

Depth Optimized Quantum Circuits for HIGHT and LEA

Kyungbae Jang, Yujin Oh, Minwoo Lee, Dukyoung Kim,
Hwajeong Seo

Contents

Our Contribution

Preliminaries

Proposed Quantum Circuits for HIGHT and LEA

Evaluation

Conclusion

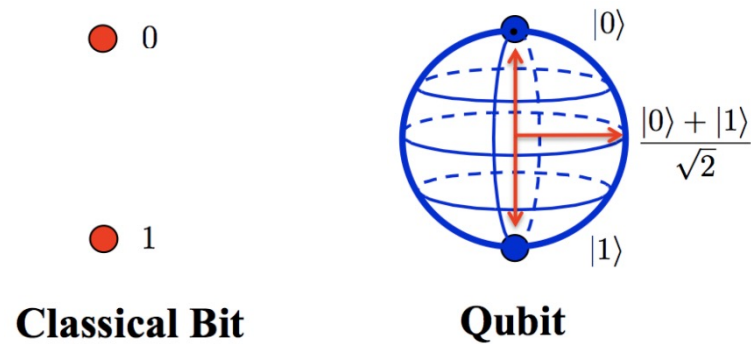
Contributions

- **Depth-optimized quantum circuits for LEA and HIGHT**
 - **We achieve depth reductions of 48% and 74% for HIGHT and LEA, respectively.**
- **Multiple methods for effectively reducing circuit depth are gathered in this work.**
 - **The implementation methods can be adopted for generic quantum circuit implementations.**
- **The required quantum complexities for HIGHT and LEA are redefined in this work.**
 - **Post-quantum security level for HIGHT and LEA are re-evaluated.**

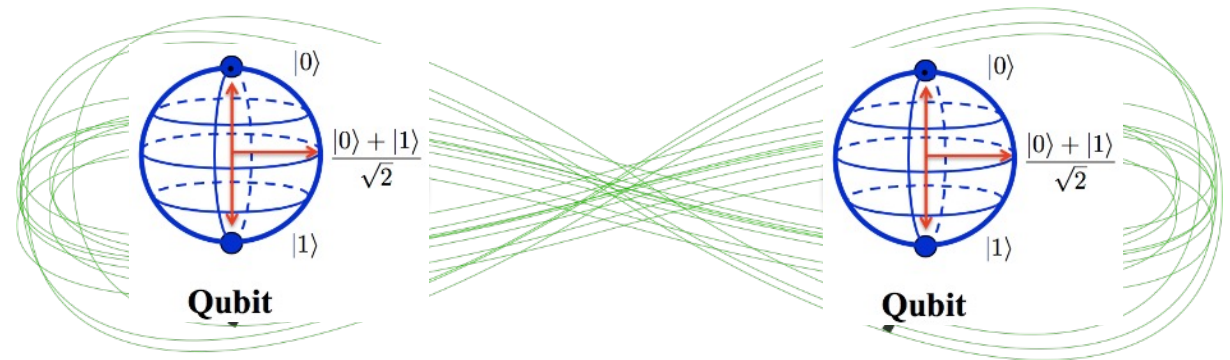
Quantum Computing

- Qubit (Quantum bit)

- Superposition

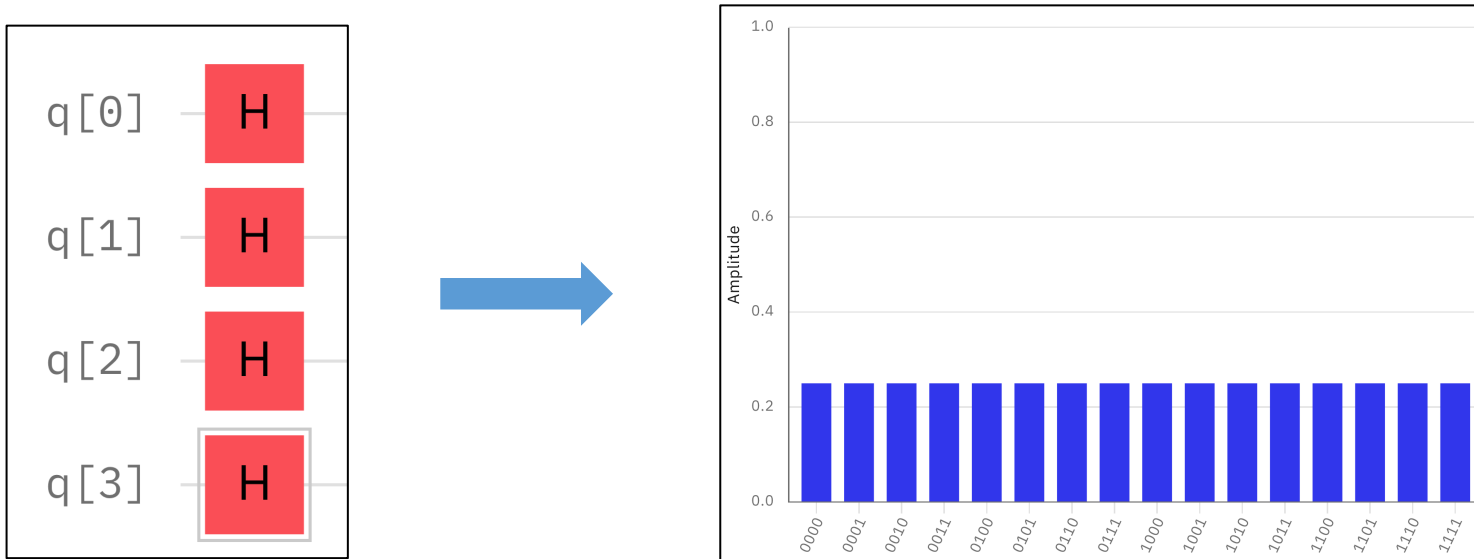


- Entanglement



Quantum Computing

- n -qubit with superposition state?

