

양자 프로그래밍 툴 Depth 최적화

장경배

<https://youtu.be/htlvll4htek>

A Framework for Depth-Efficient Quantum Implementations of Linear Layers

Kyungbae Jang¹, Anubhab Baksi², and Hwajeong Seo¹

¹ Hansung University, Seoul, South Korea

² Lund University, Sweden

starj1023@gmail.com, anubhab.baksi@eit.lth.se, hwajeong84@gmail.com

Abstract. Quantum depth plays a critical role in determining the performance of quantum implementations, yet quantum programming tools often fail to produce depth-optimal compilations of linear layers. In this work, we present a systematic and automated framework that reorders quantum gate sequences of linear layers to obtain depth-efficient quantum implementations. Our method consistently produces circuits with lower depth compared to prior implementations.

We apply the framework to a range of cryptographic operations, including the AES MixColumn, internal layers of the AES S-box, binary field squaring, and modular reduction in binary field multiplication. In all these cases, our method achieves meaningful reductions in quantum depth—for example, lowering the depth of the AES MixColumn and S-box circuits.

This work explores optimal quantum circuit designs for quantum programming tools, improves the accuracy of quantum resource estimation for cryptanalysis, and supports more realistic evaluations of post-quantum security.

Case Study: AES MixColumn

- 다음 두 논문에선 동일한 AES MixColumn (32×32)을 구현, 하지만 **다른 Depth를 보고**
 - PQCrypto'16 (M. Grassl et al, 전 NIST 표준)
 - 32 Qubits
 - 277 CNOTs
 - **39 Depth**
 - EUROCRYPT'20 (S. Jaques et al. 현 NIST 표준) →
 - 32 Qubits
 - 277 CNOTs
 - **111 Depth**

```
def Euro_plu_mix(eng, word):  
    CNOT | (word[7], word[0]);  
    CNOT | (word[8], word[0]);  
    CNOT | (word[9], word[0]);  
    CNOT | (word[15], word[0]);  
    CNOT | (word[17], word[0]);  
    CNOT | (word[25], word[0]);  
    CNOT | (word[9], word[1]);  
    CNOT | (word[10], word[1]);  
    CNOT | (word[18], word[1]);  
    CNOT | (word[26], word[1]);
```



Note that [GLRS16] describes the same technique, while achieving a significantly smaller design than the one we obtain.

양자 프로그래밍 툴 Depth 최적화

- 하나의 큐비트의 반복적인 사용은 다른 연산과의 병렬화를 방해

⇒ 양자 프로그래밍 툴이 최적의 Depth를 찾을 수 없음 (Qcrypton, ProjectQ, Q#, Qiksit, Cirq)

```
7  circuit = QuantumCircuit(9, 9)
8
9  # Depth Test
10 circuit.cx(0, 8)
11 circuit.cx(0, 4)
12 circuit.cx(0, 5)
13 circuit.cx(0, 6)
14 circuit.cx(0, 7)
15
16 circuit.cx(1, 4)
17 circuit.cx(1, 5)
18 circuit.cx(1, 6)
19 circuit.cx(1, 7)
20
21 circuit.cx(2, 4)
22 circuit.cx(2, 5)
23 circuit.cx(2, 6)
24 circuit.cx(2, 7)
25 circuit.cx(2, 8)
26
27 circuit.x(2)
28 circuit.x(4)
29 circuit.x(6)
30 circuit.x(8)
```

Qiskit Information	
Qubits	9
X	4
CX	14
Circuit Depth	9
CX Depth	8

양자 프로그래밍 툴 Depth 최적화

- 하나의 큐비트의 **반복적인 사용을 회피: Compiler-friendly**
⇒ 양자 프로그래밍 툴이 더 낮은 Depth를 찾음 (9 → 6)
(암호 안전성 분석 측면에서, 올바른 복잡도를 도출하는 것 또한 유의미함)

