ORACLES AND CHOICE SEQUENCES FOR TYPE-THEORETIC PRAGMATICS

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October 8, 2015

joint work with Darryl McAdams

INTRODUCTION

$[\![$*A woman walked in.*$]\!]$

$\triangledown$

$$[\![\textit{A woman walked in.}]\!]$$

$$\bigtriangledown$$

$$(\Sigma p \in Woman)$$

$$[\![\textit{A woman walked in.}]\!]$$

$$\triangledown$$

$$(\Sigma p \in Woman)\, WalkedIn(p)$$

$$[\![\textit{She sat down}]\!]$$
$$\triangledown$$

$$[\![\textit{She sat down}]\!]$$
$$\triangledown$$
$$\mathit{SatDown}(???)$$

$$[\![\textit{A woman walked in. She sat down}]\!]$$
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$$\triangledown$$

$$(\Sigma x \in (\Sigma p \in Woman)\, WalkedIn(p))$$

$$[\![\textit{A woman walked in. She sat down}]\!]$$
$$\bigtriangledown$$
$$(\Sigma x \in (\Sigma p \in Woman)\, WalkedIn(p))\, SatDown(???)$$

$$[\![\textit{A woman walked in. She sat down}]\!]$$

$$\triangledown$$

$$(\Sigma x \in (\Sigma p \in Woman)\, WalkedIn(p))\, SatDown(\pi_1(x))$$

$$[\![\textit{Every farmer who owns a donkey beats it.}]\!]$$

$$\triangledown$$

$$(\Pi p \in (\Sigma x \in \mathit{Farmer})\,(\Sigma y \in \mathit{Donkey})\,\mathit{Owns}(x; y))$$

$$[\![\textit{Every farmer who owns a donkey beats it.}]\!]$$

$$\triangledown$$

$$(\Pi p \in (\Sigma x \in \mathit{Farmer})\,(\Sigma y \in \mathit{Donkey})\,\mathit{Owns}(x; y))\,\mathit{Beats}(???; ???)$$

$$[\![\textit{Every farmer who owns a donkey beats it.}]\!]$$

$$\triangledown$$

$$(\Pi p \in (\Sigma x \in \mathit{Farmer})\,(\Sigma y \in \mathit{Donkey})\,\mathit{Owns}(x; y))\,\mathit{Beats}(\pi_1(p); \pi_1(\pi_2(p)))$$

TWO THINGS TO DEAL WITH

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- terms for presuppositions

- terms for presuppositions
- resolution of presuppositions

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TWO THINGS TO DEAL WITH

- terms for presuppositions (this talk)
- resolution of presuppositions

THE **require** ORACLE: STATICS

require — FORMAL RULES

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$$\mathbf{require} : (0;1) \qquad \text{(operator)}$$

$$\mathbf{require}\ x : A\ \mathbf{in}\ N \stackrel{\text{def}}{=} \mathbf{require}(A; x.N) \qquad \text{(notation)}$$

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$$\frac{\Gamma \vdash M \in A \quad \Gamma, x : A \vdash N \in B}{\Gamma \vdash \textbf{require}\ x : A\ \textbf{in}\ N \in B} \qquad (\textit{require})$$

$$[\![\textit{A woman walked in. She sat down}]\!]$$

$$\triangledown$$

$$(\Sigma x \in (\Sigma p \in \mathit{Woman})\, \mathit{WalkedIn}(p))\, \mathit{SatDown}(???)$$

require — EXAMPLES

$$[\![\textit{A woman walked in. She sat down}]\!]$$

$$\triangledown$$

$$(\Sigma x \in (\Sigma p \in Woman)\, WalkedIn(p))\ \mathbf{require}\ y : Woman\ \mathbf{in}\ SatDown(y)$$

The meaning of a sentence is a logical proposition.