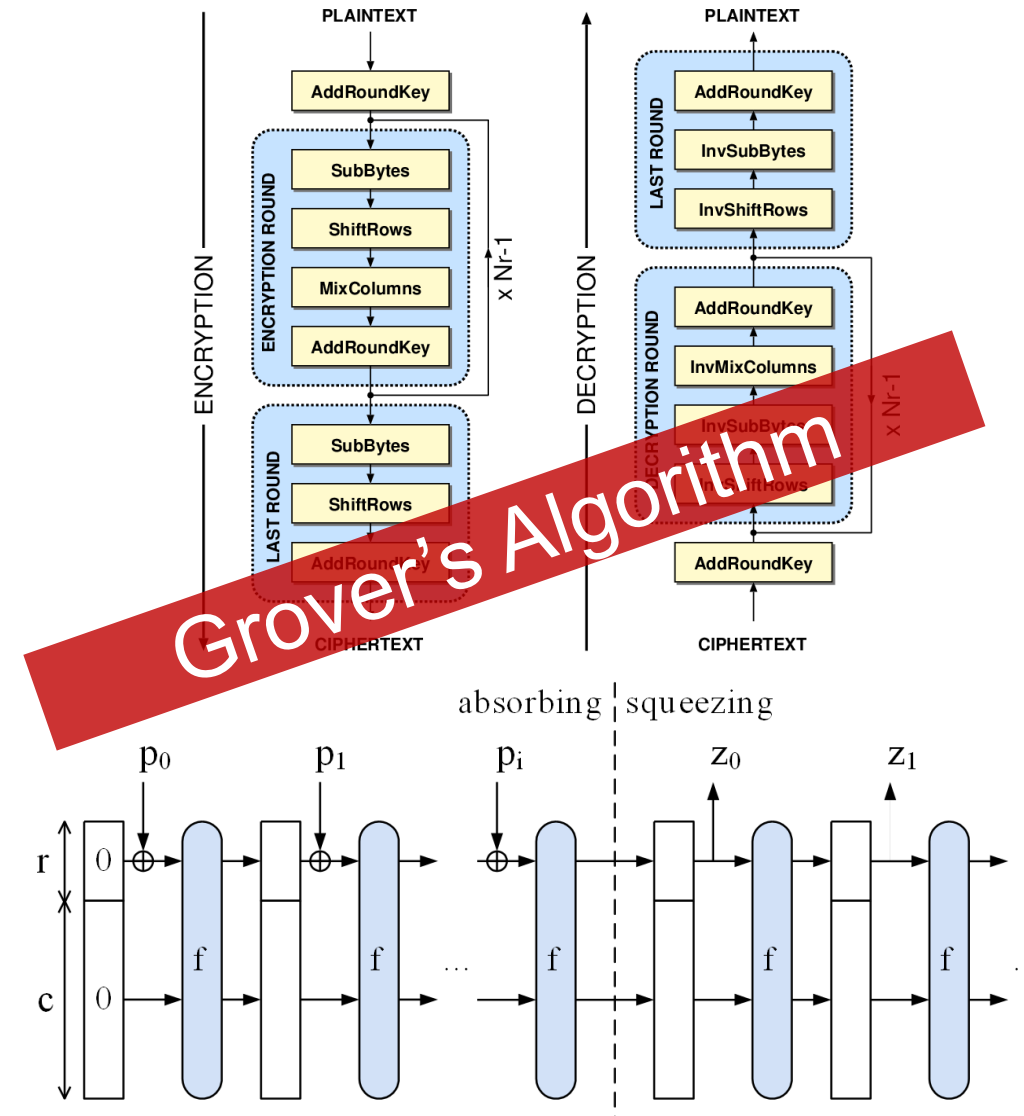
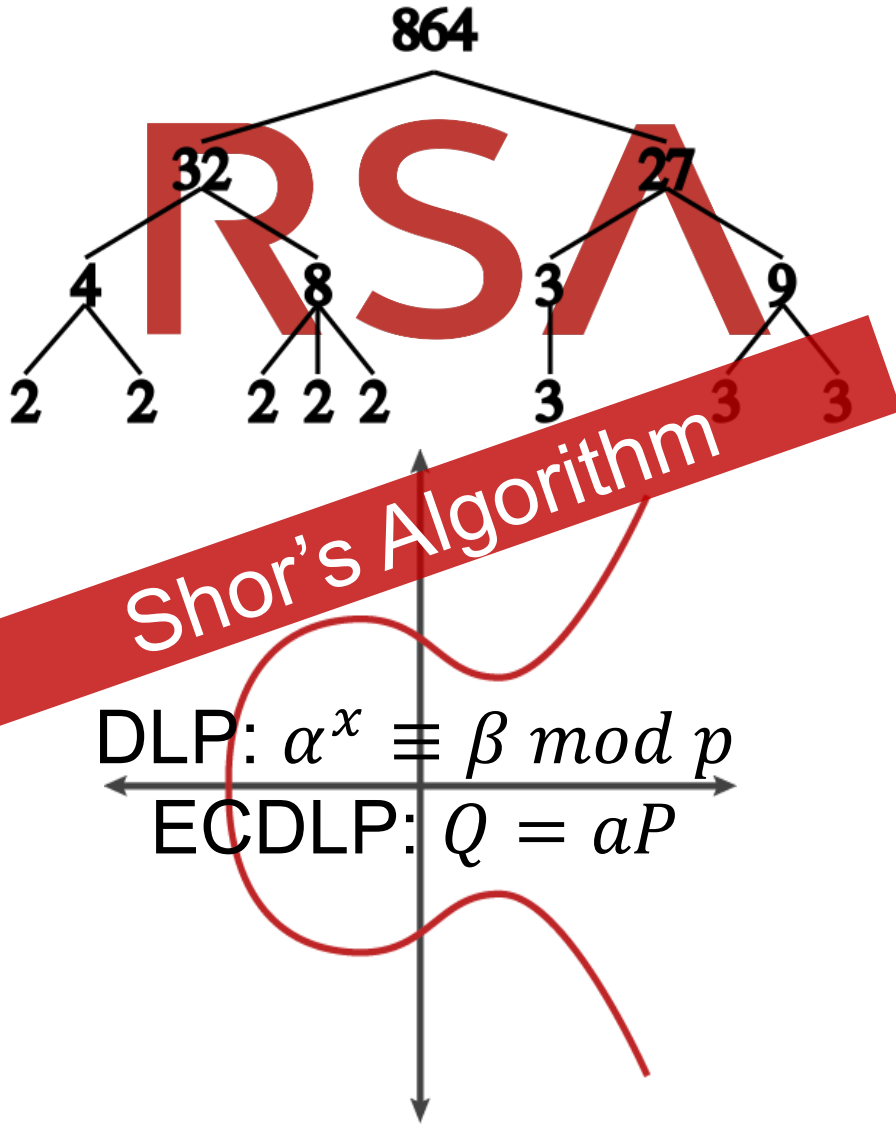


