VorVQ

Geometric Descriptions of Voronoi Tessellations of a Vector $$\operatorname{\textbf{Quantizer}}$$

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Preface

VorVQ exposes a collection of tools for computing quantities which describe the geometry of the cells of the first and second-order Voronoi tessellations induced by the learned prototypes (i.e., codebook vectors) of a vector quantizer. These

geometric descriptions are formed via computation of:

- Half-plane representations of the closed, convex polytopes comprising the Voronoi tessellation (i.e., the Voronoi cells, which we also call Voronoi polytopes in what follows).
- The Chebyshev centers [1] of each Voronoi polytope.
- The adjacency matrix of the Delaunay triangulation [2] of the prototype vectors in ℝ^d, for arbitrary dimension d. The well-known Qhull library (available for R in the geometry package) computed Delaunay triangulations for low (≤ 5) dimension; to our knowledge VorVQ is the only package to allow equivalent computation in high dimension. + The adjacency matrix of the Gabriel graph ([3]), which is a known sub-graph of the Delaunay triangulation. The R package cccd also exposes this functionality; it is re-cast here with support for parallel computation to expedite the calculation of the Delaunay triangulation.
- The Maximum Volume Inscribed ([4]) and Dikin ([5]) ellipsoids for each Voronoi polytope, which are useful for approximating the orientation and size (volume) of each Voronoi cell.

Calculation within VorVQ is implemented in C++ (via Rcpp and RcppArmadillo) with multi-threaded support (via RcppParallel).

VorVQ requires a valid installation of the Gurobi optimization library. Free academic licenses are available at Gurobi's website.

1 Background

1.1 Voronoi Tessellations

Vector quantization [6] is a common task in data compression whereby data vectors x are represented by a discrete set of n_W prototype vectors $W = \{w_j \in \mathbb{R}^d\}_{j=1}^{n_W}$ according to some assignment rule $j^* = BMU1(x)$, which returns the index j^* of x's **Best Matching Unit** in W. Likewise, an ancillary assignment rule $k^* = BMU2(x)$ returns the second Best Matching Unit to x (again, from W). The function BMU can be any valid distance; here, we only consider the most common case where the BMU is selected via Euclidean distance:

$$j^* = BMU1(x) = \arg\min_j \ ||x-w_j||_2$$

and

$$k^* = BMU2(x) = \arg\min_{k,k \neq j^*} \ ||x-w_k||_2.$$

The region $V_j \subset \mathbb{R}^d$ for which j is BMU1 is the first-order Voronoi cell generated by w_j :

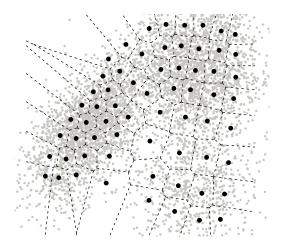
$$V_j = \{x : BMU1(x) = j\}$$

The collection of cells $\{V_j\}$ tessellates (or tiles) \mathbb{R}^d and is known as the first-order Voronoi tessellation. Similarly, the region $V_{jk} \subset \mathbb{R}^d$ for which j is

BMU1 and k is BMU2 is the **second-order Voronoi cell** generated by the interaction of w_i and w_k :

$$V_{jk}=\{x\,:\,BMU1(x)=j\quad\&\quad BMU2(x)=k\}.$$

Again, the collection $\{V_{jk}\}$ produces a (different) tessellation of \mathbb{R}^d . An example of the learned prototypes (black points) of synthetic two-dimensional data (gray points) and the induced first-order Voronoi tessellation (dashed lines) is given below:



One could, of course, extend this logic to produce higher-order $(3,4,\ldots,n_W-1)$ tessellations of \mathbb{R}^d (see [7] for further information on higher-order Voronoi tessellations). In VorVQ we only consider the first (Vor1) and second (Vor2) order cases, as they have been shown to faciliate manifold inference from learned vector quantizers ([8], [9]).

1.2 Voronoi Half-Plane Representations

The definitions of each Vor1 and Vor2 cell above specify convex polyhedra in \mathbb{R}^d . As such, each V_j or V_{jk} can be described geometrically using the half-plane representation [10] of their associated polyhedra, which is of the form $Ax \leq b$ for some coefficient matrix A and upper bound vector b. For example, transforming

the Vor1 definition above to its half-plane representation is straightforward:

$$\begin{split} V_j &= \{x \,:\, ||x-w_j||_2 \leq ||x-w_k||_2\} \quad \forall\, k \neq j \\ &= \{x \,:\, (w_k-w_j)^T x \leq \frac{1}{2} (||w_k||_2^2 - ||w_j||_2^2)\} \quad \forall\, k \neq j \\ &= \{x \,:\, A_j x \leq b_j\} \end{split}$$

where A_j is a $(n_W-1)\times d$ matrix and b_j is a (n_W-1) length vector. If global dimension-wise lower and upper bounds of the tessellated region are appended to the above system of inequalities then each Voronoi cell is guaranteed to be a closed polytope.

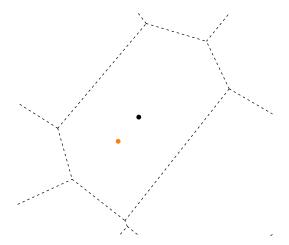
The form $V_j = \{A_j x \leq b_j\}$ is more amenable to performing optimization over each Voronoi cell. Various different objective functions, constrained to V_j , give rise to most of the computation on offer by VorVQ.

1.3 Chebyshev Centers

The computation of the ellipsoidal polytope approximators discussed in the ellipsoid section requires specifying an interior point to each Voronoi polytope. While Vor1 cells have a natural, known, interior point (the prototype w_j which generates each cell), Vor2 cells do not. There are myriad ways of finding interior points in polytopes; VorVQ provides functionality to compute each Voronoi cell's **Chebyshev center**, which is a relatively lightweight method for generating an interior point as the center of the largest sphere which can be inscribed in the polytope. Formally, center c is returned by the following linear program [1]:

$$\begin{aligned} & \text{maximize} & & r \\ & \text{subject to:} \\ & a_i^T c + r ||a_i|| \leq b_i \, \forall \, i \end{aligned}$$

where a_i and b_i are the *i*-th row of A and b as previously defined. Note that the above LP is only bounded if $Ax \leq b$ represents a closed polytope, which is ensured by appending the global bounds of the tessellated region to the system of linear inequalities. A Voronoi polytope's Chebyshev center (orange) is compared to its prototype location (black) for a two-dimensional Voronoi polytope below.