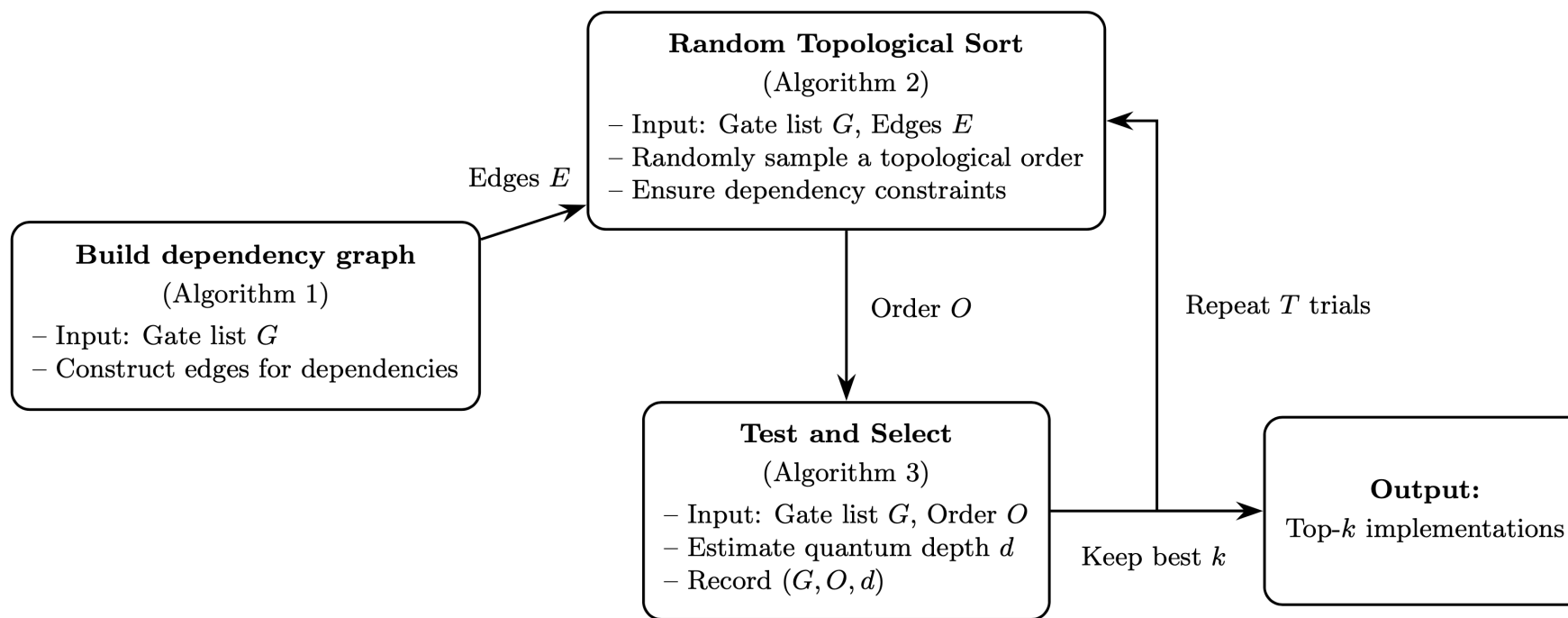
Reordered

```
32 # Depth Test (Reordered)
33 circuit.cx(0, 5)
34 circuit.cx(1, 4)
35 circuit.x(6)
36 circuit.cx(2, 7)
37 circuit.cx(0, 7)
38 circuit.cx(2, 6)
39 circuit.cx(1, 7)
40 circuit.cx(2, 5)
41 circuit.x(8)
42 circuit.cx(0, 4)
43 circuit.cx(2, 4)
44 circuit.cx(2, 8)
45 circuit.cx(0, 6)
46 circuit.x(4)
47 circuit.cx(0, 8)
48 circuit.cx(1, 6)
49 circuit.x(2)
50 circuit.cx(1, 5)
```