We can prepare 2^n states (as probability) at once!

With proper quantum algorithm? (Shor, Grover, Simon etc...)

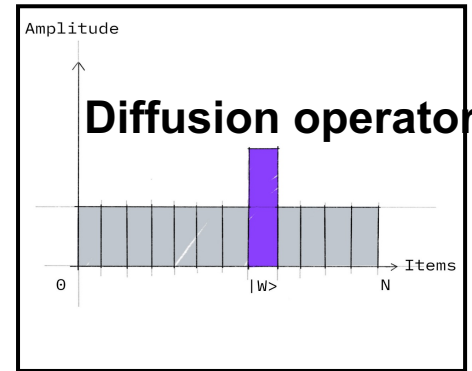
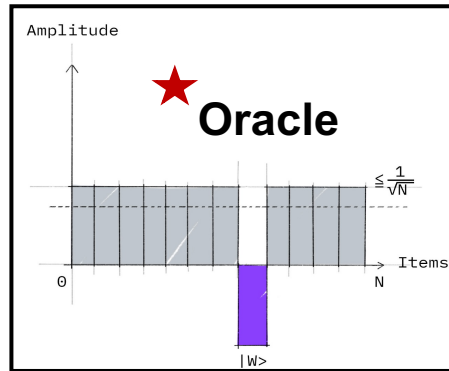
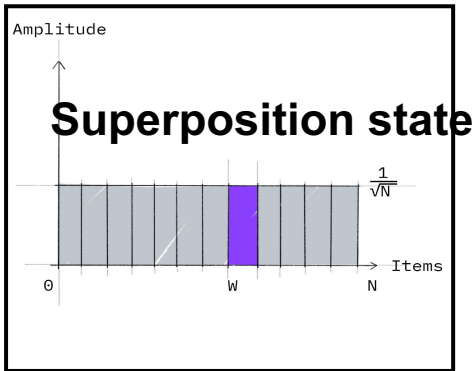
→ Meaningful result can be achieved

Cryptosystems in Quantum World



Grover's Algorithm

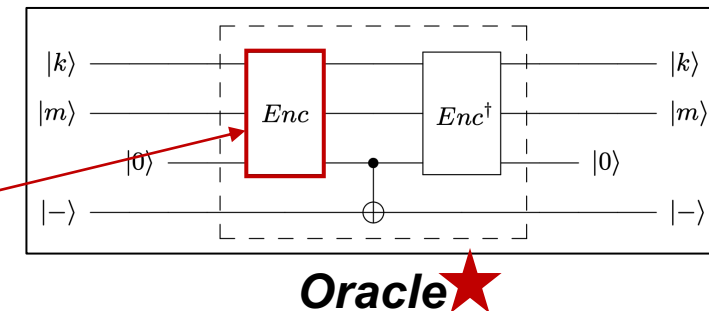
- Search complexity for N data elements
 - Classical: $O(N)$ → **Quantum (Grover): $O(\sqrt{N})$**