1.4 Prototype Proximity Graphs

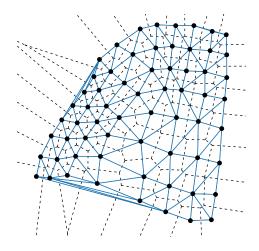
In general, a proximity graph of a vector quantizer's prototypes is a graph whose vertices are the prototypes and whose edges represent some type of neighbor relationship between prototypes. Probably the most famous proximity graph is the Delaunay triangulation [2] which connects prototypes j and k if (and only if) their Voronoi cells V_j and V_k intersect (share a face) in \mathbb{R}^d . Various methods exist to compute a Delaunay graph in low dimension (e.g., see the <code>geometry</code> package), but such methods typically involve storing a list of the vertices of the Delaunay "triangles" (simplexes, in higher dimension) and become infeasible as dimension grows. To avoid this bottleneck altogether <code>VorVQ</code> relies on the linear programming method of [10] to identify whether an edge exists between prototypes j and k in the Delaunay adjacency matrix, which we denote DADJ. Briefly, recall that the linear system defining V_j has components

$$A = \begin{pmatrix} (w_1 - w_j)^T \\ (w_2 - w_j)^T \\ \vdots \\ (w_{j-1} - w_j)^T \\ (w_{j+1} - w_j)^T \\ \vdots \\ (w_{n_W} - w_j)^T \end{pmatrix} \qquad b = \frac{1}{2} \times \begin{pmatrix} ||w_1||^2 - ||w_j||^2 \\ ||w_2||^2 - ||w_j||^2 \\ \vdots \\ ||w_{j-1}||^2 - ||w_j||^2 \\ ||w_{j+1}||^2 - ||w_j||^2 \\ \vdots \\ ||w_{n_W}||^2 - ||w_j||^2 \end{pmatrix}$$

Let a_k^T denote the the row of A involving w_j and w_k , and note that, if V_j and V_k share a face, a_k will be orthogonal to it. Consider the primal / dual LPs

$$\begin{array}{cccc} & \text{Primal} & & & \text{Dual} \\ \max_{x} & a_k^T x & & \min_{y} & b^T y \\ \text{s.t.} & Ax \leq b & & \text{s.t.} & A^T y = a_k \\ & & & y \geq 0 \end{array}$$

The primal optimum x^* will lie on the shared face between V_j and V_k if it exists; the dual optimum y_k^* corresponding to the k-th row of A will be nonzero (denoting the constraint is binding) if this is the case. Thus $DADJ_{jk} = DADJ_{kj} = 1$ if the dual optimum in the above is strictly positive: $y_k^* > 0$. By convention, $DADJ_{jk} = 0$ implies V_j and V_k are not adjacenct. Relying on symmetry, $n_W \times (n_W - 1)/2$ linear programs must be constructed and solved to determine all entries of DADJ. This is cumbersome but still feasible as each LP test can be performed in parallel. The Delaunay triangulation of the first-order Voronoi tessellation given above is:

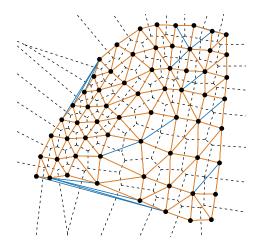


Although inherently parallelizable, the computational effort required to compute the full Delaunay adjacency matrix can be quite large. As discussed above a LP must be solved for each unique pair of first-order Voronoi cells in the tessellation, which is $\mathfrak{G}(n_W^2)$ separate optimizations. More, the complexity of each LP grows with the data dimension d. VorVQ utilizes parallel computation and the Gurobi solver which is, to our knowledge, the fastest LP solver available, but the cumulative time required to compute DADJ can still be improved.

To achieve this, VorVQ employs an intelligent search strategy to avoid solving the Delaunay LP for all $n_W(n_W-1)/2$ possible adjacencies. This strategy relies heavily on the Gabriel graph [3] of VQ prototypes, whose adjacency matrix we denote by GADJ. The Gabriel graph is another proximity graph which places an edge between vertices w_j and w_k IFF the ball whose diameter has endpoints w_j and w_k contains no other vertices w_l . That is, if w_{jk} is the midpoint between w_j and w_k , and $d(w_j, w_k)$ is their Euclidean distance, then

$$GADJ_{jk} = \begin{cases} 1 & \#\{l: d(w_l, \bar{w_{jk}}) < \frac{1}{2}d(w_j, w_k)\} = 0\\ 0 & \text{else} \end{cases}$$

The Gabriel graph of the learned prototypes of the synthetic two-dimensional data is shown below (in orange edges), overlain on the Delaunay graph from above (in blue).



It is clear that many, but not all, Delaunay edges exist in the Gabriel graph; what is not clear from the above, but has been proven [11], is that Gabriel graphs are proper sub-graphs of Delaunay graphs. Further, they are guaranteed to be connected. VorVQ exploits these two properties to effectively "seed" the Delaunay graph computation in two different ways. First, as a Delaunay sub-graph, any Gabriel edge is necessarily a Delaunay edge (so Gabriel edges do not need to be tested via the Delaunay LP). Second, as a sub-triangulation, we can intelligently focus attention on (and only test) special pairs of vertices for Delaunay adjacency. These "special pairs" are those vertices which, if connected, would help complete an existing partial triangle (simplex) in the GADJ; such

pairs necessarily are within geodesic distance = 2 on GADJ. After identification and Delaunay testing of such pairs, this process can be repeated until there are no un-tested vertex pairs of geodesic distance = 2 in the updated DADJ. Pseudo-code for this intelligent searching is given below:

Algorithm 1: GADJ Assisted Delaunay Computation

```
Input : GADJ adjacency matrix

Output: DADJ adjacency matrix

Set DADJ = GADJ;

Set isTested = GADJ;

Compute GD = geodesic distance matrix of prototypes, using DADJ;

Set TestSet = \{(j,k): \text{isTested}_{jk} = 0 \& GD_{jk} = 2\};

while \#\{TestSet\} \neq 0 do

Perform Delaunay LP test \forall (j,k) \in \text{TestSet};

Set DADJ_{jk} = 1 for each adjacency found in TestSet;

Set isTested<sub>jk</sub> = 1 \forall (j,k) \in \text{TestSet};

Update GD;

Update GD;

Update TestSet;
```

In experiments the above algorithm has resulted in a 20-90% reduction in the computational effort expended to compute the full DADJ, excluding the costing of computing the initial GADJ. Luckily, the latter also has a parallel implementation in VorVQ making the initialization step very fast for any set of prototypes of reasonable size for a vector quantizer.

