

TODO

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Outline

- 1 Automatic Differentiation
- 2 Implicit QR Iteration
- 3 Shift Optimization Results
- 4 References

Automatic Differentiation (AD)

Big Picture

- At its core, computer code consists of arithmetic and logical operations
- AD applies elementary differentiation rules to the arithmetic operations

Can overload operators to achieve this.

Examples:

$$a + b \rightarrow a + b, \partial a + \partial b$$

$$a \cdot b \rightarrow a \cdot b, \partial a \cdot b + a \cdot \partial b$$

$$\sin(a) \rightarrow \sin(a), \cos(a) \cdot \partial a$$

...

Forward Mode

- Calculate effect on output due to perturbed inputs
- Input perturbation direction is provided to seed the calculation
- Notation: \dot{w} is the derivative w.r.t. input perturbation

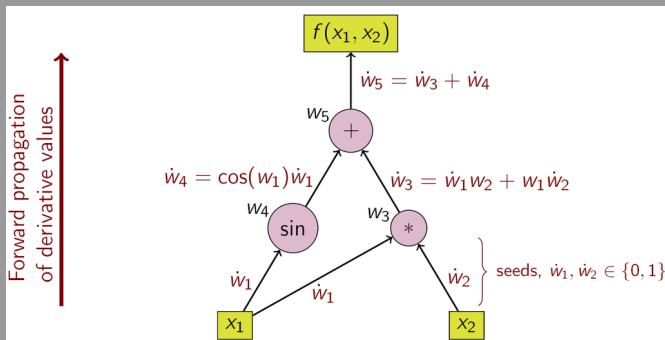


Figure: The subscripts represent intermediate steps¹

¹Figure from wikipedia's Automatic Differentiation page

Reverse Mode

- Calculate what input perturbation creates output perturbation
- Output perturbation direction is provided to seed the calculation
- Notation: \bar{w} is the derivative w.r.t. input perturbation

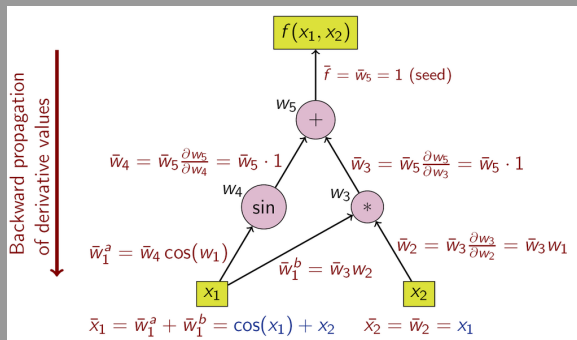


Figure: The subscripts represent intermediate steps²

²Figure from wikipedia's Automatic Differentiation page

Abstract Formulation

- An algorithm can be viewed as a map

$$f : \mathbb{R}^m \rightarrow \mathbb{R}^n$$

- Forward Mode calculates directional derivatives of f (the directions are the seeds)
- Reverse Mode calculates the gradient of f (if seeded with 1, otherwise is a scaled gradient)

Forwards vs. Backwards

For an algorithm $\vec{f} : \mathbb{R}^m \rightarrow \mathbb{R}^n$

Forward Mode

simple memory access

1-2 function evaluations for $\nabla \vec{f} \cdot \vec{d}$

m sweeps to compute $\nabla \vec{f}$

Backwards Mode

complicated memory access

1-10 function evaluations for ∇f_i

n sweeps to compute $\nabla \vec{f}$

Will focus on forward mode

Schur Decomposition AD

For $A \in \mathbb{R}^{n \times n}$, the Schur decomposition is

$$\begin{aligned} A &= QSQ' \\ Q'Q &= I \end{aligned}$$

with S upper triangular³

Thus

$$\begin{aligned} \dot{A} &= \dot{Q}SQ' + Q\dot{S}Q' + QS\dot{Q}' \\ Q'\dot{A}Q &= PS + SP' + \dot{S} \\ Q'\dot{A}Q &= PS - SP + \dot{S} \end{aligned}$$

