

# Supplementary Material

## Supplementary material for the paper titled ‘A Novel Method for Finding Differential-Linear Distinguishers: Application to Midori64, CRAFT, and Skinny64’

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### A Formation of Linear Constraints

The differential pattern and difference of a word can be described by  $(a_0, a_1, b_0, \dots, b_3)$  and constrained as

$$\left\{ \begin{array}{l} -a_0 - b_1 \geq -1 \\ a_1 - b_3 \geq 0 \\ a_1 - b_2 \geq 0 \\ a_1 - b_1 \geq 0 \\ a_1 - b_0 \geq 0 \\ -a_1 + b_0 + b_1 + b_2 + b_3 \geq 0 \\ a_0, a_1, b_0, b_1, b_2, b_3 \text{ are binary variables.} \end{array} \right. \quad (1)$$

For XOR operation, the propagation can be described by  $(a_0^{X_1}, a_0^{X_2}, a_0^Y, b_i^{X_1}, b_i^{X_2}, b_i^Y)$  for each  $i \in \{0, 1, 2, 3\}$  and linearly constrained as

$$\left\{ \begin{array}{l} -a_0^{X_1} - a_0^{X_2} - b_i^{X_1} - b_i^{X_2} - b_i^Y \geq -2 \\ a_0^{X_1} + a_0^{X_2} - a_0^Y \geq 0 \\ -a_0^{X_1} + a_0^Y - b_i^{X_1} + b_i^{X_2} + b_i^Y \geq 0 \\ b_i^{X_1} + b_i^{X_2} - b_i^Y \geq 0 \\ -a_0^{X_2} + a_0^Y + b_i^{X_1} - b_i^{X_2} + b_i^Y \geq 0 \\ a_0^{X_1}, a_0^{X_2}, a_0^Y, b_i^{X_1}, b_i^{X_2}, b_i^Y \text{ are binary variables.} \end{array} \right. \quad (2)$$

For a single S-box of Midori64, the propagation can be described by  $(a_0^X, a_1^X, a_0^Y, a_1^Y)$  and  $(b_0^X, \dots, b_3^X, b_0^Y, \dots, b_3^Y, p_0, p_1)$ , and constrained simultaneously as the fol-

lowing two systems of linear inequalities.

$$\begin{cases}
 -a_0^Y - a_1^Y \geq -1 \\
 -a_0^X + a_0^Y \geq 0 \\
 -a_1^X + a_0^Y + a_1^Y \geq 0 \\
 a_0^X + a_1^X - a_0^Y \geq 0 \\
 a_1^X - a_1^Y \geq 0 \\
 -a_0^X - a_1^X \geq -1 \\
 a_0^X, a_1^X, b_0^Y, b_1^Y \text{ are binary variables.}
 \end{cases} \quad (3)$$