1.5 Polytope Approximating Ellipsoids

As Voronoi polytopes are often highly irregular (e.g., they have drastically varying number of faces, or no observable symmetry), a study of their geometry usually proceeds by investigating a polytope approximator. Ellipsoids are a common approximator as their geometry is well understood and their functional form easy to manipulate. VorVQ offers two such approximators: the **Dikin** ellipsoid and **Maximum Volume I**inscribed **E**llipsoid (MVIE), each denoted by $\mathscr E$ and defined by a center vector c and a $d \times d$ rotation E of the following form:

$$\mathcal{E} = \{x \,:\, (x-c)^T E^{-1} (x-c) \leq 1\}.$$

1.5.1 Dikin Ellipsoids

The so-called Dikin ellipsoid $\mathscr{C}_{\mathfrak{D}}$, introduced by Dikin [5] (in Russian) but also discussed in [1], was originally put forth as a tool for an early interior point method for the solution of linear programs. Consider an arbitrary LP with objective g(y) and feasible region $P := Ay \leq b$, where $y \in \mathbb{R}^d$, $A \in \mathbb{R}^{m \times d}$ with rows a_i^T and $b \in \mathbb{R}^m$. At iteration k, given $y_k \in P$, the method computes the Dikin ellipsoid centered at y_k , $\mathscr{C}_{\mathfrak{D}}(y_k, E_{\mathfrak{D}}(y_k))$, and updates $y_{k+1} = \arg\min_{y \in \mathscr{C}_{\mathfrak{D}}} g(y)$

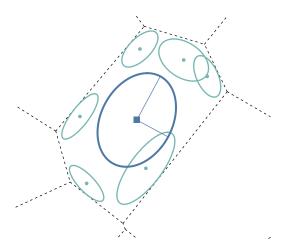
until convergence. The ellipsoidal transformation $E_{\mathfrak{D}}(y)$ is given by the Hessian of the logarithmic barrier function

$$\text{barrier}(y) = -\sum_{i=1}^m \log(b_i - a_i^T y)$$

which accumulates the log distances (i.e., slacks) to each of the hyperplanes defining P (and we note for completeness is the objective function whose argmin over y is the analytic center of P). Thus

$$E_{\mathfrak{D}}(y) = \frac{\partial^2 \mathrm{barrier}(y)}{\partial y^2} = \sum_{i=1}^m \left(\frac{1}{b_i - a_i^T y}\right)^2 a_i a_i^T = A^T \operatorname{diag}(1/(b - Ay))^2 \, A.$$

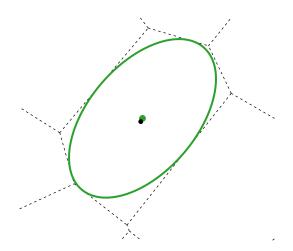
The Dikin ellipsoid approximates the local geometry of the polytope around y: bounding hyperplanes that are closer to y exert greater influence on its curvature and rotation than do those that are farther away. Several Dikin ellipsoids within a single first-order Voronoi cell are shown below: the blue ellipsoid is centered at the prototype generating the cell, while the teal are centered at various other points within the polytope showing how the Dikin ellipsoid conforms to local geometry.



1.5.2 Maximum Volume Inscribed Ellipsoids

Predictably, the MVIE $\mathscr{E}_{\mathcal{F}}$ is *inscribed* in a polytope $P := Ay \leq b$ with maximum volume. Determining the MVIE is not trivial, requiring sophisticated

interior point methods in convex optimization. In [4], Zhang and Gao develop a practical interior-point method for computation of MVIEs in arbitrary dimension. Beginning with any interior point $c_{\mathcal{F}}$ as initial center, the algorithm first computes the Dikin ellipsoid as an initial transformation $E_{\mathcal{F}}$. The remaining computational effort iteratively updates $c_{\mathcal{F}}$ and $E_{\mathcal{F}}$ to maximize $\log \det E_{\mathcal{F}}$ subject to $\mathcal{E}_{\mathcal{F}} \subset P$. Upon convergence, the optimal $\mathcal{E}_{\mathcal{F}}$ can be taken as a suitable ellipsoidal approximation of P, as the MVIE shown below (in green, with starting center point in black) indicates:



VorVQ computes the MVIEs for all cells in the (first or second-order) Voronoi tessellation using a version of Zhang's code translated to C++ from Matlab and adapted for parallel implementation.

2 Installation

VorVQ is available for Unix-like systems from github (not CRAN), installable via

```
devtools::install_github("somdisco/VorVQ")
```

Windows users can clone the repository and modify their Makevars files accordingly. The DESCRIPTION file in the package source lists the following 3rd-party R package dependencies:

```
RcppParallel,
Rdpack,
rcdd,
igraph
```

All Voronoi calculations are made efficient through use of the templated C++ class object VOR. For those who may desire custom or expanded functionality, the class template is implemented in header-only fashion in the inst/include directory of the package source.

Note that VorVQ requires a valid Gurobi installation and license on the user's local system to solve the various linear programs that arise in Voronoi computation. Gurobi can be obtained (free for academic users) by following its installation guide).

! Important Installation Note!

3 A Complete Example

3.1 Initialization

We demonstrate the calculations on offer by VorVQ with a simple, two-dimensional synthethic dataset known as CMIX9 which ships with the package for test purposes. The data are a sample (N=5000) from the 3-component Gaussian mixture defined in Table 1 of [12]. To facilitate demonstration the data were learned by a 10x10 Self-Organizing Map; the results of this learning (the learned prototypes, and the CADJ matrix) are available in the list variable CMIX9.

VorVQ is built around the internal C++ class VOR, which is a header-only class (found in the /inst/include/ source package) envisioned to be entirely portable for future integration with other projects. This section will demonstrate how to interact with the class fields and methods to compute Voronoi-related quantities (a complete list of the fields and methods of this class exposed to the user can be found with ?VOR). All methods are documented in usual R manner, with help functionality available via ?<method_name>.

To begin, an instance of the VOR class must be initialized with the prototype and CADJ matrix:

```
library(VorVQ)
## Loading required package: Rcpp
vor = VOR$new()
vor$initialize_VOR(CMIX9$W, CMIX9$CADJ)
## Initializing Voronoi object
## ++ storing W and dimensions ... done
## ++ setting global bounds = +/- 5% of W range ... done
      call $set_bounds to change
## ++ storing CADJ ... done
## ++ setting active vor1 indices (default to non-empty vor1 cells) ... done
      call $set_vor1_active to change
## ++ setting active vor2 indices (default to non-empty vor2 cells) ... done
      call $set_vor2_active to change
## Calculating Gabriel graph edges ...
## 15%
              30%
                         45%
                                    60%
                                                75%
## 90%
              100%
## Computation time: 3.08948e-05 minutes
```

Note that when calling library(VorVQ) a new Gurobi environment is also initialized, which requires a valid Gurobi license as discussed in the Installation section.

