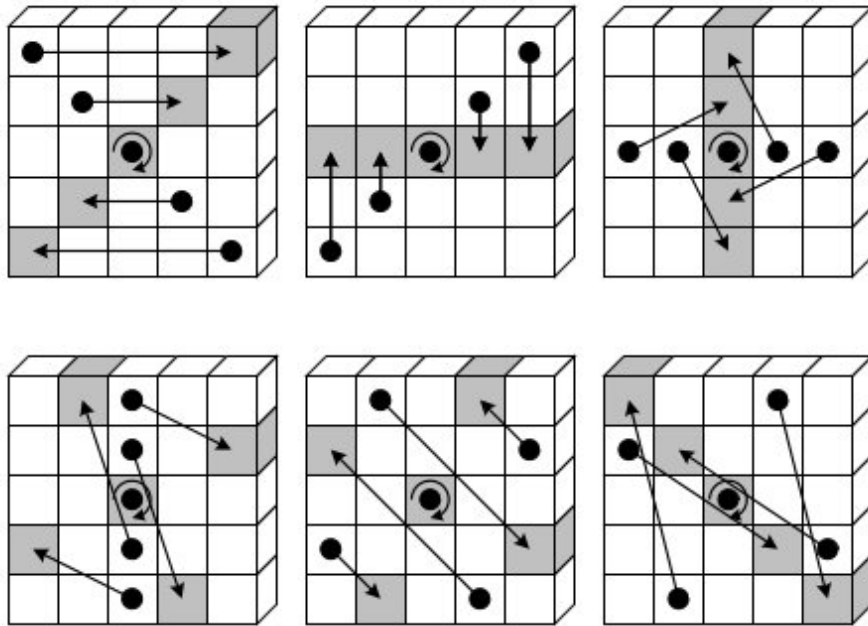


Figure 3.12:  $\pi$  applied to a single slice [5]

---

**Algorithm 3.6**  $\rho$  transformation KECCAK [5]

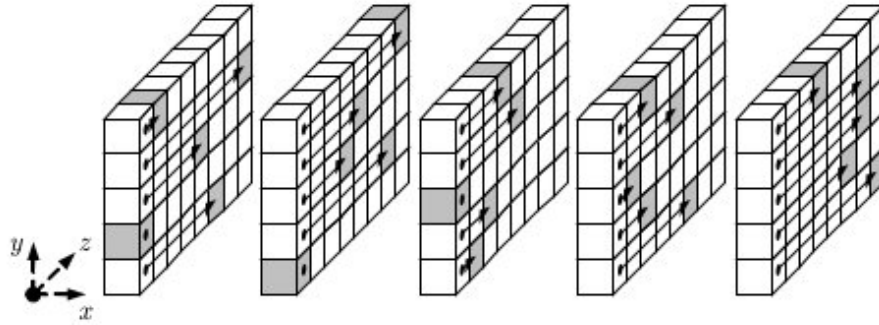
---

```

1:  $A[0, 0] = a[0, 0]$ 
2:  $\begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$ 
3:
4: for  $t = 0$  to 23 do
5:    $A[x, y] = ROT(a[x, y], (t + 1)(t + 2)/2)$ 
6:    $\begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 2 & 3 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}$ 
7:
8: end for

```

---

Figure 3.13:  $\rho$  transformation applied to lanes [5]

	$x = 2$	$x = 4$	$x = 0$	$x = 1$	$x = 2$
$y = 2$	153	231	3	10	171
$y = 1$	55	276	36	300	6
$y = 0$	28	91	0	1	190
$y = 4$	120	78	210	66	253
$y = 3$	21	136	105	45	15

Table 3.13: Offsets for  $\rho$  transformation [5]

Function	Output size	Security strength in bits		
		Collision	Preimage	2nd Preimage
SHA-1	160	$< 80$	160	$160 - L(M)$
SHA-224	224	112	224	$\min(224, 256 - L(M))$
SHA-512/224	224	112	224	224
SHA-256	256	128	256	$256 - L(M)$
SHA-512/256	256	128	256	256
SHA-384	384	192	384	384
SHA-512	512	256	512	$512 - L(M)$
SHA3-224	224	112	224	224
SHA3-256	256	128	256	256
SHA3-384	384	192	384	384
SHA3-512	512	256	512	512
SHAKE128	$d$	$\min(d/2, 128)$	$\geq \min(d, 128)$	$\min(d, 128)$
SHAKE256	$d$	$\min(d/2, 256)$	$\geq \min(d, 256)$	$\min(d, 256)$

Table 3.14: Security strength of SHA-1, SHA-2, SHA-3 functions [12]

### 3.4 SHA-3

As per the latest proposed draft, SHA-3 contains 4 cryptographic hash function variations of Keccak for digest sizes 224, 256, 384 and 512; along with two extendable output functions(XOF) [12]. XOF are hash functions, whose output can be extended to desired length.

The Keccak variants are denoted by KECCAK[c] where c is the capacity of the sponge construction used for Keccak function. The message that is to be hashed is concatenated with two bits, and the output digest length is mentioned as parameters to the Keccak function for SHA-3.

$$\text{SHA3-224}(M) = \text{KECCAK}[448](M \parallel 01, 224)$$

$$\text{SHA3-256}(M) = \text{KECCAK}[512](M \parallel 01, 256)$$

$$\text{SHA3-384}(M) = \text{KECCAK}[768](M \parallel 01, 384)$$

$$\text{SHA3-512}(M) = \text{KECCAK}[1024](M \parallel 01, 512)$$

The capacity of each of the hash function is twice that of the digest length. The input string is concatenated with two bits 01, to support domain separation. The XOF are called



SHAKE128 and SHAKE256, where the suffix 128 and 256, denote security strength of the functions. SHAKE nomenclature comes from the combination of words, SHA and Keccak. The XOF can be defined as

$$\text{SHAKE128}(M, d) = \text{KECCAK}[256](M \parallel 1111, d)$$

$$\text{SHAKE256}(M, d) = \text{KECCAK}[512](M \parallel 1111, d)$$

Here the  $d$  denotes digest size, since the output size can be varied as per parameter  $d$ . There are two capacities available for SHAKE, that of being 256 and 512, that give the security margin of 128 and 256 respectively. The message before being provided as input to SHAKE is concatenated with 4 bits of 1. The first two bits of 1 are for domain separation, while the rest 2 bits of 1 are for compatibility with Sakura coding, that enables tree hashing, in which parallel processing can be applied.

## Chapter 4

# Related work and hypothesis based on Hill Climbing to find near collisions

### 4.1 Rotational cryptanalysis of round-reduced KECCAK

Rotational cryptanalysis[17] is used to follow relation between states  $(A, A^\leftarrow)$  of KECCAK-f[1600] in course of their transformation, and thus derive a distinguisher.

**Definition 4.1.1.** A pair of two 1600-bit states  $(A, A^\leftarrow)$  is called rotational pair when each lane in state  $A^\leftarrow$  is created by bitwise rotation of operation of corresponding lane in state  $A$ . The operation moves the bit from position  $(x, y, z)$  to the position  $(x, y, z + n)$ , where  $z + n$  is done on modulo 64.  $x$  and  $y$  values range from 0 to 4, and  $z$  value ranges from 0 to 63.  $n$  is the rotational number and is same for every lane. Thus rotational pairs are  $\forall(x, y, z) : A_{(x,y,z)} = A_{(x,y,z+n)}^\leftarrow$  [18].

**Definition 4.1.2.** Set  $S_n$  is a set of  $2^{1600}$  pairs of states which are created by an operation of KECCAK-f[1600] applied to all possible rotational pairs [18].

**Definition 4.1.3.** Probability  $p_{(x,y,z)}^n$  is the probability for pair of states  $(A, A^\leftarrow)$  randomly selected from the set  $S_n$  we have  $A_{(x,y,z)} = A_{(x,y,z+n)}^\leftarrow$  [18].

**Definition 4.1.4.** Given probability distribution  $\mathcal{D}_n$  that assigns probability  $\frac{1}{n!}$  for each  $p \in \mathcal{P}_n$ . A permutation is called random if it is chosen according to uniform distribution  $\mathcal{D}_n$  [18].

It is assumed that random permutation  $p_{(x,y,z)}^n$  follows binomial distribution  $\mathcal{B}(t, s)$  where  $t$  is trials and  $s$  is success probability that is equal to 0.5. Experimental results for a chosen  $p_{(x,y,z)}$  are compared to follow distribution  $\mathcal{B}(t, s)$ . The experimental values are supposed to fall within range of  $0.5t \pm 2\sigma$  with 95% confidence interval.

The probability change through steps of Keccak, can be derived from analysis of bitwise operation. It is assumed that corresponding bits from  $(A, A^\leftarrow)$  are equal, both combinations ('00' or '11') or opposite combinations have same probability to be actual values. Operations like NOT, or rotation in Keccak do not affect the probabilities, so only probabilities for AND and XOR are considered.

