

Algoritmi trattati a lezione

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```

proc InsertionSort (A)
  for (j = 2 to A.length)
    {
      key = A[j]
      i = j - 1
      while ((i > 0) and (A[i] > key))
        {
          A[i + 1] = A[i]
          i = i - 1
        }
      A[i + 1] = key
    }

```

```

proc RecursiveBinarySearch (A, low, high, k)
  if (low > high)
    then return nil
  mid = (high + low)/2
  if (A[mid] = k)
    then return mid
  if (A[mid] < k)
    then return RecursiveBinarySearch(A, mid + 1, high, k)
  if (A[mid] > k)
    then return RecursiveBinarySearch(A, low, mid - 1, k)

```

```

proc SelectionSort (A)
  for (j = 1 to A.length - 1)
    {
      min = j
      for (i = j + 1 to A.length)
        if (A[i] < A[min])
          then min = i
      SwapValue(A, min, j)
    }

```

```

proc Merge (A, p, q, r)
  n1 = q - p + 1
  n2 = r - q
  let L[1, ..., n1] and R[1, ..., n2] be new array
  for (i = 1 to n1) L[i] = A[p + i - 1]
  for (j = 1 to n2) R[j] = A[q + j]
  i = 1
  j = 1
  for (k = p to r)
    if (i ≤ n1)
      then
        if (j ≤ n2)
          then
            if (L[i] ≤ R[j])
              then CopyFromL(i)
            else CopyFromR(j)
          else CopyFromL(i)
        else CopyFromR(j)

```

```

proc MergeSort (A, p, r)
  if (p < r)
    then
      {
        q = [(p + r)/2]
        MergeSort(A, p, q)
        MergeSort(A, q + 1, r)
        Merge(A, p, q, r)
      }

```

```

proc Partition (A, p, r)
  x = A[r]
  i = p - 1
  for (j = p to r - 1)
    if (A[j] ≤ x)
      then
        {
          i = i + 1
          SwapValue(A, i, j)
        }
  SwapValue(A, i + 1, r)
  return i + 1

```

```

proc QuickSort (A, p, r)
  if (p < r)
    then
      {
        q = Partition(A, p, r)
        QuickSort(A, p, q - 1)
        QuickSort(A, q + 1, r)
      }

```

```

proc RandomizedPartition (A, p, r)
  s = p ≤ Random() ≤ r
  SwapValue(A, s, r)
  x = A[r]
  i = p - 1
  for j = p to r - 1
    if (A[j] ≤ x)
      then
        {
          i = i + 1
          SwapValue(A, i, j)
        }
  SwapValue(A, i + 1, r)
  return i + 1

```

```

proc RandomizedQuickSort (A, p, r)
  if (p < r)
    then
      {
        q = RandomizedPartition(A, p, r)
        RandQuickSort(A, p, q - 1)
        RandQuickSort(A, q + 1, r)
      }

```

```

proc Empty (S)
  if (S.top = 0)
    then return true
  return false

```

```

proc Push (S, x)
  if (S.top = S.max)
    then return "overflow"
  S.top = S.top + 1
  S[S.top] = x

```

```

proc Pop (S)
  if (Empty(S))
    then return "underflow"
  S.top = S.top - 1
  return S[S.top + 1]

```

```

proc Enqueue (Q, x)
  if (Q.dim = Q.length)
    then return "overflow"
  Q[Q.tail] = x
  if (Q.tail = Q.length)
    then Q.tail = 1
  else Q.tail = Q.tail + 1
  Q.dim = Q.dim + 1

```

```

proc Dequeue (Q)
  if (Q.dim = 0)
    then return "underflow"
  x = Q[Q.head]
  if (Q.head = Q.length)
    then Q.head = 1
  else Q.head = Q.head + 1
  Q.dim = Q.dim - 1
  return x

```

```

proc Empty (S)
{
  if (S.head = nil)
  then return true
  return false
}

```

```

proc Push (S, x)
{
  Insert(S, x)
}

```

```

proc Pop (S)
{
  if (Empty(S))
  then return "underflow"
  x = S.head
  Delete(S, x)
  return x.key
}

```