Initialization stores the number of prototypes in the tessellation, the data dimension, and the input prototype and CADJ matrices, which can be accessed via

```
str(vor$W)
## num [1:100, 1:2] -1.976 -1.741 -1.453 -1.158 -0.741 ...
# Data dimension
vor$d
## [1] 2
# Number of prototypes
vor$nW
## [1] 100
# The prototype and CADJ matrices, copied from those input
str(vor$W)
## num [1:100, 1:2] -1.976 -1.741 -1.453 -1.158 -0.741 ...
str(vor$CADJ)
## num [1:100, 1:100] 0 26 0 0 0 0 0 0 0 ...
```

Global, dimension-wise bounds are required to ensure the tessellated region is bounded (which, in turn, ensures the Voronoi polytopes are closed). During initialization, default bounds = +/-5% of the dimension-wise range of the prototypes (in W) are saved to the field 1b and ub:

```
vor$1b
## [,1]
```

```
## [1,] -2.171620

## [2,] -2.233343

vor$ub

## [,1]

## [1,] 2.139453

## [2,] 2.511854
```

If different bounds are desired they can be set by passing new d-dimensional lower and upper bound vectors to the method set_bounds. Here, we simply copy the existing bounds and invoke the set_bounds method for demonstration:

```
newlb = vor$lb
newub = vor$ub
vor$set_bounds(newlb, newub)
```

Note that the range spanned by 1b and ub *must include* the apparent range of the prototypes in W. When calling set_bounds manually a check of such will be performed, and an error returned if any prototype vector in W is found to be outside the given bounds.

From the CADJ weights we can infer which of the first and second-order Voronoi cells of the tessellation by W contain training data. We call such cells "active", and a default list of the active indices is computed and stored during initialization.

```
str(vor$vor1_active)
## num [1:100, 1] 1 2 3 4 5 6 7 8 9 10 ...
str(vor$vor2_active)
## num [1:464, 1:2] 2 11 1 3 12 2 4 5 12 3 ...
```

The (1-based) indices in the above refer to the rows of W which generate the respective Voronoi cells. Computation of first and second-order Voronoi cell centers, MVIEs and Dikin ellipsoids is restricted to these "active cell" lists, which can be changed via set_vor1_active and set_vor2_active. For ex., to re-set the active lists to the defaults computed above:

```
vor1_indices = vor$vor1_active
vor$set_vor1_active(vor1_indices)

vor2_indices = vor$vor2_active
vor$set_vor2_active(vor2_indices)
```

Note that only those first-order Voronoi cells i and j which are Delaunay adjacency define a valid second-order Voronoi cell. If the Delaunay adjacency matrix DADJ has already been calculated (see below), the vor2_indices list (as input to set_vor2_active) will be checked to ensure all given vor2_indices are Delaunay adjacent. If any are not an error will be returned.

The expedite the calculation of the Delaunay adjacency (as discussed in section

Background), the Gabriel proximity graph of the the prototypes is automatically computed and stored in the field GADJ during initialization, which can be accessed via

```
str(vor$GADJ)
## num [1:100, 1:100] 0 1 0 0 0 0 0 0 0 0 ...
```

To compute a stand-alone Gabriel graph from a given prototype matrix, see <code>?Gabriel_ADJ</code>.

Several parameters are needed to control the methods used in the VorVQ package. Default parameter values, which can be viewed via method get_params(), are set during initialization:

```
vor$get_params()
## $parallel
## [1] TRUE
##
## $MVIE_maxiter
## [1] 100
##
## $MVIE_tol
## [1] 1e-04
##
## $MVIE_fix_c0
## [1] FALSE
##
## $seed_DADJ
## [1] TRUE
```

The list elements above indicate whether parallel computation is allowed (=T), and certain parameters required for MVIE computation: the max. number of iterations allowed, the convergence tolerance, and whether updating of the ellipsoid center is allowed (MVIE_fix_c0 = F) or prohibited (MVIE_fix_c0 = T). To set any of these parameters, pass a list object to the set function whose names match that of the parameter of interest. Partial lists are allowed. For ex., to disable and then re-enable parallel computation:

```
# Disable parallel
vor$set_params(list(parallel = F))
# Enable parallel
vor$set_params(list(parallel = T))
```

3.2 Delaunay Adjacency

The Delaunay adjacency of first-order Voronoi cells is computed via the method calc_DADJ and stored in the field DADJ:

```
# Calculate the Delaunay adjacency
vor$calc DADJ()
## Testing Delaunay graph edges ...
## Round 1 (456 edges)
                                                      75%
                                                                 100%
                               25%
                                          50%
## Round 2 (45 edges)
                              25%
                                         50%
                                                     75%
                                                                 100%
## Tested 501 of 4950 possible edges
## Computation time: 0.00135781 minutes
# View its structure
str(vor$DADJ)
   num [1:100, 1:100] 0 1 0 0 0 0 0 0 0 0 ...
```

If a DADJ is available from an external source, it can be set (and, hence, avoid its potentially expensive re-calculation) via the <code>set_DADJ</code> method. Here, we extract the previously calculated <code>DADJ</code> and pass it back to the <code>VOR</code> object for demonstration of setting from an externally sourced adjacency matrix:

```
newDADJ = vor$DADJ
vor$set_DADJ(newDADJ)
```

Note that while DADJ calculation for a vector quantizer in high dimension with many prototypes can be an expensive, it is a required step for all other computations exposed by VorVQ. Despite this, we have avoided its calculation during the initialize_VOR method to allow a user to supply an externally-computed DADJ, if one is available. To compute a stand-alone Delaunay graph from a given prototype matrix, see ?Delaunay_ADJ.

