Quantum Information Set Decoding

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https://youtu.be/YmyEfLjExt8





Information Set Decoding(ISD)

- 코드기반 암호에서 가장 효율적이라고 알려진 공격법
 - 주어진 암호문 c 와 공개키G'을 이용하여 원본메세지 m 을 복구함
- 1. n bit 의 길이 암호문 c 에서 k bit 의 벡터 c_k 를 랜덤하게 선택
 - → 오류 위치를 모르는 상태에서 자신이 선택한 k bit 벡터에 오류가 포함되지 않아야 함
- 2. 선택한 열의 index에 맞춰 G' 으로부터 G'_k 를 뽑아낸다. 이때, G'_k 는 invertible
- 3. $c + c_k G'_k^{-1} G$ 의 weight가 t 와 같거나 더 적은지 확인한다. 그렇지 않으면 1단계 부터 다시 반복
- 4. 앞의 조건이 만족한다면 원본메세지 $m = c_k G'_k^{-1}$ 로 복구가 가능하다.

Information Set Decoding Example(ISD)

1. n-bit 의 길이 암호문 c 에서 k-bit 의 벡터 c_k 를 랜덤하게 선택 c_k 를 관덤하게 오류가 포함되지 않아야 함

밑의 결과 벡터의 weight를 확인한다. (t or less)

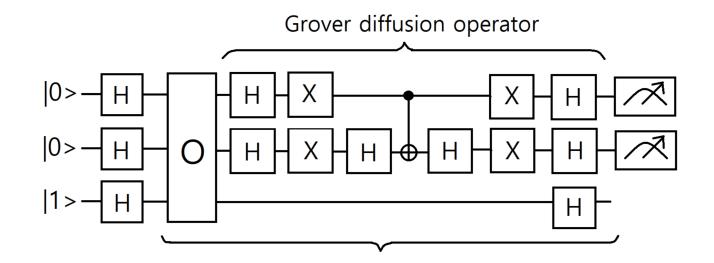
$$c + c_k G'_k^{-1} G = 0 \ 1 \ 1 \ 0 \ 1 \ 1 \ 0$$
 + 0 1 1 0 0 1 0 = 0 0 0 0 1 0 0

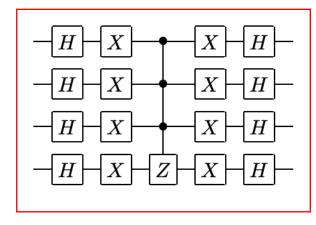
마지막으로 최종 메세지 $m = c_k G'_k^{-1} = 1101$

n = 1024, k = 524 의 Goppa 코드에 대하여이 경우의 수는 $2^{80.7}$ 결과의 작업 계수

Grover

- 정렬되지 않은 데이터베이스로부터 특정 데이터를 찾는 양자 알고리즘
- Classic 한 알고리즘에선 최대 N 번의 시도가 필요하지만, Grover 알고리즘은 \sqrt{N} 번이면 가능하다.





diffusion operator

Grover

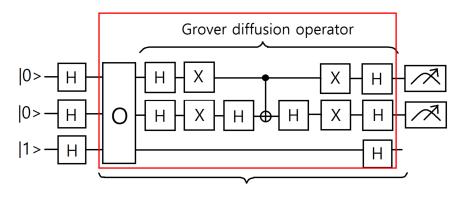
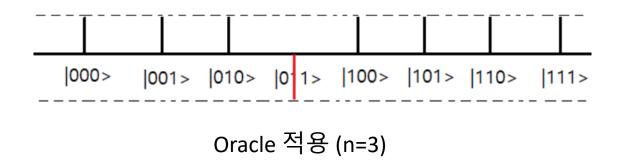
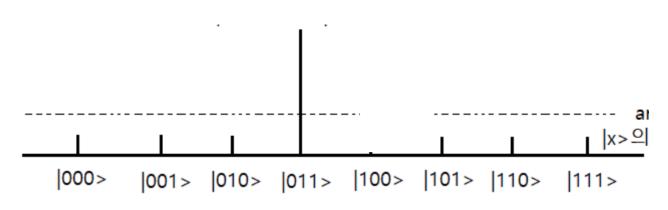


Fig. Grover algorithm (n=2)





Diffusion operator (n=3) 적용

Grover

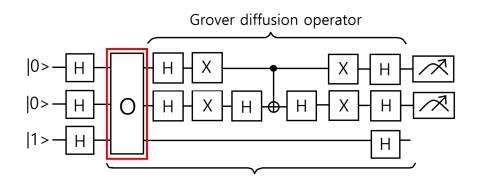
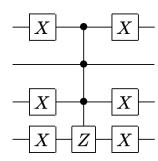


Fig. Grover algorithm (n=2)

Oracle

The oracle function performs a phase flip on the marked state. The phase flip inverts the amplitude α_{0010} of the state, making it $-\frac{1}{4}=-0.25$. Below the oracle for the state $|0010\rangle$ is shown. The corresponding QISKit code can be found in Appendix A.2, and a list of all oracles in Appendix B.1.