Where $P = Q'\dot{Q}$, and the orthonormality of Q yields that $P = -P'$

³except 2x2 blocks on diagonal

Schur Decomp AD Algorithm

- Treat \dot{A} as the seeded input, define $B = Q'\dot{A}Q$
- Noting S, \dot{S} are upper triangular and P is skew symmetric, note

$$\begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix} = \begin{bmatrix} P_{11} & -P'_{21} \\ P_{21} & P_{22} \end{bmatrix} \begin{bmatrix} S_{11} & S_{12} \\ 0 & S_{22} \end{bmatrix} - \begin{bmatrix} S_{11} & S_{12} \\ 0 & S_{22} \end{bmatrix} \begin{bmatrix} P_{11} & -P'_{21} \\ P_{21} & P_{22} \end{bmatrix} + \quad (1)$$

- The (21) block yields $B_{21} = P_{21}S_{11} - S_{22}P_{21} + 0$, use triangular sylvester equation solver⁴
- Recursively solve the (11) and (22) blocks
- Compute $\dot{Q} = QP$ and $\dot{S} = B - Pt + SP$

⁴e.g. Octave's `sy1`. Or use the same splitting trick

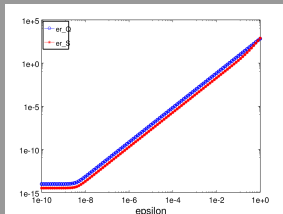
Schur Decomp AD Performance

Given a Schur decomposition $A = Q * S * Q'$, and seed \dot{A} :

$$A + \varepsilon \dot{A} = (Q + \varepsilon \dot{Q})(S + \varepsilon \dot{S})(Q + \varepsilon \dot{Q})' + O(\varepsilon^2)$$

$$(Q + \varepsilon \dot{Q})'(Q + \varepsilon \dot{Q}) = I + O(\varepsilon^2)$$

Following is the difference (error) observed as measured in the Frobenius norm.



On 10 random matrices in $\mathbb{R}^{400 \times 400}$, octave's built in schur function took 3.78 seconds (cumulative), and the schurAd algorithm spent 3.47 seconds (cumulative) calculating the derivatives.

Implicit multishift QR and Enhancements

Multishift QR

Multishift QR sweep:

function MQRSWEEP($H, \vec{\sigma}$)

form householder vector from 1st column of $\prod_{i=1}^s (H - \sigma_i I)$

Apply householder transformation to H using this vector

Reduce H to hessenberg form

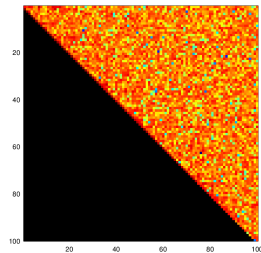
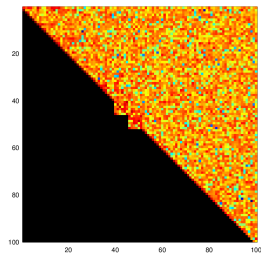
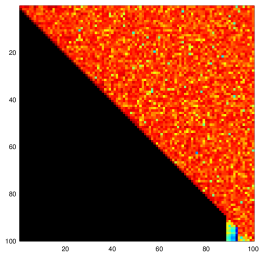
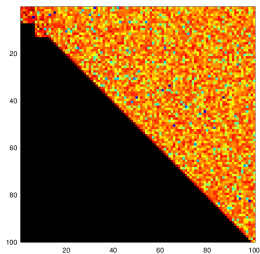
Deflate H if small entry on subdiagonal

return H , deflated eigenvalues

end function

In practice for $s > 5$, add the shifts in 5 at a time rather than all at once.

One 'Sweep'



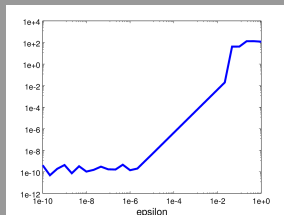
QRI AD Performance

Consider $f : \mathbb{R}^s \rightarrow \mathbb{R}^{n \times n}$, where the inputs are the shifts $\vec{\sigma}$, and the output is the matrix after a sweep of QRI.

Then we have in the direction \vec{d}

$$f(\vec{\sigma}) + \varepsilon \dot{f}(\vec{\sigma}) = f(\vec{\sigma} + \varepsilon \vec{d}) + O(\varepsilon^2)$$

Below is the difference (error) observed as measured in the Frobenius norm.



On 10 random matrices in $\mathbb{R}^{400 \times 400}$, using 20 shifts, it took 8.74 seconds (cumulative) to run a sweep of QRI, and 22.47 seconds (cumulative) to run a sweep of QRI with 1 directional derivative.