3.3 Voronoi Polytope Definitions

With DADJ set, we can now retrieve the half-plane representation of any first or second-order Voronoi polytope. DADJ is used to ensure only non-redundant constraints are included in the system Ax <= b defining each polytope. To guarantee the polytope is actually a polytope instead of a polyhedron (i.e., that it is closed), the global bounds stored in 1b and ub are appended to the system. For example, to obtain the polytope definition for the first-order Voronoi cell generated by the prototype 1 (in the 1st row of W):

```
vor$get_vor1_polytope(1)
## $A
##
               [,1]
                          [,2]
##
  [1,] 0.23451935 -0.0529177
  [2,] 0.05103292 0.4011529
  [3,] 0.34175500 0.3073637
        1.00000000
                    0.0000000
  [5,] 0.00000000
                    1.0000000
  [6,] -1.00000000
                    0.0000000
## [7,] 0.00000000 -1.0000000
##
```

```
## $b

## [1,] -0.3798787

## [2,] -0.4326060

## [3,] -0.8864174

## [4,] 2.1394531

## [5,] 2.5118537

## [6,] 2.1716202

## [7,] 2.2333428
```

and to obtain the definition of the second-order polytope defined by prototypes 1 and 2:

```
vor$get_vor2_polytope(1,2)
## $A
##
                [,1]
                            [,2]
##
    [1,] 0.23451935 -0.0529177
##
    [2,] 0.05103292
                      0.4011529
    [3,] -0.18348643
                      0.4540706
##
         0.34175500
                      0.3073637
##
    [4,]
##
    [5,]
         0.10723566
                      0.3602814
##
    [6,]
         1.00000000
                      0.0000000
    [7,] 0.00000000
##
                      1.0000000
##
    [8,] -1.00000000 0.0000000
    [9,] 0.00000000 -1.0000000
##
##
## $b
##
                [,1]
    [1,] -0.37987869
##
    [2,] -0.43260597
##
##
    [3,] -0.05272729
##
    [4,] -0.88641741
##
    [5,] -0.50653872
##
    [6,] 2.13945314
##
    [7,]
         2.51185371
##
    [8,]
         2.17162020
    [9,]
         2.23334278
```

Again, note that second-order polytopes only exist for pairs of Delaunay adjacent prototypes. In this example, first-order cells generated by prototypes 1 and 99 are not adjacenct in \mathbb{R}^2 , so their second-order polytope cannot be formed:

```
vor$get_vor2_polytope(1,99)
## Error in vor$get_vor2_polytope(1, 99): DADJ(iidx, jidx)!=1: Voronoi cells iidx and jidx a
```

The definitions of first and second-order Voronoi polytopes are required to compute the Chebyshev centers, Dikin ellipsoids, and MVIEs; however, as they can

be large in high-dimensional settings they are never stored but instead generated as needed with the above methods. To return stand-alone polytope definitions given a prototype and Delaunay adjacency matrix, see ?vor1_polytope and ?vor2_polytope.

3.4 Chebyshev Centers

Computation of the Dikin and MVIE ellipsoids requires the specification of an interior point to each Voronoi polytope. For first-order cells this interior point is taken as the prototype which generated each cell, but interior points for second-order cells need to be supplied by the user. For ease, VorVQ provides the methods calc_vor1_centers and calc_vor2_centers to compute the Chebyshev center of each first and second-order Voronoi polytope:

```
vor$calc_vor1_centers()
## Calculating Chebyshev centers for vor1 cells ...
## 15%
              30%
                         45%
                                     60%
                                                75%
## 90%
              100%
## Computation time: 0.000150192 minutes
str(vor$vor1 centers)
    num [1:100, 1:2] -2.009 -1.763 -1.329 -1.123 -0.689 ...
vor$calc_vor2_centers()
## Calculating Chebyshev centers for vor2 cells ...
## 9%
             19%
                        29%
                                    39%
                                               49%
## 59%
              69%
                          79%
                                     89%
                                                99%
## 100%
## Computation time: 0.000287662 minutes
str(vor$vor2_centers)
   num [1:464, 1:2] -1.91 -2.08 -2.02 -1.37 -1.68 ...
```

After calculation, the centers are stored in the rows of the matrix fields $vor1_centers$ and $vor2_centers$ (each has d columns). Row i of $vor1_centers$ lists the center for the first-order Voronoi polytope defined by the i-th element of the $vor1_active$ list, and row i of $vor2_centers$ lists the center for the second-order Voronoi polytope defined by the i-th row of the $vor2_active$ list.

There are, of course, other ways of obtaining polytope interior points. VorVQ allows the user to specify their own centers via the set_vor1_centers and set_vor2_centers methods. Here, we copy the previously computed center matrices, clear their respective fields in the VOR object, and re-set them:

```
# Extract existing centers
vor1_centers = vor$vor1_centers
vor2_centers = vor$vor2_centers
```

```
# Clear the centers fields
vor$clear_vor1_centers()
str(vor$vor1_centers)
   num[0 , 0 ]
vor$clear_vor2_centers()
str(vor$vor2_centers)
## num[0 , 0 ]
# Re-assign
vor$set_vor1_centers(vor1_centers)
## Checking interior points for vor1 cells ...
                                     60%
## 15%
              30%
                         45%
                                                75%
## 90%
              100%
## Computation time: 1.37452e-05 minutes
str(vor$vor1_centers)
   num [1:100, 1:2] -2.009 -1.763 -1.329 -1.123 -0.689 ...
vor$set_vor2_centers(vor2_centers)
## Checking interior points for vor2 cells ...
## 9%
             19%
                        29%
                                   39%
                                               49%
## 59%
              69%
                         79%
                                    89%
                                                99%
## 100%
## Computation time: 1.77836e-05 minutes
str(vor$vor2_centers)
## num [1:464, 1:2] -1.91 -2.08 -2.02 -1.37 -1.68 ...
```

Obviously, the supplied centers must be interior to each respective polytope. The $set_vor\{1/2\}_centers$ methods will verify this by extracting the definition of each polytope and verifying that $Ac \leq b$, for each center c. If any supplied centers are not interior, an error will be returned.

See ?polytope_Chebyshev_center for computing Stand-alone Chebyshev centers, given an input half-plane representation of the polytope of the form $Ax \leq b$.

3.5 Dikin Ellipsoids

The Dikin ellipsoids for each Voronoi polytope listed in vor1_active and vor2_active are computed via the methods:

```
# Vor1 Dikin

vor$calc_vor1_Dikin()

## Calculating Dikin ellipsoids for vor1 cells ...

## 15% 30% 45% 60% 75%

## 90% 100%

## Computation time: 2.34927e-05 minutes
```

```
# Vor2 Dikin
vor$calc_vor2_Dikin()

## Calculating Dikin ellipsoids for vor2 cells ...

## 9% 19% 29% 39% 49%

## 59% 69% 79% 89% 99%

## 100%

## Computation time: 1.98246e-05 minutes
```

The Dikin ellipsoid, computed at a particular interior point c to each polytope, is defined by its ellipsoidal rotation matrix E (see Background). For each first-order Dikin ellipsoid, method ${\tt calc_vor1_Dikin}$ will take c to be the corresponding center found in the rows of ${\tt vor1_centers}$, if this field has been set by either ${\tt calc_vor1_centers}$ or ${\tt set_vor1_centers}$; otherwise, c is taken to be the corresponding prototype vector from ${\tt W}$. The second-order Dikins are computed at whatever values are found in ${\tt vor2_centers}$ (so this field must be set prior to calculation of second-order Dikin ellipsoids, see above).

