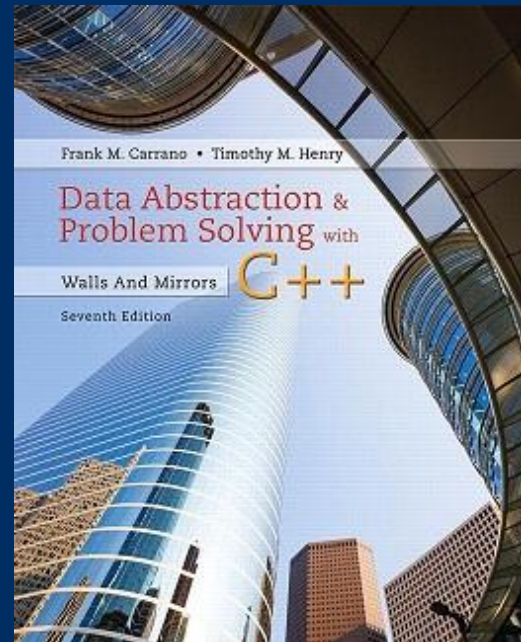


Chapter 17

Heaps

CS 302 - Data Structures

M. Abdullah Canbaz





Reminders

- Assignment 5 is available
 - Due April 11th at 2pm
- TA
 - Athanasia Katsila,
Email: akatsila [at] nevada {dot} unr {dot} edu,
Office Hours: Tuesday, 10:30 am - 12:30 pm at SEM 211
- Unfortunately No!
 - Please take quizzes on time!
- No 70% rule!



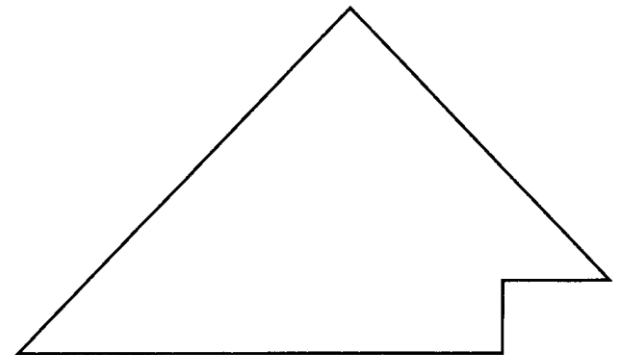
The ADT Heap

- A heap is a complete binary tree that either is
 - Empty or ...
 - Whose root contains a value \geq each of its children and has heaps as its subtrees
- It is a special binary tree ... different in that
 - It is ordered in a weaker sense
 - it will always be a complete binary tree

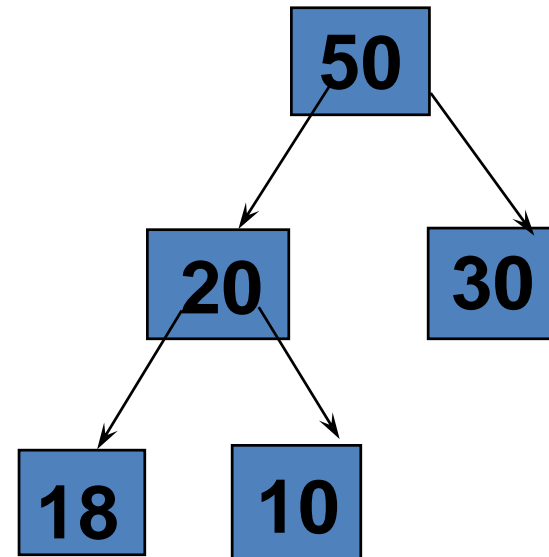
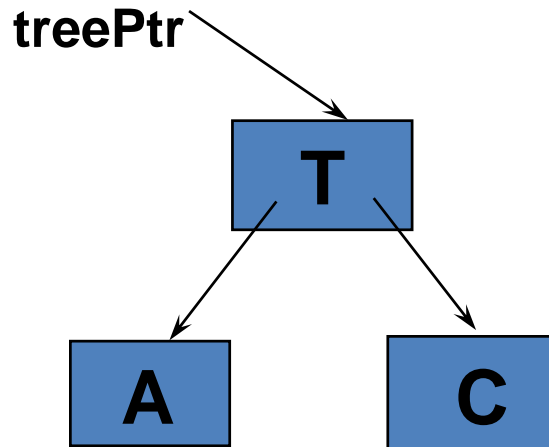
What is a Heap?

