## CHAPTER 1: BASIC CONCEPTS IN GRAPH THEORY

## **Discrete Mathematics 2**

Lecturer: Nguyen Kieu Linh

Posts and Telecommunications Institute of Technology

Hanoi, 2023 http://www.ptit.edu.vn



#### Contents

- 1 Graph Definitions
  - Graphs
  - Undirected graph
  - Directed graph
- 2 Basic terminologies in undirected graphs
- 3 Basic terminologies in directed graphs
- 4 Some special types of graphs



#### Contents

- 1 Graph Definitions
  - Graphs
  - Undirected graph
  - Directed graph
- 2 Basic terminologies in undirected graphs
- 3 Basic terminologies in directed graphs
- 4 Some special types of graphs



## **Graphs**

#### Definition

A graph  $G = \langle V, E \rangle$  consists of V, a nonempty set of vertices (or nodes) and E, a set of edges. Each edge has either one or two vertices associated with it, called its endpoints. An edge is said to connect its endpoints.

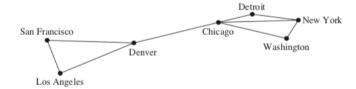


Figure: A Computer Network.



#### Contents

- 1 Graph Definitions
  - Graphs
  - Undirected graph
  - Directed graph
- 2 Basic terminologies in undirected graphs
- 3 Basic terminologies in directed graphs
- 4 Some special types of graphs



## Some types of Graphs

- \* Undirected graph
  - \* Undirected Simple Graph
  - ⋆ Undirected Multigraph
  - ⋆ Undirected Pseudograph
- \* Directed graph
  - ⋆ Directed Simple Graph
  - ⋆ Directed Multigraph



## **Undirected Simple Graph**

## Undirected simple graph $G = \langle V, E \rangle$ :

- $\star V$  is the set of vertices
- ★ E is the set of edges, consisting of unordered pairs of two distinct vertices in V
- \* There is at most one edge connecting two vertices.

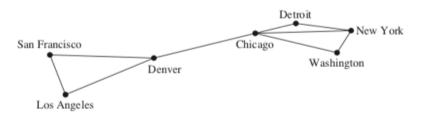


Figure: Undirected Simple Graph.



## **Undirected Multigraph**

## Undirected multigraph $G = \langle V, E \rangle$ :

- $\star$  V is the set of vertices
- ★ E is the set of edges, consisting of unordered pairs of two distinct vertices in V
- \*  $e_1, e_2 \in E$  are called multiple edges if they connect the same two vertices,  $e_1 = (u, v), e_2 = (v, u)$ .



Figure: Undirected multigraph.



## **Undirected Pseudograph**

## Undirected Pseudograph $G = \langle V, E \rangle$ :

- $\star$  V is the set of vertices
- $\star$  E is the set of edges, consisting of unordered pairs of two vertices (maybe the same) in V
- \* The graph includes edges that connect a vertex to itself. Such edges are called *loops*, and sometimes we may even have more than one loop at a vertex.

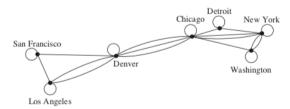


Figure: Undirected Pseudograph.



#### Contents

- 1 Graph Definitions
  - Graphs
  - Undirected graph
  - Directed graph
- 2 Basic terminologies in undirected graphs
- 3 Basic terminologies in directed graphs
- 4 Some special types of graphs



## Directed Simple Graph

## Directed Simple Graph $G = \langle V, E \rangle$ :

- $\star$  V is the set of vertices
- ★ E is the set of *directed edges* (or arcs, arrows), consisting of ordered pairs of two distinct vertices in V
- \* There is at most one directed edges (or arcs) from a vertex u to another one v. The directed edge associated with the ordered pair (u, v) is said to start at u and end at v.



Figure: Directed Simple Graph.



## Directed Multigraph

## Directed Multigraph $G = \langle V, E \rangle$ :

- $\star$  V is the set of vertices
- ★ E is the set of <u>directed edges</u> (or <u>arcs</u>), consisting of ordered pairs of two distinct vertices in V
- ★  $e_1, e_2 \in E$  are called *multiple directed edges* if they connect the same two vertices,  $e_1 = (u, v), e_2 = (v, u)$ .



Figure: Directed Simple Graph.