$$\begin{cases}
 b_2^X - b_1^Y - b_2^Y - b_3^Y - p_0 + 3p_1 \geq 0 \\
 -2b_0^X - b_1^X - b_3^X + 2b_1^Y + 6b_2^Y + 2b_3^Y + 5p_0 - 6p_1 \geq -4 \\
 b_1^X - 2b_2^X + b_3^X + b_1^Y - 2b_2^Y + b_3^Y - 3p_0 + 2p_1 \geq -2 \\
 4b_0^X + 3b_1^X + 6b_2^X - b_3^X - 8b_0^Y - 2b_1^Y - 2b_2^Y - 6b_3^Y + 5p_0 + 7p_1 \geq -1 \\
 +4b_0^X + 5b_1^X + 3b_2^X + 5b_3^X + 3b_0^Y - 4b_1^Y - 3b_2^Y - 4b_3^Y + 2p_0 + p_1 \geq 0 \\
 b_1^X - 2b_2^X + b_3^X + 3b_0^Y + 3b_1^Y - b_2^Y + 3b_3^Y - 2p_0 - 2p_1 \geq -2 \\
 8b_0^X - b_1^X - 2b_2^X + 2b_3^X - 5b_0^Y - 4b_1^Y - 7b_2^Y - b_3^Y + 9p_0 + 8p_1 \geq -3 \\
 -7b_0^X + 4b_1^X + b_2^X - 2b_3^X + b_0^Y - 8b_1^Y - b_2^Y + 4b_3^Y + 6p_0 + p_1 \geq -9 \\
 -3b_0^X - 4b_1^X + b_2^X - 2b_3^X + b_0^Y + 6b_1^Y + 4b_2^Y - 5b_3^Y + 7p_0 - 5p_1 \geq -9 \\
 -b_1^X - b_2^X - b_3^X + b_0^Y - b_1^Y + b_2^Y - b_3^Y - 2p_0 + 2p_1 \geq -3 \\
 -b_0^X - 3b_1^X - 2b_2^X + 3b_3^X - 4b_0^Y + 3b_1^Y - b_2^Y - 2b_3^Y + 2p_0 + 2p_1 \geq -6 \\
 b_1^X - 2b_2^X - b_3^X + 2b_0^Y - b_1^Y - 2b_2^Y - 3b_3^Y + 3p_0 + 3p_1 \geq -3 \\
 2b_0^X + b_1^X + 3b_2^X + b_3^X - 3b_0^Y - 2b_1^Y + b_2^Y - 2b_3^Y + p_0 + 2p_1 \geq 0 \\
 b_0^X - b_1^X + b_2^X - b_3^X + b_0^Y + 2b_1^Y - b_2^Y + 2b_3^Y - 2p_0 \geq -2 \\
 b_1^X - b_2^X + b_3^X - p_0 \geq -1 \\
 -b_0^X - b_1^X - 2b_2^X - 3b_3^X - 2b_0^Y - 2b_1^Y + 3b_2^Y + 4b_3^Y + 2p_0 - p_1 \geq -7 \\
 -4b_0^X - 5b_1^X - 2b_2^X + b_3^X + 6b_0^Y - 2b_1^Y + 5b_2^Y + 4b_3^Y + 3p_0 - 7p_1 \geq -11 \\
 3b_0^X - 2b_1^X + 2b_2^X + b_3^X - b_0^Y - 2b_1^Y - 2b_2^Y + b_3^Y + 2p_1 \geq -2 \\
 b_0^X + 2b_1^X + 3b_2^X + 2b_3^X - 2b_0^Y + b_1^Y + b_3^Y + 2p_0 - 2p_1 \geq 0 \\
 -b_1^X - 2b_2^X - 2b_3^X + b_0^Y + 3b_1^Y + 2b_2^Y + p_0 - 3p_1 \geq -5 \\
 -b_1^X - b_3^X - b_1^Y - b_3^Y - p_0 \geq -4 \\
 -b_0^X + b_1^X - b_3^X - b_0^Y - b_2^Y + b_3^Y \geq -3 \\
 -2b_0^X - 2b_1^X - b_2^X + 3b_3^X - b_0^Y + 3b_1^Y - 2b_2^Y - b_3^Y + p_1 \geq -5 \\
 -b_0^X - b_3^X + b_0^Y + b_1^Y + b_2^Y - p_1 \geq -2 \\
 b_0^X + b_2^X - b_3^X + b_1^Y - b_2^Y - p_0 \geq -2 \\
 b_i^X, b_i^Y, p_0, p_1 \text{ are binary variables.}
 \end{cases} \quad (4)$$

## B The descriptions of Midori64, CRAFT and Skinny64

### B.1 Midori64

The state of Midori64 is a  $4 \times 4$  matrix, arranged as follows:

$$\begin{pmatrix} s_0 & s_4 & s_8 & s_{12} \\ s_1 & s_5 & s_9 & s_{13} \\ s_2 & s_6 & s_{10} & s_{14} \\ s_3 & s_7 & s_{11} & s_{15} \end{pmatrix},$$

where  $s_i$  represents the cell with a bit size of 4.

For the round function  $F = \text{AK} \circ \text{MC} \circ \text{SC} \circ \text{SB}$ . The details about four components of  $F$  are clarified as below:

- **SubCell (SB)**: Each cell in the state goes through the same S-box, the truth table of this S-box is listed as Table 1.