A heap is a binary tree that satisfies these special **SHAPE** and **ORDER** properties:

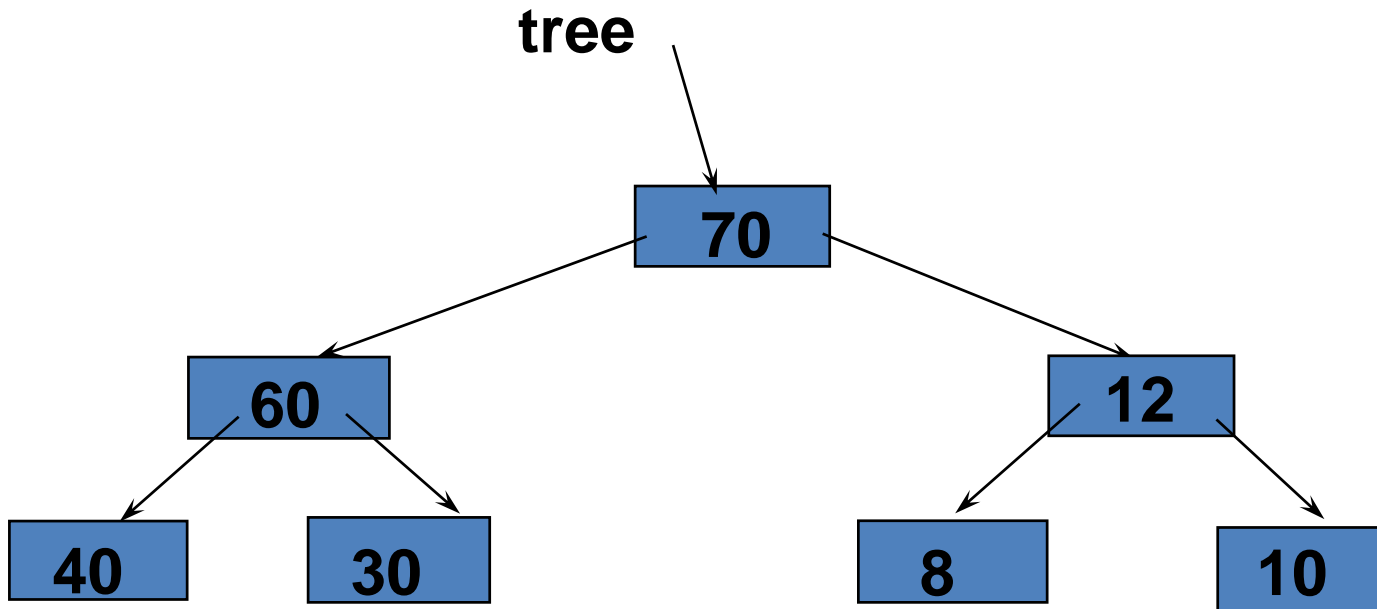
- Its shape must be a complete binary tree.
- For each node in the heap, the value stored in that node is greater than or equal to the value in each of its children.



