# Efficient Key Management Scheme for Health Blockchain

Paper review

https://youtu.be/NXwijDHKAYU

- Is consists of
  - Consensus Mechanism
  - Digital Signature
  - Hash Chains
  - Shared Database

- Provides
  - Non-Repudiation
  - Integrity
  - Distributed Storage
  - Time-based Traceability

- Can be used in
  - Healthcare
  - Fintech
  - Computational Law
  - Audit
  - Notarization

- Has
  - Speed of recording Issue
  - Efficiency of consensus Issue
  - Privacy of data Issue

## Privacy of Data

Zero-Knowledge Proof

- Encrypt and Store
  - Key Management Issue

## Key Management Issue

One key for all blocks

One key for a block

Considering application scenario

### Health Blockchain

Blockchain handling health data

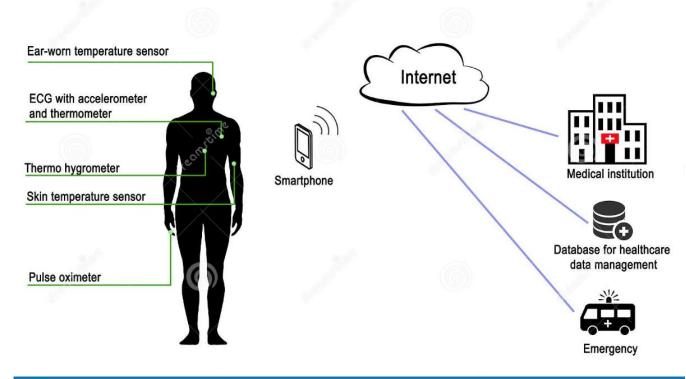
Has a number of private data

→ Must solve privacy issue

## Body Sensor Networks (BSNs)

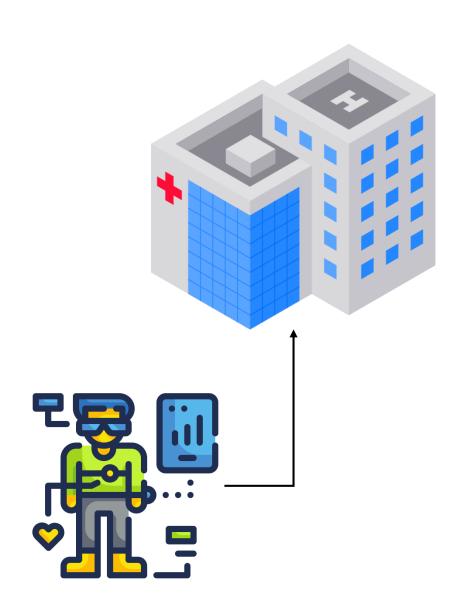
BAN, WBAN, MBAN

#### Wearable Wireless Body Area Network



# Original BSNs Problems

- Monopoly Problem
- Vulnerability Problem
- Privacy Problem
- Integrity Problem



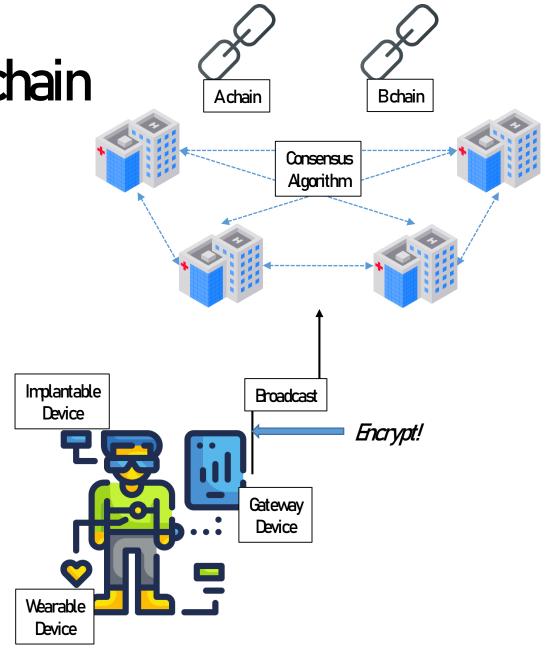
Proposed BSNusing Blockchain

Monopoly Problem

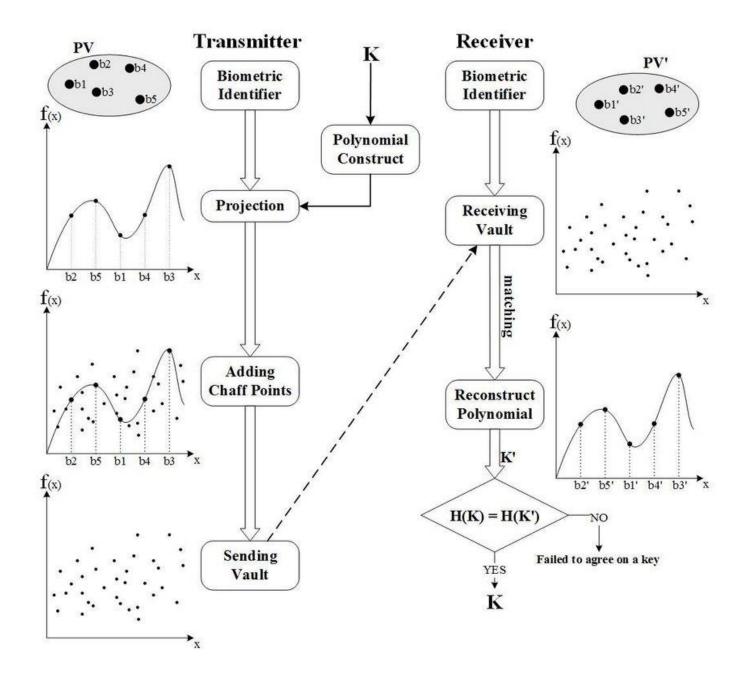
Vulnerability Problem

Privacy Problem

Integrity Problem



# Fuzzy Vault



# Fuzzy Vault using PPG (photoplethysmography) signal

#### Production of PPG vector.

- I. Collect PPG signals
- II. Encoding signals into vectors using FFT (fast Fourier transform)

$$F_S = \langle f_S^1, f_S^2, \dots, f_S^a \rangle, F_r = \langle f_r^1, f_r^2, \dots, f_r^a \rangle$$

#### II. Creating Polynomial.

- I. Create p(x) with a public order a
- II. Coefficients are produced from a random number  $\rightarrow$  common key

$$Coef. = e_a, e_{a-1}, \dots, e_1, e_0$$
,  $Key = e_a \parallel e_{a-1} \parallel, \dots, \parallel e_1 \parallel e_0$ 

# Fuzzy Vault using PPG (photoplethysmography) signal

#### III. Vault Production.

- I. Compute a set  $D = \{f_S^i, p(f_S^i)\}, 1 \le i \le a$  predefined
- II. Build a chaff points set  $C = \{c_i, d_i\}, 1 \le i \le W$  using random numbers
- III. Vault  $R = D \cup C$

#### IV. Vault Transmission.

I. Send  $R \parallel T(K,R)$  where T is MAC function

#### V. Opening Vault.

- I. Reconstruct p(x) using  $f_r^i$ , Lagrangian Interpolation