**Lemma 4.1.1.** *Given input bits  $a, b$  and output bit  $out$ ; with  $p_a$  and  $p_b$  defined as per definition 4.1.3, then for AND operation  $P_{out} = \frac{1}{2}(p_a + p_b - p_a p_b)$  [18]*

**Lemma 4.1.2.** *Given input bits  $a, b$  and output bit  $out$ ; with  $p_a$  and  $p_b$  defined as per definition 4.1.3, then for XOR operation  $P_{out} = p_a + p_b - 2p_a p_b$  [18]*

A 4 round rotational was built, that found some values deviate from the 0.5 like  $p_{(4,4,14)}^{54} = 0.5625$ . 10,000 random samples of rotational pairs are ran on 4 round KECCAK-f[1600]. The mean from that sample is 5682 which is beyond range from mean of  $\mathcal{B}(10000, 0.5)$  which is  $5000 \pm 2.5$ . Thus demonstrating a distinguisher for 4 rounds. After 4 rounds  $p_{(x,y,z)}^n = 0.5$ , and hence the distinguisher cannot be directly extended. But instead probability that relation between two pairs of states  $(A_{(x,y,z)}, A_{(x,y,z+n)}^\leftarrow)$  and  $(A_{(x,y',z)}, A_{(x,y'',z+n)}^\leftarrow)$  are observed, that should follow distribution  $\mathcal{B}(10000, 0.5)$ . Values for  $p_{(2,1,37)}^{63}$  and  $p_{(2,2,37)}^{63}$  have the highest deviation from 0.5 at end of fourth round. The probability that they are in the same relation is given by  $p_{(x,y,z)}^n p_{(x,y'',z)}^n + (1 - p_{(x,y,z)}^n)(1 - p_{(x,y,z)}^n)$  which comes to 0.499024. The  $\rho$  and  $\pi$  steps in Keccak only change the position of the probability to bit pairs  $(A_{(1,2,43)}, A_{(1,2,44)}^\leftarrow)$  and  $(A_{(2,0,16)}, A_{(2,0,17)}^\leftarrow)$ . The bias is observed by generating sufficient number of samples 'm' based on Chernoff bound based on inequality

$$m \geq \frac{1}{(P_c - 0.5)^2} \ln \frac{1}{\sqrt{\epsilon}}$$

where  $\epsilon$  is error set to 0.05, thus value of m turns out to be 403,000,000. Based on this value algorithm 4.1 is implemented.

The mean for  $\mathcal{B}(403000000, 0.5)$  with the standard deviation has range of  $201,500,000 \pm 2.10037$ , but experimentally from implementing algorithm 4.1, it comes to around 201,450,503, thus concluding this as a distinguisher.

---

**Algorithm 4.1** Find pair probability [18]

---

```

1: Generate 403,000,000 rotational pairs.
2: for all 403,000,000 rotational pairs do
3:   Run Keccak for 5 rounds on states  $A$  and  $A^{\leftarrow}$ .
4:   if  $(A_{(1,2,43)} \oplus A_{(1,2,44)}^{\leftarrow} \oplus A_{(2,0,16)} \oplus A_{(2,0,17)}^{\leftarrow} = 0)$  then
5:     mean := mean + 1
6:   end if
7: end for
8: return mean

```

---

For the preimage attack an unknown message with cyclical pattern like that of 4 0's followed by 4 1's alternatively, of 512 bits is chosen. There are 256 possible messages of the cyclic pattern. For the preimage attack, a rotational counterpart of the preimage is searched for, that would reduce the complexity from random search. Algorithm 4.2 describes the steps.

---

**Algorithm 4.2** Preimage for 3 round Keccak for unknown cyclic input [18]

---

```

1: Guess first 8 lanes of  $A^{\leftarrow}$ 
2: Run Keccak-f[1600] for 3 rounds on state  $A^{\leftarrow}$ 
3: for  $n := 0$  to  $n < 64$  do
4:   candidate := true
5:   for 10 sets of coordinates  $(x, y, z)$  being on list created on precomputation do
6:     if  $(p_{(x,y,z)}^n = 1)$  and  $(A_{(x,y,z)} = A_{(x,y,z)}^{\leftarrow})$  then
7:       candidate := false
8:     end if
9:     if  $(p_{(x,y,z)}^n = 0)$  and  $(A_{(x,y,z)} \neq A_{(x,y,z)}^{\leftarrow})$  then
10:      candidate := false
11:    end if
12:  end for
13:  if candidate = true then
14:    Rotate the guessed state by  $n$  bits.
15:    Verify input to Keccak, that runs for 3 rounds.
16:  end if
17: end for

```

---

For a given state there are 64 possible rotational pairs, derived from length of lane. We are searching for preimage of 512 bits, thus the probability of guessing the rotational

counterpart  $A^\leftarrow$  is  $2^{-512} \cdot 64 = 2^{-506}$ , or rather  $2^{506}$  guesses. There are  $2^{256}$  messages possible of the cyclic pattern that we consider here. There are 10 sets of  $(x, y, z)$  coordinates for each of the rotational number. The probability of the candidate having  $p_{(x,y,z)}^n$  similar to that on list is  $2^{-10}$ . So  $2^{512}/2^{10} = 2^{502}$  number of checks are required at most to find the candidate. Thus the total work is equivalent to  $2^{506}$  calls to KECCAK-512 for 3 rounds to find the preimage.

The above method for finding the preimage cannot be directly extended to 4 rounds since  $\iota$  operation in Keccak renders  $p_{(x,y,z)}^n \neq 0, 1$ . To overcome this, the rotational state is ran on modified Keccak-512 for four rounds which does not implement  $\iota$  function. Algorithm 4.3 is for finding preimage.

---

**Algorithm 4.3** Preimage for 4 round Keccak for unknown cyclic input without  $\iota$  function in Keccak [18]

---

```

1: Choose 512 bits at random for state  $A^\leftarrow$ 
2: Run KECCAK-f[1600] without the  $\iota$  transformation is run on them.
3: for for  $n := 0$  to  $n < 64$  do
4:   candidate := true
5:   for 9 sets of coordinates  $(x, y, z)$  that are in list created from precomputation do
6:     if  $(p_{(x,y,z)}^n = 0)$  and  $(A_{(x,y,z)} \neq A_{(x,y,z+n)}^\leftarrow)$  then
7:       candidate := false
8:     end if
9:     if  $(p_{(x,y,z)}^n = 1)$  and  $(A_{(x,y,z)} = A_{(x,y,z+n)}^\leftarrow)$  then
10:      candidate := false
11:    end if
12:  end for
13:  if candidate = true then
14:    Rotate the guessed state by  $n$  bits.
15:    Verify input to Keccak, that runs for 4 rounds of modified Keccak-512.
16:  end if
17: end for

```

---

Just like the 3 round preimage finding method, the above method for finding the preimage for 4 rounds has complexity of  $2^{506}$  calls to Keccak-512.

## 4.2 Finding near collisions with Hill Climbing

A generic algorithm applied to find collisions, in reduced rounds of some SHA-3 competitors was Hill Climbing [28]. Near collisions in which more than 75% of the bits were same for two different messages, were found for reduced rounds of BLAKE-32, Hamsi-256 and JH. Near collision results are important for knowing the security margins. In some cases, output of hash functions may be truncated for compatibility or efficiency purposes. In such cases near collisions could be improved to obtain collisions.

A  $\epsilon/n$  bit near collision for hash function  $h$  and two messages  $M_1$  and  $M_2$ , where  $M_1 \neq M_2$  can be defined as

$$HW(h(M_1, CV) \oplus h(M_2, CV)) = n - \epsilon$$

where HW is the Hamming weight, and CV is the chaining value, and  $n$  is the hash size in bits.

$\epsilon/n$	Complexity ( $\approx$ )
128 / 256, 256 / 512, 512 / 1024	$2^4$
151 / 256, 287 / 512, 553 / 1024	$2^{10}$
166 / 256, 308 / 512, 585 / 1024	$2^{20}$
176 / 256, 323 / 512, 606 / 1024	$2^{30}$
184 / 256, 335 / 512, 623 / 1024	$2^{40}$
191 / 256, 345 / 512, 638 / 1024	$2^{50}$
197 / 256, 354 / 512, 651 / 1024	$2^{60}$

Table 4.1: Approximate complexity to find a  $\epsilon/n$ -bit near collision by generic random search [28]

Hill Climbing starts with a random candidate, and then choosing a random successor that has a better fit to the solution. In practice for message  $M$  and chaining value  $CV$

$$HW(h(M, CV) \oplus h(M, CV + \delta)) = n/2,$$

can be considered secure, where  $\delta$  is  $n$ -bit vector with small Hamming weight. However, if the diffusion for the hash function  $h$  is not proper, then we obtain a lower Hamming weight.

In such situation a correlation between two chaining values differing in small weight  $\delta$  can obtain near collisions, with hill climbing algorithm.

Here, the aim of hill climbing algorithm will be to minimize the function

$$f_{M_1, M_2}(x) = HW(h(M_1, x) \oplus h(M_2, x))$$

where  $x \in \{0, 1\}^n$ , where  $M_1$  and  $M_2$  are message blocks. CV is chosen as any random chaining value. Then the set of k-bit neighbours for the CV, will be

$$S_{CV}^k = \{x \in \{0, 1\}^n \mid HW(CV \oplus x) \leq k\}$$

where

$$size\ of\ S_{CV}^k = \sum_{i=0}^k \binom{n}{i}.$$

The k-opt condition can be defined as

$$f_{M_1, M_2}(CV) = \min_{x \in S_{CV}^k} f_{M_1, M_2}(x)$$

We can now describe algorithm 4.3, that is used in hill climbing algorithm to find the nearest match.

---

**Algorithm 4.4** Hill Climbing algorithm ( $M_1, M_2, k$ ) [28]

---

- 1: Randomly select CV
  - 2:  $f_{best} = f_{M_1, M_2}(CV)$
  - 3:
  - 4: **while** (CV is not k-opt) **do**
  - 5:     CV = x such that  $x \in S_{CV}^k$  with  $f(x) < f(best)$
  - 6:      $f_{best} = f_{M_1, M_2}(CV)$
  - 7:
  - 8: **end while**
  - 9: **return** (CV,  $f_{best}$ )
- 

Given two message  $M_1$  and  $M_2$ , and a randomly chosen chaining value CV, the  $f_{M_1, M_2}(CV)$  is obtained. The set  $S_{CV}^k$  is searched for a better fit CV, and if found is updated. And the search is repeated again in the k-bit neighbourhood of new CV.