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What we want:

$$(\Sigma x \in (\Sigma p \in Woman)\, WalkedIn(p))\, \mathbf{require}\; y : Woman\; \mathbf{in}\; SatDown(y)$$

$$\sim$$

$$(\Sigma x \in (\Sigma p \in Woman)\, WalkedIn(p))\, \pi_1(x)$$

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What we want:

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$$\sim$$

$$(\Sigma x \in (\Sigma p \in \mathit{Woman})\, \mathit{WalkedIn}(p))\, \pi_1(x)$$

$$\boxed{\text{where } M \sim N \stackrel{\text{def}}{=} (M \preceq N) \wedge (N \preceq M)}$$

EVERY GRAMMATICAL SENTENCE HAS A MEANING

…but only some of them denote propositions (types)!

require—NAÏVE DYNAMICS

$$\frac{M \in A \quad [M/x]\,N \Downarrow N'}{\mathbf{require}\ x : A\ \mathbf{in}\ N \Downarrow N'} \tag{??}$$

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Can the above be made precise? There are two problems:

1. circularity
2. non-determinism

(HOLD THAT THOUGHT)

A POSITIVE EXAMPLE

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$$[\![\textit{The President ran a marathon}]\!]$$

$$\triangledown$$

$$\llbracket \textit{ The President ran a marathon } \rrbracket$$

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$$\textbf{require } x : \mathit{President} \textbf{ in } (\Sigma y \in \mathit{Marathon})\, \mathit{Ran}(x; y)$$

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$$\Downarrow$$
$$(\Sigma y \in \mathit{Marathon})\,\mathit{Ran}(\mathit{Obama}; y)$$

A NEGATIVE EXAMPLE

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$$[\![\textit{The unicorn ran a marathon}]\!]$$

$$\triangledown$$

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$$\triangledown$$

$$\textbf{require } x : \mathit{Unicorn} \textbf{ in } (\Sigma y \in \mathit{Marathon})\, \mathit{Ran}(x; y)$$

$$\llbracket\ \textit{The unicorn ran a marathon}\ \rrbracket$$
$$\triangledown$$
$$\mathbf{require}\ x : \mathit{Unicorn}\ \mathbf{in}\ (\Sigma y \in \mathit{Marathon})\, \mathit{Ran}(x; y)$$

(not a proposition)

Is **require** COMPUTATIONALLY EFFECTIVE?

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$$\boxed{w \Vdash \mathscr{J}_\alpha}$$

THE CREATING SUBJECT

All mathematics is a mental construction performed by an idealized subject, subject to the following observations about knowledge:

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2. at a point in time, the subject knows whether or not it has experienced a judgment (decidability)

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Corollary
The meaning of a judgment $\mathscr{J}$ must be explained in terms of its forcing condition, $\boxed{w \Vdash \mathscr{J}}$*, for any stage/world* w*.*

...

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Remark

Contra Dummett, I by no means take the above as requiring that the following shall be true in a constructive metatheory, divorced from time:

$$\forall w. \forall \mathscr{J}. (w \Vdash \mathscr{J}) \vee \neg(w \Vdash \mathscr{J}) \qquad \text{(Dummett's infelicity)}$$

The above is impossible in a Beth model.

logical consequence $\Rightarrow$ semantic consequence

Brouwer?, Martin-Löf, Sundholm $\Rightarrow$ *Brouwer?, Heyting, Allen, Zeilberger*

MEANING THEORY WITHOUT UNIFORMITY CONSTRAINTS

logical consequence ⇒ semantic consequence

Brouwer?, Martin-Löf, Sundholm ⇒ *Brouwer?, Heyting, Allen, Zeilberger*

proof conditions ⇒ assertion conditions

Dummett, Martin-Löf, Sundholm ⇒ *Brouwer, Heyting, Van Atten*

MEANING THEORY WITHOUT UNIFORMITY CONSTRAINTS

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Dummett, Martin-Löf, Sundholm ⇒ *Brouwer, Heyting, Van Atten*

global meaning explanation ⇒ **local meaning explanation**

Husserl, Dummett, Martin-Löf ⇒ *Brouwer, Beth, Kripke, Grothendieck, Lawvere, Joyal*

assertion acts (judgments) are intensional (local)