**Table 1.** The truth table of S-box used in Midori64 and CRAFT.

$x$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	c	a	d	3	e	b	f	7	8	9	1	5	0	2	4	6

- **ShuffleCell (SC)**: The cells in the state are permuted as follows:

$$\mathcal{P} = [0, 10, 5, 15, 14, 4, 11, 1, 9, 3, 12, 6, 7, 13, 2, 8].$$

- **MixColumn (MC)**: The state matrix is multiplied by a binary almost MDS matrix  $M$ , as shown below:

$$M = \begin{pmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}.$$

- **AddRoundKey (AK)**: The  $r$ th round-key is XOR-ed with the state.

At the last round, the state only goes through the S-box and is XOR-ed with the last round-key.

### B.2 CRAFT

The 64-bit internal state of CRAFT is viewed as a  $4 \times 4$  square array of 16 4-bit cells as follows:

$$\begin{pmatrix} s_0 & s_1 & s_2 & s_3 \\ s_4 & s_5 & s_6 & s_7 \\ s_8 & s_9 & s_{10} & s_{11} \\ s_{12} & s_{13} & s_{14} & s_{15} \end{pmatrix},$$

where  $s_i$  denote the  $i$ th cell.

The round function of CRAFT is  $F = \text{SB} \circ \text{PN} \circ \text{ATK}_i \circ \text{ARC}_i \circ \text{MC}$ , where  $i \in \{0, 1, \dots, 30\}$ . The five transformations involved are explained below.

- **MixColumn (MC)**: Each column of the state is multiplied by the following binary matrix:

$$M = \begin{pmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}.$$

- **AddConstants<sub>i</sub> (ARC<sub>i</sub>)**: The state cells  $s_4$  and  $s_5$  are XOR-ed with two 4-bit constants  $a_i$  and  $b_i$  in the  $i$ th round, respectively.
- **AddTweakey<sub>i</sub> (ATK<sub>i</sub>)**: The state is XOR-ed with the round tweakkey.
- **PermuteNibbles (PN)**: An involutory permutation  $\mathcal{P}$  is applied to the cell positions in the state.

$$\mathcal{P} = [15, 12, 13, 14, 10, 9, 8, 11, 6, 5, 4, 7, 1, 2, 3, 0].$$

- **SubBox (SB)**: Each cell of the state goes through the same S-box. The truth table of the S-box is given in Table 1.

Additional, the last round excludes the **SB** operation.

### B.3 Skinny64

Similar to CRAFT, the internal state of **Skinny64** is splited to 16 cells with a bit size of 4, and arranged as a square array as

$$\begin{pmatrix} s_0 & s_1 & s_2 & s_3 \\ s_4 & s_5 & s_6 & s_7 \\ s_8 & s_9 & s_{10} & s_{11} \\ s_{12} & s_{13} & s_{14} & s_{15} \end{pmatrix},$$

where  $s_i$  denote the  $i$ th cell.

The components of the round function of **Skinny64** are explained briefly as follows.

- **SubCell (SC)**: Each cell in the state goes through the same S-box. The truth table of **Skinny64**'s S-box is listed as the following table.

$x$	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
$S(x)$	c	6	9	0	1	a	2	b	3	8	5	d	4	e	7	f

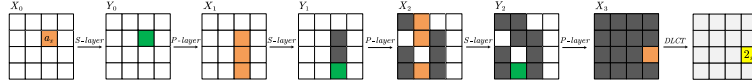
- **AddConstants (AC)**: The state  $s_0$  and  $s_4$  are XOR-ed with two constants  $c_0$  and  $c_1$ , respectively.
- **AddRoundTweakey (ART)**: The first two rows of the state array are XOR-ed with the round tweakkey.
- **ShiftRows (SR)**: The  $i$ th row of the state array is right rotated by  $i$  positions for  $i \in \{0, 1, 2, 3\}$ .
- **MixColumn (MC)**: Each column of the state array is multiplied by the following binary matrix  $M$ :

$$M = \begin{pmatrix} 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}.$$

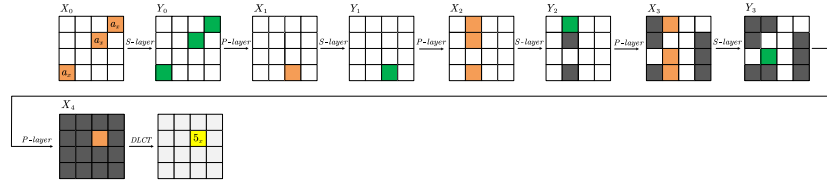
## C The DL Distinguishers of Midori64, CRAFT, and Skinny64