Qiskit Information	
Qubits	9
X	4
CX	14
Circuit Depth	9 → 6
CX Depth	8 → 6

Compiler-friendly 최적화 프레임워크

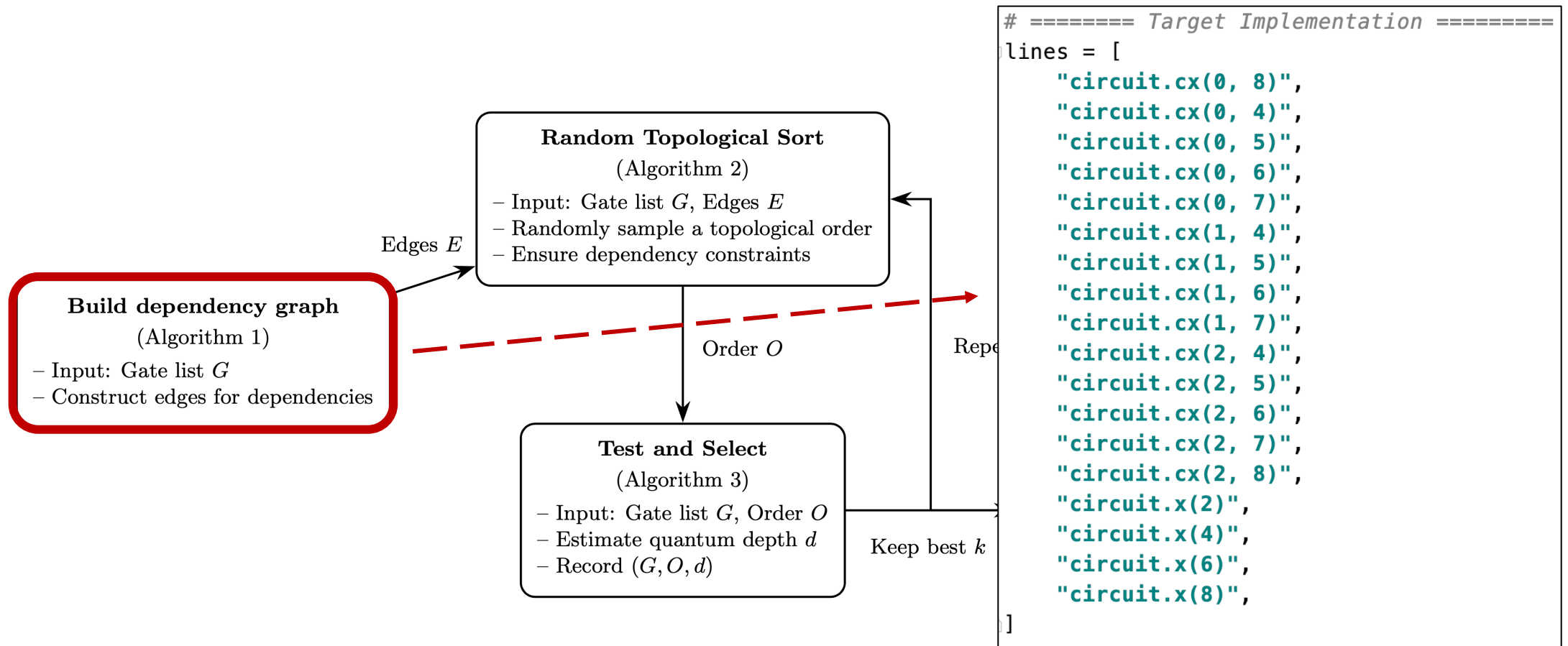
- 양자 프로그래밍 툴 용 Compiler-friendly 최적화 프레임워크 개발