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$\vert_x \mathscr{J}(x)$ (general judgment)

$\mathscr{J}_2 \ (\mathscr{J}_1)$ (hypothetical judgment)

$M \Downarrow N$ (evaluation)

BETH-KRIPKE SEMANTICS FOR ASSERTIONS

assertion acts (judgments) are intensional (local)

$\mid_x \mathscr{J}(x)$	(general judgment)
$\mathscr{J}_2\ (\mathscr{J}_1)$	(hypothetical judgment)
$M \Downarrow N$	(evaluation)
$A\ \mathit{type}$	(typehood)

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$$A\ \mathit{type} \qquad \text{(typehood)}$$

$$A\ \mathit{verif} \qquad \text{(verification)}$$

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$A\ \mathit{verif}$	(verification)
$A\ \mathit{true}$	(truth)
$M = N \in A$	(membership)

assertion acts (judgments) are intensional (local)

$w \Vdash \mid_x \mathscr{J}(x)$	(general judgment)
$w \Vdash \mathscr{J}_2 \ (\mathscr{J}_1)$	(hypothetical judgment)
$w \Vdash M \Downarrow N$	(evaluation)
$w \Vdash A \ \mathit{type}$	(typehood)
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SEMANTICS OF HIGHER-ORDER ASSERTIONS

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$$w \Vdash \mathscr{J}_2 \ (\mathscr{J}_1) \qquad \text{(hypothetical judgment)}$$

SEMANTICS OF HIGHER-ORDER ASSERTIONS

$$w \Vdash |_x \, \mathscr{I}(x) \iff \cdots$$

$$w \Vdash \mathscr{I}_2 \; (\mathscr{I}_1) \iff \cdots$$

$$\begin{aligned} w \Vdash \,\mid_x \mathscr{J}(x) &\iff \forall u \geq w. \forall x \in \mathscr{D}_u.\ u \Vdash \mathscr{J}(x) \\ w \Vdash \mathscr{J}_2\ (\mathscr{J}_1) &\iff \forall u \geq w.\ u \Vdash \mathscr{J}_1 \Rightarrow u \Vdash \mathscr{J}_2 \end{aligned}$$

$$
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w \Vdash |_x\, \mathscr{J}(x) &\iff \forall u \geq w . \forall x \in \mathscr{D}_u .\ u \Vdash \mathscr{J}(x) \\
w \Vdash \mathscr{J}_2\ (\mathscr{J}_1) &\iff \forall u \geq w .\ u \Vdash \mathscr{J}_1 \Rightarrow u \Vdash \mathscr{J}_2
\end{aligned}
$$

where $\mathscr{D}_w$ is the species of constructions that have been effected by stage w

The meaning of a proposition/type is an intensional (world-indexed) specification of verification acts, i.e. a local meaning explanation for $\boxed{w \Vdash P\ \mathit{verif}}$ (and its synthesis).

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For a type A, implicit in the explanation of $w \Vdash A\ \mathit{verif}$ is a $\mathbb{W}$-indexed family of PERs $\mathscr{V}[\![A]\!]_w \subseteq \mathscr{D}_w \times \mathscr{D}_w$ whose members reflect the computational content (extension) of verification acts.

Truth (justification) consists in recognizing the effectiveness of a procedure for verification.

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In the model, this corresponds to the inevitability of verification (i.e. a bar, in which verification occurs at all nodes):

$$w \Vdash A \; true \iff \exists \mathfrak{B} \; \mathbf{bars} \; w. \forall u \in \mathfrak{B}. \; u \Vdash A \; verif \qquad \text{(due to Dummett)}$$

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1. **Canonical membership** reflects **verification**

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2. **Membership** reflects **justification**

$$w \Vdash M = N \in A \quad \bowtie \quad w \Vdash A \; \mathit{true}$$

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1. **Canonical membership** reflects **verification**
2. **Membership** reflects **justification**
3. **Computation** reflects **the recognition of a bar**

$$w \Vdash M = N \in A \quad \rtimes\!\ltimes \quad w \Vdash A\ \mathit{true}$$

The analytic judgments of type theory are reflections on mathematical activity.