There are two ways of choosing the next best CV, one by choosing the first chaining value that has a lower  $f$  value, the greedy way. And another by choosing the best chaining value amongst  $S_{CV}^k$ , which is steepest ascent. The algorithm terminates once we get k-opt chaining value.

## 4.3 Tabu Search, Simulated Annealing and Random search

### 4.3.1 Simulated Annealing

---

**Algorithm 4.5** Simulated Annealing Algorithm for obtaining near collisions

---

```

1: function SIMULATED-ANNEALING( $M_1, M_2, CV, \text{schedule}$ )
2:    $\text{current} \leftarrow CV$ 
3:   for  $t = 1$  to  $\infty$  do
4:      $T \leftarrow \text{schedule}(t)$ 
5:     if  $T = 0$  then
6:       return  $\text{current}$ 
7:     end if
8:      $\text{next} \leftarrow$  a randomly selected successor from set  $S_{\text{current}}^k$ 
9:      $\Delta E \leftarrow f_{M_1, M_2}(\text{current}) - f_{M_1, M_2}(\text{next})$ 
10:    if  $\Delta E > 0$  then
11:       $\text{current} \leftarrow \text{next}$ 
12:    else
13:       $\text{current} \leftarrow \text{next}$ , with probability  $e^{\Delta E/T}$ 
14:    end if
15:  end for
16: end function

```

---

The problem with hill climbing, is that it can get locked in the local maxima, and fail to get the global maxima. This is due to hill climbing not taking a downhill or a step with lower value. However, if hill climbing is tweaked to combine with random walk, then the problem of local maxima can be avoided. Simulated annealing picks a random successor, and accepts it if the value is higher than previous. However, if the successor has a lower value, then it is accepted with a probability less than 1. The probability has an exponential decrease proportional to the decreased value of the move, and the temperature. Thus at higher temperature or at the initial stages, a downhill successor is more likely to be



accepted, than in the later stages [21].

### 4.3.2 Tabu Search

---

**Algorithm 4.6** Tabu Search for obtaining near collisions [8]

---

```

1: function TABU-SEARCH( $TabuList_{size}, M_1, M_2, CV$ )
2:    $S_{best} \leftarrow CV$ 
3:    $TabuList \leftarrow \text{null}$ 
4:   while  $S_{best}$  not k-opt do
5:      $CandidateList \leftarrow \text{null}$ 
6:      $S_{neighbourhood} \leftarrow S_{S_{best}}^k$ 
7:     for  $S_{candidate} \in S_{best_{neighbourhood}}$  do
8:       if ( $\neg \text{ContainsAnyFeatures}(S_{candidate}, TabuList)$ ) then
9:          $CandidateList \leftarrow S_{candidate}$ 
10:      end if
11:    end for
12:     $S_{candidate} \leftarrow \text{LocateBestCandidate}(CandidateList)$ 
13:    if  $\text{Cost}(S_{candidate}) \leq \text{Cost}(S_{best})$  then
14:       $S_{best} \leftarrow S_{candidate}$ 
15:       $TabuList \leftarrow \text{featureDifferences}(S_{candidate}, S_{best})$ 
16:      while  $TabuList > TabuList_{size}$  do
17:         $\text{DeleteFeature}(TabuList)$ 
18:      end while
19:    end if
20:  end while
21:  return  $S_{best}$ 
22: end function

```

---

Tabu search implements the neighbourhood search for the solutions, until the termination condition. The algorithm uses a fixed amount of memory, to keep note of states, visited some fixed amount of time in past. The idea behind keeping the state, is to restrict the search, to states that have not been visited previously. The algorithm can be tweaked, to accept moves in tabu list through aspiration criteria, or inferior moves just to explore new possible states. Tabu search has been applied to mostly combinatorial optimization problems [11, 20].

---

**Algorithm 4.7** Random selection from k-bit neighbourhood of  $CV$ 


---

```

1: function RANDOM-SELECTION( $M_1, M_2, CV, \text{number\_of\_trials}$ )
2:    $\text{current} \leftarrow CV$ 
3:    $\text{trial} \leftarrow 0$ 
4:   while  $\text{trial} < \text{number\_of\_trials}$  do
5:      $\text{next} \leftarrow$  randomly selected candidate from  $S_{\text{current}}^k$ 
6:     if  $f_{M_1, M_2}(\text{next}) - f_{M_1, M_2}(\text{current})$  then
7:        $\text{current} \leftarrow \text{next}$ 
8:     end if
9:   end while
10:  return  $\text{current}$ 
11: end function

```

---

## 4.4 Hypothesis

### HYPOTHESIS

- Reduced state Keccak, has better resistance to near collisions than BLAKE and Grøstl. For the attack algorithms hill climbing, simulated annealing, tabu search and random selection.
- Simulated annealing and tabu search, are better at finding near collisions compared to hill climbing and random selection.
- State size has no effect on efficiency of Keccak permutation rounds, for reduced number of rounds.

As per the press release from NIST, one of the reasons for choosing Keccak, was that it had a large security margin. All the five finalists from SHA-3 competition were found to be secure and have good security margins. However there has not been much study, on the comparative security margins for the candidate's reduced versions. Hill climbing has been shown as good generic greedy algorithm to find near collisions for reduced versions of some SHA-3 candidates. A generic algorithm does not exploit the inner permutations or construction, of a hash function. Rather takes a good guess approach, to what the solution can be depending on the fitness of the candidate solution. Thus making it an ideal tool to

test on any hash function, irrespective of its design. In addition to hill climbing, it would also be interesting to observe the success other variations of generic algorithms like Tabu search, or Simulated Annealing will be able to get, and compare those results against a random search. In theory for an ideal hashing function the performance of the generic algorithm will be equivalent to the random search algorithm on an average.

Comparative studies on SHA-3 candidates have been using the statistical test suites provided by NIST to check any deficiencies [9] [15]. Other than particular attacks like zero-sum property has been tested on Keccak and Blue Midnight Wish [1].

# Chapter 5

## Research Approach and Methodology

### 5.1 Experiment Structure

The experiment is designed to find the number of attempts it takes to find near collision amongst the three candidate hashing algorithm BLAKE, Keccak and Grøstl when subjected to attacks from following algorithms hill climbing, simulated annealing, tabu search and random selection. The idea is to take a seed message and update it, so we get two different input message with tiny difference. Then try to minimize our cost function, which is minimizing the hamming weight of the string obtained by bitwise XOR of the two message digests, obtained by feeding the two input message that are padded by the same chaining value. The minimization of the of the cost function is obtained by the collision finding the algorithm, which they acheive by selecting the suitable chaining value.

The experiment is conducted for a number of trials, where the digest lengths are varied at the standard bit lengths of 224, 256, 384 and 512 as standardized by SHA-3. The full version of SHA-3 finalist hashing algorithms have been found to be practically secure, so we do the experiments on the reduced versions of these algorithms. A hashing algorithm can be reduced in ways like reducing the digest size, the internal state size, or the number of permutation rounds. We choose to vary the permutation rounds in the hash functions for our experiment. The factors that are kept common for comparison are the digest lengths, message pairs and permutation rounds. For each trial and each hashing algorithm, the chaining value is randomly chosen and then worked upon by the collision finding algorithm.

### 5.1.1 Input

The string "The quick brown fox jumps over the lazy dog", was chosen as the root seed message. This seed string contains all letters from English alphabet, and is neither too small or large. Pairs were made from this seed string, by toggling a bit, from the ASCII/UTF-8 bit representation of this string. The toggling of the bits are divided into 3 parts - starting, middle, and end. In starting part the most significant bits of the string are toggled, while in the end part the least significant bits are toggled. In the middle part, the bits toggled equally from the most significant and least significant side of the bit that is in middle of the string. For example let bit representation of a string be 0110001000011000, then in starting section there will be strings generated of kind 1110001000011000, 0010001000011000, 0100001000011000 and so on. Only 1 bit from the seed string is toggled, starting from the first bit, then the second bit. This process is repeated till the number of bits asked to be toggled. In this case of experiment we updated 20 bits from the seed string, thus generating 20 strings from seed string each having a bit difference from the seed.

The generated strings are paired up with the seed string and are written to a file. Each file has a pair of strings written to it, separated by newline. The text files holding these pairs are named the number, derived from the order in which the bit was updated for the generated string. For example the input file 1.txt will have the entry of seed string and the generated string that has the first bit toggled. Following are the contents of the file 1.txt in our case.

```
54686520717569636b2062726f776e20666f78206a756d7073206f76657220746865206c617a7920646f67
d4686520717569636b2062726f776e20666f78206a756d7073206f76657220746865206c617a7920646f67
```

Notice that the first line is the hexadecimal representation of the seed string "The quick brown fox jumps over the lazy dog", and the second line has the first bit toggled, from the seed string and again represented in hexadecimal format. In similar way rest of the 20 files are created and named for the bits updated from the most significant bit onwards. These

files are then stored in the folder "Start", which in turn is stored in the folder called "Input" that holds all the input strings.

Similarly in the "Input" folder two more folders "Middle" and "End" are created, and each filled with 20 text files named or numbered 1 to 20. In the "Middle" folder are files that have two lines of string that have one bit difference in the middle section of the string. The bits updated are equally distributed around most and least significant between the middle bit in the seed string. For example in case of two bits being toggled, for the "Middle" section for seed string in bit form 0110001000011000 you will get two strings like 0110001100011000, and 0110001010011000. The bit toggling starts from the most significant bit selected from the middle bit to least significant bit, that is from left to right. So in the example provided if the seed string is 0110001000011000, then file Input/Middle/1.txt will contain

```
0110001000011000
0110001100011000
```

an file Input/Middle/2.txt will contain.

```
0110001000011000
0110001010011000
```

Please note that the above mentioned fragments are examples, and not actual contents. The actual contents are going to be hexadecimal representation of bit value, of the seed string "The quick brown fox jumps over the lazy dog", and the hexadecimal representation of bits of the updated string, as shown for the file 1.txt in "Start" folder category.