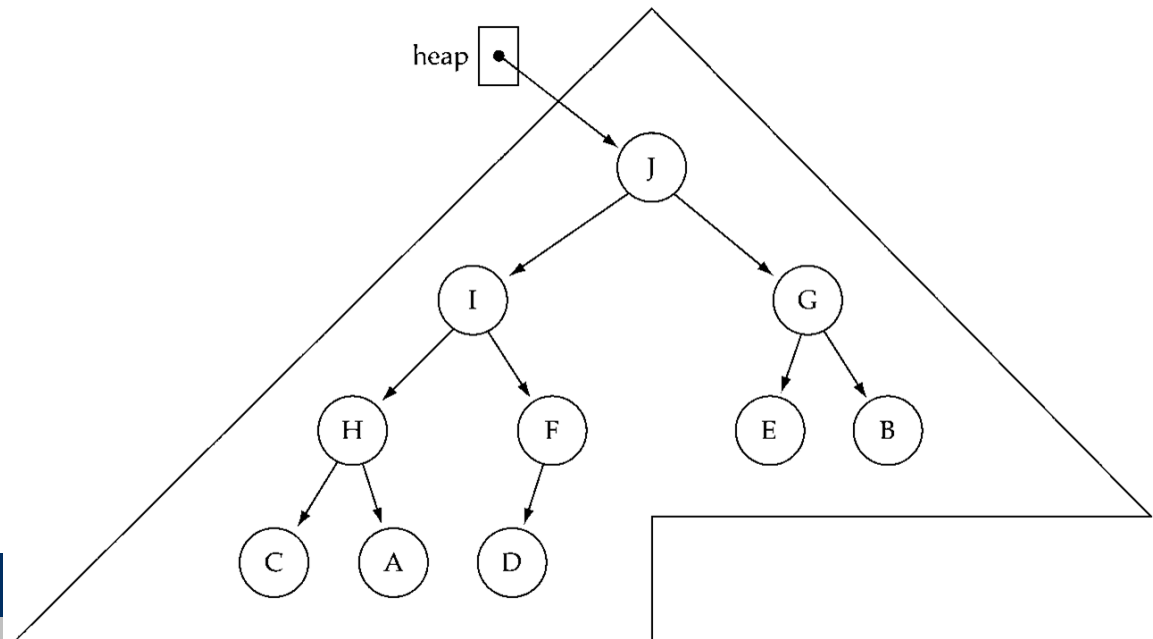
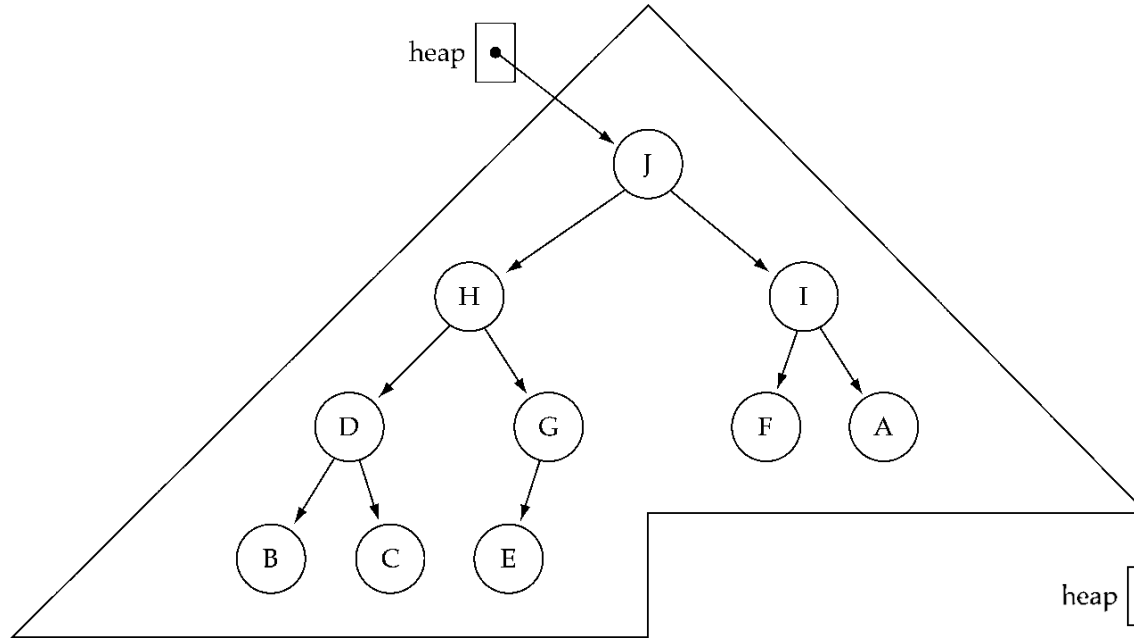
Are these Both Heaps?