1. **Canonical membership** reflects **verification**
2. **Membership** reflects **justification**
3. **Computation** reflects **the recognition of a bar**

$$\bigwedge \left\{ \begin{array}{l} w \Vdash M \Downarrow M' \\ w \Vdash N \Downarrow N' \\ \mathscr{V}[\![A]\!]_w(M', N') \end{array} \right\} \bowtie \quad \exists \mathfrak{B}\ \textbf{bars}\ w. \forall u \in \mathfrak{B}.\ u \Vdash A\ \mathit{verif}$$

INTUITIONISTIC VS CONSTRUCTIVE VALIDITY

a statement is intuitionistically valid if it is forced at all nodes

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Intuitionism subsumes constructivism, but goes much further by allowing the observation of non-constructive objects (Fourman)

THE **require** ORACLE: DYNAMICS

choice sequences (streams of objects) may be propounded over time based on the previous experience of the creating subject.

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Example:

$$\alpha(i) = \begin{cases} 0 & i \Vdash A\ \mathit{true} \\ 1 & \neg(i \Vdash A\ \mathit{true}) \end{cases} \tag{KS}$$

Let $\mathscr{K}_A : \mathbf{FinSet}^{\mathbb{W}^{\mathrm{op}}}$ be the sheaf of constructions of A *true* effected "so far" for each canonical proposition A.

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We now can give a precise, but non-deterministic, dynamics to **require**:

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We now can give a precise, but non-deterministic, dynamics to **require**:

$$\frac{w \Vdash A \Downarrow A' \quad M \in \mathscr{K}_{A'}(w) \quad w \Vdash [M/x]N \Downarrow N'}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \Downarrow N'} \tag{*}$$

We need a way to deterministically choose a representative of $\mathscr{K}_A(w)$. First, let $\varkappa_A$ be the choice sequence of lists given by enumerating $\mathscr{K}_A(w)$ at each stage w, in order of time.

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Idea: reformulate Type Theory relative to a choice sequence of "choosers".

A spread direction $\mathfrak{S}$ is a restriction on choice sequences which is defined by a condition on their finite approximations (prefixes, neighborhoods), subject to the following laws:

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1. the empty neighborhood shall be admitted

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2. if a neighborhood is admitted, so shall all its subneighborhoods

$$\vert_{\vec{u},m}\ \mathfrak{S}(\vec{u})\ \left(\mathfrak{S}(\vec{u} \frown m)\right)$$

SPREADS: SETS OF CHOICE SEQUENCES

A spread direction $\mathfrak{S}$ is a restriction on choice sequences which is defined by a condition on their finite approximations (prefixes, neighborhoods), subject to the following laws:

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$$|_{\vec{u},m}\ \mathfrak{S}(\vec{u})\ \left(\mathfrak{S}(\vec{u} \frown m)\right)$$

3. a neighborhood may always be refined within the spread

$$|_{\vec{u}}\ \mathfrak{S}(\vec{u} \frown m)\ \left(\mathfrak{S}(\vec{u})\right)$$

A spread direction for index-choosers:

$$\frac{}{\mathfrak{S}(\langle\rangle)} \qquad \frac{\mathfrak{S}(\vec{u}) \quad |_n\ \rho(n) < n \ \ (n \in \mathbb{N}^+)}{\mathfrak{S}(\vec{u} \frown \rho)} \qquad \text{(spread law)}$$

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Reformulate type theory relative to an arbitrary $\alpha \in \mathfrak{S}$! For instance:

$$\frac{w \Vdash M \Downarrow_\alpha M' \quad w \Vdash N \Downarrow_\alpha N' \quad \mathscr{V}[\![A]\!]_w^\alpha (M', N')}{w \Vdash M = N \in_\alpha A}$$

Deterministic choice for $\varkappa_A$:

$$\frac{w \Vdash A \Downarrow_\alpha A' \quad |\varkappa_{A'}(w)| = \ell \quad \mathsf{hd}(\alpha)(\ell) = i \quad \varkappa_{A'}(w)(i) = M}{w \Vdash \varkappa_A \ni_\alpha M}$$