#### Convention

- We will focus on Undirected Simple Graph and Directed Simple Graph
- \* "Undirected Graph" means "Undirected Simple Graph"
- \* "Directed Graph" means "Directed Simple Graph"



#### Contents

- 1 Graph Definitions
  - Graphs
  - Undirected graph
  - Directed graph
- 2 Basic terminologies in undirected graphs
- 3 Basic terminologies in directed graphs
- 4 Some special types of graphs



## Vertex Degree

#### Definition

Two vertices u and v in an undirected graph G are called *adjacent* (or neighbors) in G if u and v are endpoints of an edge e of G. Such an edge e is called *incident* with the vertices u and v and e is said to connect u and v.

#### Definition

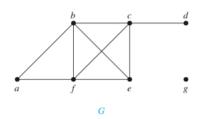
The degree of a vertex in an undirected graph is the number of edges incident with it, except that a loop at a vertex contributes twice to the degree of that vertex. The degree of the vertex v is denoted by deg(v).

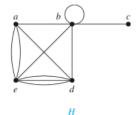


## Vertex Degree

## **Example**

What are the degrees in the graphs G and H displayed in following figure?





- \* A vertex of degree zero is called *isolated*.
- \* A vertex is *pendant* if and only if it has degree one.



## Vertex Degree

#### **Theorem**

\* Let  $G = \langle V, E \rangle$  be an undirected graph with m edges. Then

$$2m = \sum_{v \in V} deg(v)$$

\* An undirected graph has an even number of vertices of odd degree.

## Example

How many edges are there in a graph with 10 vertices each of degree six?



#### Path and Circuit

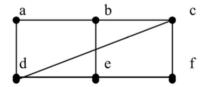
- \* A path with length n from vertex u to vertex v in undirected graph  $G = \langle V, E \rangle$  is sequence  $x_0, x_1, \ldots, x_{n-1}, x_n$  in which n is a positive integer,  $x_0 = u, x_n = v, (x_i, x_{i+1}) \in E$ ,  $i = 0, 1, 2, \ldots, n-1$ .
- \* The above path can be represented as a sequence of edges  $(x_0, x_1)(x_1, x_2), \ldots, (x_{n-1}, x_n)$ .
- $\star$  u is the starting point and v is the ending point of the path
- \* A circuit is a path ending at the starting point (u = v)
- \* A path or a circuit is said to be *simple* if there is no repetition of edges
- \* A cycle is a simple circuit with no repeated vertices other than the first and last ones



#### Path and Circuit

## Example

- $\star$  a, d, c, f, e is a simple path with length 4
- $\star$  d, e, c, b is not a path because (e, c) is not an edge
- $\star$  b, c, f, e, b is a circuit with length 4
- \* Path with length 5 : a, b, e, d, a, b is not simple because (a, b) appears twice





## **Connected Graph**

#### **Definition**

- \* An undirected graph is said to be *connected* if there is a path between every pair of vertices
- \* If is not connected, G consists of several connected subgraphs (two subgraphs do not share any vertex)
  - Each such subgraph is called a *connected component* of *G*.
  - An undirected graph is connected if and only if it has only one connected component

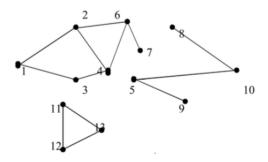
#### **Theorem**

In an undirected graph, if there exist a vertex  $v \in V$  such that there is a path from v to all the other vertices of V, the graph is connected.



## Example

How many connected components are there in graph G?





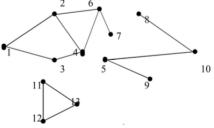
# Bridge and Cut Vertex

#### **Definition**

In an undirected graph, a *bridge* is an edge of the graph whose deletion increases its number of connected components. A *cut vertex* is a vertex whose deletion (with its boundary edges) increases its number of connected components.

#### Example

Find the bridges and cut vertices in the graph below



Discrete Mathematics 2 - Nguyen Kieu Linh ----



#### Contents

- 1 Graph Definitions
  - Graphs
  - Undirected graph
  - Directed graph
- 2 Basic terminologies in undirected graphs
- 3 Basic terminologies in directed graphs
- 4 Some special types of graphs



## In-degree and Out-degree

#### Definition

When (u, v) is an edge of the graph G with directed edges, u is said to be adjacent to v and v is said to be adjacent from u. The vertex u is called the *initial vertex of* (u, v), and v is called the terminal or end vertex of (u, v).

#### Definition

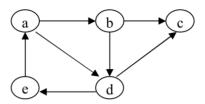
In a graph with directed edges the *in-degree* of a vertex v, denoted by  $\deg^- v$ , is the number of edges with v as their terminal vertex. The *out-degree* of v, denoted by  $\deg^+(v)$ , is the number of edges with v as their initial vertex.



## In-degree and Out-degree

## Example

Find the in-degree and out-degree of each vertex in the following graph with directed edges.