In this section, we will illustrate the DL distinguishers of Midori64, CRAFT and Skinny64, where the notation like ' $a_x$ ' is the hexadecimal representation of the 4-bit vector (1, 0, 1, 0). The white (resp. black) word indicates its differential pattern is Z (resp. U\*). The orange (resp. green) word means the pattern is N before (resp. after) going through the S-box. Besides, the yellow (resp. gray) word in the last matrix represents the non-zero (resp. zero) output mask.

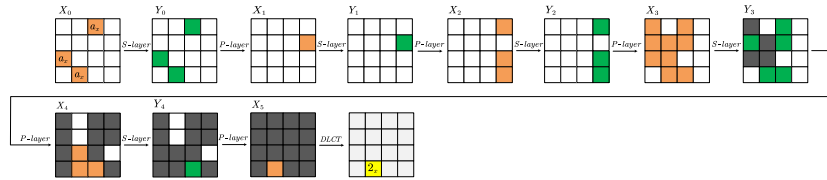
### C.1 The DL distinguishers of Midori64



**Fig. 1.** The 4-round DL distinguisher of Midori64.

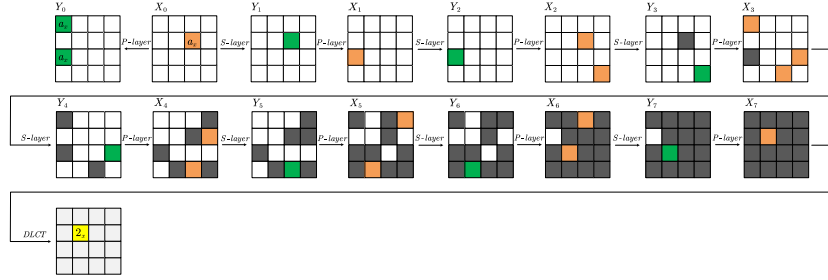


**Fig. 2.** The 5-round DL distinguisher of Midori64.

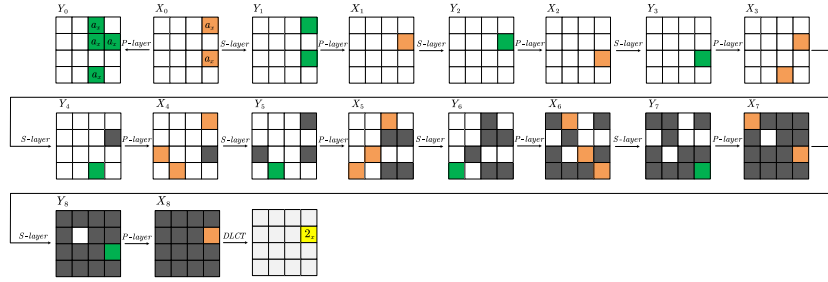