```

proc Empty (Q)
{
  if (Q.head = nil)
  then return true
  return false
}

```

```

proc Enqueue (Q, x)
{
  Insert(Q, x)
}

```

```

proc Dequeue (Q)
{
  if (Empty(Q))
  then return "underflow"
  x = Q.tail
  Q.tail = x.prev
  Delete(Q, x)
  return x.key
}

```

```

proc Enqueue (A, i, priority)
{
  if (i > A.length)
  then return "overflow"
  A[i] = priority
}

```

```

proc DecreaseKey (A, i, priority)
{
  if ((A[i] < priority) or (A[i].empty = 1))
  then return "error"
  A[i] = priority
}

```

```

proc ExtractMin (A)
{
  MinIndex = 0
  MinPriority = ∞
  for (i = 1 to A.length)
  {
    if ((A[i] < MinPriority) and (A[i].empty = 0))
    then
      {
        MinPriority = A[i]
        MinIndex = i
      }
  }
  if (MinIndex = 0)
  then return "underflow"
  A[MinIndex].empty = 1
  return MinIndex
}

```

```

proc Parent (i)
{
  return ⌊ i/2 ⌋
}

```

```

proc Left (i)
{
  return 2 · i
}

```

```

proc Right (i)
{
  return 2 · i + 1
}

```

```

proc MinHeapify (H, i)
{
  l = Left(i)
  r = Right(i)
  smallest = i
  if ((l ≤ H.heapsize) and (H[l] < H[i]))
  then smallest = l
  if ((r ≤ H.heapsize) and (H[r] < H[smallest]))
  then smallest = r
  if smallest ≠ i
  then
    {
      SwapValue(H, i, smallest)
      MinHeapify(H, smallest)
    }
}

```

```

proc BuildMinHeap (H)
{
  H.heapsize = H.length
  for (i = ⌊ H.length/2 ⌋ downto 1) MinHeapify(H, i)
}

```

```

proc HeapSort (H)
{
  BuildMaxHeap(H)
  for (i = H.length downto 2)
  {
    {
      SwapValue(H, i, 1)
      H.heapsize = H.heapsize - 1
      MaxHeapify(H, 1)
    }
}

```

```

proc Enqueue (Q, priority)
{
  if (Q.heapsize = Q.length)
  then return "overflow"
  Q.heapsize = Q.heapsize + 1
  Q[heapsize] = ∞
  DecreaseKey(Q, Q.heapsize, priority)
}

```

```

proc DecreaseKey (Q, i, priority)
{
  if (priority > Q[i])
  then return "error"
  Q[i] = priority
  while ((i > 1) and (Q[Parent(i)] > Q[i]))
  {
    SwapValue(Q, i, Parent(i))
    i = Parent(i)
  }
}

```

```

proc ExtractMin (Q)
{
  if (Q.heapsize < 1)
  then return "underflow"
  min = Q[1]
  Q[1] = Q[Q.heapsize]
  Q.heapsize = Q.heapsize - 1
  MinHeapify(Q, 1)
  return min
}

```

```

proc Enqueue (Q, i, priority)
{
  if (i > Q.length)
  then return "overflow"
  Q[i] = priority
}

```

```

proc DecreaseKey (Q, i, priority)
{
  if ((Q[i] < priority) or (Q[i].empty = 1))
  then return "error"
  Q[i] = priority
}

```

```

proc ExtractMin (Q)
{
  MinIndex = 0
  MinPriority = ∞
  for (i = 1 to Q.length)
  {
    if ((Q[i] < MinPriority) and (Q[i].empty = 0))
    then
      {
        MinPriority = Q[i]
        MinIndex = i
      }
  }
  if (MinIndex = 0)
  then return "underflow"
  Q[MinIndex].empty = 1
  return MinIndex
}

```

```

proc CountingSort (A, B, k)
{
  let C[0, ..., k] new array
  for (i = 0 to k) C[i] = 0
  for (j = 1 to A.length) C[A[j]] = C[A[j]] + 1
  for (i = 1 to k) C[i] = C[i] + C[i - 1]
  for (j = A.length downto 1)
  {
    B[C[A[j]]] = A[j]
    C[A[j]] = C[A[j]] - 1
  }
}

```