< 프레임워크 Overview >

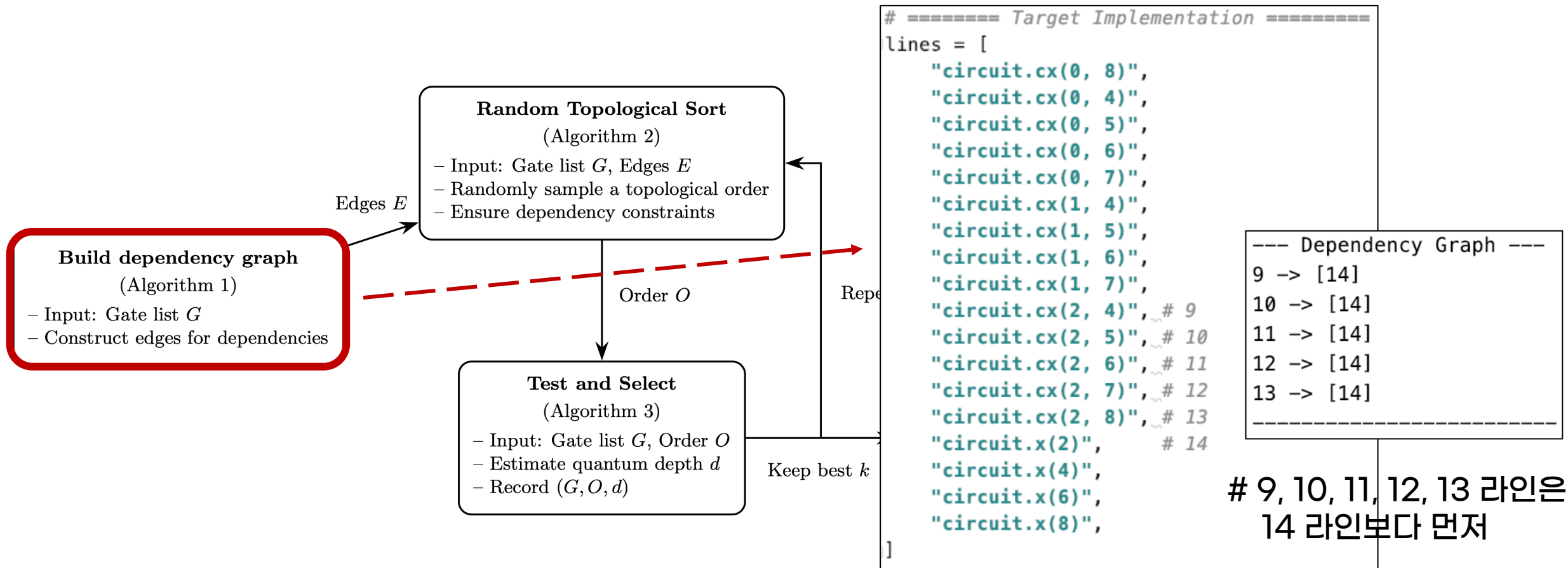
Compiler-friendly 최적화 프레임워크

- 양자 프로그래밍 툴 용 Compiler-friendly 최적화 프레임워크 개발
 - 최적화하고 싶은 **양자 구현 입력**



Compiler-friendly 최적화 프레임워크

- 양자 프로그래밍 툴 용 Compiler-friendly 최적화 프레임워크 개발
 - 해당 양자 구현의 연산들 간의 **의존 그래프** 생성



Algorithm 1: Build dependency graph

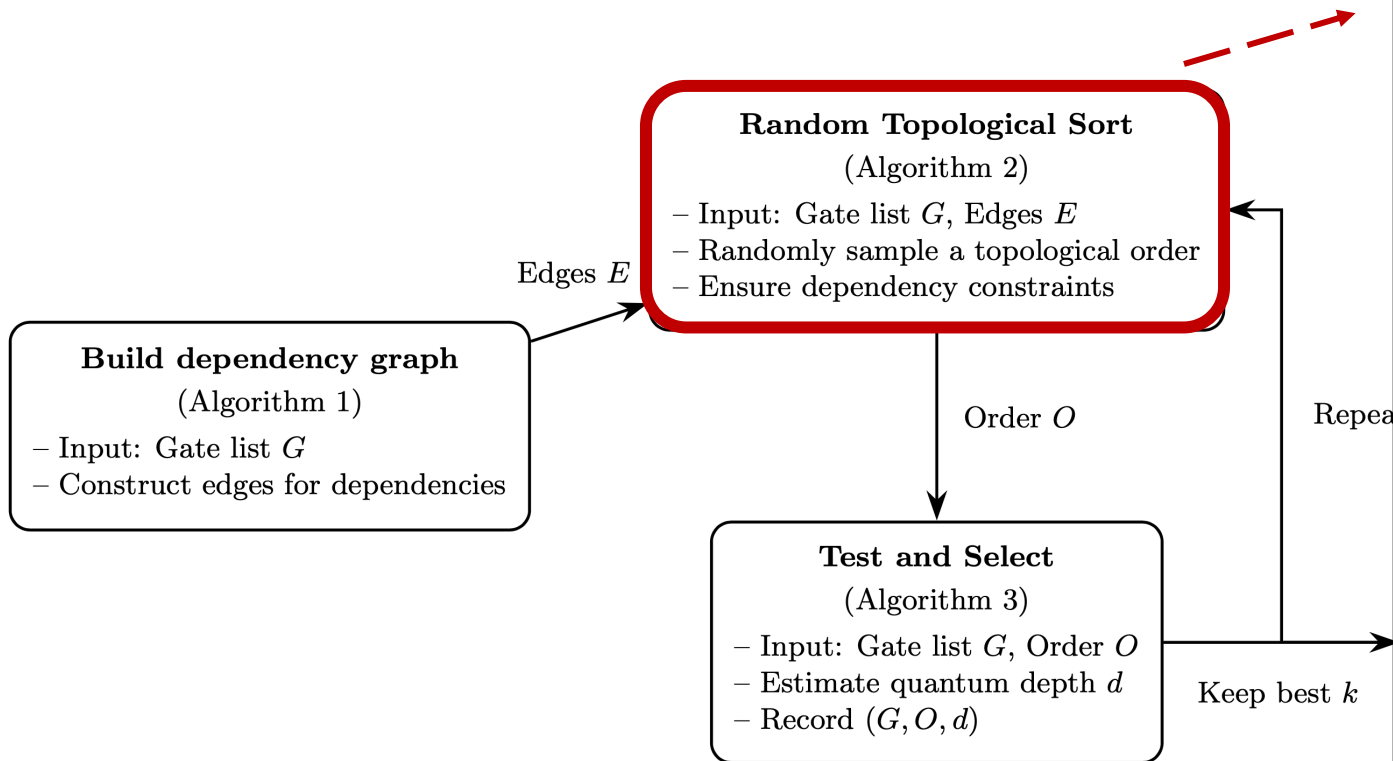
Input: Gate list $G = \{g_1, \dots, g_n\}$ with $g = \text{CNOT}(ctrl, tgt)$ or $\text{X}(tgt)$

