

# **Distributed Directory Service**

17.03.2019

# **Group 7**

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# **Overview**

The project aims to realize a distributed directory service. We aim to deploy a subset of the functionality described in the OSI LDAP standard, but hope to get down the exact implementation working according to our own terms.

# **User Features**

We plan to implement the following user features for the three user categories:

# 1. Directory Admins:

- I. Be able to specify and modify directory schema.
- II. Manage access controls over the directory data.

# 2. Publishers:

- I. Modify existing data
- II. Add entries
- III. Delete entries

#### 3. Subscribers:

- I. Read data of an object, if they know its name
- II. Search for objects, filtered by certain rules, within a specified scope

# **System Features**

We expect the following system features out of our service:

### 1. Fault Model:

- I. One crash fault tolerance
- II. One link failure tolerance:

Hopefully, we will be able to achieve at least a somewhat degraded service. We just don't want to let the link failure go unhandled. Publishers:

#### 2. Consistency Model:

→ Eventual Consistency with a time constraint:

The published updates must become consistent across the database within \*some\* fixed number of read requests.

# **Design Questions**

- 1. Where to store complete directory:
  - a. All eggs in one basket (single node):

#### Pro:

- I. Easy to maintain
- II. Don't need to go any further to search for data

#### Con:

- I. Highly resource intensive. Impractical for large directories.
- II. Longer search time as no true parallelization (can only achieve parallel search via threading)
- b. Distribute among multiple nodes:

#### Pro:

- I. Distributed Load. Truly parallel search for data possible.
- II. Faster search compared to single node if there is lots of data
- III. Scalable

#### Con:

- I. Need to reroute requests, probably need to remember what data might be where
- II. For smaller data, may perform slower due to rerouting overhead
- III. DIT structure change will become very complicated

**CHOSEN OPTION:** Distribute tree data over multiple nodes.

#### **NEED TO ALSO THINK:**

How to partition the tree and what servers must be connected to what servers? Should the network structure be static or dynamic (depending on request frequencies)?

# 2. How to reroute requests:

# a. Chaining:

#### Pro:

- I. Less work for DUA
- II. Access is secure amongst directory servers

#### Con:

- I. More work for us
- II. Need to remember forwarded query contexts until solved
- III. Requesting faraway information repeatedly via the same chain is as inefficient as things can be

#### b. Referrals:

#### Pro:

- I. Less work for us
- II. Requesting faraway information repeatedly becomes fast as you know where to go once you ask

#### Con:

- I. More work for DUA
- II. Users get to know more about the directory structure. Might introduce some security issues.

CHOSEN OPTION: Some hybrid of the two would be preferable.

#### 3. Who detects the fault?

#### a. Centralized Node:

#### Pro:

- I. It knows about everything. Can handle issues like link failure better.
- II. Can resolve issues that might affect multiple nodes at one.

#### Con:

- I. Too much communication traffic to central node.
- II. Everything needs to go to one guy, so fault resolution will be slower.
- III. Tougher to prove if it can handle everything.

# b. Local neighbours:

# Pro:

- I. No overhead of getting all information anywhere. Everything is locally resolved.
- II. If we can prove it work for one neighbour pair, the unit can repeated and is guaranteed to handle all sorts of faults.
- III. More scalable.

#### Con:

- I. Link failures are tough to handle as both neighbours believe the other is dead while both are alive.
- II. May not be able to handle more widely disruptive faults using only local resolutions.

**CHOSEN OPTION: Leaning toward local fault resolution.** 

Will try to figure out how to handle link failures.

# 4. How many replicas?

#### a. Two:

#### Pro:

- I. Can handle single crash fault.
- II. Consistency issues are easier to resolve.

#### Con:

I. Data access points may be too far from where data is kept. So, longer chains, and more latency.

# b. Many:

#### Pro:

I. Data access will be faster if data is distributed evenly around the globe.

## Con:

I. Consistency gets tougher to maintain across multiple replicas.

CHOSEN OPTION: Will try to keep few (three-four) replicas, but get the data dynamically closer to where they are being accessed.

# 5. How to maintain data consistency?

a. Passive Replication:

## Pro:

- I. Consistency is always maintained.
- II. Simpler to implement.

#### Con:

I. Writes are slower.

# b. Active Replication:

#### Pro:

- I. Flexibility. Lots of ways to maintain consistency.
- II. Can be made as fast as needed.

#### Con:

I. Implementation is more complex.

CHOSEN OPTION: Passive Replication as we expect very few writes.