```

proc RadixSort (A, d)
{
  for (i = 1 to d) AnyStableSort(A) on digit i
}

```

```

proc ListInsert (L, x)
{
  x.next = L.head
  if (L.head ≠ nil)
  then L.head.prev = x
  L.head = x
  x.prev = nil
}

```

```

proc ListSearch (L, k)
{
  x = L.head
  while (x ≠ nil) and (x.key ≠ k) x = x.next
  return x
}

```

```

proc ListDelete (L, x)
{
  if (x.prev ≠ nil)
  then x.prev.next = x.next
  else L.head = x.next
  if (x.next ≠ nil)
  then x.next.prev = x.prev
}

```

```

proc MakeSet (S, S, x, i)
{
  S[i].set = x
  S.head = x
  S.tail = x
}

```

```

proc Union (x, y)
{
  S1 = x.head
  S2 = y.head
  if (S1 ≠ S2)
  then
    {
      S1.tail.next = S2.head
      z = S2.head
      while (z ≠ nil)
      {
        z.head = S1
        z = z.next
      }
      S1.tail = S2.tail
    }
}

```

```

proc FindSet (x)
{ return x.head.head }

```

```

proc MakeSet (S, S, x, i)
{
  S[i].set = x
  S.head = x
  S.tail = x
  S.rank = 0
}

```

```

proc Union (x, y)
{
  S1 = x.head
  S2 = y.head
  if (S1 ≠ S2)
  then
    {
      if (S2.rank > S1.rank)
      then
        {
          Stemp = S1
          S1 = S2
          S2 = Stemp
        }
      S1.tail.next = S2.head
      z = S2.head
      while (z ≠ nil)
      {
        z.head = S1
        z = z.next
      }
      S1.tail = S2.tail
      S1.rank = S1.rank + S2.rank
    }
}

```

```

proc MakeSet (x)
{
  x.p = x
  x.rank = 0
}

```

```

proc Union (x, y)
{
  x = Findset(x)
  y = Findset(y)
  if (x.rank > y.rank)
  then y.p = x
  if (x.rank ≤ y.rank)
  then
    {
      x.p = y
      if (x.rank = y.rank)
      then y.rank = y.rank + 1
    }
}

```

```

proc FindSet (x)
{
  if (x ≠ x.p)
  then x.p = FindSet(x.p)
  return x.p
}

```

```

proc TreeInOrderTreeWalk (x)
{
  if (x ≠ nil)
  then
    {
      TreeInOrderTreeWalk(x.left)
      Print(x.key)
      TreeInOrderTreeWalk(x.right)
    }
}

```

```

proc TreePreOrderTreeWalk (x)
{
  if (x ≠ nil)
  then
    {
      Print(x.key)
      TreeInOrderTreeWalk(x.left)
      TreeInOrderTreeWalk(x.right)
    }
}

```

```

proc TreePostOrderTreeWalk (x)
{
  if (x ≠ nil)
  then
    {
      TreeInOrderTreeWalk(x.left)
      TreeInOrderTreeWalk(x.right)
      Print(x.key)
    }
}

```

```

proc BSTTreeSearch (x, k)
{
  if ((x = nil) or (x.key = k))
  then return x
  if (k ≤ x.key)
  then return BSTTreeSearch(x.left, k)
  else return BSTTreeSearch(x.right, k)
}

```

```

proc BSTTreeMinimum (x)
{
  if ((x.left = nil))
  then return x
  return BSTTreeMinimum(x.left)
}

```

```

proc BSTTreeSuccessor (x)
{
  if (x.right ≠ nil)
  then return BSTTreeMinimum(x.right)
  y = x.p
  while ((y ≠ nil) and (x = y.right))
  {
    x = y
    y = y.p
  }
  return y
}

```

```

proc BSTTreeInsert (T, z)
{
  y = nil
  x = T.root
  while (x ≠ nil)
  {
    y = x
    if (z.key ≤ x.key)
    then x = x.left
    else x = x.right
  }
  z.p = y
  if (y = nil)
  then T.root = z
  if ((y ≠ nil) and (z.key ≤ y.key))
  then y.left = z
  if ((y ≠ nil) and (z.key > y.key))
  then y.right = z
}

```