In the similar manner, files are created for the "End" category. Least significant bits are toggled, one by one proceeding towards the significant bits. Say the seed string is 0110001000011000 then in file 1.txt the seed will be paired with string 0110001000011001, and in file 2.txt it will be paired with 0110001000011010 and so on. The files for input can be found along with the source code implementation git repository.

### 5.1.2 Output

The output has detailed folder structure, due to breadth of the experiment parameters and numbers noted down for the same. It has the following structure

Output/length\_of\_chain\_value/collision\_algorithm/digest\_size/SHA3\_finalist\_algorithm  
/number\_of\_rounds/Category\_of\_toggled\_input/files\_named\_as\_in\_input

For example if the experiment is conducted on input 1.txt on Input/Start category, with a chaining value length of 32 bits. The collision algorithm Hill Climbing in ran on SHA-3 finalist algorithm BLAKE, that hashes both the input strings concatenated with the chaining value for a digest size of 224 bits, running only 2 rounds of permutation. Then the output file 1.txt is created in the location as per the above given structure in path Output/32/HillClimbing/224/BLAKE/2/Start/1.txt.

Each file has 8 entries in it, which are

1. The number of times near collisions were found and not found, from the number of trials.
2. The cumulative total number of iterations that it took for the algorithm to find the near collision.
3. The cumulative total number of iterations ran, from the trials when it did not find near collision.
4. Average number of iterations for when near collisions were found and not found.
5. Total cumulative iterations altogether for the experiment for all trials, and average iterations.

Instead of parameter time, we note down the average of iterations of operations over number of trials, that it takes for the collision finding algorithm, to find near collisions. The output files can be found with source code implementation in git repository.

### **5.1.3 Rational for the experiment structure, parameters and collected data**

#### **Choosing iterations over execution time and noting collisions**

After creation of the output files, the average iterations for each case of input that is start, end or middle over the pairs starting from 1 to 20 are entered manually into a excel sheet. Thus we get an average of iterations for all the hashing algorithms, over all the collision finding algorithms, for all the digest size and number of permutation round for input type having strings differing by a bit at a certain point.

The iteration is a whole number starting from zero, unique to a collision finding algorithm instance. Each time in the main algorithm loops over to find a collision based on existing chaining value, by choosing one of its neighbours; the iteration is incremented. Sometimes the iteration is also incremented for processes inside the loop like finding the best neighbour from the neighbourhood in case of tabu search, which is also considered as part of algorithm operation. Iterations show how many operations were required before success was obtained. This is important in our experiment in hill climbing, where we run the algorithm till we find near collision.

The near collision in this case is defined as having 65% bits as same. So if a collision algorithm gets 65% bits of two hash digests to be similar then it is noted as a success. This is a small margin, and randomly 50% of the bits can be guessed correctly. However, getting more than 50% of the bits to be similar, by non random methods with finite attempts can be marked as a success.

#### **Why have two message pairs in the way they are?**

Having two different inputs, with slight differences gives us the opportunity to determine the diffusion properties of the hashing algorithms. Hashing algorithms ideally should have diffusion properties that distribute small differences in text over the ciphertext so as to make those small updates undetectable from the given two message digests. The ease in finding a near collision in two different messages, enables to derive conclusions on the strength of



diffusion properties of the hash function.

### **Limiting neighbourhood number to 2**

The set of k-bit neighbourhood for the chaining value is defined as

$$S_{CV}^k = \{x \in \{0, 1\}^n \mid HW(CV \oplus x) \leq k\}$$

where

$$size\ of\ S_{CV}^k = \sum_{i=0}^k \binom{n}{i}.$$

Thus the size of neighbourhood is summation of the combinatorial series of chaining value bit length, to selections starting from 1 to k, where k is the at most number of bits you want to be updated in the chaining value to be included in the neighbourhood set.

For the experiment we are keeping the k-bit of neighbourhood at 2. Increasing the k-bit neighbourhood above 3, may get better solutions, but then the neighbourhood set size grows too fast. Searching in such a huge neighbourhood does not leave any of the collision methods feasible.

### **Breaking down the input, into categories**

Purpose of dividing the bit difference of the messages into three categories of start, middle and end; is to test if the hashing algorithms have any particular vulnerability when messages are updated at select points. Different hashing algorithms have different approaches to substitution and permutation, and that could leave gaps in how the diffusion of the input message is achieved in digest. Achieving collision should be equally hard irrespective of the point where the message has been changed from the original. This classification of message pair is to find if weakness in diffusion in hashing algorithms is in any related to position of bit change in the message.

The experiment is more on lines of black box testing, where we just evaluate the output rather than examining how the input is processed. Thus we need a plethora of test cases

to cover most possibilities before concluding any weakness in one hash function compared to others. The hash functions are compared across the digest lengths of 224, 256, 384 and 512 bits, which have been standardized by NIST.

### **Achieving reduced versions of algorithms**

For our experiments, we achieve the reduction in the hashing algorithm by reducing the number of permutation rounds in a function. There are other ways to achieve reduction in a hash function, like reducing the internal state size, or reducing the number of output digest bits. NIST has specifically increased the digest sizes from SHA-1 standard of 128 bits to new 4 different bit sizes of 224, 256, 384 and 512; considering future security considerations from increasing computational power. Thus we chose not to reduce the number of output bits. Internal state size reduction was another way, but different algorithms have different state sizes. The distribution of input and other constant parameters is also different, in each of the hashing algorithm. Thus for reducing the state size, a need in reducing the word size or reducing other constants would be required, which may diminish the confusion and diffusion properties for hashing algorithms in ways unfair for experimental standardization. Thus we decided to reduce the hash function, by reducing the number of rounds.

Although reducing the number of permutation rounds is not the perfect way of reducing a hash function, given that creators of hashing functions do a trade off on a number of factors like word size, state size, S-box, the construction model etc. Thus suggesting a recommended number of permutation rounds for the hashing function, that would make sure of security of the hashing function considered holistically. However, reducing the number of permutation rounds is easy to achieve, and gives a uniform parameter amongst the hashing function to compare. That is, are the permutation functions involved in the hashing good enough properties of confusion and diffusion in minimal application.

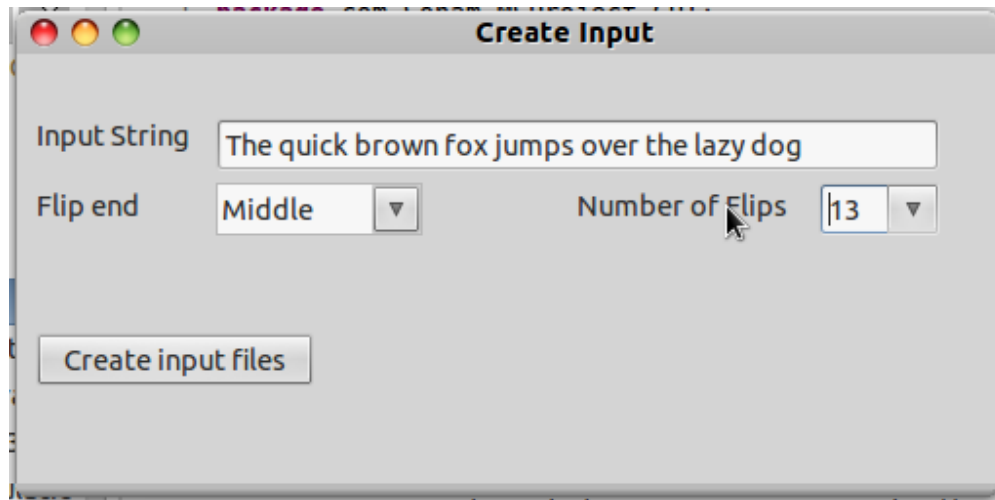


Figure 5.1: Screen shot of GUI screen input, used to create the input files.

## 5.2 Implementation

### 5.2.1 Input Creation

To create the input, a GUI was created in which the seed text could be inserted and then choose the input case, like if you want to toggle the bits at the start, middle or end of the string; and number of bits you want to toggle. On click of the create input file button, files with the bits flipped to the number provided paired with seed string will be created in respective folder.

The click button on the GUI shown in figure 5.1 calls the `createFile()` function in class `CreateInputFile`, which is responsible for file creation. Details for both the class are shown in figure 5.2.

