Quantum Information Set Decoding (Error-Free)

장경배

https://youtu.be/_cXkz7ql9ds





- Grover 알고리즘을 활용한 **코드기반암호** 분석의 경우 **대칭키 암호** 분석과 다름
- 대칭키 암호의 경우, n-bit **키의** Search space는 2^n , Grover 적용시 $\sqrt{2^n}$
- 코드기반암호의 경우 특정 무게의 n-bit 벡터를 찾는 경우, 실제 Search space가 2^n 이 아님

→ 이 경우 Grover가 비효율적으로 적용되어, **거의 동일한 성능**을 보여줌

McEliece	Workload Cryptanalysis		Minimal	Quantum-
parameters	(in binary operations)		number	computer
m,t	classic	quantum	of	bit
	computer		Qubits	security ⁷
11,32	2^{91}	2^{86}	25	80
11,40	2^{98}	2^{94}	50	88
12, 22	2^{93}	2^{87}	29	80
12,45	2^{140}	2^{133}	28	128

• Grover가 ISD(Prange)에 완벽히 적용될 수 있음을 제시, 복잡도는 절반으로 감소

Grover vs. McEliece

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Abstract. This paper shows that quantum information-set-decoding attacks are asymptotically much faster than non-quantum information-set-decoding attacks.

Key words. code-based cryptography, post-quantum cryptography

Author(s)	Year	$\max_{0 \le R \le 1} \alpha(R, \omega_{\text{GV}})$ to 4 dec. places
Prange [23]	1962	0.1207
Dumer [11]	1991	0.1164
MMT [18]	2011	0.1114
BJMM [4]	2012	0.1019
MO [19]	2015	0.0966

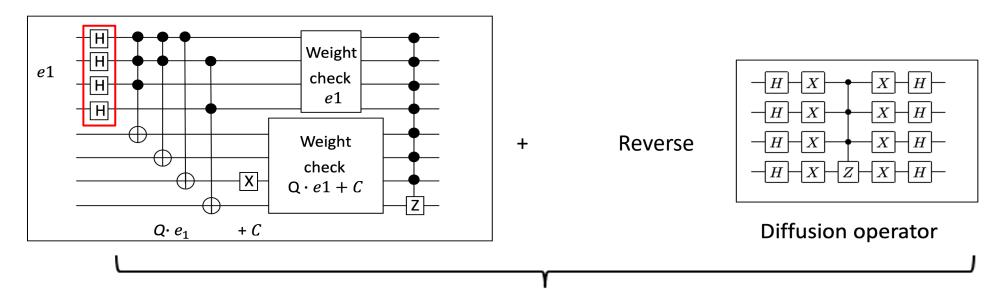
Author(s)	Year	Ingredients	$\max_{0 \le R \le 1} \alpha(R, \omega_{\text{GV}})$ to 5 dec. places
Bernstein [5]	2010	Prange+Grover	0.06035
This paper	2017	Shamir-Schroeppel+Grover+QuantumWalk	0.05970
This paper	2017	MMT+"1+1=0"+Grover+QuantumWalk	0.05869

• Error-free information set (Bernstein)

In this paper I will be satisfied with a limited class of b-bit-to-1-bit circuits,

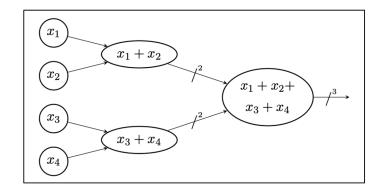
Hamming weight

• 현재 회로에서는 고려 X



n 번 반복하여 관측 확률 증가 (Grover Search space → e1-qubit)

• Weight check 에 대부분의 양자 자원이 사용됨



• Weight check 한번에 대한 비용이며, 추가로 Grover iteration의 배수만큼 더 소모됨

Gate counts:

Allocate: 6978

CCX : 12158

CX: 24318

Deallocate: 6978

Measure: 2720

Depth : 28215.