**Fig. 3.** The 6-round DL distinguisher of Midori64.

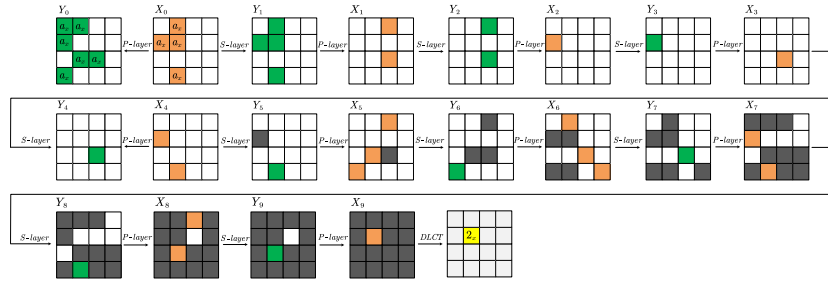
## C.2 The DL distinguishers of CRAFT



**Fig. 4.** The 8-round DL distinguisher of CRAFT.



**Fig. 5.** The 9-round DL distinguisher of CRAFT.



**Fig. 6.** The 10-round DL distinguisher of CRAFT.

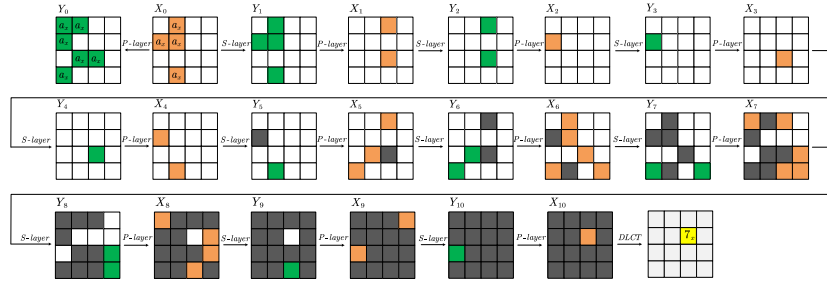


Fig. 7. The 11-round DL distinguisher of CRAFT.

### C.3 The DL distinguishers of Skinny64

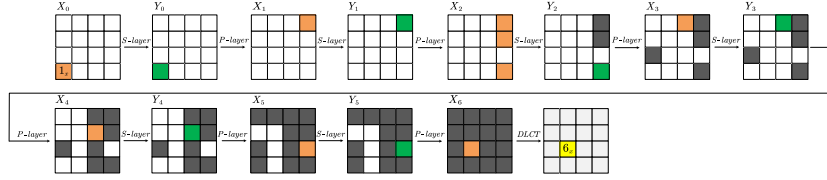


Fig. 8. The 7-round DL distinguisher of Skinny64.

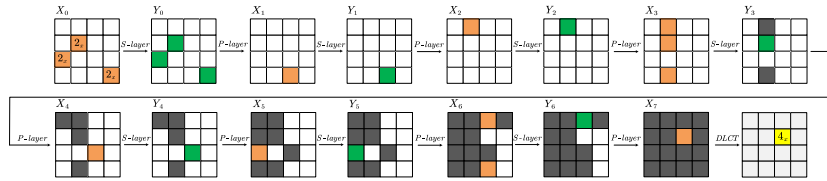


Fig. 9. The 8-round DL distinguisher of Skinny64.

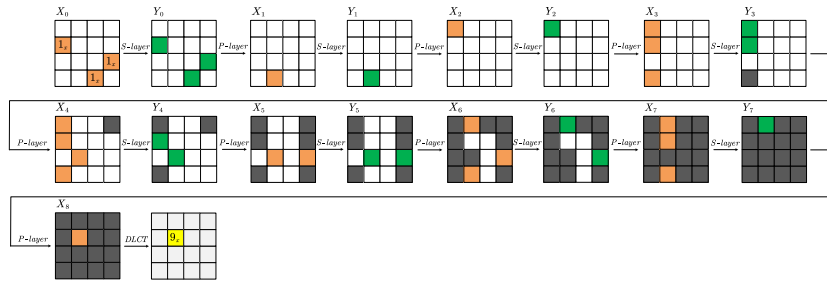
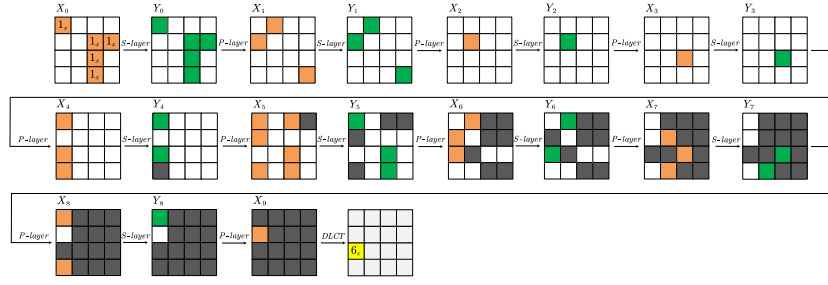


Fig. 10. The 9-round DL distinguisher of Skinny64.



**Fig. 11.** The 10-round DL distinguisher of Skinny64.