- II. Recovery K'
- III. Validate K' using T

## Fuzzy Vault using PPG signal problem

• v+1 feature points for vth-order polynomial

Relation between parameters

- Revised.
  - Key encoding
  - LOTR (Lower-order twice reconstruction)

# Proposed BSNagain

#### I. Initialization period

- I. Some biosensor can generate key
- II. Set specific node A

#### II. When gateway device needs to encrypt

- I. Gateway gives order and asks A to generate key
- II. A makes pre-key  $k_a \parallel k_{a-1} \parallel$ , ...,  $\parallel k_1$ ,  $\parallel k_0$  and encodes it into codeword  $e_i$
- III. A uses  $e_i$  as a coefficients to construct ath order p(x)

## Proposed BSNagain

#### III. Making Vault

- A collects PPG signals from adjacent nodes.
- A encodes these signals into a vector  $F_s$  using FFT
- III. A makes vault  $R = D \cup C$

## IV. Encrypting Vault Random function Pre-distributed key

- A calculates  $K^* = F(k^*, K)$  as the encryption key
- A generates random number r
- III. Do  $M = E(r \oplus k^*, R)$

## Proposed BSNagain

#### V. Sending to G

I. A sends  $K^* \parallel M \parallel r \parallel H(K^*) \parallel ID_A$  to Gateway device G

#### VI. Broadcasting

- I. G uses  $K^*$  to encrypt physiological data Mp
- II. G broadcasts e(Mp) and  $M \parallel r \parallel H(K^*) \parallel ID_A \parallel B_A$
- III. G deletes the  $K^*$

## Proposed BSNagain (Recovering)

e(Mp) and  $M \parallel r \parallel H(K^*) \parallel ID_A \parallel B_A$ 

#### I. Searching block

- I. User uses gateway to point out block on chain
- II. G searches block by index  $B_A$
- III. G sends  $M \parallel r \parallel H(K^*)$  to A

#### II. Once A received

- I. Decrypt M using  $r, k^*$  to get vault R
- II. Collect signals from adjacent biosensors
- III. Encoding these signals and recovery p(x) using LOTR and interpolation

#### III. Verifying

- I. A decodes coefficients from p(x) using RS code
- II. A checks  $H(K^*) = H(K^{*'})$

## Security and performance analysis

I. Attacking the blockchain

$$e(Mp)$$
 and  $M \parallel r \parallel H(K^*) \parallel ID_A \parallel B_A$   
 $R = D \cup C$ 

- II. Attacking the BSN
  - I. On or into the human body
  - II. Increase order



- III. Performance
  - I. BSN nodes are in charge of generation, backup, recovery