```
Section 1

Execution
```

The code can be executed with ./run.sh <num time steps>

```
Section 2

Explanation of the code
```

The main code is in src.c which performs one run of message parsing by 3 different ways. Following is the general schema for each run:

- 1. Initialize data as a dynamic (but contiguous) block of size N*N with random **double** in the range [0,1].
- 2. Repeat the following steps num_time_steps times:
 - Compute the stencil for the data.
 - Send/Recv in all possible directions by a specific methodology.
 - Compute the final halo (only the boundaries).
- 3. Free the data, and related data structures.

Schema 1: General Structure of the code

```
1 double stime = MPI_Wtime();
 2 for(int t=0; t<num_time_steps; t++){</pre>
        // Perform stencil computation
        compute stencil(data, side len);
 4
5
        // Send to possible directions
 6
 7
       if(can transfer('u', my rank, cluster len, p)){
 8
            // If it is possible to transfer upwards, send.
 9
       if(can transfer('d', my rank, cluster len, p)){
10
11
            // If it is possible to transfer downwards, send.
12
       if(can_transfer('l', my_rank, cluster_len, p)){
13
            // If it is possible to transfer leftwards, send.
14
15
       if(can transfer('r', my rank, cluster len, p)){
16
            /// If it is possible to transfer rightwards, send.
17
18
19
20
        // Recieve from possible directions
       if(can_transfer('u', my_rank, cluster_len, p)){
21
22
            // If it is possible to transfer upwards, receive.
23
       if(can_transfer('d', my_rank, cluster_len, p)){
24
            // If it is possible to transfer downwards, receive.
25
26
27
       if(can transfer('l', my rank, cluster len, p)){
28
           // If it is possible to transfer leftwards, receive.
29
```

```
30
       if(can_transfer('r', my_rank, cluster_len, p)){
31
           // If it is possible to transfer rightwards, receive.
32
33
34
       // Get the final averages for time t
35
       compute halo(data, recv data, side len, my rank, cluster len, p);
36 }
37 double ftime = MPI Wtime();
38
39 double time = ftime - stime;
40 MPI Reduce (&time, &maxTime, 1, MPI DOUBLE, MPI MAX, 0, MPI COMM WORLD);
41
42 // print only by rank 0
43 if (!my_rank) printf ("%lf\n", maxTime);
```

Following the above schema, the three strategies for message send/receive can be implemented inside each of the **if** condition. Following are three excerpts from our code for the three given strategies.

Part 1: Single MPI_Send/Recv

```
1 // sending upwards
2 if(can_transfer('u', my_rank, cluster_len, p)){
      int target = my rank - cluster len;
4
       for(int i=0; i<side len; i++){
5
           MPI Send(&data[0*side len + i], 1, MPI DOUBLE, target, i, MPI COMM WORLD);
6
7
8
9 // recv from below
10 if(can_transfer('d', my_rank, cluster_len, p)){
       int source = my_rank + cluster_len;
12
       MPI Status st;
13
       for(int i=0; i < side_len; i++){
14
           MPI Recv(&recv data[1][i], 1, MPI DOUBLE, source, i, MPI COMM WORLD, &st);
15
       }
16 }
```

Here we need to make side len number of MPI Send() MPI Recv() calls for transfer in a single direction.

Part 2: MPI_Send/Recv using MPI_Pack/Unpack

```
1 // sending upwards
 2 if(can transfer('u', my rank, cluster len, p)){
       int target = my rank - cluster len, upos = 0;
 4
       for(int j=0; j < side len; <math>j++){
           \label{eq:mpi_ack} $$ MPI\_Pack (\&data[0*side\_len + j], 1, MPI\_DOUBLE, s\_ubuf, side\_len*8, \&upos, MPI\_COMM\_WORLD); $$
 5
 6
       MPI Send(s_ubuf, side_len*8, MPI_PACKED, target, 1, MPI_COMM_WORLD);
 7
 8 }
10 // recv from below
11 if(can transfer('d', my rank, cluster len, p)){
       int source = my rank + cluster len, curr pos = 0;
12
13
       MPI Status st;
       MPI Recv(r dbuf, side len*8, MPI PACKED, source, 1, MPI COMM WORLD, &st);
14
```

```
MPI_Unpack(r_dbuf, side_len*8, &curr_pos, recv_data[1], side_len, MPI_DOUBLE, MPI_COMM_WORLD);

16 }
```

Here we need to make a single MPI_Send() MPI_Recv() call for transfer in a single direction with side_len number of MPI_Pack()/MPI_Unpack().

Part 3: MPI_Send/Recv using MPI_Datatype Vector

```
1 // define datatypes for row and col
2 MPI Datatype row vector, col vector;
3 MPI Type vector (side len, 1, 1, MPI DOUBLE, &row vector);
4 MPI Type commit (&row_vector);
5 MPI Type vector (side len, 1, side len, MPI DOUBLE, &col vector);
6 MPI Type commit (&col vector);
  // sending upwards (same for row and col)
9 if(can transfer('u', my rank, cluster len, p)){
      int target = my rank - cluster len;
10
11
       MPI Send(&data[0], 1, row vector, target, 1, MPI COMM WORLD);
12 }
13
14 // recv from below (row)
15 if(can_transfer('d', my_rank, cluster_len, p)){
      int source = my rank + cluster len;
17
       MPI Status st;
18
       MPI Recv(recv data[1], 1, row vector, source, 1, MPI COMM WORLD, &st);
19 }
20
21 // recv from left
22 if(can_transfer('l', my_rank, cluster_len, p)){
23
      int source = my rank -1;
24
       MPI Status st;
25
       MPI Recv(recv data[2], side len, MPI DOUBLE, source, 4, MPI COMM WORLD, &st);
26 }
```