- Grover's key search for Symmetric key ciphers (k -bit key)
 - Prepare k -qubit in a superposition state (by using *Hadamard* gates)

$$|\psi\rangle = H^{\otimes k} |0\rangle^{\otimes k} = \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}} \right) = \frac{1}{2^{k/2}} \sum_{x=0}^{2^k-1} |x\rangle$$

- Implement a **quantum circuit for the target cipher**, then encrypt $\sqrt{2^k}$ times
 → **Optimization target** ★



Maximum Depth (MAXDEPTH)

- In Grover's search (single instance), numerous quantum queries are performed in sequential
 - **Total depth = Time-complexity.**
- NIST suggests the parameter, namely **MAXDEPTH (Maximum Depth)**

Level 1: 2^{40} Depth	MAXDEPTH	Cycle time (faster →)		
		$1\mu\text{s}$	200ns	1ns
Level 3: 2^{64} Depth	2^{40}	12.7 days	2.55 days	18.3 mins
	2^{48}	8.92 years	1.78 years	3.26 days
Level 5: 2^{96} Depth	2^{56}	2,280 years	457 years	2.28 years
	2^{64}	585,000 years	117,000 years	585 years

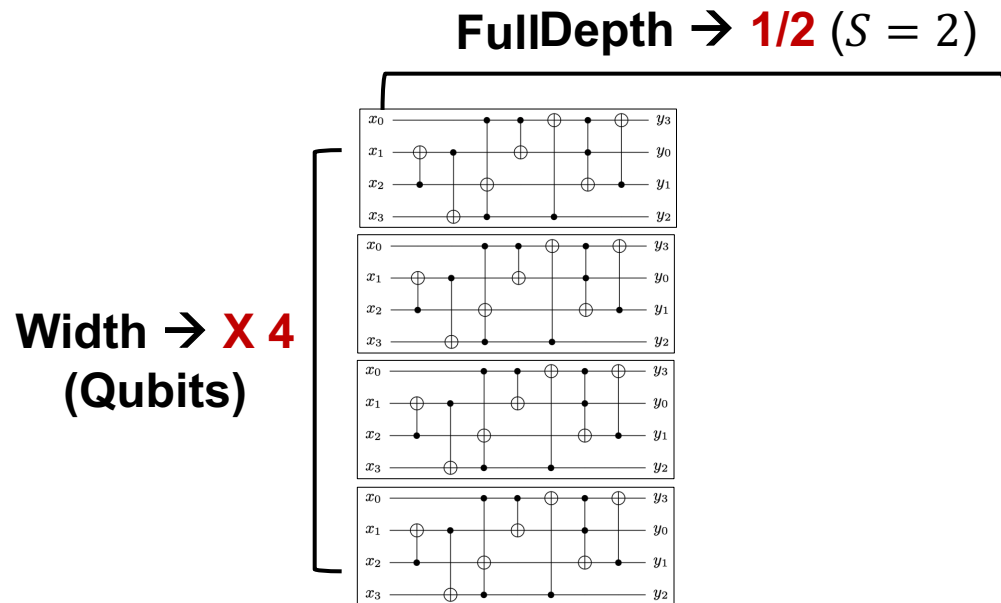
2.55 days: near-term and plausible

- If we do not satisfy the MAXDEPTH, Parallelization of Grover's search is required
→ Unfortunately, **Parallelization of Grover's search is poor**

Grover parallelization

- **Poor performance** of Grover parallelization
 - If we operate Grover Instances of S in parallel, depth is only reduced by \sqrt{S} .
→ The **DW-cost (Depth \times Width)** is transformed to the **D^2W -cost (Depth² \times Width)**

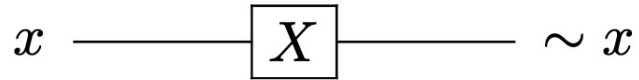
Example) if we want to **reduce the depth by half (i.e., 1/2)**, width is **increased by a factor of 4 (S=4)**



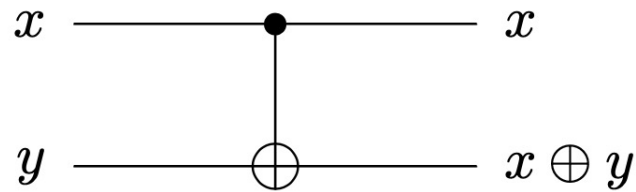
This is why we should optimize the depth for Grover's key search !!

Basic Quantum Gates

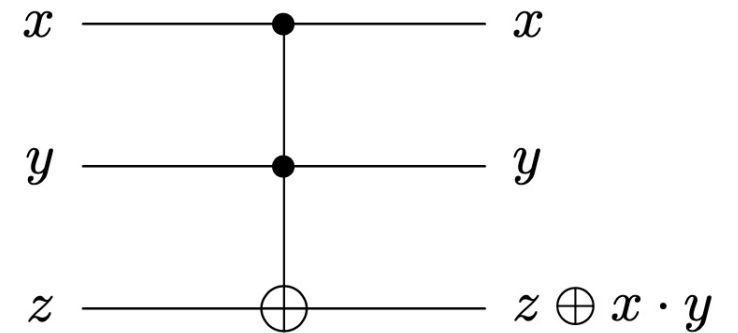
- The **NOT (X) gate** replaces **classical NOT operation**
- The **CNOT gate** replaces **classical XOR operation**
- The **Toffoli gate** replaces **classical AND operation**