The rotations defining the Dikin ellipsoids for each active Voronoi polytope are stored in the slices of the arrays $vor1_Dikin_E$ (which has dimensions $d \times d \times length(vor1_active)$) and $vor2_Dikin_E$ (which has dimensions $d \times d \times nrow(vor2_active)$).

```
# Vor1 Dikin
str(vor$vor1_Dikin_E)
## num [1:2, 1:2, 1:100] 0.01182 -0.00156 -0.00156 0.02077 0.0341 ...
# Vor2 Dikin
str(vor$vor2_Dikin_E)
## num [1:2, 1:2, 1:464] 0.01172 0.00206 0.00206 0.02736 0.00627 ...
```

See ?Dikin_ell for computing a stand-alone Dikin ellipsoid given a polytope definition $Ax \leq b$ and an interior point c.

3.6 MVIEs

The Maximum Volume Inscribed Ellipsoids for each Voronoi polytope listed in vor1_active and vor2_active are computed via the methods:

```
# Vor1 MVIE
vor$calc_vor1_MVIE()
## Calculating max vol incribed ellipsoids for vor1 cells ...
## 15% 30% 45% 60% 75%
## 90% 100%
## Computation time: 4.11339e-05 minutes

# Vor2 MVIE
vor$calc_vor2_MVIE()
```

```
## Calculating max vol incribed ellipsoids for vor2 cells ...
## 9% 19% 29% 39% 49%
## 59% 69% 79% 89% 99%
## 100%
## Computation time: 0.000110632 minutes
```

The MVIE for each polytope is defined by both an interior point c and an ellipsoidal rotation matrix E (see Background). To initialize computation of each first-order MVIE, method calc_vor1_MVIE will take the starting c to be the corresponding center found in the rows of vor1_centers, if this field has been set by either calc_vor1_centers or set_vor1_centers; otherwise, the starting c is taken to be the corresponding prototype vector from w. The starting interior points for second-order MVIEs are taken from whatever center values are found in vor2_centers (so this field must be set prior to calculation of second-order MVIEs, see above).

The MVIE centers resulting from the above computations for each active Voronoi polytope are stored in the rows of the matrix fields

```
# Vor1
str(vor$vor1_MVIE_c)
## num [1:100, 1:2] -2.051 -1.751 -1.351 -1.141 -0.834 ...

# Vor2
str(vor$vor2_MVIE_c)
## num [1:464, 1:2] -1.85 -2.06 -2.05 -1.41 -1.7 ...
```

while the rotation matrices are stored in the slices of the array fields

```
# Vor 1
str(vor$vor1_MVIE_E)
## num [1:2, 1:2, 1:100] 0.0144 0.0311 0.0311 0.2919 0.0407 ...
# Vor 2
str(vor$vor2_MVIE_E)
## num [1:2, 1:2, 1:464] 0.0183 0.049 0.049 0.3943 0.0127 ...
```

The status codes resulting from calling Zhang and Gao's interior-point solution to the MVIE problem for each Voronoi polytope are stored in

```
# Vor 1
str(vor$vor1_MVIE_status)
## int [1:100, 1] 0 0 0 0 0 0 0 0 0 0 ...

# Vor 2
str(vor$vor2_MVIE_status)
## int [1:464, 1] 0 0 0 0 0 0 0 0 0 0 ...
```

See ?max_vol_inscr_ell for interpretation of these status codes (0 indicates successful termination). Additional information about each MVIE is also computed and stored. For instance, the log-determinant of each rotation is stored in the elements of the vector fields

```
# Vor 1
str(vor$vor1_MVIE_logdet)
## num [1:100, 1] -5.73 -4.15 -4.16 -6.74 -3.43 ...

# Vor 2
str(vor$vor2_MVIE_logdet)
## num [1:464, 1] -5.34 -9.05 -5.98 -5.4 -11.62 ...
```

while the vector fields

```
# Vor 1
str(vor$vor1_MVIE_volratio)
## num [1:100, 1] 0.01103 0.02432 0.02413 0.00664 0.03478 ...
# Vor 2
str(vor$vor2_MVIE_volratio)
## num [1:464, 1] 0.015663 0.002445 0.011376 0.015148 0.000678 ...
```

store the ratios of the volumes of each active Voronoi MVIE to the sum of volumes of all active first or second-order MVIEs.

The volume of a general d-dimensional ellipsoid $\mathcal E$ with center c, rotation E, and radius r defined by

$$(x-c)^T E^{-1}(x-c) \le r^2$$

is given by

$$\operatorname{Vol}(\mathscr{E}) = V_d |E|^{1/2} r^d$$

where V_d is the volume of the d-dimensional unit hypersphere. By convention (see Background), the MVIE rotations E are constructed such that r=1 in the above. Since both V_d and |E| can grow quite large in higher-dimensional settings, only the log-determinant of each MVIE is stored (to avoid numerical overflow, in worst-case scenarios). Obviously, the actual first and second-order Voronoi MVIE volumes can be re-constructed easily from the *_MVIE_logdet fields and the above formula, if desired. The *_MVIE_volratio fields are provided as a convenience to the user to quickly determine the proportional volume of each MVIE (relative to all calculated), without having to explicitly exponentiate the potentially large values in *_MVIE_logdet to compute their relative proportions.

See ?max_vol_inscr_ell to compute a stand-alone MVIE given a polytope half-plane definition $Ax \leq b$ and a starting interior point c_0 . This help documentation also explains the meaning of the MVIE parameters (gettable / settable by methods get_params and set_params, see above).

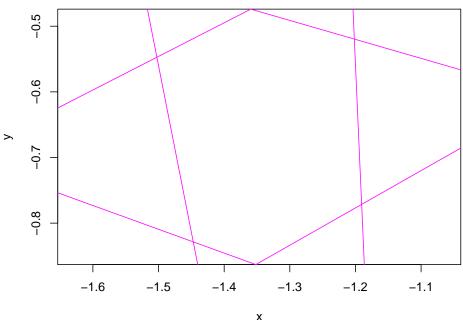
A final note: MVIE computation can be expensive, depending on the data dimension d and the complexity of the Voronoi tessellation (i.e., the size of each of the linear inequality systems $Ax \leq b$ definining each Voronoi polytope). It is highly recommended to allow these computations to be performed in parallel (which is the default selection specified during initialize_VOR).

3.7 Visualizations

VorVQ also provides visualizations of Voronoi regions if the data dimension d=2. Each of the following functions allow many different customizations outlined in their respective $?vis_**$ documentation.