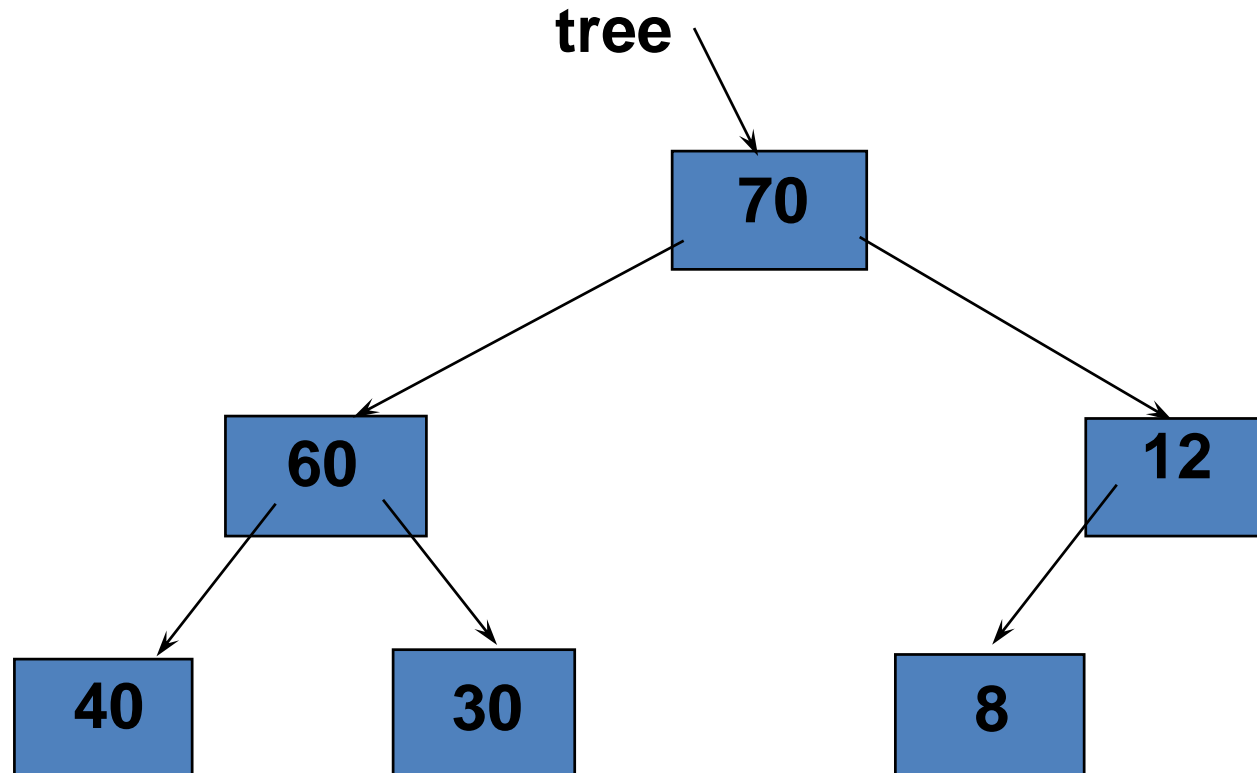
Is this a Heap?