(a) X (NOT) gate



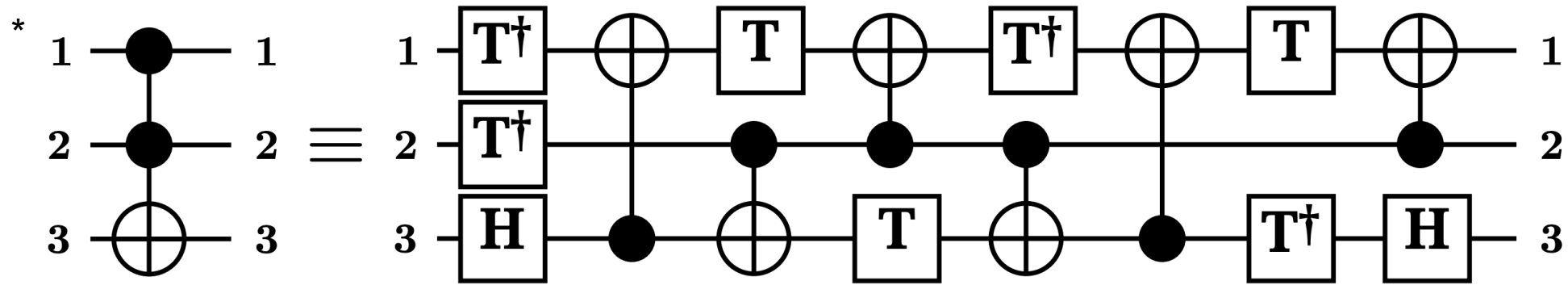
(b) CNOT gate



(c) Toffoli (CCNOT) gate

Toffoli gate

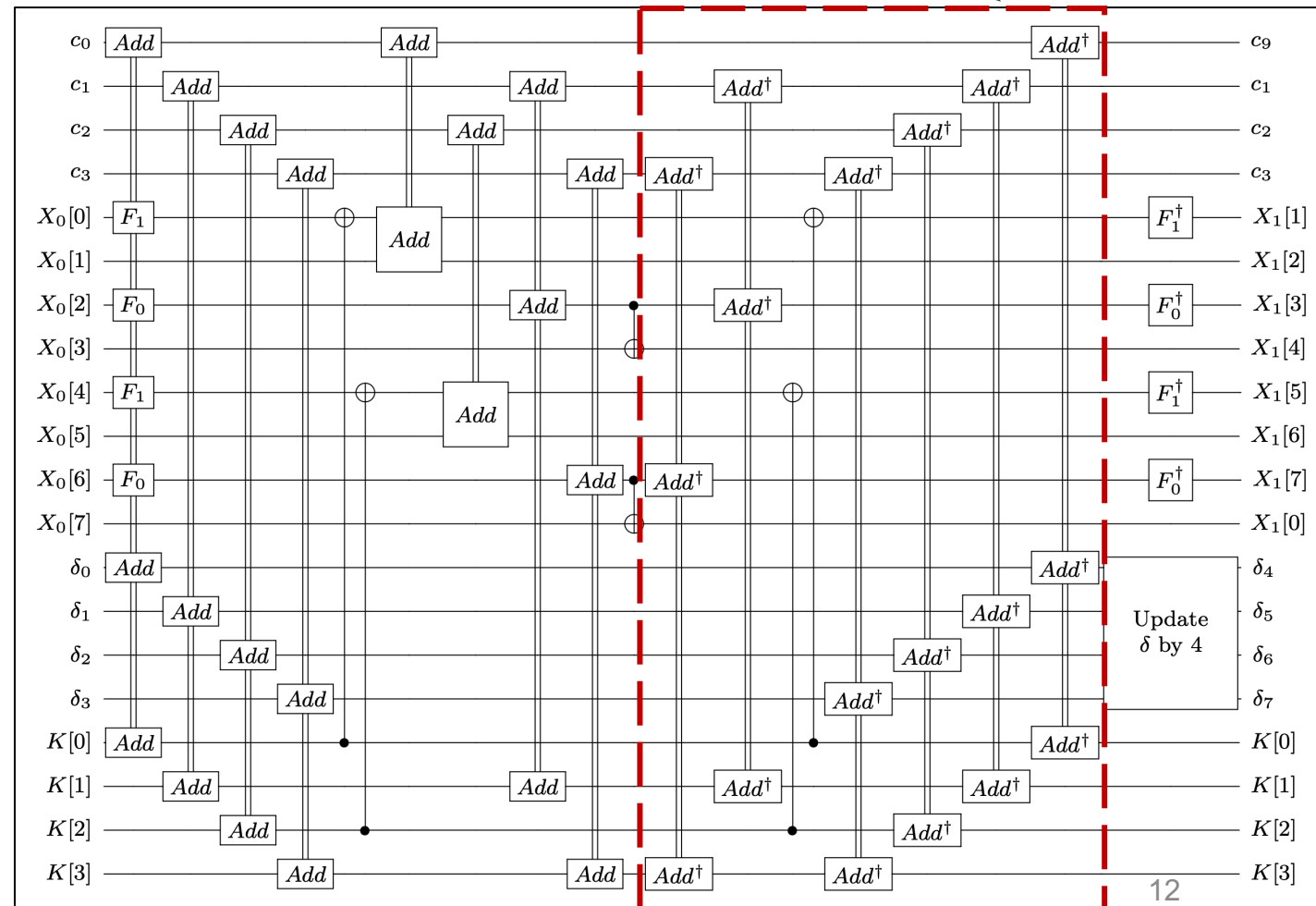
- Actually, the Toffoli gates are **more complex** than other quantum gates



(a) Toffoli gate (T -depth 4, total depth 8).