### 5.2.2 Hash function implementation

Each of the SHA-3 finalist hashing algorithm, were implemented in their own individual Java class. For the purpose of the experiment, we only required the control over what input string went to the hashing algorithm, the number of permutation rounds, and digest size in

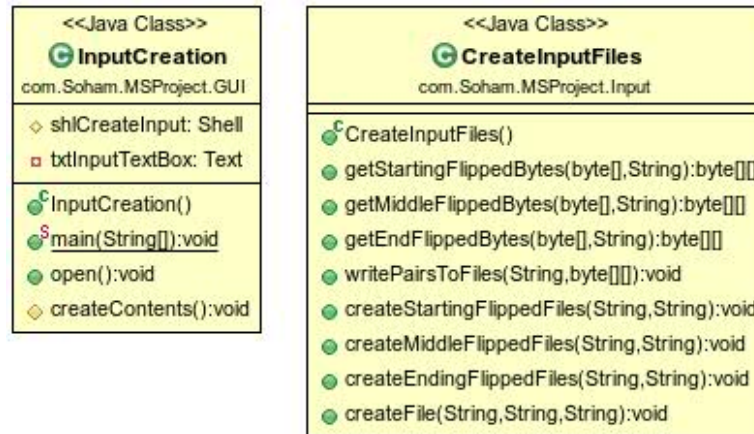


Figure 5.2: Class diagram of the input creation.

bits. Thus all the SHA-3 algorithms implement the interface "Hash" that standardizes the call to each of finalist algorithm. The input string to the hash function, is the hexadecimal string representation of the ASCII bit value of the input string. The digest size is an integer from the following four values 224, 256, 384 and 512. And the rounds can be anything between 1 and maximum number of rounds for a given hash. For example in Keccak the number of rounds is 24, and if 25 is input in rounds, then the hashing will be done for 24 rounds. So a safety check for a upward limit has not been built. By default, if you put zero, then the recommended values of permutation rounds for the respective algorithm are used. The class diagram for the SHA-3 hash function implementation is shown in figure 5.3.

### 5.2.3 Experiment with different collision methods

From the experiment selection dropdowns shown in figure 5.4, we can choose from the following factors

1. The length of the chaining value that will be padded to the message, in bits.
2. The algorithm for finding collision.
3. The length of message digest in bits.

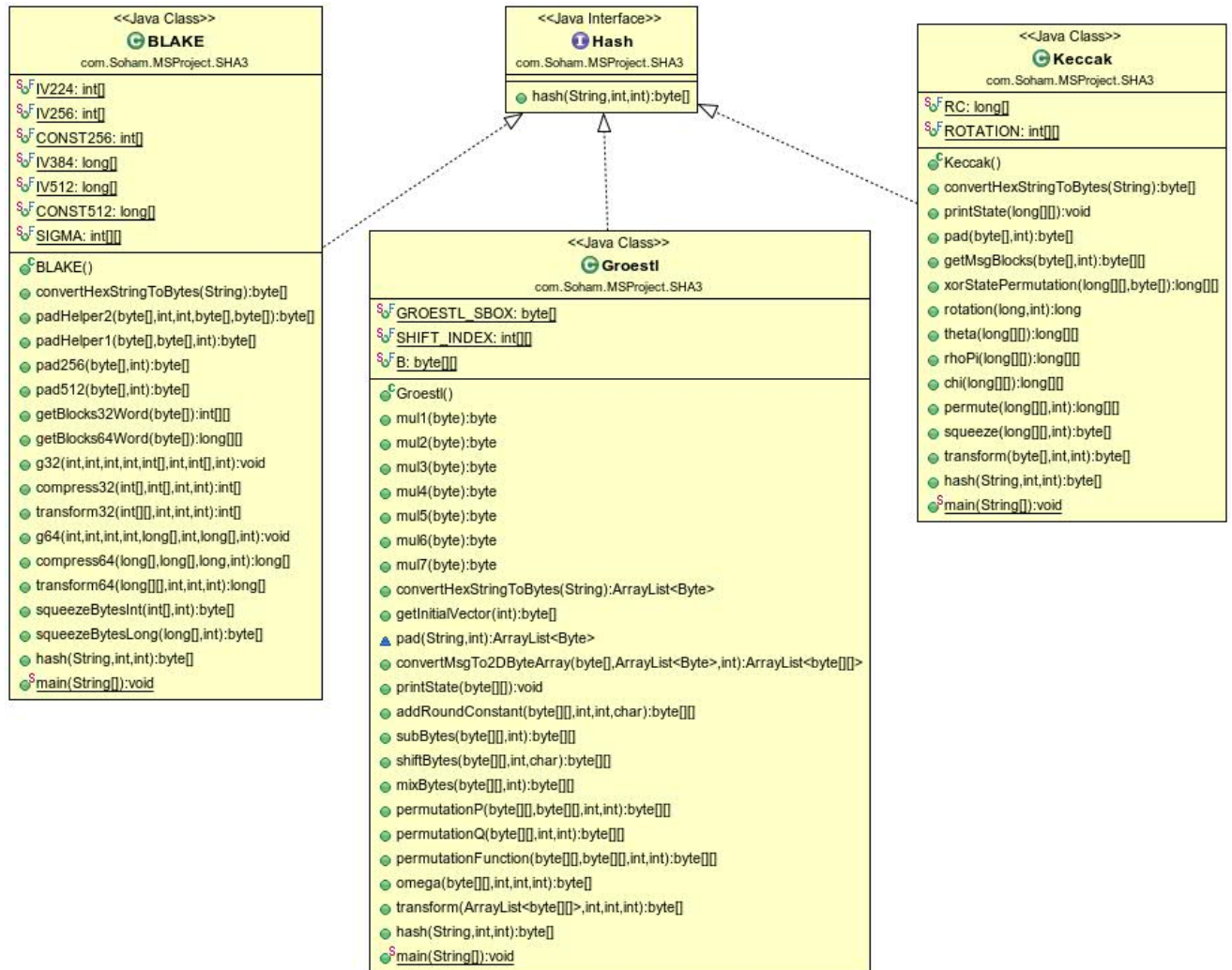


Figure 5.3: Class diagram of the classes for the 3 hash functions implemented.

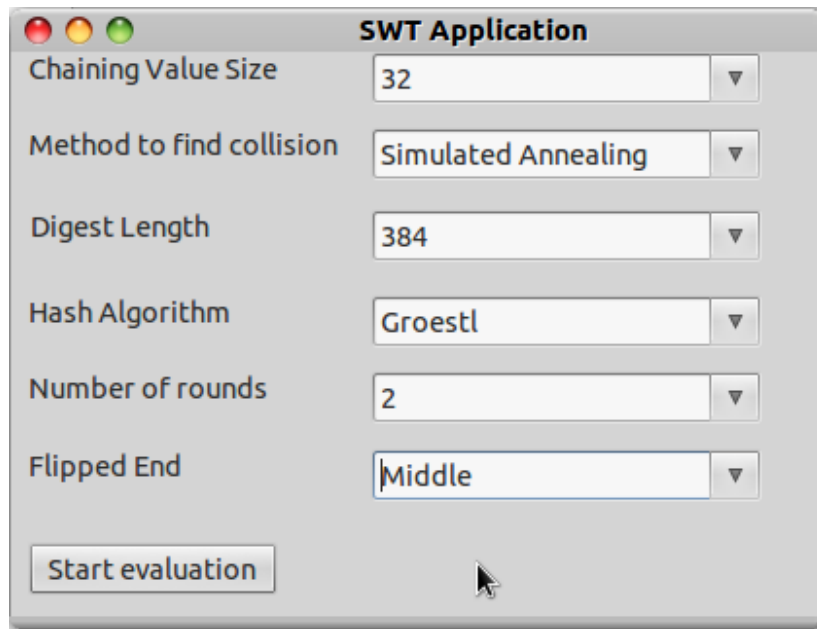


Figure 5.4: Class diagram of the classes for the 3 hash functions implemented.

4. The SHA-3 finalist algorithm, for hashing.
5. Number of permutation rounds, for the hashing algorithm.
6. The input case from start, middle or end, that you want to experiment with.

The classes that instantiate the experiment, and their relation with the GUI class is shown in figure 5.5. It is the responsibility of the "Experiment" class, to go over the input files of the said category, and then provide the collision finding algorithms with the input message pairs, and instantiating them.

The Experiment class is the one that calls FindCollision algorithm, which is the parent class, that generates the random chaining value as per the bit length provided. Evaluates the cost function, and generating neighbourhood solution chaining values for the collision finding purpose. The relationship between the classes is shown in the UML class diagram shown in figure 5.6.

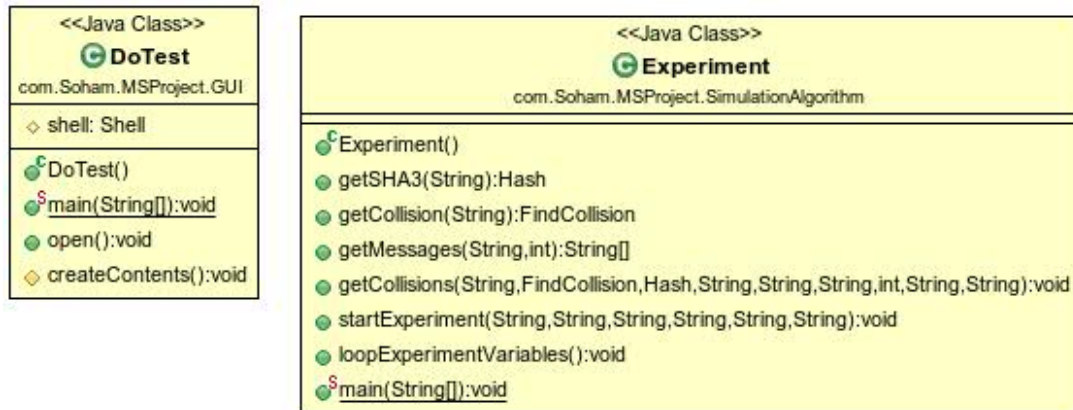


Figure 5.5: Class diagram of the GUI class and the initiation of experiment.

## 5.2.4 Testing the implemented code

Code implementing hashing involves permuting and substituting large number of bits, resulting in subtle bugs to magnify and distort the output results. As a consequence we have written, unit test cases to make sure, that our code works in individual and parts, so that the correctness of the results can be guaranteed with absence of bugs or bias in the code. InputTest class is created to make sure, that the toggling of the bits is done as expected. ExperimentTest class makes sure, that the correct parameters are called and passed like the correct hashing algorithm, with the expected digest size and rounds etc. The SHA3Test class makes sure that the implementation of all the SHA-3 finalist create the right message digests as expected. Finally the test class HillClimbTest makes sure that neighbourhood test for the chaining value is done properly. The class diagram for the test is shown in figure 5.7.

