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DSAI Theory Endsem

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Finding Imp. vertex {

(a) For every vertex $v \in V$ {

Do BFS (v)

If (the BFS tree of v
contains all other
vertices) {

graph has Important
vertex v

return v ;

}

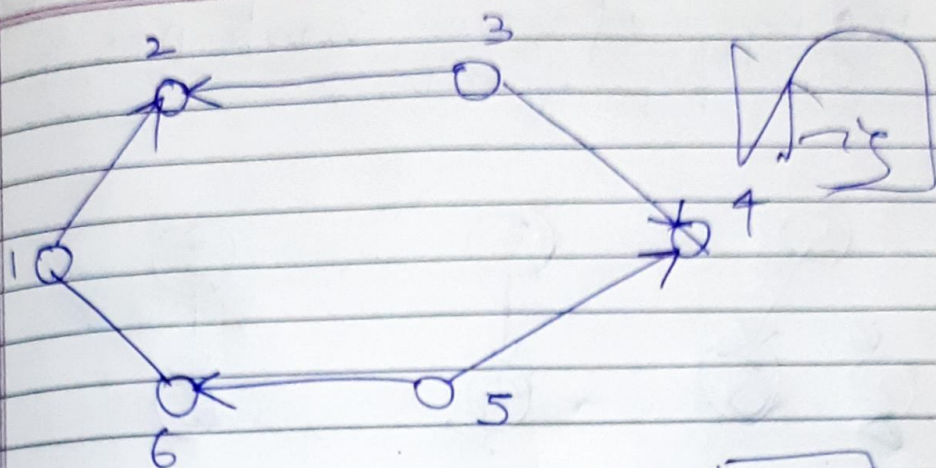
~~else~~

}

~~graph has no~~

graph has no Important vertex.

}



Algorithm for finding if DAG has an Imp. vertex or not.

Step 1. Make the DFS tree of the DAG

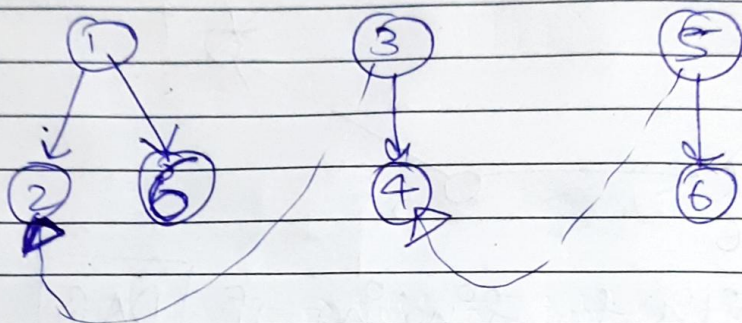
Step 2. If there are no components in the DFS tree i.e. it is a simply connected acyclic graph (tree) then the DAG has an Imp. vertex.

Step 3: However if DAG has more than one components then for an important vertex to exist

(i) there should be a back edge b/w components

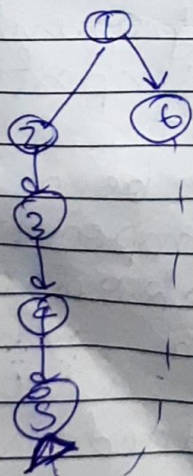
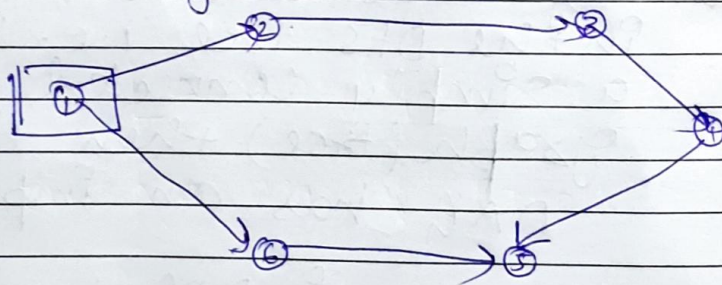
(ii) and the back edge b/w ~~from~~ within a component should end at a vertex that ~~has a~~ indeed has a back edge to prev component.

For eg. if you make DFS of
the Ag given.



This graph does not have
an imp. vertex.

However if.



this has an
imp. vertex.

Note! Running time = $O(n+m)$

of DFS and even if we
check all non tree
edges

~~hasImpVertex (root)~~

hasImpVertex ()

Do DFS (u)

if (DFS do not have any components)

then exist an imp vertex

else

(if there are no back edges)
there do not exist
imp vertex.

}

if there are no back edges
within components &
there does not
exist imp vertex

}

elif (back edges b/w
every component and back
with component end at
vertex which has
back edge to prev component)
there exist imp vertex

}

}

Important

→ pass head by reference (not by value)

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value

(4)

1. Put (head, int n) {
 create a new node v with value n;
 temp = head;
 tempPrev = head;
 *

 while (temp → value < n) {
 temp = temp → next;

 }

 while (tempPrev → next → value < n) {
 tempPrev = tempPrev → next;

 }

 tempPrev → next = v;
 v → next = temp;

 → * if (head == NULL) {
 head = v;
 v → next = NULL

 }

 if (temp → next == NULL) {
 temp → next = v;
 v → next = NULL;

 }

}

0	1	2	3	4	5	6	7	8	9	10	11	12
1	-2	0	2	3	4	5	6	7	9	12	13	14

↑
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~~if false~~

identical with (arr, n) {

int s = 0;

int e = n - 1;

int m;

~~while~~ while (true) {

if (s > e)

return false // or return -1

if (arr[m] > m)

e = m - 1

else if (arr[m] < m)

s = m + 1

else if (arr[m] == m)

return true; // return m;

}

}

$O(\log n) \rightarrow$ we can make a Binary Tree for this


```

4  mandepth (root) {
    if (root == NULL)
    if (root == NULL) {
        return 0;
    }
    if (root->left == NULL OR
        root->right == NULL) {
        return 0;
    }
    ansL = mandepth(root->left);
    ansR = mandepth(root->right);
    return ansL;
    return max{ansL, ansR} + 1;
}

```

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The algorithm goes over every vertex
its running time $= O(n)$