Grover Oracle + ISD

Challenge (Classic Mceliece encoding)

$$C_0 = He^{^{\scriptscriptstyle au}}$$
 라는 신드롬 계산 식에서 $\,C_0\,$ 와 $\,H\,$ 가 주어진다 해도

low – weight 벡터 e 를 찾아내기 매우 어려움 \rightarrow Finding low-weight codeword problem

Challenge

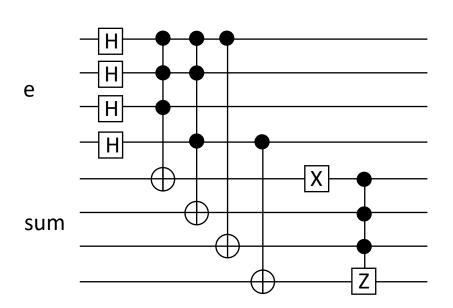
 $C_0 = He^{^{\scriptscriptstyle au}}$ 라는 신드롬 계산 식에서 C_0 와 H 가 주어진다 해도

low – weight 벡터 e 를 찾아내기 매우 어려움 \rightarrow Finding low-weight codeword problem

ISD Challenge

$$1110$$
 1100
 1000
 0101
 $C_k = 0111$
 $C_k = 0111$

 $c_k = H_k e$ 를 만족하는 특정 weight의 벡터 e 찾기



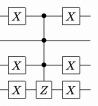
$$c_k = 0.1.1.1$$

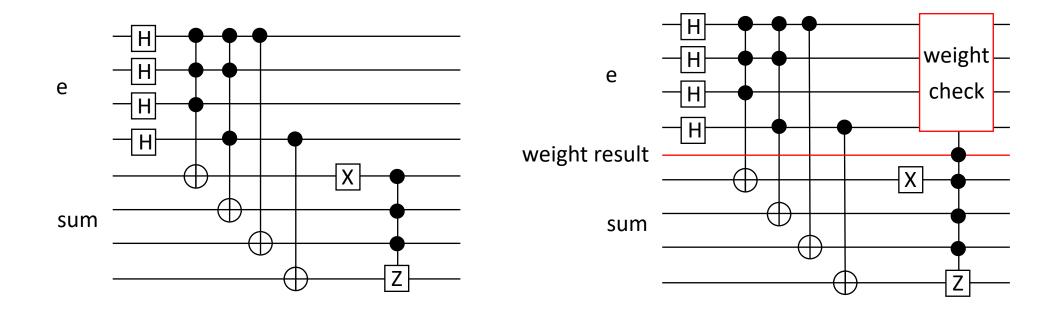
$$c_k = 0 1 1 1$$

$$c_k H_k^{-1} = e$$

Oracle

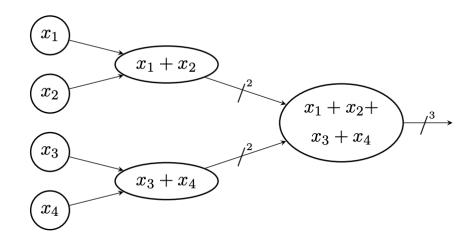
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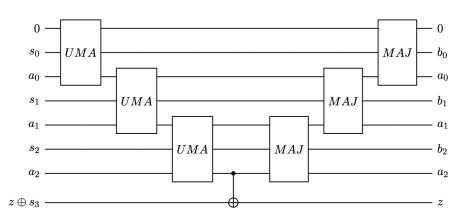


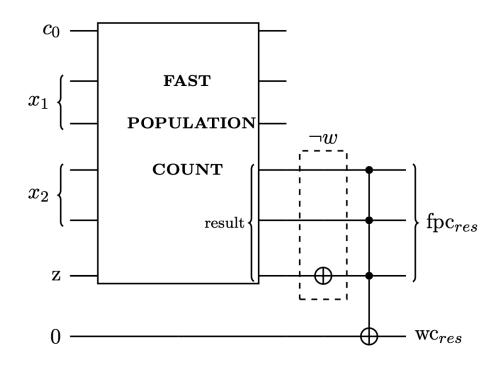


 $c_k = H_k e$ 를 만족하는 특정 weight의 벡터 e 찾기

Hamming weight 확인 (4-bit)