AD Definitions

- the shifts $\vec{\sigma}$ form a natural set of inputs
- for convenience, $f : \mathbb{R}^s \rightarrow \mathbb{R}^s$
- define $J_{i,j}(\vec{\sigma}) = \frac{\partial f_i}{\partial \sigma_j}(\vec{\sigma})$, so $J \in \mathbb{R}^{s \times s}$
- Assume we want to find $\vec{\sigma}$ such that $f(\vec{\sigma}) = \vec{0}$
- Newton's method provides an update procedure $\vec{\sigma}^{(n+1)} = \vec{\sigma}^{(n)} - J^{-1}(\vec{\sigma}^{(n)})f(\vec{\sigma}^{(n)})$

AD Variation on standard QRI

Standard Algorithm:

repeat

$$\sigma \leftarrow \text{eig}(H_{n-s:n, n-s:n})$$

$$H \leftarrow \text{MQRSSweep}(H, \sigma)$$

until $\text{length}(H) < 2$

AD Variation:

Choose the output of f to be the s bottommost subdiagonals

$$\sigma \leftarrow \text{eig}(H_{n-s:n, n-s:n})$$

repeat

$$\dot{H} \leftarrow \text{MQRSSweep}(H, \sigma) \text{ for seeded } I_{s \times s}$$

Deflate H if possible

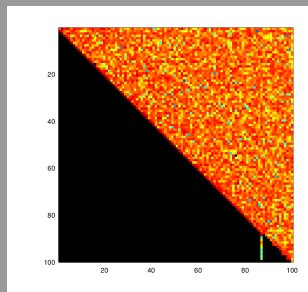
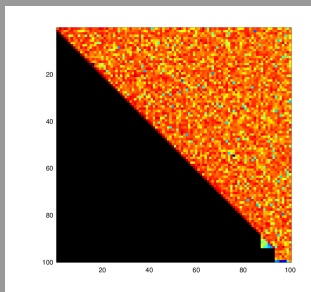
Form $J(\vec{\sigma})$

$$\vec{\sigma} \leftarrow \vec{\sigma} - J^{-1}(\vec{\sigma})f(\vec{\sigma})$$

until $\text{length}(H) < 2$

Aggressive Early Deflation

- drive the bulges through the matrix via QRI
- perform a Schur decomposition on the bottom portion containing the bulges, introducing a spike
- deflate matrix if spike values are below a tolerance

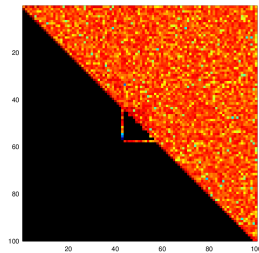
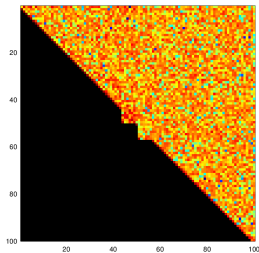
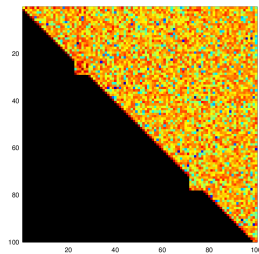
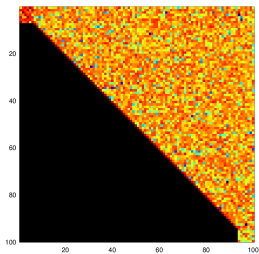


For the AD version, the outputs are the bottom s values on the spike.

Middle Deflation

- ① Bulges are driven from both ends
 - ② When they meet, a Schur decomposition is performed, generating 2 spikes
 - ③ If the tips of the spikes have enough zeros, can split matrix in the middle
- Performing this in the middle of QRI was proposed in Braman
 - the outputs for the AD algorithm will be the tips of the spikes

One 'Sweep' to middle



Implementation Details

- Schur decomposition is not unique since the eigenvalues can be reordered. Thus they were reordered to yield minimum value in the spike tips
- For aggressive early deflation, a second spike on the deflated matrix without another QR iteration can yield further deflatable eigenvalues []
- For all the AD methods, there is freedom to choose a larger or smaller set of outputs, but the optimization strategy would have to change.
- AD allows for further derivatives to be calculated in the same manner

Bibliography