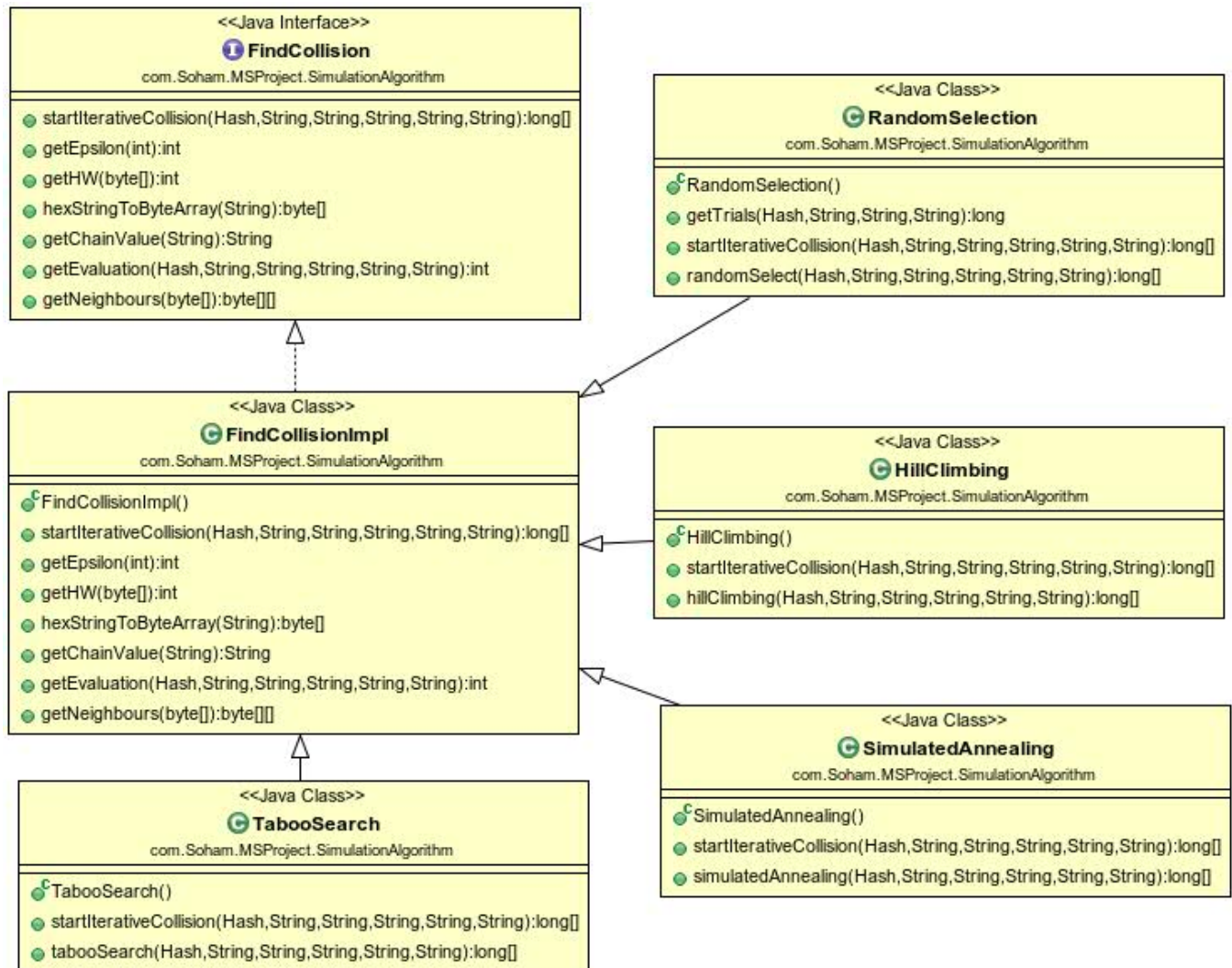


Figure 5.6: Class diagram of the classes for finding collisions.



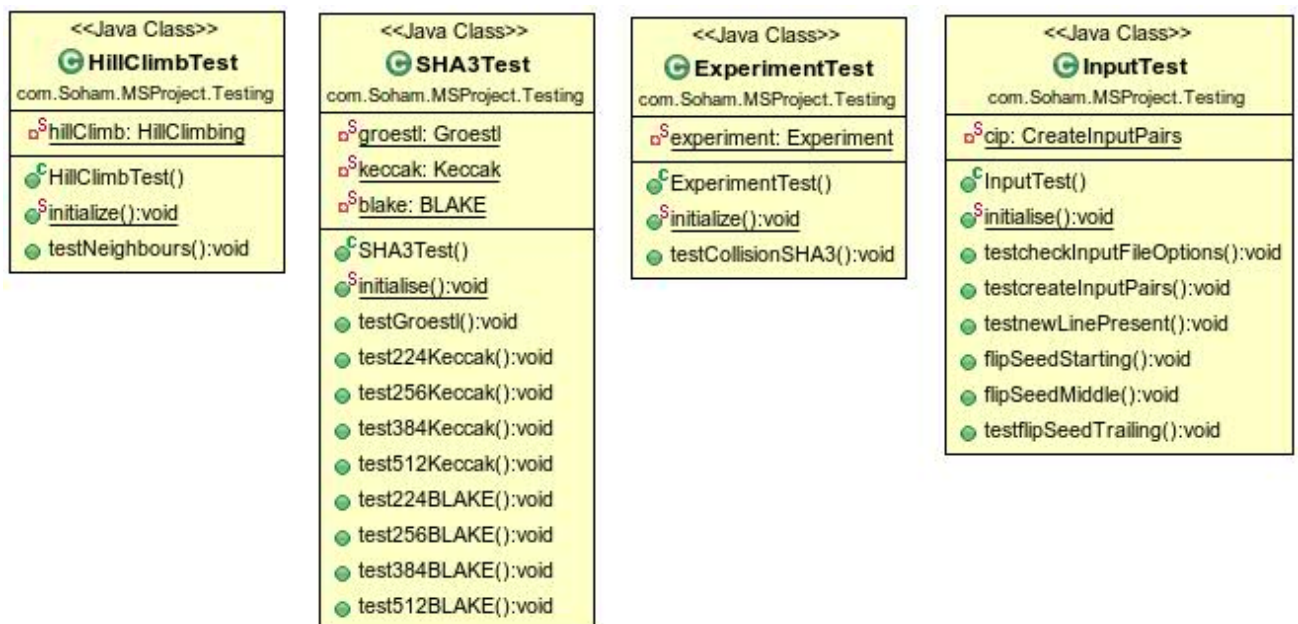


Figure 5.7: Class diagram of the classes used for testing, using JUnit4.

# Chapter 6

## Observations

The experiment first subjected all the 3 SHA-3 finalist algorithm to hill climbing, simulated annealing, and random selection; for permutation rounds of 1, 2, 3. The chaining value length was kept at 32 bits; and all the possible four message digest sizes were evaluated. Tabu search algorithm was applied only to Grøstl for rounds 1 and 2, for digest size of 224 bits. At first time, each collision finding algorithm was given 128 trials of experiment, for each hashing algorithm for a particular digest size and permutation round and input class. The experiment, was again conducted as mentioned; but with chaining value of length 64 bits. Since tabu search had previously proved to be expensive, with absence of encouraging results; it was discontinued. The experiment on random selection was not done, since hill climbing and simulated annealing took much longer time, than it took for the chaining value of 64 bits, and even that was discontinued. Since the results from the 64 bit length chaining value weren't that encouraging, hence we stuck to chaining value of bit length 32. We this time repeated our experiment, but with 256 trials instead of 128 trials, and experimented with rounds of 3 and 4. All the following observations and numbers can be found on the git repository for the source code implementation of the experiment.

### 6.1 Observations on programming implementations

Java comes as not the best choice in implementing algorithms related to hashing. Although since this experiment required GUI, and platform compatibility which made Java ideal tool. The SHA-3 finalist algorithms have word size or state size parameters of 32 or 64

bits, which exploit the computer architecture of having these word sizes built in or used, and using them as single data entities. In languages like C, there are primitive data types of size 4 bytes or 8 bytes, which can be easily integrated into the hashing algorithm. Java also has similar word size, but by default the data type is always signed. This creates problems while bit shifting or implementing some of the modulus operations. The leading bit is zero and during bit shifts if the bit is toggled to 1, then the data in it will be treated as a negative number, due to which some of the optimizations cannot be performed easily.

While implementing Keccak care should be taken, that bytes should be arranged in little endian rather than the usual big endian format. This makes applying the permutation on the lanes straightforward mapping of algebraic functions to Java.

## 6.2 Average iterations

Table 6.1 to 6.3 show the average iterations for the hill climbing algorithm applied to the SHA-3 finalist algorithm, for a chaining value of length 32 bits. For most of the algorithms except for the BLAKE for 1 permutation round the number of iterations increase as the digest size is increased. Also the number of iterations also seem to increase when number of rounds are increased, but at round 3 and 4 they are almost similar. For 32 bit chaining value the number of iterations for hill climbing algorithm applied to all SHA-3 finalist algorithm, remains around 900. For Keccak in hill climbing, when the permutation round was set at 1, the average iterations it took was around 532, lower than other SHA-3 finalist for same number of rounds. The number of iterations for random selection, simulated annealing for 32 bit chaining value were fixed at 1024 iterations. For the 64 bit length chaining value, for the simulated annealing experiment the number of iterations were set to 11 times the number of the digest size in bit length. So for digest length 224, 256, 384 and 512 we set iterations at 2464, 2816, 4224, 5632 respectively.