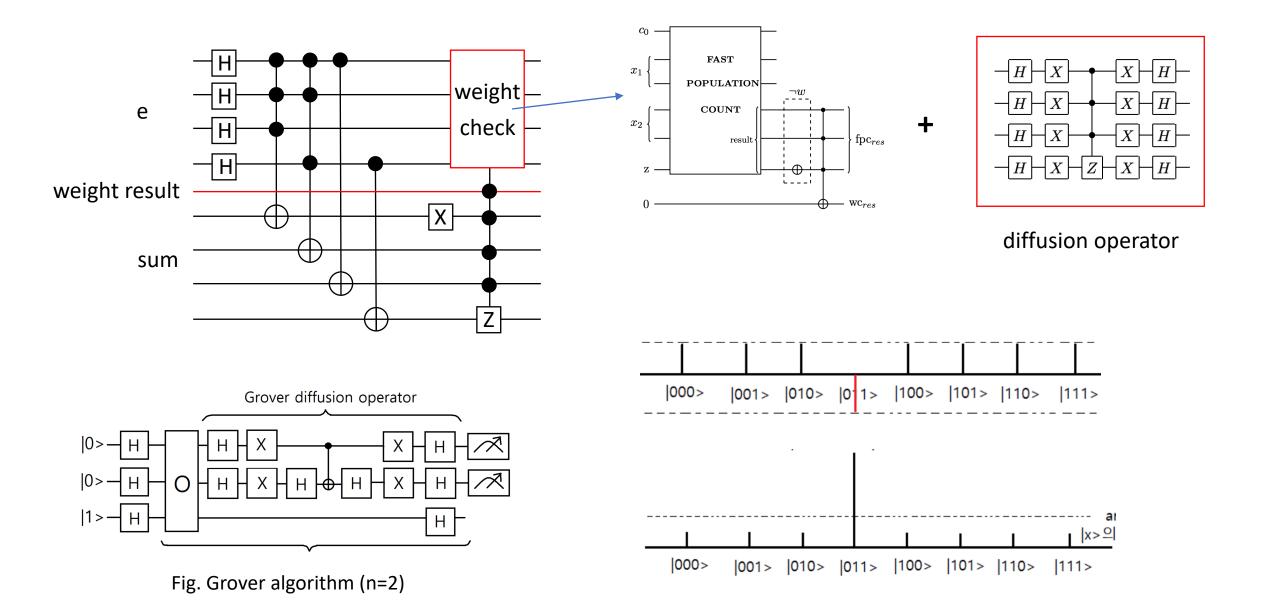






weight 가 011 (3) 인지 확인하는 회로

Grover Oracle + Diffusion

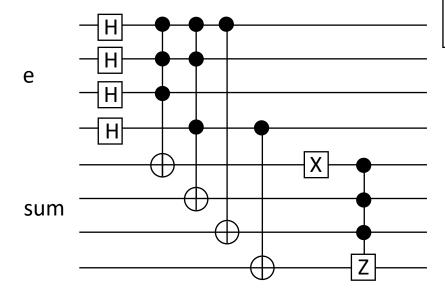


ISD Challenge

```
1110
1100
1000
0101
```

$$c_k = 0 1 1 1$$

```
H_k
```



```
mat = eng.allocate_qureg(4) # vector e
  target = eng.allocate_qureg(4) # syndrome s

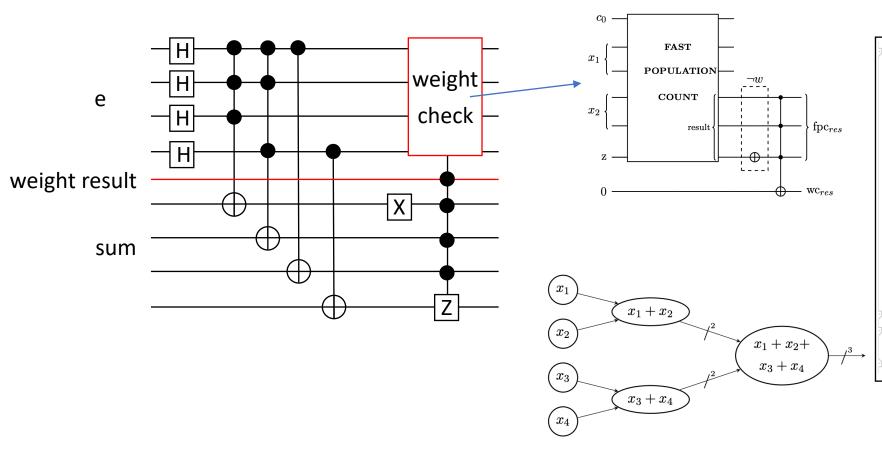
c = eng.allocate_qubit() # ripple carry bit
  z = eng.allocate_qureg(3) # 4-bit vector for weight check

w_res = eng.allocate_qubit() # weight result
  n = 15

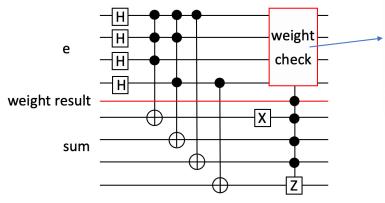
All(H) | mat

with Loop(eng, n):
  # oracle
  Set_matrix(eng, mat, target, c, z, w_res)
```

```
def Set_matrix(eng, mat, target,c, z,w_res):
   with Compute(eng):
       CNOT | (mat[0], target[0]) #1
       CNOT | (mat[1], target[0]) #0
       CNOT | (mat[2], target[0]) #0
       CNOT | (mat[0], target[1]) #1
       CNOT | (mat[1], target[1]) #0
       CNOT | (mat[3], target[1]) #1
      CNOT | (mat[0], target[2]) #1
       CNOT | (mat[3], target[3]) #1
      X | (target[0])
      H_weight(eng, mat, c, z, w_res)
   with Control(eng, w_res):
        with Control(eng, target[0:-1]):
            Z | target[-1]
   Uncompute(eng)
```



```
def H_weight(eng, vector, c, z, w_res):
   Toffoli | (vector[0], vector[1], z[0])
   CNOT | (vector[0], vector[1]) # z1, vector[1] -> 0 1
   Toffoli | (vector[2], vector[3], z[1])
   CNOT | (vector[2], vector[3]) # z2, vector[3] -> 0 1
   MAJ(eng, vector[1], vector[3], c)
   MAJ(eng, z[0], z[1], vector[1])
   CNOT | (z[0], z[2])
   UMA(eng, z[0], z[1], vector[1])
   UMA(eng, vector[1], vector[3], c)
                                          #z2 z1 v[3]
   X | z[2] # set weight condition : 011
   with Control(eng, z[2]):
       with Control(eng, z[1]):
           with Control(eng, vector[3]):
               X | w_res
```