```

proc BSTTreeDelete (T, z)
{
  if (z.left = nil)
  then BSTTransplant(T, z, z.right)
  if ((z.left ≠ nil) and (z.right = nil))
  then Transplant(T, z, z.left)
  if ((z.left ≠ nil) and (z.right ≠ nil))
  then
    {
      y = BSTTreeMinimum(z.right)
      if (y.p ≠ z)
      then
        {
          BSTTransplant(T, y, y.right)
          y.right = z.right
          y.right.p = y
        }
      BSTTransplant(T, z, y)
      y.left = z.left
      y.left.p = y
    }
}

```

```

proc BSTTreeTransplant (T, u, v)
{
  if (u.p = nil)
  then T.root = v
  if ((u.p ≠ nil) and (u = u.p.left))
  then u.p.left = v
  if ((u.p ≠ nil) and (u = u.p.right))
  then u.p.right = v
  if (v ≠ nil)
  then v.p = u.p
}

```

```

proc BSTTreeLeftRotate (T, x)
{
  y = x.right
  x.right = y.left
  if (y.left ≠ T.nil)
  then y.left.p = x
  y.p = x.p
  if (x.p = T.nil)
  then T.root = y
  if ((x.p ≠ T.nil) and (x = x.p.left))
  then x.p.left = y
  if ((x.p ≠ T.nil) and (x = x.p.right))
  then x.p.right = y
  y.left = x
  x.p = y
}

```

```

proc RBTreeInsert (T, z)
{
  y = T.nil
  x = T.root
  while (x ≠ T.nil)
  {
    y = x
    if (z.key < x.key)
    {
      then x = x.left
    }
    else x = x.right
  }
  z.p = y
  if (y = T.nil)
  {
    then T.root = z
  }
  if ((y ≠ T.nil) and (z.key < y.key))
  {
    then y.left = z
  }
  if ((y ≠ T.nil) and (z.key ≥ y.key))
  {
    then y.right = z
  }
  z.left = T.nil
  z.right = T.nil
  z.color = RED
  RBTreeInsertFixup(T, z)
}

```

```

proc RBTreeInsertFixup (T, z)
{
  while (z.p.color = RED)
  {
    if (z.p = z.p.p.left)
    {
      then RBTreeInsertFixUpLeft(T, z)
    }
    else RBTreeInsertFixUpRight(T, z)
  }
  T.root.color = BLACK
}

```

```

proc RBTreeInsertFixupLeft (T, z)
{
  y = z.p.p.right
  if (y.color = RED)
  {
    then
    {
      z.p.color = BLACK
      y.color = BLACK
      z.p.p.color = RED
      z = z.p.p
    }
    else
    {
      if (z = z.p.p.right)
      {
        then
        {
          z = z.p
          LeftRotate(T, z)
        }
        z.p.color = BLACK
        z.p.p.color = RED
        TreeRightRotate(T, z.p.p)
      }
    }
  }
}

```

```

proc RBTreeInsertFixupRight (T, z)
{
  y = z.p.p.left
  if (y.color = RED)
  {
    then
    {
      z.p.color = BLACK
      y.color = BLACK
      z.p.p.color = RED
      z = z.p.p
    }
    else
    {
      if (z = z.p.p.left)
      {
        then
        {
          z = z.p
          TreeRightRotate(T, z)
        }
        z.p.color = BLACK
        z.p.p.color = RED
        TreeLeftRotate(T, z.p.p)
      }
    }
  }
}

```

```

proc BTreeSearch (x, k)
{
  i = 1
  while ((i ≤ x.n) and (k > x.keyi)) i = i + 1
  if ((i ≤ x.n) and (k = x.keyi))
  {
    then return (x, i)
  }
  if (x.leaf = true)
  {
    then return nil
  }
  DiskRead(x.ci)
  return BTreeSearch(x.ci, k)
}

```

```

proc BTreeCreate (T)
{
  x = Allocate()
  x.leaf = true
  x.n = 0
  DiskWrite(x)
  T.root = x
}

```