Output: Directed edge set $E \subseteq \{(i \rightarrow j) \mid 1 \leq i < j \leq n\}$

```
1:  $E \leftarrow \emptyset$ 
2: for each qubit label  $q$ :  $Writes[q] \leftarrow []$ ,  $Reads[q] \leftarrow []$ 
3: for  $i \leftarrow 1$  to  $n$  do ▷ collect per-qubit reads/writes
4:   if  $g_i$  is CNOT then
5:      $Reads[ctrl(g_i)].append(i)$ 
6:      $Writes[tgt(g_i)].append(i)$ 
7:   else if  $g_i$  is X then
8:      $Writes[tgt(g_i)].append(i)$ 
9:   end if
10: end for
11: for  $j \leftarrow 1$  to  $n$  do ▷ RAW: all prior writes  $\rightarrow$  current read
12:   if  $g_j$  is CNOT then
13:      $c \leftarrow ctrl(g_j)$ 
14:      $P \leftarrow \{i \in Writes[c] \mid i < j\}$ 
15:     if  $P \neq \emptyset$  then
16:        $E \leftarrow E \cup \{(i \rightarrow j) \mid i \in P\}$ 
17:     end if
18:   end if
19: end for
20: for  $j \leftarrow 1$  to  $n$  do ▷ WAR: all prior reads  $\rightarrow$  current write
21:   if  $g_j$  is CNOT or X then
22:      $q \leftarrow tgt(g_j)$ 
23:      $R \leftarrow \{i \in Reads[q] \mid i < j\}$ 
24:     if  $R \neq \emptyset$  then
25:        $E \leftarrow E \cup \{(i \rightarrow j) \mid i \in R\}$ 
26:     end if
27:   end if
28: end for
29: return  $E$ 
```