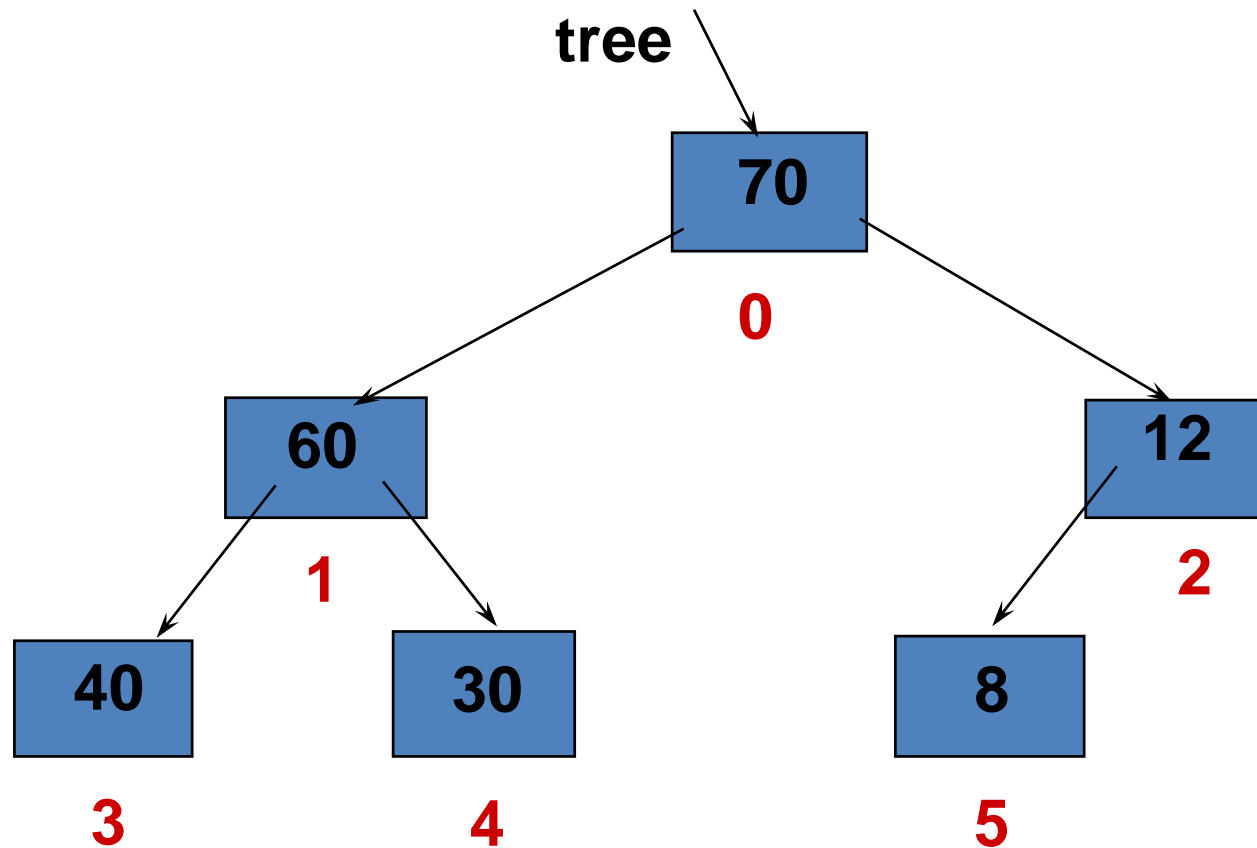


Unique?