```

proc BTreeSplitChild (x, i)
{
  z = Allocate()
  y = x.ci
  z.leaf = y.leaf
  for j = 1 to t - 1 z.keyj = y.keyj+t
  if (y.leaf = false)
  {
    then for j = 1 to t z.cj = y.cj+t
  }
  y.n = t - 1
  for j = x.n + 1 downto i + 1 x.cj+1 = x.cj
  x.ci+1 = z
  for j = x.n downto i x.keyj+1 = x.keyj
  x.keyi = y.keyt
  x.n = x.n + 1
  DiskWrite(y)
  DiskWrite(z)
  DiskWrite(x)
}

```

```

proc BTreeInsert (T, k)
{
  r = T.root
  if (r.n = 2 · t - 1)
  {
    then
    {
      s = Allocate()
      T.root = s
      s.leaf = false
      s.n = 0
      s.c1 = r
      BTreeSplitChild(s, 1)
      BTreeInsertNonFull(s, k)
    }
    else BTreeInsertNonFull(r, k)
  }
}

```

```

proc BTreeInsertNonFull (x, k)
{
  i = x.n
  if (x.leaf = true)
  {
    then
    {
      while ((i ≥ 1) and (k < x.keyi))
      {
        {
          x.keyi+1 = x.keyi
          i = i - 1
        }
        x.keyi+1 = k
        x.n = x.n + 1
        DiskWrite(x)
      }
    }
    else
    {
      while ((i ≥ 1) and (k < x.keyi)) i = i - 1
      i = i + 1
      DiskRead(x.ci)
      if (x.ci.n = 2 · t - 1)
      {
        then
        {
          BTreeSplitChild(x, i)
          if (k > x.keyi)
          {
            then i = i + 1
          }
          BTreeInsertNonFull(x.ci, k)
        }
      }
    }
  }
}

```

```

proc HashInsert (T, k)
{
  let x be a new node with key k
  {
    i = h(k)
    ListInsert(T[i], x)
  }
}

```

```

proc HashSearch (k)
{
  i = h(k)
  return ListSearch(T[i], k)
}

```

```

proc HashDelete (k)
{
  i = h(k)
  x = ListSearch(T[i], k)
  ListDelete(T[i], x)
}

```

```

proc OaHashInsert (T, k)
{
  i = 0
  repeat
  {
    {
      j = h(k, i)
      if (T[j] = nil)
      {
        then
        {
          T[j] = k
          return j
        }
        else i = i + 1
      }
    }
  }
  until (i = m)
  return "overflow"
}

```

```

proc OaHashSearch ( $T, k$ )
{
   $i = 0$ 
  repeat
  {
     $j = h(k, i)$ 
    if ( $T[j] = k$ )
    {
      then return  $j$ 
    }
     $i = i + 1$ 
  }
  until ( $(T[j] = \text{nil})$  or  $(i = m)$ )
  return nil
}

```

```

proc HashComputeModulo ( $w, B, m$ )
{
  let  $d = |w|$ 
   $z_0 = 0$ 
  for ( $i = 1$  to  $d$ )  $z_{i+1} = ((z_i \cdot B) + a_i) \bmod m$ 
  return  $z_d + 1$ 
}

```

```

proc BreadthFirstSearch ( $G, s$ )
{
  for ( $u \in G.V \setminus \{s\}$ )
  {
     $u.color = \text{WHITE}$ 
     $u.d = \infty$ 
     $u.\pi = \text{nil}$ 
  }
   $s.color = \text{GREY}$ 
   $s.d = 0$ 
   $s.\pi = \text{nil}$ 
   $Q = \emptyset$ 
  Enqueue( $Q, s$ )
  while ( $Q \neq \emptyset$ )
  {
     $u = \text{Dequeue}(Q)$ 
    for ( $v \in G.Adj[u]$ )
    {
      if ( $v.color = \text{WHITE}$ )
      {
        then
        {
           $v.color = \text{GRAY}$ 
           $v.d = u.d + 1$ 
           $v.\pi = u$ 
          Enqueue( $Q, v$ )
        }
      }
    }
     $u.color = \text{BLACK}$ 
  }
}

```