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Deterministic dynamics for **require**:

$$\frac{w \Vdash \varkappa_A \ni_\alpha M \quad w \Vdash [M/x]N \Downarrow_{\mathsf{tl}(\alpha)} N'}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \Downarrow_\alpha N'} \qquad (\text{for } \alpha \in \mathfrak{S})$$

Theorem

The following rule is valid in our semantics:

$$\frac{w \Vdash A\ \mathit{true} \quad w \Vdash x : A \vdash_\alpha N \in B}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\frac{w \Vdash A\ \mathit{true} \quad w \Vdash x : A \vdash_\alpha N \in B}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\frac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \frac{}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

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$$\frac{\dfrac{}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \Downarrow_\alpha N'} \quad \dfrac{}{\mathscr{V}[\![B]\!]^\alpha_w (N', N')}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \in_\alpha B}\ \mathit{require}$$

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$$\dfrac{\dfrac{\dfrac{}{w \Vdash \varkappa_A \ni_\alpha M} \quad \dfrac{}{[M/x]N \Downarrow_\alpha N'}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \Downarrow_\alpha N'} \qquad \dfrac{}{\mathscr{V}[\![B]\!]^\alpha_w (N', N')}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \in_\alpha B}\ \mathit{require}$$

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$$\frac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \frac{\overline{w \Vdash \rvert_{y,z}\, [y/x]N = [z/x]N \in_\alpha B\ \left(y = z \in_\alpha A\right)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\frac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \overline{[M/x]N \Downarrow_\alpha N'}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \Downarrow_\alpha N'} \qquad \overline{\mathscr{V}\,[\![B]\!]_w^\alpha\,(N', N')}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\dfrac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \dfrac{\dfrac{\overline{\forall u{\geq}w.\forall y,z{\in}\mathscr{D}_u.\ u{\Vdash}y{=}z{\in_\alpha}A \Rightarrow u{\Vdash}[y/x]N{=}[z/x]N{\in_\alpha}B}\ \mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \left(y = z \in_\alpha A\right)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\dfrac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \dfrac{}{[M/x]N \Downarrow_\alpha N'}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \Downarrow_\alpha N'} \qquad \dfrac{}{\mathscr{V}\,[\![B]\!]^\alpha_w\,(N', N')}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\frac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \cfrac{\cfrac{\overline{\forall u \geq w. \forall y,z \in \mathscr{D}_u.\ u \Vdash y = z \in_\alpha A \Rightarrow u \Vdash [y/x]N = [z/x]N \in_\alpha B}\ \mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \left(y = z \in_\alpha A\right)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\cfrac{\cfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\ \mathscr{F}(w, M, M) \qquad \cfrac{}{w \Vdash M \in_\alpha A}}{}$$

$$\cfrac{\cfrac{\cfrac{\cfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \cfrac{}{[M/x]N \Downarrow_\alpha N'}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \Downarrow_\alpha N'} \qquad \cfrac{}{\mathscr{V}\,[\![B]\!]^\alpha_w\,(N', N')}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\dfrac{}{w \Vdash A\ \mathit{true}}\,\mathscr{D} \qquad \dfrac{\dfrac{\dfrac{}{\forall u \geq w. \forall y,z \in \mathscr{D}_u.\ u \Vdash y = z \in_\alpha A \Rightarrow u \Vdash [y/x]N = [z/x]N \in_\alpha B}\,\mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \big(y = z \in_\alpha A\big)}}{w \Vdash x : A \vdash_\alpha N \in B}\,\mathscr{E}$$

$$\dfrac{\dfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\ \mathscr{F}(w,M,M) \qquad \dfrac{\dfrac{}{w \Vdash \varkappa_A \ni_\alpha M}}{w \Vdash M \in_\alpha A}}{}$$