The average iterations for tabu search for Grøstl message for digest length 224 and permutation round 1 and 2 was 335254.443. Also there were no collisions found with this method. Being so computationally expensive and with no positive results, the experiment

Digest Size	Rounds			
	1	2	3	4
224	803.6331	891.3276	886.56465	883.1557
256	804.5944	891.35455	892.13813	885.4943
384	1137.8868	898.2596	899.3996	902.6081
512	1118.9453	902.0659	903.68877	906.0339

Table 6.1: Average iterations over all input cases for Hill Climbing for BLAKE for chaining value of bit length 32

Digest Size	Rounds			
	1	2	3	4
224	875.893	888.3345	886.71898	885.8383
256	889.72437	887.639	891.2398	889.7382
384	872.1038	897.5908	898.66531	897.1204
512	896.17804	904.18335	904.1414	902.55646

Table 6.2: Average iterations over all input cases for Hill Climbing for Grøstl for chaining value of bit length 32

with tabu search was discontinued.

As the number of bits in chaining value is increased from 32 to 64, the number of iterations increases to around 3500. The increase in number of iterations is more than three and half times, than what is observed in 32 bit chaining value.

## 6.3 Near collisions

If the collision algorithm was able to find a chaining value by iteratively going through its neighbourhood, that got the two message digests to agree on more than 65% of the bits, then it was noted as a collision. In the tables starting from 6.7 to 6.something we have shown the number of input pairs in a input case of start, middle or end that showed collision, and the maximum number of trials in a pair that collision was obtained. In the tables below, the letters S, M, E stand for start, middle and end input case. The number of

	Rounds			
Digest Size	1	2	3	4
224	532.2517	880.6999	883.23035	888.5428
256	532.341	888.98126	889.24855	895.0305
384	533.9329	892.61993	899.02125	904.8013
512	533.70807	900.46466	905.87	902.2507

Table 6.3: Average iterations over all input cases for Hill Climbing for Keccak for chaining value of bit length 32

	Rounds		
Digest Size	1	2	3
224	4190.473	3465.0413	3483.752
256	4264.3496	3456.9246	3431.7415
384	3885.6829	3515.1746	3528.4922
512	3984.2366	3535.423	3559.4246

Table 6.4: Average iterations over all input cases for Hill Climbing for BLAKE for chaining value of bit length 64

	Rounds		
Digest Size	1	2	3
224	3687.7493	3468.3171	3469.366
256	3714.0242	3479.8074	3473.6194
384	3594.1716	3522.1423	3512.8254
512	3581.0823	3544.302	3559.9834

Table 6.5: Average iterations over all input cases for Hill Climbing for Grøstl for chaining value of bit length 64

Digest Size	Rounds		
	1	2	3
224	2118.8818	3535.5684	3457.5586
256	2118.3499	3557.9949	3466.1624
384	2134.1736	3676.9707	3521.3984
512	2139.1523	3764.485	3562.9753

Table 6.6: Average iterations over all input cases for Hill Climbing for Keccak for chaining value of bit length 64

cases of collision, followed by maximum number of trials in a case where a collision was obtained, and separated by slash are entered into the tables. For tabu search there were no collisions found in two rounds of Grøstl that it was implemented on.

### 6.3.1 Hill Climbing

In the collision tables for 32 bit chaining value trials for round 1 and 2 were 128, while for round 3 and 4 were 256. So the numbers for round 3 and 4 should reflect more accuracy. Although there would be very little inconsistency when comparing numbers from round 1 and 2; and round 3 and 4. Since, round 1 and 2 show more near collision figures as compared to rounds 3 and 4. For the 64 bit chaining value, the number of trials was set at 128.

For BLAKE in round 1, collisions can be found in each instance and trial. However for rounds 2, 3, 4 the message pairs that show collision and for trials reduced to less than 20, for digest lengths of 224 and 256. For digest lengths of 384, 512 in higher number of rounds, collisions are not obtained. The aforementioned observations hold true for chaining values of lengths 32 and 64 both.

Grøstl follows a similar pattern with as that of BLAKE, though for digest size 256 and round 1 it seems to be less collision prone, but experimental factors like choosing a chaining value could be a reason for the outlier. As the digest size increases with the number of rounds, the collision decreases. Rounds 3 and 4 have no collision for BLAKE

Digest size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	2/50	19/128	20/128	2/2	8/1	3/2	8/3	12/2	14/3	12/2	7/2	6/1
256	0/0	19/128	20/128	1/1	1/1	2/1	5/2	2/1	3/1	1/1	0/0	2/1
384	18/128	20/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0
512	18/128	20/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.7: Collisions and maximum trials a input pair had collision for BLAKE with Hill Climbing algorithm for 32 bit chaining value.

Digest Size	Rounds								
	1			2			3		
	S	M	E	S	M	E	S	M	E
224	16/128	20/128	20/128	10/3	16/4	12/6	15/3	12/3	16/4
256	15/128	20/128	20/128	5/2	4/1	0/0	3/1	8/2	5/1
384	20/128	20/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0
512	19/128	20/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.8: Collisions and maximum trials a input pair had collision for BLAKE with Hill Climbing algorithm for 64 bit chaining value.

Digest size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	9/4	1/2	0/0	4/2	6/2	6/2	11/2	10/3	8/2	13/3	9/2	11/4
256	1/2	0/0	0/0	3/1	1/1	2/1	3/1	4/1	3/1	2/1	2/1	3/1
384	0/0	0/0	12/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0
512	8/17	0/0	14/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.9: Collisions and maximum trials a input pair had collision for Grøstl with Hill Climbing algorithm for 32 bit chaining value.

Digest Size	Rounds								
	1			2			3		
	S	M	E	S	M	E	S	M	E
224	17/10	16/6	10/7	13/4	10/4	17/4	15/4	18/4	16/3
256	9/5	12/3	4/2	4/2	3/1	7/2	7/1	4/1	4/1
384	1/128	0/0	13/128	0/0	0/0	0/0	0/0	0/0	0/0
512	8/30	0/0	20/128	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.10: Collisions and maximum trials a input pair had collision for Grøstl with Hill Climbing algorithm for 64 bit chaining value.

at digest sizes greater than 256.

Keccak from the numbers seems to be least effective if only single or double permutation rounds are used. Collision can be found in Keccak even for 3 rounds, at digest size 384. Keccak only seems secure after getting more than 2 rounds for digest size 512 or rounds more than 3 for digest size 384.

### 6.3.2 Simulated Annealing

Similar to Hill Climbing in simulated annealing, collision tables for 32 bit chaining value trials for round 1 and 2 were from 128 trials, while for round 3 and 4 were 256 trials. The number of instances and trials where collisions were found in simulated annealing as compared to those in hill climbing, are reduced. With higher digest sizes higher than



Size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	all	all	all	all	all	all	11/3	12/3	17/5	12/3	11/3	10/2
256	all	all	all	all	all	all	2/1	4/1	5/2	4/1	1/1	3/1
384	all	all	all	all	all	all	0/0	0/0	1/1	0/0	0/0	0/0
512	all	all	all	all	all	all	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.11: Collisions and maximum trials a input pair had collision for Keccak with Hill Climbing algorithm for 32 bit chaining value. The word "all" stands for number 20/128.

Digest Size	Rounds								
	1			2			3		
	S	M	E	S	M	E	S	M	E
224	20/128	20/128	20/128	20/128	20/128	20/128	18/7	18/8	19/6
256	20/128	20/128	20/128	20/128	20/128	20/128	7/3	7/2	6/2
384	20/128	20/128	20/128	20/128	20/128	20/128	1/1	0/0	0/0
512	20/128	20/128	20/128	20/128	20/128	20/128	0/0	0/0	0/0

Table 6.12: Collisions and maximum trials a input pair had collision for Keccak with Hill Climbing algorithm for 64 bit chaining value.

Digest size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	0/0	19/128	20/128	0/0	0/0	0/0	1/1	0/0	0/0	0/0	1/1	0/0
256	0/0	18/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0	1/1	0/0	1/1
384	10/116	9/109	20/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0
512	14/110	16/104	20/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.13: Collisions and maximum trials a input pair had collision for BLAKE with Simulated Annealing algorithm for 32 bit chaining value.

Digest size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	1/1	0/0	0/0	1/1	0/0	0/0	1/1	1/1	0/0	0/0	0/0	0/0
256	0/0	0/0	0/0	0/0	0/0	0/0	1/1	0/0	0/0	0/0	1/1	0/0
384	0/0	0/0	12/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0
512	1/1	0/0	14/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.14: Collisions and maximum trials a input pair had collision for Grøstl with Simulated Annealing algorithm for 32 bit chaining value.

256 and rounds higher than 2, simulated annealing similarly failed to find any success. However for hash algorithms with rounds 3 and 4 with digests sizes 224 and 256; random search seems to perform better than simulated annealing.

### 6.3.3 Random Selection

Random selection was not applied with chaining value of 64 bit length, after not so optimum results obtained from its' use in hill climbing algorithm; and simulated annealing applied on Grøstl. Random selection also had similar success like hill climbing and simulated annealing for digests obtained from rounds above 2, and digests lengths above 256.

Digest size	Rounds								
	1			2			3		
	S	M	E	S	M	E	S	M	E
224	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0
256	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0
384	1/128	0/0	12/128	0/0	0/0	0/0	0/0	0/0	0/0
512	3/1	16/104	16/128	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.15: Collisions and maximum trials a input pair had collision for Grøstl with Simulated Annealing algorithm for 64 bit chaining value.