HIGHT quantum circuit

- We present a **Shallow architecture** for HIGHT
 - In the previous implementation (QIP'22), there is an overhead for the **reverse operation**
- In quantum implementations, the reverse operation is often utilized to initialize ancilla qubits and reuse them.
- In our Shallow architecture, **there is no depth overhead for the reverse operation.**



HIGHT quantum circuit

- In the previous work, **the subsequent round function is delayed** until the completion of the reverse operation of the current round function.
 - Since the current and subsequent round functions **share the ancilla qubits each other**.

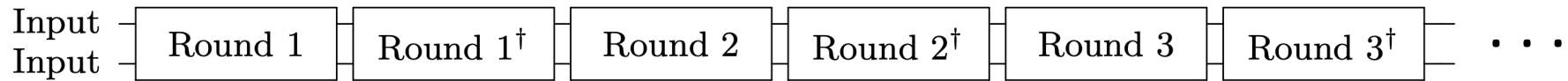


Fig. 2: The regular architecture adopted in [13]

- In the shallow architecture, the reverse operation of **the current round function is performed simultaneously with the subsequent round function** (i.e., in parallel).
 - we run **two sets of ancilla qubits** by allocating additional ancilla qubits for the subsequent round function.

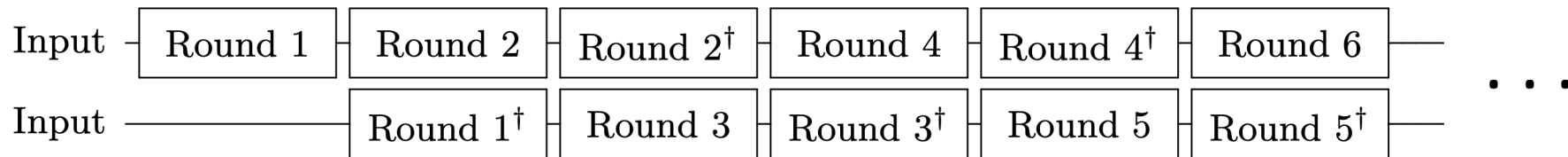


Fig. 3: The shallow architecture adopted in this work.

HIGHT quantum circuit

In HIGHT, linear layer operations which called $F_0(X)$ and $F_1(X)$ are given by:

$$F_0(x) = (x \lll 1) \oplus (x \lll 2) \oplus (x \lll 7)$$

$$F_1(x) = (x \lll 3) \oplus (x \lll 4) \oplus (x \lll 6)$$

In the previous work, **in-place implementation** was presented

→ **low qubit count** **but high circuit depth.**

We present an **out-of-place implementation**

→ **reduce the depth** **but increases the qubit count** (but we reuse them)

Operation	Source	#CNOT	#Qubit (reuse)	Depth
F_0	[14] and [13]	21	8	15
F_0	Ours	24	16 (8)	3
F_1	[14] and [13]	24	8	17
F_1	[14] and [13]	24	16 (8)	3

HIGHT quantum circuit

- We **effectively reduce the Toffoli/full depth** by allocating additional ancilla qubits.
 - **48% depth reduction**
 - All of the trade of metrics; $TD-M$, $FD-M$, TD^2-M , FD^2-M are optimized.