$$\dfrac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\,\mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \dfrac{}{[M/x]N \Downarrow_\alpha N'}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \Downarrow_\alpha N'} \qquad \dfrac{}{\mathscr{V}[\![B]\!]^\alpha_w (N', N')}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\frac{}{w \Vdash A\ \mathit{true}}\,\mathscr{D} \qquad \frac{\dfrac{\dfrac{}{\forall u{\geq}w.\forall y,z{\in}\mathscr{D}_u.\ u{\Vdash}y{=}z{\in_\alpha}A \Rightarrow u{\Vdash}[y/x]N{=}[z/x]N{\in_\alpha}B}\,\mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \big(y = z \in_\alpha A\big)}}{w \Vdash x : A \vdash_\alpha N \in B}\,\mathscr{E}$$

$$\frac{\dfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\,\mathscr{F}(w,M,M) \qquad \dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\,\mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha {\color{darkred}M}}}{w \Vdash M \in_\alpha A}}{}$$

$$\frac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\,\mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha {\color{darkred}M}} \qquad \dfrac{}{[M/x]N \Downarrow_\alpha {\color{darkred}N'}}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \Downarrow_\alpha {\color{darkred}N'}} \qquad \dfrac{}{\mathscr{V}[\![B]\!]^\alpha_w (N', N')}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\dfrac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \dfrac{\dfrac{\dfrac{}{\forall u{\geq}w.\forall y,z{\in}\mathscr{D}_u.\ u{\Vdash}y{=}z{\in_\alpha}A \Rightarrow u{\Vdash}[y/x]N{=}[z/x]N{\in_\alpha}B}\ \mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \left(y = z \in_\alpha A\right)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\dfrac{\dfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\ \mathscr{F}(w,M,M) \qquad \dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M}}{w \Vdash M \in_\alpha A}}{w \Vdash [M/x]N \in_\alpha B}$$

$$\dfrac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \dfrac{}{[M/x]N \Downarrow_\alpha N'}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \Downarrow_\alpha N'} \qquad \dfrac{}{\mathscr{V}[\![B]\!]^\alpha_w(N',N')}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\dfrac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \dfrac{\dfrac{\dfrac{}{\forall u{\geq}w.\forall y,z{\in}\mathscr{D}_u.\ u{\Vdash}y{=}z{\in_\alpha}A \Rightarrow u{\Vdash}[y/x]N{=}[z/x]N{\in_\alpha}B}\ \mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \left(y = z \in_\alpha A\right)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\dfrac{\dfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\ \mathscr{F}(w, M, M) \quad \dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M}}{w \Vdash M \in_\alpha A}}{\overline{w \Vdash [M/x]N \in_\alpha B}}$$

$$\dfrac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \quad \dfrac{}{[M/x]N \Downarrow_\alpha N'}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \Downarrow_\alpha N'} \quad \dfrac{}{\mathscr{V}\,[\![B]\!]_w^\alpha\,(N', N')}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\dfrac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \dfrac{\dfrac{\overline{\forall u{\geq}w.\forall y,z{\in}\mathscr{D}_u.\ u{\Vdash}y{=}z{\in_\alpha}A \Rightarrow u{\Vdash}[y/x]N{=}[z/x]N{\in_\alpha}B}\ \mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \big(y = z \in_\alpha A\big)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\dfrac{\dfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\ \mathscr{F}(w, M, M) \qquad \dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M}}{w \Vdash M \in_\alpha A}}{\dfrac{\overline{w \Vdash [M/x]N \Downarrow_\alpha N'}\ \mathscr{G} \qquad \overline{\mathscr{V}[\![B]\!]^\alpha_w (N', N')}\ \mathscr{H}}{w \Vdash [M/x]N \in_\alpha B}}$$

