

LED 암호 구현

송민호

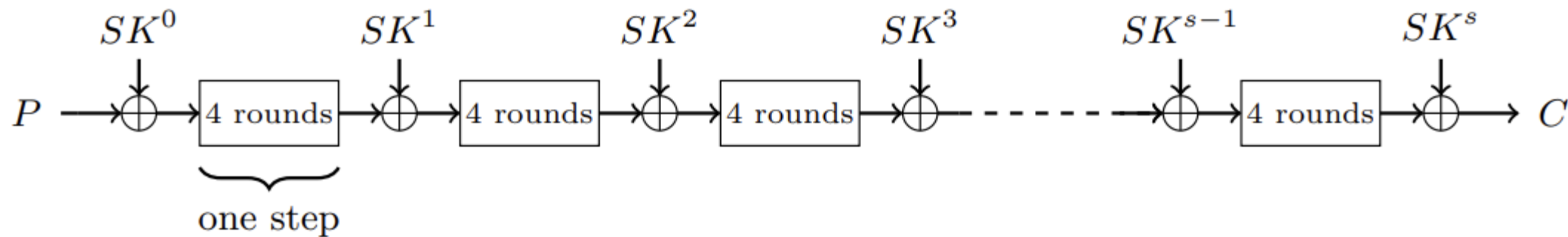
유튜브 주소: <https://youtu.be/PYDjJXXBX3U>

LED

- 2011년 CHES에서 발표한 경량 암호
- 블록 길이 64-bit, 키 길이 64-bit, 80-bit, 96-bit, 128-bit
 - 키 길이가 64-bit인 경우, 8 라운드 연산
 - 그 외에는 12라운드 연산
- 4-bit를 이용한 4x4 배열을 4-bit 단위로 연산

LED

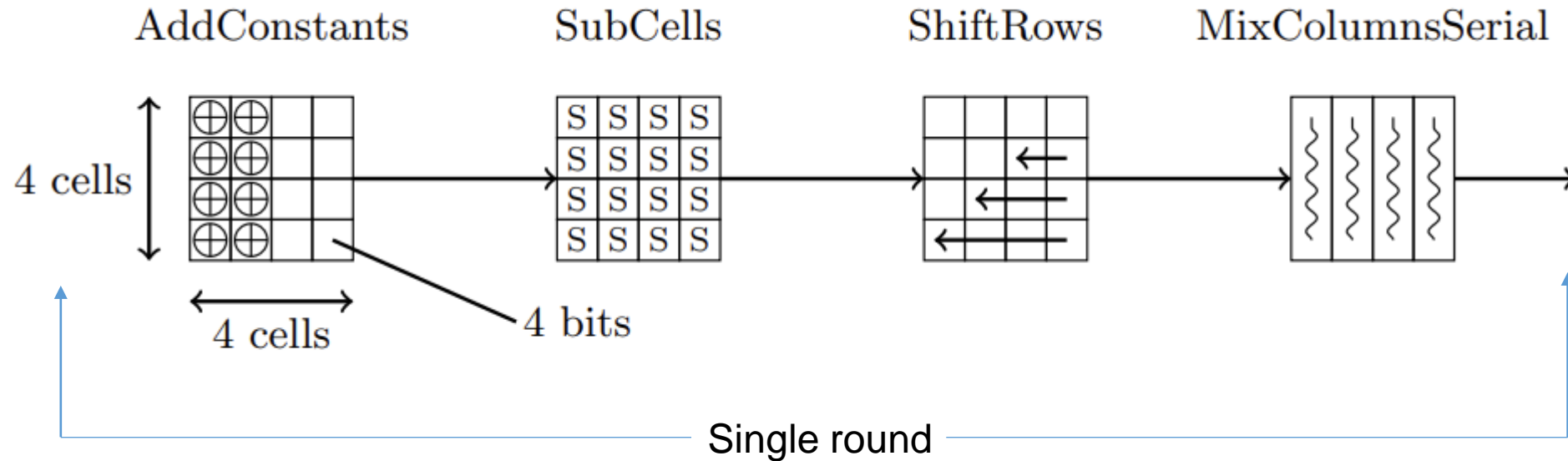
- 연산 순서



```
for  $i = 0$  to  $s - 1$  do {  
    addRoundKey (STATE,  $SK^i$ )  
    step (STATE)  
}  
addRoundKey (STATE,  $SK^s$ )
```