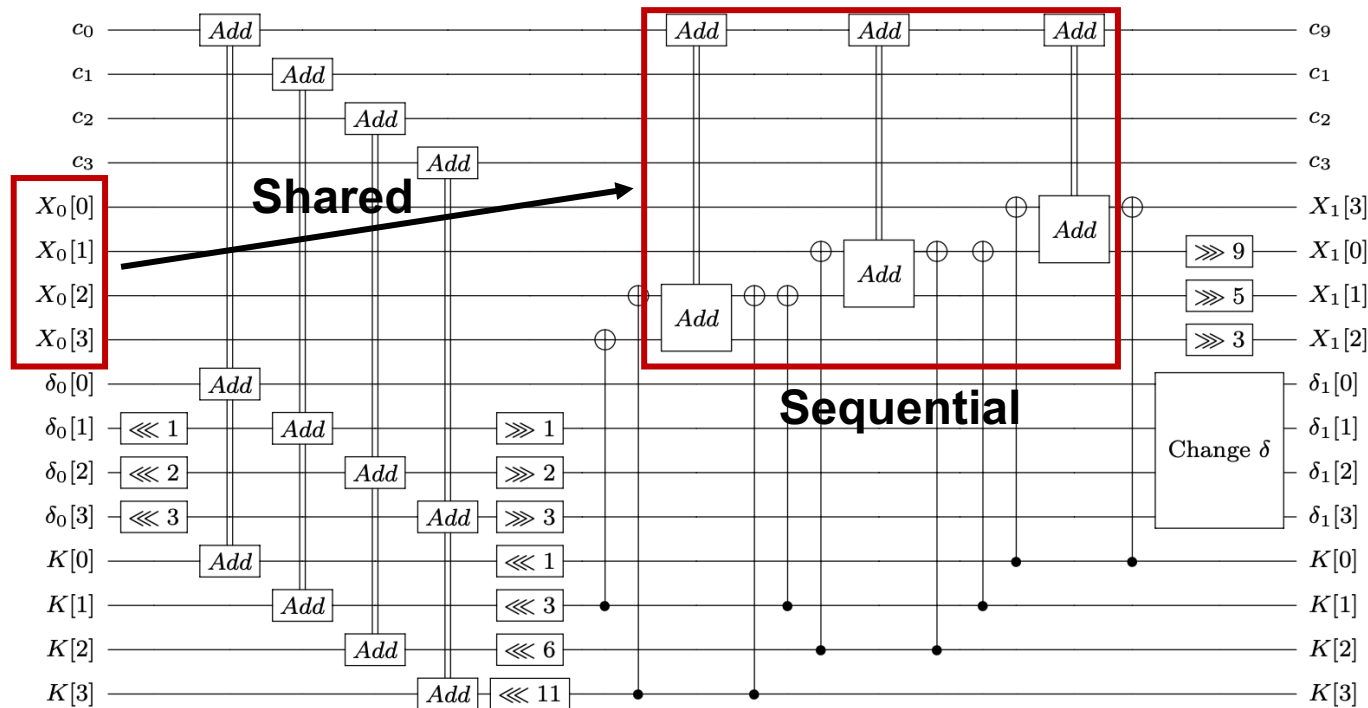
Source	#CNOT	#1qCliff	#T	Toffoli depth (TD)	#Qubit (M)	Full depth (FD)	$TD-M$	$FD-M$	TD^2-M	FD^2-M
[14]	64,799	13,444	50,176	.	201	68,415	.	$1.639 \cdot 2^{23}$.	$1.711 \cdot 2^{39}$
[13]	57,558	16,144	40,540	1,664	228	14,058	$1.447 \cdot 2^{18}$	$1.528 \cdot 2^{21}$	$1.176 \cdot 2^{29}$	$1.311 \cdot 2^{35}$
Ours	57,440	16,598	40,422	832	296	7,308	$1.879 \cdot 2^{17}$	$1.031 \cdot 2^{21}$	$1.527 \cdot 2^{27}$	$1.84 \cdot 2^{33}$

TABLE II: Quantum resources required for implementations of HIGHT.

LEA quantum circuit

Parallel Additions for Round Function

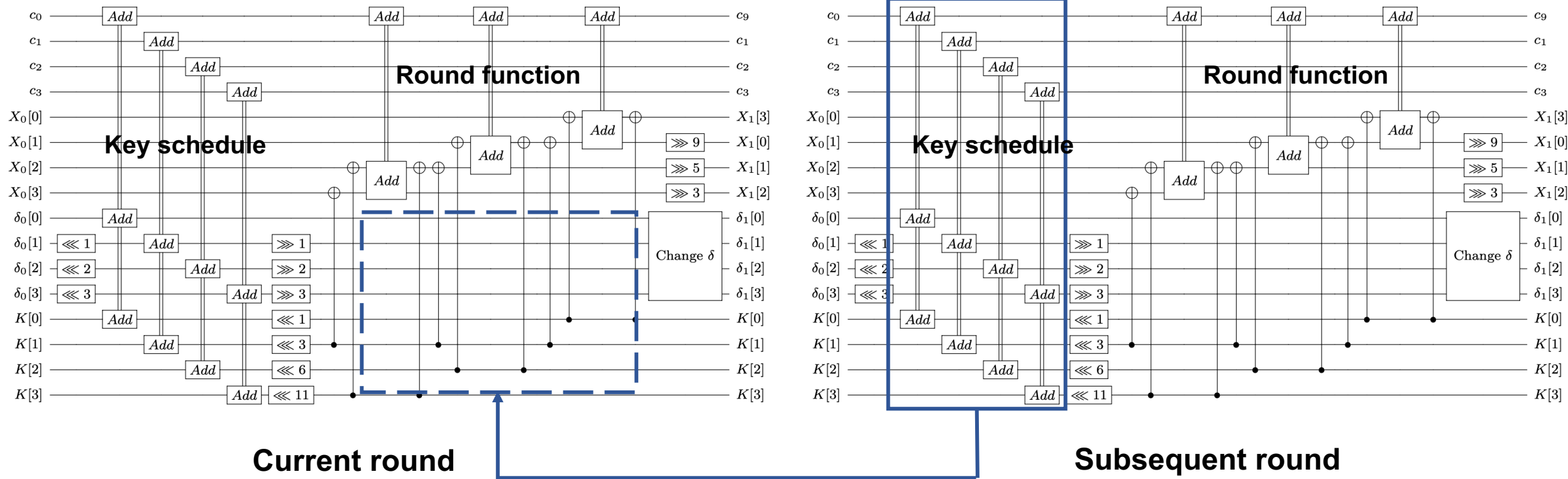
- In the previous implementation, **sequential additions** are performed.
 - Since the **inputs ($X_0[0] \sim [3]$) are shared** in the three additions.



- We perform the **three additions in parallel**.
 - To enable this, **we copy inputs ($X_0[0] \sim [3]$) before the additions**.

LEA quantum circuit

- The subsequent key schedule and the current round function can be executed in parallel.
 - As we did before, to enable this, we allocate additional ancilla qubits.