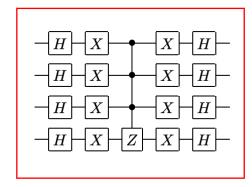
```
FAST
x_1
         POPULATION
            COUNT
x_2
                 result
```

```
def ISD(eng):
    mat = eng.allocate_gureg(4) # vector e
    target = eng.allocate_qureg(4) # syndrome s
    c = eng.allocate_qubit() # ripple carry bit
    z = eng.allocate_qureg(3) # 4-bit vector for weight check
    w_res = eng.allocate_qubit() # weight result
    n = 15
    All(H) | mat
    with Loop(eng, n):
        # oracle
        Set_matrix(eng, mat, target, c, z, w_res)
       # Diffusion
        with Compute(eng):
            All(H) | mat
           # All(H) | target
           All(X) \mid mat
            # All(X) | target
        with Control(eng, mat[0:-1]):
            # with Control(eng, target[0:-1]):
            Z \mid mat[-1]
                                                                            Uncompute(eng)
       Uncompute(eng)
```

```
fpc_{res}
```

def Set_matrix(eng, mat, target,c, z,w_res):

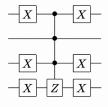
```
with Compute(eng):
   CNOT | (mat[0], target[0]) #1
   CNOT | (mat[1], target[0]) #0
   CNOT | (mat[2], target[0]) #0
   CNOT | (mat[0], target[1]) #1
   CNOT | (mat[1], target[1]) #0
  CNOT | (mat[3], target[1]) #1
  CNOT | (mat[0], target[2]) #1
  CNOT | (mat[3], target[3]) #1
  X | (target[0])
  H_weight(eng, mat, c, z, w_res)
with Control(eng, w_res):
    with Control(eng, target[0:-1]):
        Z | target[-1]
```



diffusion operator

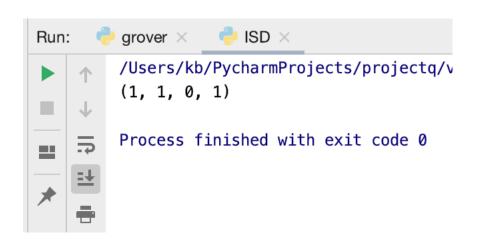
Oracle

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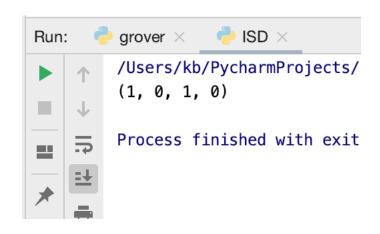


$$G' = S G P = \begin{cases} 1 & 1 & 1 & 1 & 0 & 0 & 0 & m = 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 & 1 & 0 & c = 0 & 1 & 1 & 0 & 1 & 1 & 0 \end{cases} \qquad c_k = 0 1 1 1 \qquad G'_k = \begin{cases} 1110 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 1000 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1001 & 0 & 0 & 0 & 1 & 0 & 0 \end{cases}$$

결과:1101 (weight = 3, n=3)



결과: n=1 → 1101이 자주 나오나, 다른 값도 나옴



감사합니다