Here we need to make a single MPI_Send() MPI_Recv() call for transfer in a single direction. Note that while receiving a row, we can safely use the row_vector datatype, while we need to use MPI_DOUBLE when recv for a column is made because we want to fit it in a horizontal buffer.

```
Section 3

Issues faced during testing/development
```

Following are some mentionable bugs we faced during code development/testing:

- 1. During MPI_Send()/MPI_Recv() when only one single MPI_DOUBLE was sent/recv, we had a problem in our implementation where our send/recv buffers were overlapping.
- 2. During MPI_Pack()/MPI_Unpack(), we were not resetting the value of position variable, thus we were getting segmentation faults
- 3. During MPI_Unpack() we realized that we were losing the data when using MPI_Isend(), this was because the call was non-blocking and therefore we were upacking before the buffer was filled. This could be fixed in two ways, first using MPI_Wait() for the specific call, and then unpacking. Second, to use a blocking recv so that the unpacking buffer is filled with the sender's data. We used the latter in our implementation (in all of the send/recv calls blocking send/recv are used).

- 4. While using MPI_Type_vector(), we faced a segfault because we were receiving the col_vector datatype in a horizontal buffer, since the layout is preserved, the recv call was trying to fill the recv buffer out of bounds thus producing the error. We fixed this by using MPI_DOUBLE and receiving one of it at a time.
- 5. While setting up the NodeAllocator, we faced a package not found error, so we had to manually install psutil. Since at that time, the CSE cluster was not connected to the internet, we had to use this script to manually connect to the internet: https://github.com/satendrapandeymp/Auth_IITK/blob/master/auth-iitk.py.

Section 4

Observations

The following 4 plots are corresponding to the number of processes involved, i.e. 16, 36, 49, 64.

- In all of the cases, multiple send/recv takes the highest time. This is due to the overhead for multiple send/recv calls.
- Using MPI_Type_vector takes the smallest time because of minimum overhead of type creation. In pack/unpack, there is an extra overhead due to packing/unpacking, while in send/recv, overhead is introduced due to multiple send/recv calls.

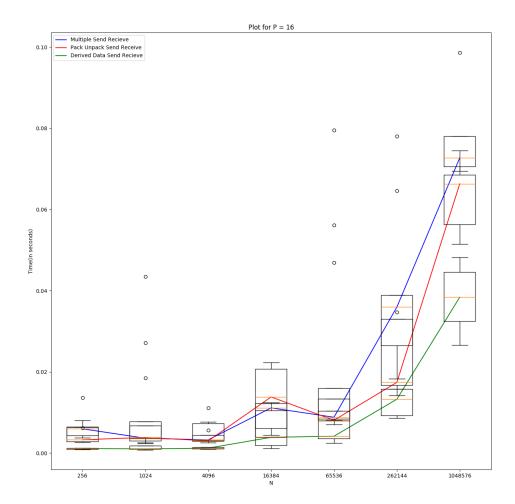


Figure 1: Number of Processes = 16

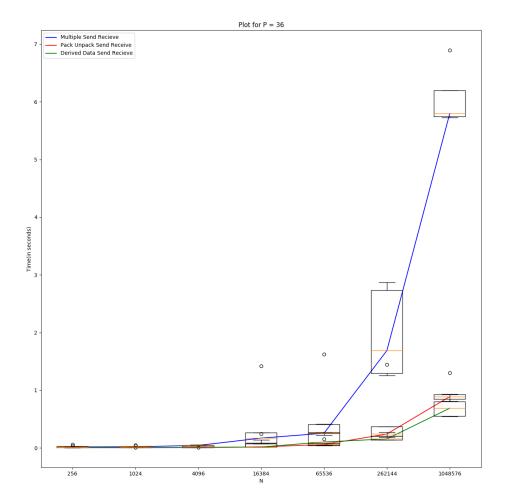


Figure 2: Number of Processes = 36

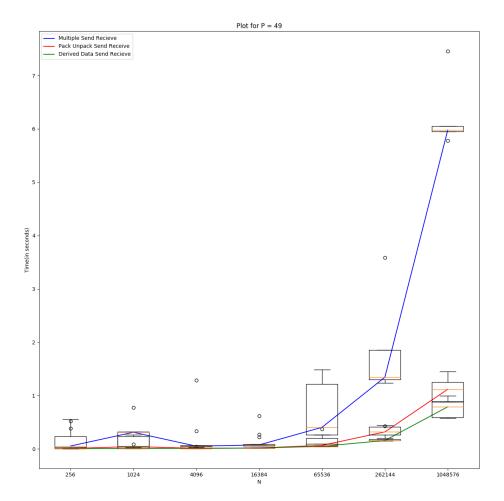


Figure 3: Number of Processes = 49

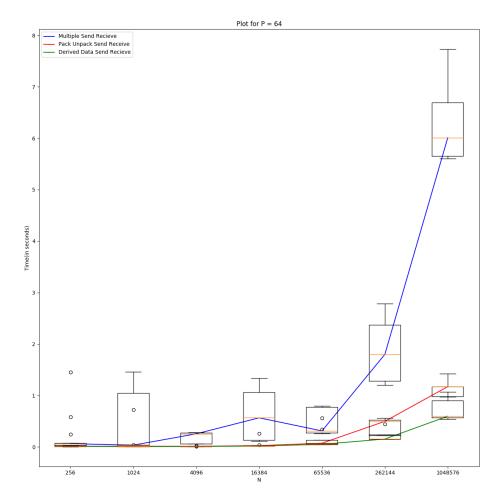


Figure 4: Number of Processes = 64