LEA quantum circuit

- We significantly reduce the Toffoli/full depth by allocating additional ancilla qubits.
 - **74% depth reduction**
 - Due to the **significant increases in qubit count**, the trade of metrics, **$TD-M, FD-M$, increases.**
 - However, **thanks to the depth optimization**, the trade-off metrics for parallelization, **TD^2-M, FD^2-M , are optimized.**

Cipher	Source	#CNOT	#1qCliff	#T	Toffoli depth (TD)	#Qubit (M)	Full depth (FD)	$TD-M$	$FD-M$	TD^2-M	FD^2-M
LEA-128	[14]	94,104	30,592	71,736	.	289	82,825	.	$1.427 \cdot 2^{24}$.	$1.803 \cdot 2^{40}$
	[13]	94,104	31,588	71,736	5856	388	47,401	$1.083 \cdot 2^{21}$	$1.096 \cdot 2^{24}$	$1.549 \cdot 2^{33}$	$1.586 \cdot 2^{39}$
	Ours	94,104	31,588	71,736	1,464	2,695	12,326	$1.881 \cdot 2^{21}$	$1.98 \cdot 2^{24}$	$1.345 \cdot 2^{32}$	$1.49 \cdot 2^{38}$
LEA-192	[14]	138,852	45,758	107,604	.	353	124,181	.	$1.306 \cdot 2^{25}$.	$1.238 \cdot 2^{42}$
	[13]	138,852	47,748	107,604	6832	518	55,301	$1.688 \cdot 2^{21}$	$1.707 \cdot 2^{24}$	$1.407 \cdot 2^{34}$	$1.441 \cdot 2^{40}$
	Ours	138,852	47,748	107,604	1,708	3,209	14,298	$1.307 \cdot 2^{22}$	$1.367 \cdot 2^{25}$	$1.09 \cdot 2^{33}$	$1.193 \cdot 2^{39}$
LEA-256	[14]	156,672	36,753	129,024	.	417	175,234	.	$1.089 \cdot 2^{26}$.	$1.456 \cdot 2^{43}$
	[13]	158,688	54,630	122,976	7808	582	63,108	$1.083 \cdot 2^{22}$	$1.095 \cdot 2^{25}$	$1.033 \cdot 2^{35}$	$1.054 \cdot 2^{41}$
	Ours	158,688	54,630	122,976	1,952	3,657	16,257	$1.702 \cdot 2^{22}$	$1.772 \cdot 2^{25}$	$1.622 \cdot 2^{33}$	$1.758 \cdot 2^{39}$

TABLE III: Quantum resources required for implementations of LEA.

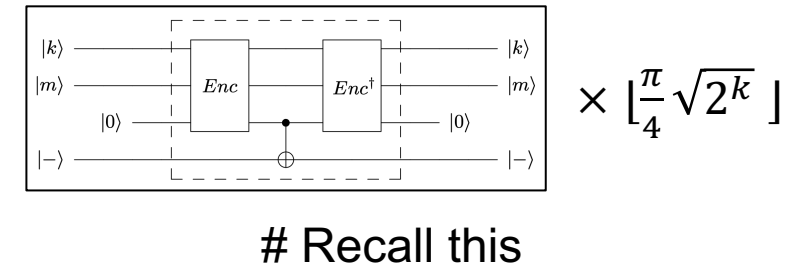
Evaluation of Post-Quantum Security

- Based on the optimized quantum circuits for Grover's key search, we estimate required resources for quantum key search.

Grover's key search (k -bit key) are estimated as follows: Quantum circuit $\times 2 \times \lfloor \frac{\pi}{4} \sqrt{2^k} \rfloor$.

Cipher	Total gates	Total depth	Complexity	NIST level
HIGHT	$1.372 \cdot 2^{81}$	$1.402 \cdot 2^{77}$	$1.924 \cdot 2^{158}$	Level 1 (2^{157})
LEA-128	$1.183 \cdot 2^{82}$	$1.182 \cdot 2^{78}$	$1.398 \cdot 2^{160}$	Level 1 (2^{157})
LEA-192	$1.763 \cdot 2^{114}$	$1.371 \cdot 2^{110}$	$1.209 \cdot 2^{225}$	Level 3 (2^{221})
LEA-256	$1.008 \cdot 2^{147}$	$1.558 \cdot 2^{142}$	$1.57 \cdot 2^{289}$	Level 5 (2^{285})

TABLE IV: Quantum resources required for Grover's key search for HIGHT and LEA.



- We evaluate the post- quantum security level suggested by NIST.
 - Level 1, 3, and 5** correspond to the attack complexity for **AES-128, -192, and -256**, respectively.
 - HIGHT and LEA achieve the appropriate post-quantum security** level according to the key size.

Conclusion

- **Multiple techniques** are gathered in this work **to effectively reduce circuit depth.**
(such as shallow architecture and copying for parallel operation)
- **Depth-optimized quantum circuits** offer optimal performance for Grover's key search.
 - We provide the **lowest quantum attack complexity** and the **best trade-off performance** for major metrics **under the depth constraint.**
- **We re-evaluate post-quantum security for HIGHT and LEA** (with NIST standard).
 - The quantum circuit of the target cipher is a **fundamental block in quantum cryptanalysis.**
 - Thus, the quantum circuits in this work can be utilized for other quantum algorithms (not only for Grover's exhaustive search).

Thank you!