• 실제 Classic McEliece 파라미터에 대한 QISD

```
3.1 Parameter set kem/mceliece348864

KEM with m=12, n=3488, t=64. Field polynomials f(z)=z^{12}+z^3+1 and F(y)=y^{64}+y^3+y+z. This parameter set is proposed and implemented in this submission.
```

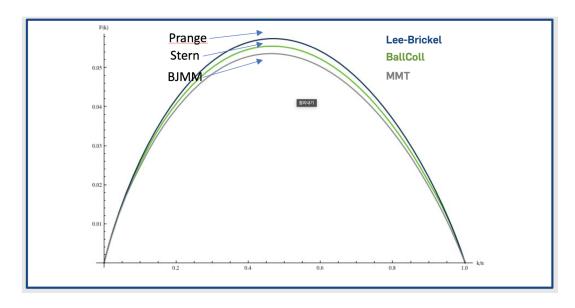
```
def QISD(eng):

# Goppa code size : (768 x 3488)
# Information set size : (768 x 2720)

e = eng.allocate_qureg(2720)
syndrome = eng.allocate_qureg(768)
temp = eng.allocate_qureg(2040)
temp2 = eng.allocate_qureg(576)
carry = eng.allocate_qureg(682)
carry2 = eng.allocate_qureg(191)

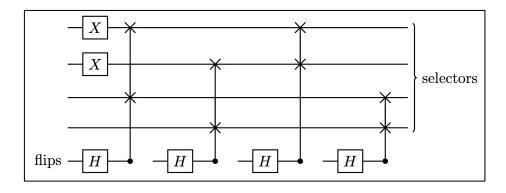
c0 = eng.allocate_qubit()
```

다양한 버전의 Infomration Set Decoding들이 있지만, 큰 성능 차이 x



- Quantum을 적용하였을 때 자원 측면에서의 효율성?
 - Lee-Brickel은 Weight check가 2번인 반면, Prange에서는 1번

• 특정 Hamming weight를 고려하는 중첩 상태의 Input 설계



< 중첩 상태의 4-qubit vector, **Hamming weight = 2 >**

• Search space 감소

```
\begin{array}{c}
1100 \\
1010 \\
2^4 \rightarrow 1001 \\
0110 \\
0101 \\
0011
\end{array}
```

```
def ISD(eng):
    vector = eng.allocate_qureg(4) # vector e
   flips = eng.allocate_qureg(4) # vector e
    target = eng.allocate_qureg(4) # syndrome s
   X | vector[0]
   X | vector[1] Hamming weight
   All(H) | flips
    for i in range(1): Iteration 횟수
        with Compute(eng):
            with Control(eng, flips[0]):
                Swap | (vector[0], vector[2])
           with Control(eng, flips[1]):
                Swap | (vector[1], vector[3])
           with Control(eng, flips[2]):
                Swap | (vector[0], vector[1])
           with Control(eng, flips[3]):
                Swap | (vector[2], vector[3])
```

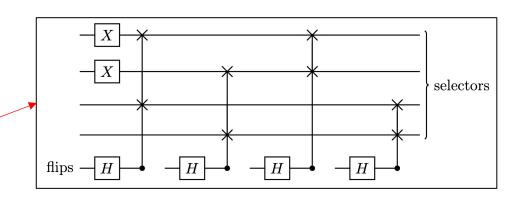
```
CNOT | (vector[0], target[2])
CNOT | (vector[1], target[2])

CNOT | (vector[2], target[1])
CNOT | (vector[3], target[1])

CNOT | (vector[2], target[2])

X | target[0]
X | target[1]
X | target[3]  # answer list : 1000, 0100, 0011,

with Control(eng, target[0:-1]):
        Z | target[-1]
Uncompute(eng)
```

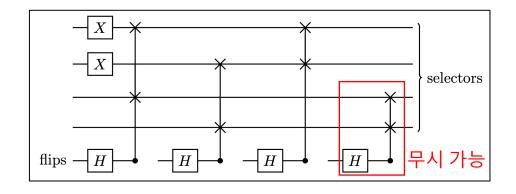


임의로 **0010**을 찾는 Oracle 설계, **Hamming weight 고려** Index 순서 (0123)