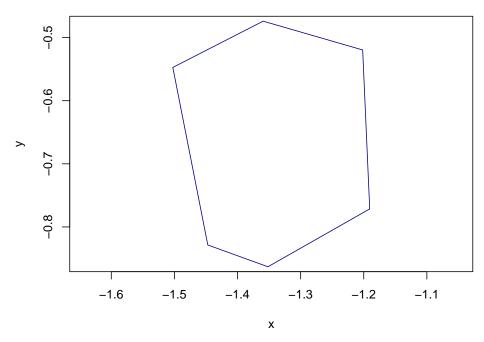
We start by visualizing the boundaries of the half-planes defining the polytope. While any polytope definition would suffice we select the definition of the polytope corresponding to prototype 13 as an example:

```
poly = vor$get_vor1_polytope(13)
vis_polytope_constraints(A = poly$A, b = poly$b)
```



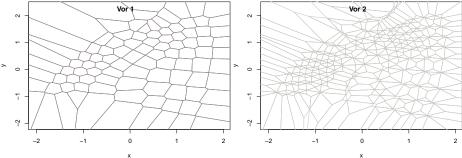
The polytope itself (with redundant sections of the half-plane boundaries removed) can be visualized via

```
vis_polytope(A = poly$A, b = poly$b, vertex.cex = 0)
```



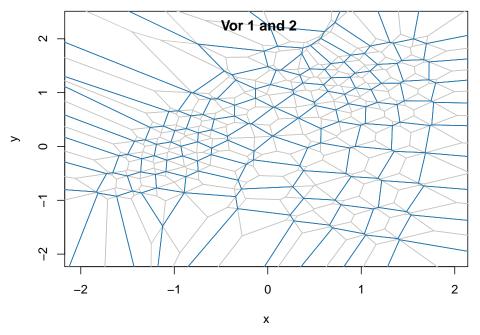
To visualize all (active) first or second-order Voronoi cells in a tessellation:

```
vis_vor1_polytopes(VOR = vor, add = F, vertex.cex = 0)
title("Vor 1")
vis_vor2_polytopes(VOR = vor, add = F, vertex.cex = 0)
title("Vor 2")
```



The vis_vor* visualizations are designed to allow layering, just like base R plots. Here, we plot both first (blue) and second-order (gray) Voronoi tessellations together:

```
vis_vor2_polytopes(VOR = vor, add = F, vertex.cex = 0)
vis_vor1_polytopes(VOR = vor, add = T, vertex.cex = 0, edge.col = tableau20("blue"))
title("Vor 1 and 2")
```



User-supplied text annotations can be placed at the centers of each Voronoi cell with the vis_vor_annotate functions. By default, the functions called without any text supplied will annotate each Voronoi cell with its id:

```
vis_vor1_polytopes(VOR = vor, add = F, vertex.cex = 0)
vis_vor1_annotate(VOR = vor, add = T, text.col = tableau20("blue"))
title("Vor 1")

# Vor 2
vis_vor2_polytopes(VOR = vor, add = F, vertex.cex = 0)
vis_vor2_annotate(VOR = vor, add = T, text.col = tableau20("green"))
title("Vor 2")

**Or**
**The structure of the structu
```

Approximating ellipsoids (if previously calculated and stored) can also be viewed in situ:

```
# Vor 1 Dikin
vis_vor1_polytopes(VOR = vor, add = F, vertex.cex = 0)
vis_vor1_Dikin(VOR = vor, add = T)
# Vor 2 Dikin
vis_vor2_polytopes(VOR = vor, add = F, vertex.cex = 0)
vis_vor2_Dikin(VOR = vor, add = T)
# Vor 1 MVIE
vis_vor1_polytopes(VOR = vor, add = F, vertex.cex = 0)
vis_vor1_MVIE(VOR = vor, add = T)
# Vor 2 MVIE
vis_vor2_polytopes(VOR = vor, add = F, vertex.cex = 0)
vis_vor2_MVIE(VOR = vor, add = T)
```

3.8 Saving & Loading

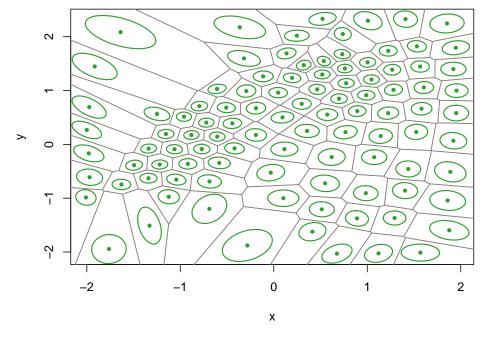
VOR objects can be saved to disk for future access via the save method. This method requires a user-specified file name and must have extension .vor.

```
vor$save("cmix9.vor")
```

While the file has extension .vor it is saved in R's binary .rds format as a list. R's readRDS command *can* be used to return this list to the user's environment

in a new R session, but the pointers to the C++ object are destroyed on exit of the first session. The load method will properly restore a previously saved VOR object to memory. Once loaded all object methods are available. For example, we can load the previously saved VOR object, re-compute its centers, and re-plot its first-order Dikin ellipsoids

```
vor2 = VorVQ::VOR$new()
vor2$load("cmix9.vor")
# Re-compute centers
vor2$calc_vor1_centers()
## Calculating Chebyshev centers for vor1 cells ...
              30%
## 15%
                         45%
                                     60%
                                                75%
## 90%
              100%
## Computation time: 6.08267e-05 minutes
# Check that it is the same
vis_vor1_polytopes(VOR = vor2, add = F, vertex.cex = 0)
vis_vor1_Dikin(VOR = vor2, add = T)
```



3.9 Performing All Voronoi Calculations

If the results of all Voronoi calculations exposed by VorVQ are desired the wrapper method calc_all can be called just after VOR object initialition. Internally, this wrapper calls the following previously discussed methods, in the given order:

```
calc_DADJ()
set_vor2_centers()
calc_vor1_Dikin()
calc_vor2_Dikin()
calc_vor1_MVIE()
calc_vor2_MVIE().
```