LED

- One step
 - Single round 4번 반복



LED

- LED 64/64
- Key 64-qubit
- Block 64-qubit
 - 첫 번째 행 b0 = 16-qubit

```
k = eng.allocate_qureg(64)
b0 = eng.allocate_qureg(16)
b1 = eng.allocate_qureg(16)
b2 = eng.allocate_qureg(16)
b3 = eng.allocate_qureg(16)

addRoundKey(eng, b0, b1, b2, b3, k)

cnt = 0

for i in range(0, rounds):
    for j in range(0, 4):
        AddConstants(eng, b0, b1, b2, b3, cnt)
        SubCells_layer(eng, b0, b1, b2, b3)
        ShiftRows(eng, b0, b1, b2, b3)
        MixColumns(eng, b0, b1, b2, b3)
        cnt += 4

    addRoundKey(eng, b0, b1, b2, b3, k)
```

addRoundKey

- Key와 평문을 XOR 연산
- 키 스케줄이 존재하지 않음
 - 매 라운드마다 주어진 키 그대로 사용

$$\begin{bmatrix} k_0 & k_1 & k_2 & k_3 \\ k_4 & k_5 & k_6 & k_7 \\ k_8 & k_9 & k_{10} & k_{11} \\ k_{12} & k_{13} & k_{14} & k_{15} \end{bmatrix}$$

for 64-bit keys

```
def addRoundKey(eng, b0, b1, b2, b3, k):  
    for i in range(16):  
        CNOT | (k[i+48], b0[i])  
        CNOT | (k[i+32], b1[i])  
        CNOT | (k[i+16], b2[i])  
        CNOT | (k[i], b3[i])
```

AddConstants

- Round Constants와 평문을 XOR 연산

```
def AddConstants(eng, b0, b1, b2, b3, cnt):  
    rc = [0x4000, 0x5100, 0x2000, 0x3100, 0x4000, 0x5300,  
          0x4300, 0x5700, 0x2300, 0x3700, 0x4700, 0x5600,  
          0x4600, 0x5700, 0x2600, 0x3700, 0x4500, 0x5700,  
          0x4700, 0x5100, 0x2700, 0x3100, 0x4600, 0x5300,  
          0x4300, 0x5500, 0x2300, 0x3500, 0x4700, 0x5200,  
          0x4200, 0x5600, 0x2200, 0x3600, 0x4500, 0x5400,  
          0x4400, 0x5100, 0x2400, 0x3100, 0x4000, 0x5200,  
          0x4200, 0x5700, 0x2200, 0x3700, 0x4500, 0x5600]  
  
    Round_constant_XOR(eng, b0, rc[cnt], 16)  
    Round_constant_XOR(eng, b1, rc[cnt+1], 16)  
    Round_constant_XOR(eng, b2, rc[cnt+2], 16)  
    Round_constant_XOR(eng, b3, rc[cnt+3], 16)
```

```
def Round_constant_XOR(eng, k, rc, bit):  
    for i in range(bit):  
        if (rc >> i & 1):  
            X | k[i]
```

1000000000000000
-> 1000000000000000

SubCells

- Sbox 연산
 - Sbox LIGHTER-R을 통해 최적화
 - ./non-lin-search -qvw -o "C56B90AD3EF84712"
-f "libraries/NCT_gc.conf" -s "LED"
 - CCNOT2 = Toffoli
 - RNOT = X

x	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
$S[x]$	C	5	6	B	9	0	A	D	3	E	F	8	4	7	1	2

```
// from : 00FF 0F0F 3333 5555
```

```
F[0] = X[1];  
F[1] = X[2];  
F[2] = X[3];  
F[3] = X[0];
```

```
F[1] = CNOT1(F[1], F[0]);  
F[3] = CCNOT2(F[1], F[0], F[3]);  
F[0] = CCNOT2(F[3], F[1], F[0]);  
F[1] = CCNOT2(F[2], F[0], F[1]);  
F[0] = CNOT1(F[0], F[3]);  
F[3] = RNOT1(F[3]);  
F[0] = CNOT1(F[0], F[1]);  
F[2] = CNOT1(F[2], F[3]);  
F[1] = CNOT1(F[1], F[2]);  
F[2] = RNOT1(F[2]);  
F[3] = CCNOT2(F[1], F[0], F[3]);
```

```
X[0] = F[1];  
X[1] = F[3];  
X[2] = F[0];  
X[3] = F[2];
```

```
//to : 9B70 E16C 32E5 59A6  
// Cost : 11  
// Logic Library : NCT_gc
```


SubCells

- Endian 정렬

```
F[0] = X[1];  
F[1] = X[2];  
F[2] = X[3];  
F[3] = X[0];
```

→

```
F[0] = X[2]  
F[1] = X[1]  
F[2] = X[0]  
F[3] = X[3]
```

```
F[1] = CNOT1(F[1], F[0]);  
F[3] = CCNOT2(F[1], F[0], F[3]);  
F[0] = CCNOT2(F[3], F[1], F[0]);  
F[1] = CCNOT2(F[2], F[0], F[1]);
```

→

CNOT | (b[2], b[1])

```
Toffoli_gate(eng, b[1], b[2], b[3])  
Toffoli_gate(eng, b[3], b[1], b[2])  
Toffoli_gate(eng, b[0], b[2], b[1])
```

```
X[0] = F[1];  
X[1] = F[3];  
X[2] = F[0];  
X[3] = F[2];
```