- 중첩 상태의 vector를 결정하는 건, flips 큐비트
 - flips 큐비트를 대상으로하여 Diffusion operator 수행
- 3개의 후보 중, Solution은 **0011**
- 4-qubit vector를 search, 반복 횟수는 1

```
def ISD(eng):
    vector = eng.allocate_qureg(4) # vector e
    flips = eng.allocate_qureg(4) # vector e
   target = eng.allocate_qureg(4) # syndrome s
   X | vector[0]
   X | vector[1]
   All(H) | flips
                                                 Oracle
   for i in range(1):
       with Compute(eng):
           with Control(eng, flips[0]):
               Swap | (vector[0], vector[2])
           with Control(eng, flips[1]):
               Swap | (vector[1], vector[3])
           with Control(eng, flips[2]):
               Swap | (vector[0], vector[1])
           with Control(eng, flips[3]):
               Swap | (vector[2], vector[3])
           CNOT | (vector[0], target[2])
           CNOT | (vector[1], target[2])
           CNOT | (vector[2], target[1])
           CNOT | (vector[3], target[1])
           CNOT | (vector[2], target[2])
           X | target[0]
           X | target[1]
           X | target[3]
                               # answer list : 1000, 0100, 0011
       with Control(eng, target[0:-1]):
           Z | target[-1]
       Uncompute(eng)
       with Compute(eng):
                                        Diffusion operator
           All(H) | flips
           All(X) | flips
       with Control(eng, flips[0:-1]):
          Z \mid flips[-1]
       Uncompute(eng)
       with Control(eng, flips[0]):
           Swap | (vector[0], vector[2])
       with Control(eng, flips[1]):
           Swap | (vector[1], vector[3])
       with Control(eng, flips[2]):
           Swap | (vector[0], vector[1])
       with Control(eng, flips[3]):
           Swap | (vector[2], vector[3])
   All(Measure) | vector
   All(Measure) | flips
```

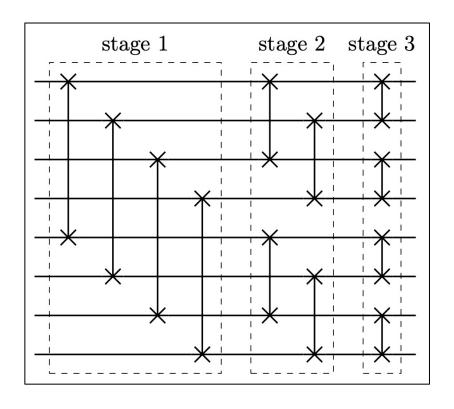
• Hamming Weight(2)에 따른 Flips 최적화 가능,



• 4-bit 벡터에서 Hamming weight가 2인 벡터들

Quantum ISD + Guessing Phase

- 6 x 8 의 Goppa code에 대한 QISD
- 8-bit 벡터를 대상으로 QISD, Weight 는 2



```
def ISD(eng):
    mat = eng.allocate_qureg(8)  # vector e, weight=2
    target = eng.allocate_qureg(6)  # syndrome s
    flips = eng.allocate_qureg(10)  # weight result

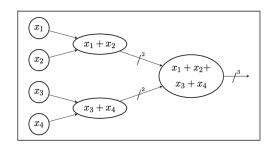
# n = 7  # optimal iteration number : pi/4 * arcsin(N^(1/2)) - (1/2)

print('Answer e = 0 0 0 0 1 1 0 0')
    X | mat[0]
    X | mat[1]

All(H) | flips
```

Conclusion

- Guessing phase를 같이 고려하면, ISD에 Grover search 알고리즘이 완벽히 적용될 수 있음
- Hamming Weight에 따른 flips 부분 최적화 가능
- Weight check에 필요한 자원들을 사용하지 않을 수 있음



```
# Goppa code size: (768 x 3488)
# Information set size: (768 x 2720)

e = eng.allocate_qureg(2720)
syndrome = eng.allocate_qureg(768)
temp = eng.allocate_qureg(2040)
temp2 = eng.allocate_qureg(576)
carry = eng.allocate_qureg(682)
carry2 = eng.allocate_qureg(191)

c0 = eng.allocate_queg(191)
```

• 더 알아봐야 함 → 완벽히 적용, 복잡도, flips 설계 및 최적화

감사합니다