Note that calc_vor1_centers is excluded from the calc_all wrapper as it is not mandatory (the prototype vectors themselves are used as interior points for first-order Voronoi cells by default). As an example, we can re-initialize the VOR object and repeat all calculations above with one call to calc_all:

```
vor = VOR$new()
vor$initialize_VOR(CMIX9$W, CMIX9$CADJ)
## Initializing Voronoi object
## ++ storing W and dimensions ... done
## ++ setting global bounds = \pm of W range ... done
      call $set bounds to change
##
## ++ storing CADJ ... done
## ++ setting active vor1 indices (default to non-empty vor1 cells) ... done
##
      call $set_vor1_active to change
## ++ setting active vor2 indices (default to non-empty vor2 cells) ... done
##
      call $set_vor2_active to change
## Calculating Gabriel graph edges ...
## 15%
              30%
                         45%
                                     60%
                                                75%
## 90%
              100%
## Computation time: 1.38954e-05 minutes
vor$calc_all()
## Testing Delaunay graph edges ...
## Round 1 (456 edges)
                              25%
                                          50%
                                                     75%
                                                                 100%
## Round 2 (45 edges)
                              25%
                                         50%
                                                                100%
                                                    75%
## Tested 501 of 4950 possible edges
## Computation time: 0.000603641 minutes
##
## Calculating Chebyshev centers for vor2 cells ...
             19%
## 9%
                        29%
                                    39%
                                               49%
## 59%
              69%
                          79%
                                     89%
                                                99%
## 100%
## Computation time: 0.000188645 minutes
##
## Calculating Dikin ellipsoids for vor1 cells ...
## 15%
              30%
                         45%
                                     60%
                                                75%
## 90%
              100%
## Computation time: 1.39063e-05 minutes
## Calculating Dikin ellipsoids for vor2 cells ...
```

```
## 9%
             19%
                         29%
                                     39%
                                                49%
## 59%
              69%
                          79%
                                      89%
                                                 99%
## 100%
## Computation time: 1.45799e-05 minutes
##
## Calculating max vol incribed ellipsoids for vor1 cells ...
## 15%
              30%
                          45%
                                      60%
                                                 75%
## 90%
              100%
##
  Computation time: 2.7756e-05 minutes
##
## Calculating max vol incribed ellipsoids for vor2 cells ...
## 9%
             19%
                         29%
                                    39%
                                                49%
## 59%
              69%
                          79%
                                      89%
                                                 99%
## 100%
## Computation time: 8.26927e-05 minutes
```

3.10 Higher-Dimensional Settings

When the data dimension d is high several of the Voronoi calculations exposed by VorVQ can be fairly time consuming, particularly the calculation of the Delaunay adjacency and MVIEs (mostly the second-order MVIEs, as there can be many thousand active second-order Voronoi cells in a higher-dimensional tessellation). To mitigate this it is recommended to allow parallel computation (which is allowed by default, see <code>?set_params</code>). To give the user an idea of how long these calculations may take, below we <code>calc_all</code> quantities for the 100-dimensional synthethic SHGR tessellation provided by VorVQ as a higher-dimensional example case (see <code>?SHGR</code> for more details):

```
# Initialize with SHGR
vor = VOR$new()
vor$initialize_VOR(SHGR$W, SHGR$CADJ)
## Initializing Voronoi object
  ++ storing W and dimensions ... done
##
  ++ setting global bounds = \pm- 5% of W range ... done
      call $set_bounds to change
##
  ++ storing CADJ ... done
  ++ setting active vor1 indices (default to non-empty vor1 cells) ... done
##
      call $set_vor1_active to change
  ++ setting active vor2 indices (default to non-empty vor2 cells) ... done
##
      call $set_vor2_active to change
## Calculating Gabriel graph edges
              20%
                          30%
                                                50%
## 10%
                                     40%
              70%
                                                100%
## 60%
                          80%
                                     90%
## Computation time: 0.000852747 minutes
```

```
# View number of active 2nd-order Voronoi cells
nrow(vor$vor2_active)
## [1] 2875
# Calc all
vor$calc_all()
## Testing Delaunay graph edges ...
## Round 1 (9807 edges)
                                25%
                                           50%
                                                       75%
                                                                   100%
## Round 2 (13371 edges)
                                 25%
                                            50%
                                                        75%
                                                                   100%
## Round 3 (23379 edges)
                                 25%
                                             50%
                                                        75%
                                                                   100%
## Round 4 (7792 edges)
                                25%
                                           50%
                                                       75%
                                                                   100%
## Round 5 (6943 edges)
                                25%
                                           50%
                                                       75%
                                                                  100%
## Round 6 (174 edges)
                                                      75%
                                                                  100%
                               25%
                                           50%
## Tested 61466 of 79800 possible edges
## Computation time: 1.58323 minutes
## Calculating Chebyshev centers for vor2 cells ...
             19%
                         29%
## 9%
                                    39%
                                                49%
## 59%
              69%
                                     89%
                                                 99%
                          79%
## 100%
## Computation time: 0.243345 minutes
##
## Calculating Dikin ellipsoids for vor1 cells ...
                                                48%
## 9%
             19%
                         29%
                                    38%
              68%
## 58%
                          77%
                                     87%
                                                 97%
## 100%
## Computation time: 0.00194854 minutes
##
## Calculating Dikin ellipsoids for vor2 cells ...
## 9%
             19%
                         29%
                                    39%
                                                49%
## 59%
              69%
                          79%
                                     89%
                                                 99%
## 100%
## Computation time: 0.0290169 minutes
##
## Calculating max vol incribed ellipsoids for vor1 cells ...
## 9%
             19%
                         29%
                                    38%
                                                48%
## 58%
              68%
                          77%
                                     87%
                                                 97%
## 100%
## Computation time: 1.30166 minutes
## Calculating max vol incribed ellipsoids for vor2 cells ...
## 4%
             9%
                        14%
                                   19%
                                               24%
## 29%
                          39%
                                     44%
                                                 49%
              34%
## 54%
              59%
                          64%
                                     69%
                                                 74%
              84%
                          89%
                                                 99%
## 79%
                                     94%
```

```
## 100%
## Computation time: 25.9597 minutes
```

The MVIE calculation for the 2,875 second-order Voronoi cells required approximately 30 minutes of parallel computation with 15 available threads on a 2.4 GHz 8-Core Intel Core i9 processor running Mac OS Catalina (10.15.6). If these processing times are prohibitive, the user may instead compute approximate MVIEs, which are optimized over only the ellipsoidal rotations (i.e., not jointly with the centers, which remain fixed at their starting values). This modification, controlled by setting the parameter MVIE_fix_c0 = T, see ?set_params) can result in significant computational savings (at the expense of volumetric precision), as shown below:

```
vor$set_params(list(MVIE_fix_c0 = T))
vor$calc_vor2_MVIE()
## Calculating max vol incribed ellipsoids for vor2 cells ...
## 4%
             9%
                        14%
                                    19%
                                               24%
## 29%
              34%
                          39%
                                                 49%
                                      44%
## 54%
              59%
                          64%
                                      69%
                                                 74%
## 79%
              84%
                          89%
                                                 99%
                                      94%
## 100%
## Computation time: 14.1547 minutes
```

References

- [1] S. Boyd, S. P. Boyd, and L. Vandenberghe, *Convex optimization*. Cambridge university press, 2004.
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