## In-degree and Out-degree

#### **Theorem**

For any directed graph  $G = \langle V, E \rangle$ , we have

$$\sum_{v \in V} \mathsf{deg}^+(v) = \sum_{v \in V} \mathsf{deg}^-(v) = |E|.$$

#### **Notation**

- \* Many properties of directed graphs do not depend on directions. In some cases, we can ignore the directions on directed edges.
- \* The undirected graph receiving by removing directions on directed edges is called the corresponding undirected graph.



#### Path and Circuit

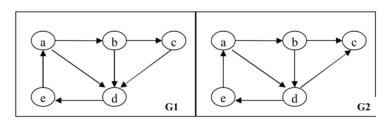
- \* A path with length n from vertex u to vertex v in directed graph  $G = \langle V, E \rangle$ . is sequence  $x_0, x_1, \ldots, x_{n-1}, x_n$  in which n is a positive integer,  $x_0 = u, x_n = v, (x_i, x_{i+1}) \in E$ ,  $i = 0, 1, 2, \ldots, n-1$ .
- \* The above path can be represented as a sequence of edges  $(x_0, x_1)(x_1, x_2), \dots, (x_{n-1}, x_n)$ .
- $\star$  u is the starting point and v is the ending point of the path
- $\star$  A *circuit* is a path ending at the starting point (u = v)
- \* A path or a circuit is said to be *simple* if there is no repetition of directed edges.



## Strongly Connected Graph, Weakly Connected Graph

#### Definition

- \* Directed graph  $G = \langle V, E \rangle$  is said to be *strongly connected* if there is a path between every pair of vertices.
- \* Directed graph  $G = \langle V, E \rangle$  is said to be weakly connected if its corresponding undirected graph is connected.





#### Orientation

#### **Definition**

An *orientation* of an undirected graph is an assignment of a direction to each edge, turning the initial graph into a directed graph. A *strong orientation* is an orientation that results in a strongly connected graph.

#### **Theorem**

For any undirected graph  $G = \langle V, E \rangle$ , there exists a strong orientation on G if and only if all its edges are not bridge.



#### Contents

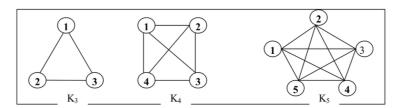
- 1 Graph Definitions
  - Graphs
  - Undirected graph
  - Directed graph
- 2 Basic terminologies in undirected graphs
- 3 Basic terminologies in directed graphs
- 4 Some special types of graphs



## Complete Graph

#### Definition

- \* Complete graph n vertices, denoted by  $K_n$ , is a simple undirected graph that exists an edge connecting between two every vertices
- \* Number of edges:  $\frac{n(n-1)}{2}$

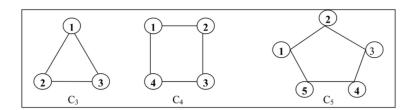


## Cycle Graph

### **Definition**

Cycle Graph  $C_n$ ,  $n \geq 3$ , consists of n vertices

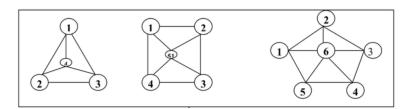
$$\{v_2, v_3\}, ..., \{v_{n-1}, v_n\}, \text{ and } \{v_n, v_1\}.$$



## Wheel Graph

#### Definition

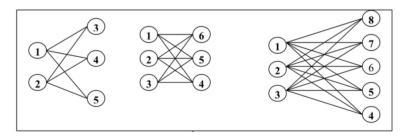
Wheel graph n vertices, denoted by  $W_n$  is a graph formed by connecting a single vertex to all vertices of a cycle graph  $C_{n-1}$ .



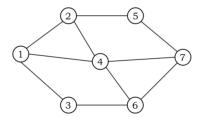
## Bipartite Graph (Bigraph)

#### Definition

Bigraph  $G = \langle V, E \rangle$  is a graph whose vertices can be divided into two disjoint sets X and Y and such that every edge connects a vertex in to one in, i.e. (x, y), in which  $x \in X, y \in Y$ 

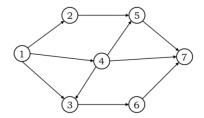


**Exercise 1.** Determine the degree of each vertex in the below undirected graph



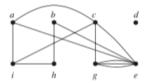


**Exercise 2.** Determine the in-degree and out-degree of each vertex in the below directed graph





# **Exercise 3.** Determine the degree of each vertex in the below undirected graph





**Exercise 4.** Determine the in-degree and out-degree of each vertex in the below directed graph