Compiler-friendly 최적화 프레임워크

- 양자 프로그래밍 툴 용 Compiler-friendly 최적화 프레임워크 개발
 - 생성된 의존 그래프 기반 **랜덤 Reordering**



Algorithm 2: Randomized topological sort

Input: Gate list $G = \{g_1, \dots, g_n\}$, dependency edges E

Output: A randomized topological order $Order$ or None

1: Compute in-degrees $\deg[j]$ for all $1 \leq j \leq n$

2: $Pool \leftarrow \{j \mid \deg[j] = 0\}$, $Order \leftarrow []$

3: **while** $Pool \neq \emptyset$ **do**

4: Randomly shuffle $Pool$

5: $u \leftarrow$ pop one from $Pool$

6: append u to $Order$

7: **for each** $(u \rightarrow w) \in E$ **do**

8: $\deg[w] \leftarrow \deg[w] - 1$

9: **if** $\deg[w] = 0$ **then**

10: add w to $Pool$

11: **end if**

12: **end for**

13: **end while**

14: **if** $|Order| = n$ **then**

15: **return** $Order$

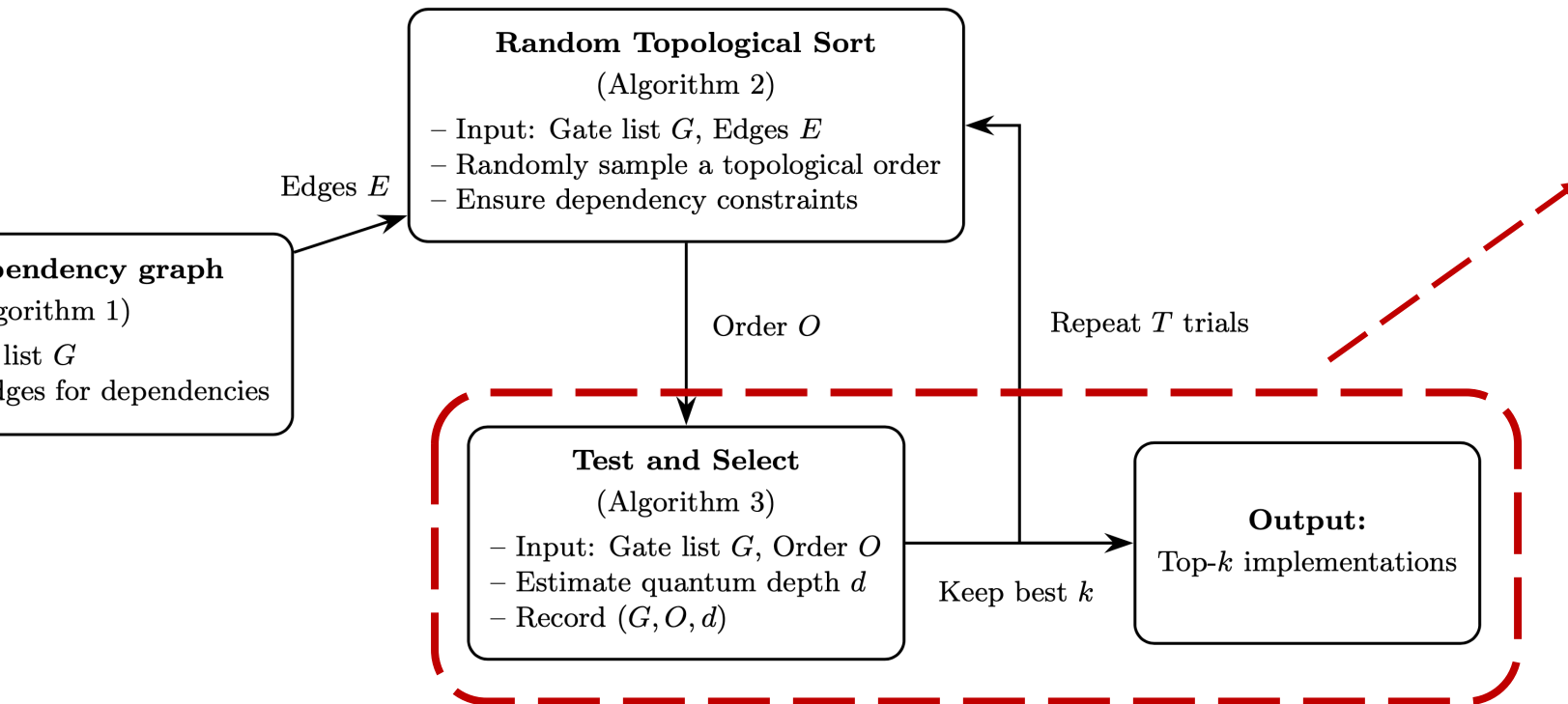
16: **else**

17: **return** None

18: **end if**

Compiler-friendly 최적화 프레임워크

- 양자 프로그래밍 툴 용 Compiler-friendly 최적화 프레임워크 개발
 - Reorder된 구현의 **Depth 추정** 및 **Top- k 분류** (1000번 시도, $k = 3$)



Collected 1000 valid orders (trials=1000).