→

```
X[2] = F[0]  
X[1] = F[3]  
X[0] = F[1]  
X[3] = F[2]
```

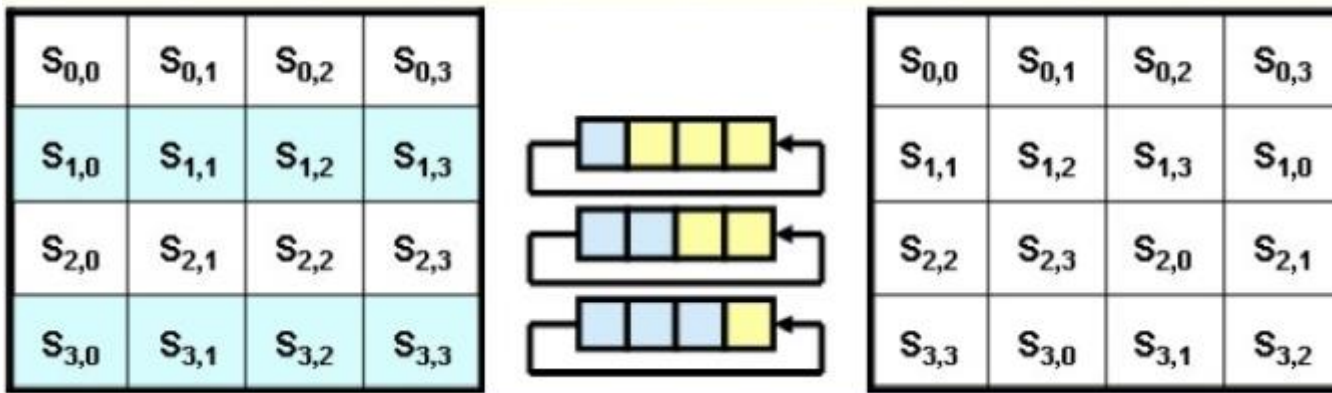
→

```
F[0] = X[2] -> F[0]  
F[1] = X[0] -> F[2]  
F[2] = X[3] -> F[3]  
F[3] = X[1] -> F[1]
```

```
out = []  
out.append(b[0])  
out.append(b[2])  
out.append(b[3])  
out.append(b[1])
```

ShiftRows

- 각 행의 Cell을 Shift



```
for i in range(4): #5
    new_b1.append(b1[12+i])
for i in range(4): #6
    new_b1.append(b1[i])
for i in range(4):
    new_b1.append(b1[4+i])
for i in range(4): #4
    new_b1.append(b1[8+i])
```

MixColumnsSerial

- 행렬곱

$$\begin{pmatrix} 4 & 1 & 2 & 2 \\ 8 & 6 & 5 & 6 \\ B & E & A & 9 \\ 2 & 2 & F & B \end{pmatrix} = M.$$

$$\begin{matrix} 4 & 1 & 2 & 2 \\ 8 & 6 & 5 & 6 \\ B & E & A & 9 \\ 2 & 2 & F & B \end{matrix} \times \begin{matrix} b0[0] & b0[1] & b0[2] & b0[3] \\ b1[0] & b1[1] & b1[2] & b1[3] \\ b2[0] & b2[1] & b2[2] & b2[3] \\ b3[0] & b3[1] & b3[2] & b3[3] \end{matrix}$$

$$(4 \times b0[0]) \oplus (1 \times b1[0]) \oplus (2 \times b2[0]) \oplus (2 \times b3[0])$$

- Reduction 사용

- 4-bit로 맞춰주기 위함
- $X^4 + X + 1$

1	0	1	0	0	0
X^5	X^4	X^3	X^2	X	1
1	0	1	0	0	0
1	0	1	1	1	0
		1	1	1	0

MixColumnsSerial

4 1 2 2	x	b0[0]	b0[1]	b0[2]	b0[3]
8 6 5 6		b1[0]	b1[1]	b1[2]	b1[3]
B E A 9		b2[0]	b2[1]	b2[2]	b2[3]
2 2 F B		b3[0]	b3[1]	b3[2]	b3[3]

(4 x b0[0]) = 0100 x (b3 b2 b1 b0)

b3	b2	b1	b0	0	0
X^5	X^4	X^3	X^2	X	1
b3	b2	b1	b0	b2	b2
b3	b2	b1	$b0 \oplus b3$	$b2 \oplus b3$	b2
		b1	$b0 \oplus b3$	$b2 \oplus b3$	b2

```
#4
CNOT | (b0[0], out[2])
CNOT | (b0[1], out[3])

CNOT | (b0[2], out[0])
CNOT | (b0[2], out[1])
CNOT | (b0[3], out[1])
CNOT | (b0[3], out[2])
```

Q & A