Size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	all	all	all	all	all	all	0/0	0/0	0/0	0/0	0/0	0/0
256	all	all	all	all	all	all	0/0	0/0	0/0	0/0	1/1	0/0
384	all	all	all	all	all	all	0/0	0/0	0/0	0/0	0/0	0/0
512	all	all	all	all	all	all	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.16: Collisions and maximum trials a input pair had collision for Keccak with Simulated Annealing algorithm for 32 bit chaining value. "all" stands for 20/128.

Digest size	Rounds								
	1			2			3		
	S	M	E	S	M	E	S	M	E
224	20/128	20/128	20/128	20/128	20/128	20/128	0/0	1/1	0/0
256	20/128	20/128	20/128	20/128	20/128	20/128	0/0	0/0	0/0

Table 6.17: Collisions and maximum trials a input pair had collision for Keccak with Simulated Annealing algorithm for 64 bit chaining value.

Size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	2/35	19/128	20/128	3/2	4/2	3/1	7/1	10/3	7/2	11/2	8/3	9/1
256	0/0	19/128	20/128	1/1	2/1	0/0	3/1	1/1	0/0	4/1	1/1	2/1
384	18/128	19/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0
512	18/128	20/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.18: Collisions and maximum trials a input pair had collision for BLAKE with random selection algorithm for 32 bit chaining value.

Size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	6/3	1/3	0/0	2/1	8/2	5/1	9/4	8/1	9/2	7/2	6/3	7/2
256	0/0	0/0	0/0	0/0	0/0	0/0	2/1	1/1	0/0	1/1	1/1	1/1
384	0/0	0/0	12/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0
512	8/20	0/0	14/128	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.19: Collisions and maximum trials a input pair had collision for Grøstl with random selection algorithm for 32 bit chaining value.

Size	Rounds											
	1			2			3			4		
	S	M	E	S	M	E	S	M	E	S	M	E
224	20/128	20/128	20/128	20/128	20/128	20/128	4/2	5/1	6/2	0/0	0/0	0/0
256	20/128	20/128	20/128	20/128	20/128	20/128	3/2	4/1	6/1	3/1	2/1	3/1
384	20/128	20/128	20/128	20/128	20/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0
512	20/128	20/128	20/128	20/128	20/128	20/128	0/0	0/0	0/0	0/0	0/0	0/0

Table 6.20: Collisions and maximum trials a input pair had collision for Keccak with random selection algorithm for 32 bit chaining value.

# Chapter 7

## Conclusions and future work

### 7.1 Validity of the hypothesis

In section 4.4 two hypothesis were suggested, both have been sort of some what disproved. When the reduced version of an algorithm is only done by reducing the number of rounds, then for just rounds 1 and 2, Keccak performs worse than BLAKE and Grøstl. Also from the numbers and taking into account the number of success in getting collisions, hill climbing performs better than simulated annealing and tabu search. Random selection also performs, some what comparable to simulated annealing, but not much. Suggesting that for reduced rounds, that random selection could also be used to find near collisions.

#### 7.1.1 Collision resistance of SHA-3 finalist in reduced rounds

For 1 round permutation, all the algorithms should almost equivalent weakness, that is collision in almost all instances that is message pairs, and trial. However for Keccak for the three collision finding algorithms hill climbing, simulated annealing, and random selection; collisions were seen in every instance and trial. This behaviour is again noticed in Keccak for permutation of round 2; however for BLAKE and Grøstl, this number of trials and instances in which collision could be found are reduced.

However as the permutation rounds are increased to 3 and 4, the resistance to collision finding algorithms for Keccak becomes equivalent to that of BLAKE and Grøstl. For digest sizes, of 224 and 256; some message pairs do show collision in trials less than 10. But for digest sizes of 384 and 512; no collisions are found.

### 7.1.2 Feasibility of the collision algorithms

From the collision algorithms; hill climbing seems to be the most optimised in terms of collision finding to time required ratio. Tabu search performs the worst, in terms of collision to CPU time required. For example using tabu search on Grøstl for message digest of 224 bits, and rounds 1 and 2; tabu search required around 334894 iterations on average, without any collisions. For the same conditions hill climbing did not require more than 900 iterations, with some success.

At lower permutation rounds like 1 and 2, for Keccak; and for round 1 for BLAKE and Grøstl random selection is as effective as hill climbing and simulated annealing. Even for higher rounds like 3 and 4; for lower digest sizes of 224 and 256, random selection was able to find collision amongst all three SHA-3 candidate algorithms in around 10 instances in trials less than 10. This performance is comparable to that of simulated annealing.

Tabu search performs poorly, given that it goes for the steepest descent approach. Tabu search tries to pick the best chaining value from the neighbourhood, where the k-bit neighbourhood, whose size is roughly equal to  $n^2$  where n is the bit length of the chaining value. Also the tabu search has to go through elements of neighbourhood comparing each of them with the tabu list, before inserting them into the candidate set where they will be considered as solution. This places additional overhead on the algorithm performance. The tabu list can be made up of states, solutions previously found, moves previously made. For our case the best element for tabu list would have been elements from previous iteration that were not selected, but this still makes our set large. Although the process can be speed up, by getting the best element from neighbourhood during neighbourhood creation, but you still have the over head of comparing the elements twice for the best element and the comparing with tabu list. The problem does not mould well, with the tabu search solution, because trying to store state, previous solution, or moves in tabu list are all expensive.

Simulated annealing though trying to avoid the local minima, does not perform any better than hill climbing, when finding near collisions. Over trials over 256 and 128, and

chaining values 32 and 64; simulated annealing found far fewer collisions than hill climbing. While hill climbing was able to find collision in instances in more than 50% of the message pairs for each of the SHA-3 algorithm for message digest of 224 and 256, for permutation rounds of 3 and 4. But simulated annealing is only able to find collision in 1 instance out of 60 for 1 trial.

## 7.2 Effect of digest size

It seems as the digest size is increased, the near collision resistance increases, only if the rounds are above 2. The input text here "The quick brown fox jumps over the lazy dog" was of 344 bits, and as per table 7.1 even with 32 bit chaining value, the entire message is treated in one part itself, as seen from table 7.1. Thus our modifications to the message pairs do make it through in one go, rather than blocks of message from the pairs which have all bits same.

The reason higher message digests seem to have more collision resistance may be due to the fact that they have more bits to shuffle around, thus making it harder for simple algorithms without any heuristics to find the collisions. Note should be made that near collisions in this case is done when bits matching are more than 65% of the digest length. So it could be that bits upto 64% might match and yet be rejected not as near collision found. Also the higher digest sizes like 384 and 512 only show collision resistance only if applied with rounds 3 and 4.

Thus the minimum security recommendation for collision avoidance would be go for digest sizes of 384 and above.

## 7.3 Effects of the number of rounds

From the observation of tables 6.7 to 6.20, we can conclude that higher the number of permutation rounds, the more is the collision resistance. Rounds 1 and 2, can be considered as not secure. It should be noted that for permutation rounds 1 and 2; Keccak seems to

Hash algorithm	Digest size			
	224	256	384	512
BLAKE	512	512	1024	1024
Grøstl	512	512	1024	1024
Keccak	1152	1088	832	576

Table 7.1: Number of input bits to one function block, in the respective SHA-3 finalist algorithm

Hash algorithm	Digest size			
	224	256	384	512
BLAKE	14	14	16	16
Grøstl	10	10	14	14
Keccak	24	24	24	24

Table 7.2: Number of permutation rounds, in the respective SHA-3 finalist algorithm

perform poorly against BLAKE and Grøstl. But as rounds are increased to 3 and 4, the performance of Keccak with respect to other algorithms is comparable.

On the basis of observation it can be stated that, at minimum 3 rounds of permutation are required, and 4 to be secure as per the results of this experiment with limited resources. Please note that as per table 7.2, the recommended number of permutation rounds for each SHA-3 finalist algorithm is different. While Grøstl has the least number of rounds, Keccak has most number of rounds. This could be one explanation of Grøstl and BLAKE's, better performance though not most ideal; than Keccak for rounds 1 and 2.

## 7.4 Chaining value length

Longer chaining value do not seem to help in getting near collisions. For perfect collisions, longer chaining values might be helpful. However for near collisions, shorter chaining values work fine. With larger chaining values, the number of times the algorithm has to loop through also increases. For increase in chaining value from 32 bits to 64 bits, the loop



execution increased by factor of 3 times more than what it required for 32 bits.

## 7.5 Bit differences in message in particular positions

All the SHA-3 algorithms displayed lack of diffusion property if any, irrespective of input bit updated in any of the three positions start, middle and end specified in the message input string. The bit update made to middle, and end to the input message are almost equivalent, since, the input message is padded with the chaining value. Thus we do not experiment with the bits updated at the end of the input string, only with bits updated at start of the string or in the middle of the string. Since in the experiment trials the chaining value is chosen randomly, hence small differences in finding collisions amongst the input, for a given hashing algorithm, digest size and round can be ignored. Based on this fact, it seems that diffusion in the given algorithm is place agnostic.

## 7.6 Future work

1. This experiment can be repeated with the two other finalist algorithm Skein and JH, alongwith Keccak.
2. Instead of reducing the number of rounds, the reduced version of the finalist algorithm could be obtained by reducing the state size, and repeating this experiment.
3. Experiments can also be done with other combinations of factors of hashing algorithms, like reducing the digest size, but increasing the number of rounds.
4. In this experiment, the near collision was benchmarked at 65% of the output bits match. Instead further work can be done to find the effort required to find perfect collisions for message pairs for reduced rounds for certain hashing algorithms.

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