```

proc DepthFirstSearch ( $G$ )
{
  for ( $u \in G.V$ )
  {
     $u.color = \text{WHITE}$ 
     $u.\pi = \text{nil}$ 
     $time = 0$ 
    for ( $u \in G.V$ )
    {
      if ( $u.color = \text{WHITE}$ )
      {
        then DepthVisit( $G, u$ )
      }
    }
  }
}

```

```

proc DepthVisit ( $G, u$ )
{
   $time = time + 1$ 
   $u.d = time$ 
   $u.color = \text{GREY}$ 
  for ( $v \in G.Adj[u]$ )
  {
    if ( $v.color = \text{WHITE}$ )
    {
      then
      {
         $v.\pi = u$ 
        DepthVisit( $G, v$ )
      }
    }
  }
   $u.color = \text{BLACK}$ 
   $time = time + 1$ 
   $u.f = time$ 
}

```

```

proc CycleDet ( $G$ )
{
   $cycle = \text{False}$ 
  for ( $u \in G.V$ )  $u.color = \text{WHITE}$ 
  for ( $u \in G.V$ )
  {
    if ( $u.color = \text{WHITE}$ )
    {
      then DepthVisitCycle( $G, u$ )
    }
  }
  return  $cycle$ 
}

```

```

proc DepthVisitCycle ( $G, u$ )
{
   $u.color = \text{GREY}$ 
  for ( $v \in G.Adj[u]$ )
  {
    if ( $v.color = \text{WHITE}$ )
    {
      then DepthVisitCycle( $G, v$ )
    }
    if ( $v.color = \text{GREY}$ )
    {
      then  $cycle = \text{true}$ 
    }
  }
   $u.color = \text{BLACK}$ 
}

```

```

proc TopologicalSort ( $G$ )
{
  for ( $u \in G.V$ )  $u.color = \text{WHITE}$ 
   $L = \emptyset$ 
   $time = 0$ 
  for ( $u \in G.V$ )
  {
    if ( $u.color = \text{WHITE}$ )
    {
      then DepthVisitTS( $G, u$ )
    }
  }
  return  $L$ 
}

```

```

proc DepthVisitTS ( $G, u$ )
{
   $time = time + 1$ 
   $u.d = time$ 
   $u.color = \text{GREY}$ 
  for ( $v \in G.Adj[u]$ )
  {
    if ( $v.color = \text{WHITE}$ )
    {
      then DepthVisitTS( $G, v$ )
    }
  }
   $u.color = \text{BLACK}$ 
   $time = time + 1$ 
   $u.f = time$ 
  ListInsert( $L, u$ )
}

```

```

proc StronglyConnectedComponents ( $G$ )
{
  for ( $u \in G.V$ )
  {
     $u.color = \text{WHITE}$ 
     $u.\pi = \text{nil}$ 
     $time = 0$ 
  }
  for ( $u \in G.V$ )
  {
    if ( $u.color = \text{WHITE}$ )
    {
      then DepthVisit( $G, u$ )
    }
  }
  for ( $u \in G.V$ )
  {
     $u.color = \text{WHITE}$ 
     $u.\pi = \text{nil}$ 
     $time = 0$ 
  }
   $L = \emptyset$ 
  for ( $u \in G^T.V$  in rev. finish time order)
  {
    if ( $u.color = \text{WHITE}$ )
    {
      then DepthVisit( $G^T, u$ )
    }
  }
  ListInsert( $L, u$ )
  return  $L$ 
}

```

```

proc MST-Prim ( $G, w, r$ )
{
  for ( $v \in G.V$ )
  {
    do
    {
       $v.key = \infty$ 
       $v.\pi = \text{nil}$ 
       $r.key = 0$ 
       $Q = G.V$ 
    }
  }
  while ( $Q \neq \emptyset$ )
  {
    do
    {
       $u = \text{ExtractMin}(Q)$ 
      for ( $v \in G.Adj[u]$ )
      {
        do
        {
          if ( $((v \in Q) \text{ and } (W(u, v) < v.key))$ )
          {
            do
            {
               $v.\pi = u$ 
               $v.key = W(u, v)$ 
            }
          }
        }
      }
    }
  }
}

```