=== Top Reordered Schedules ===

[1] Depth = 6	[2] Depth = 6	[3] Depth = 6
circuit.cx(2, 4)	circuit.cx(2, 7)	circuit.cx(0, 4)
circuit.cx(2, 8)	circuit.cx(2, 8)	circuit.cx(2, 5)
circuit.cx(0, 7)	circuit.cx(0, 4)	circuit.cx(1, 5)
circuit.x(8)	circuit.cx(1, 5)	circuit.cx(2, 6)
circuit.cx(0, 5)	circuit.cx(1, 6)	circuit.x(4)
circuit.x(6)	circuit.cx(0, 5)	circuit.cx(0, 7)
circuit.cx(2, 7)	circuit.cx(2, 6)	circuit.cx(2, 8)
circuit.cx(1, 6)	circuit.cx(0, 7)	circuit.cx(1, 4)
circuit.cx(0, 4)	circuit.cx(1, 4)	circuit.cx(1, 7)
circuit.cx(1, 5)	circuit.cx(0, 6)	circuit.cx(0, 5)
circuit.cx(2, 6)	circuit.x(6)	circuit.x(6)
circuit.cx(0, 6)	circuit.cx(2, 4)	circuit.cx(2, 4)
circuit.cx(1, 4)	circuit.cx(1, 7)	circuit.cx(2, 7)
circuit.cx(1, 7)	circuit.x(4)	circuit.cx(0, 6)
circuit.cx(0, 8)	circuit.cx(2, 5)	circuit.x(8)
circuit.x(4)	circuit.x(2)	circuit.cx(1, 6)
circuit.cx(2, 5)	circuit.x(8)	circuit.cx(0, 8)
circuit.x(2)	circuit.cx(0, 8)	circuit.x(2)

Algorithm 3: Test and select

Input: Gate list G , sample budget T , top- k

Output: Top- k schedules with minimum depth

```
1:  $E \leftarrow \mathbf{Build\ dependency\ graph}(G)$ 
2:  $Seen \leftarrow \emptyset$  ▷ set of visited orders
3:  $Cand \leftarrow \emptyset$  ▷ set of candidate schedules
4: for  $t \leftarrow 1$  to  $T$  do
5:    $O \leftarrow \mathbf{Randomized\ topological\ sort}(E)$ 
6:   if  $O = \text{None}$  or  $O \in Seen$  then
7:     continue
8:   end if
9:    $Seen \leftarrow Seen \cup \{O\}$ 
10:   $d \leftarrow \mathbf{Estimate\ depth}(G, O)$  ▷ e.g., depth_of_dag
11:   $Cand \leftarrow Cand \cup \{(O, d)\}$ 
12: end for
13: Sort  $Cand$  by depth in ascending order
14: return first  $k$  elements of  $Cand$ 
```

Real-World Examples

- **실제 양자 회로 구현들을 대상으로 테스트**
 - 선형 연산이 아니더라도, **내부 작은 선형 연산들을** 대상으로 적용하여 최적화 가능
 \Rightarrow AES S-box, Multiplication (모듈러 reduction)

Linear layer	Size (matrix)	Source	#Qubit	Full depth
AES MixColumn	32×32	Jaques et al. [9] This paper	32*	111 81
		Liu et al. [12] This paper	135	13 9
		Shi et al. [14] This paper	131	14 9
		Sun et al. [17] This paper	122	18 11
		Li et al. [10] This paper	137	11 8
		Lin et al. [11] This paper	123	16 11
AES S-box (internal linear layers)	-	Jang et al. [6] [*] This paper	76	67 64
Squaring ($\mathbb{F}_{2^{16}}$)	16×16	J ⁺ [8] (Naïve) (Compiler-friendly) This paper	32	22 14 11
Squaring ($\mathbb{F}_{2^{127}}$)	127×127	J ⁺ [8] (Naïve) (Compiler-friendly) This paper	254	175 125 125
Multiplication ($\mathbb{F}_{2^{16}}$)	-	J ⁺ [7] (adopted in [8]) This paper	243	43 41
Multiplication ($\mathbb{F}_{2^{163}}$)	-	J ⁺ [7] (adopted in [8]) This paper	13161	66 65

 **Best AES S-box**

감사합니다