$$\dfrac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \overline{[M/x]N \Downarrow_\alpha N'}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \Downarrow_\alpha N'} \qquad \overline{\mathscr{V}[\![B]\!]^\alpha_w (N', N')}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\dfrac{}{w \Vdash A \ \mathit{true}}\ \mathscr{D} \qquad \dfrac{\dfrac{\dfrac{}{\forall u{\geq}w.\forall y,z{\in}\mathscr{D}_u.\ u{\Vdash}y{=}z{\in_\alpha}A \Rightarrow u{\Vdash}[y/x]N{=}[z/x]N{\in_\alpha}B}\ \mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B \ \left(y = z \in_\alpha A\right)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\ \mathscr{F}(w, M, M) \qquad \dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A \ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M}}{w \Vdash M \in_\alpha A}}{\dfrac{}{w \Vdash [M/x]N \Downarrow_\alpha N'}\ \mathscr{G} \qquad \dfrac{}{\mathscr{V}[\![B]\!]_w^\alpha (N', N')}\ \mathscr{H}}}{w \Vdash [M/x]N \in_\alpha B}$$

$$\dfrac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A \ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \dfrac{\checkmark}{[M/x]N \Downarrow_\alpha N'}\ \mathscr{G}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \Downarrow_\alpha N'} \qquad \dfrac{}{\mathscr{V}[\![B]\!]_w^\alpha (N', N')}}{w \Vdash \mathbf{require}\ x : A\ \mathbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\dfrac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \dfrac{\dfrac{\overline{\forall u{\geq}w.\forall y,z{\in}\mathscr{D}_u.\ u{\Vdash}y{=}z{\in_\alpha}A \Rightarrow u{\Vdash}[y/x]N{=}[z/x]N{\in_\alpha}B}\ \mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \big(y = z \in_\alpha A\big)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\ \mathscr{F}(w,M,M) \qquad \dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M}}{w \Vdash M \in_\alpha A}}{\overline{w \Vdash [M/x]N \Downarrow_\alpha N'}\ \mathscr{G} \qquad \overline{\mathscr{V}[\![B]\!]^\alpha_w (N', N')}\ \mathscr{H}}}{w \Vdash [M/x]N \in_\alpha B}$$

$$\dfrac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \dfrac{\checkmark}{[M/x]N \Downarrow_\alpha N'}\ \mathscr{G}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \Downarrow_\alpha N'} \qquad \dfrac{\checkmark}{\mathscr{V}[\![B]\!]^\alpha_w (N', N')}\ \mathscr{H}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$$\dfrac{}{w \Vdash A\ \mathit{true}}\ \mathscr{D} \qquad \dfrac{\dfrac{\overline{\forall u{\geq}w.\forall y,z{\in}\mathscr{D}_u.\ u{\Vdash}y{=}z{\in_\alpha}A \Rightarrow u{\Vdash}[y/x]N{=}[z/x]N{\in_\alpha}B}\ \mathscr{F}}{w \Vdash |_{y,z}\ [y/x]N = [z/x]N \in_\alpha B\ \big(y = z \in_\alpha A\big)}}{w \Vdash x : A \vdash_\alpha N \in B}\ \mathscr{E}$$

$$\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash M \in_\alpha A \Rightarrow w \Vdash [M/x]N \in_\alpha B}\ \mathscr{F}(w, M, M) \quad \dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M}}{w \Vdash M \in_\alpha A}}{\dfrac{}{w \Vdash [M/x]N \Downarrow_\alpha N'}\ \mathscr{G} \qquad \dfrac{}{\mathscr{V}[\![B]\!]^{\alpha}_{w}(N', N')}\ \mathscr{H}}}{w \Vdash [M/x]N \in_\alpha B}$$

$$\dfrac{\dfrac{\dfrac{\dfrac{\checkmark}{w \Vdash A\ \mathit{true}}\ \mathscr{D}}{w \Vdash \varkappa_A \ni_\alpha M} \qquad \dfrac{\checkmark}{[M/x]N \Downarrow_\alpha N'}\ \mathscr{G}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \Downarrow_\alpha N'} \qquad \dfrac{\checkmark}{\mathscr{V}[\![B]\!]^{\alpha}_{w}(N', N')}\ \mathscr{H}}{w \Vdash \textbf{require}\ x : A\ \textbf{in}\ N \in_\alpha B}\ \mathit{require}$$

$\square$

Questions?