```

proc MST-Kruskal ( $G, w$ )
{
   $T = \emptyset$ 
  for ( $v \in G.V$ )
  {
    do MakeSet( $v$ )
  }
  SortNoDecreasing( $G.E$ )
  for ( $((u, v) \in G.E - \text{in order})$ )
  {
    do
    {
      if ( $(\text{FindSet}(u) \neq \text{FindSet}(v))$ )
      {
        then
        {
           $T = T \cup \{(u, v)\}$ 
          Union( $u, v$ )
        }
      }
    }
  }
  return  $A$ 
}

```

```

proc DAGShortestPath ( $G, s$ )
{
  TopologicalSort( $G$ )
  InitalizeSingleSource( $G, s$ )
  for ( $u \in G.V - \text{in order}$ )
  {
    for ( $v \in G.Adj[u]$ ) Relax( $u, v, W$ )
  }
}

```

```

proc InitalizeSingleSource ( $G, s$ )
{
  for ( $v \in G.V$ )
  {
     $v.d = \infty$ 
     $v.\pi = \text{nil}$ 
     $s.d = 0$ 
  }
}

```

```

proc Relax ( $u, v, W$ )
{
  if ( $v.d > u.d + W(u, v)$ )
  then
  {
     $v.d = u.d + W(u, v)$ 
     $v.\pi = u$ 
  }
}

```

```

proc Bellman-Ford ( $G, s$ )
{
  InitalizeSingleSource( $G, s$ )
  for  $i = 1$  to  $|G.V| - 1$ 
  {
    for  $((u, v) \in G.E)$  Relax( $u, v, W$ )
  }
  for  $((u, v) \in G.E)$ 
  {
    if ( $v.d > u.d + W(u, v)$ )
    then return false
  }
  return true
}

```

```

proc Dijkstra ( $G, s$ )
{
  InitalizeSingleSource( $G, s$ )
   $S = \emptyset$ 
   $Q = G.V$ 
  while ( $Q \neq \emptyset$ )
  {
     $u = \text{ExtractMin}(Q)$ 
     $S = S \cup \{u\}$ 
    for ( $v \in G.Adj[u]$ )
    {
      if ( $v \in Q$ )
      then Relax( $u, v, W$ )
    }
  }
}

```

```

proc ExtendShortestPaths ( $L, W$ )
{
   $n = L.rows$ 
  let  $L'$  be a new matrix
  for  $i = 1$  to  $n$ 
  {
    for  $j = 1$  to  $n$ 
    {
       $L'_{ij} = \infty$ 
      for  $k = 1$  to  $n$ 
      {
         $L'_{ij} = \min\{L'_{ij}, L_{ik} + W_{kj}\}$ 
      }
    }
  }
  return  $i + 1$ 
}

```

```

proc SlowAllPairsMatrix ( $W$ )
{
   $n = W.rows$ 
   $L^1 = W$ 
  for  $m = 2$  to  $n - 1$ 
  {
     $L^m = \text{ExtendShortestPaths}(L^{m-1}, W)$ 
  }
  return  $L^{n-1}$ 
}

```

```

proc FastAllPairsMatrix ( $W$ )
{
   $n = W.rows$ 
   $L^1 = W$ 
   $m = 1$ 
  while ( $m < n - 1$ )
  do
  {
     $L^{2 \cdot m} = \text{ExtendShortestPaths}(L^m, L^m)$ 
     $m = 2 \cdot m$ 
  }
  return  $L^m$ 
}

```

```

proc Floyd-Warshall ( $W$ )
{
   $n = W.rows$ 
   $D^0 = W$ 
  for  $k = 1$  to  $n$ 
  {
    let  $D^k$  be a new matrix
    for  $i = 1$  to  $n$ 
    {
      for  $j = 1$  to  $n$ 
      {
         $D^k_{ij} = \min\{D^{k-1}_{ij}, D^{k-1}_{ik} + D^{k-1}_{kj}\}$ 
      }
    }
  }
}

```