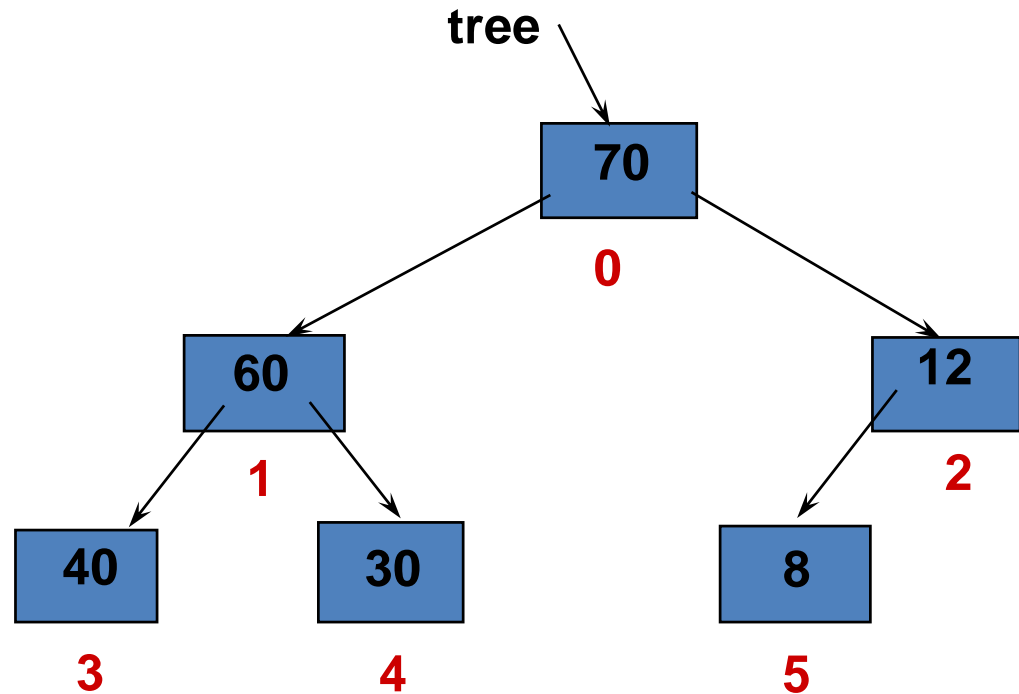
Where is the Largest Element in a Heap Always Found?



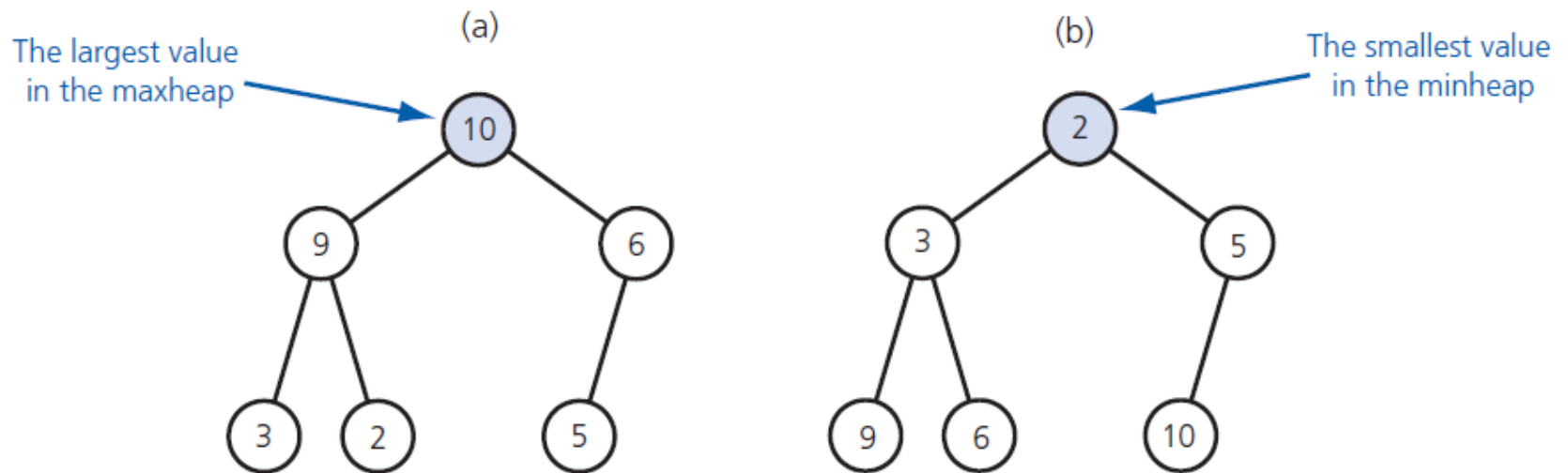


tree.nodes

[0]	70
[1]	60
[2]	12
[3]	40
[4]	30
[5]	8
[6]	



- (a) A maxheap and (b) a minheap

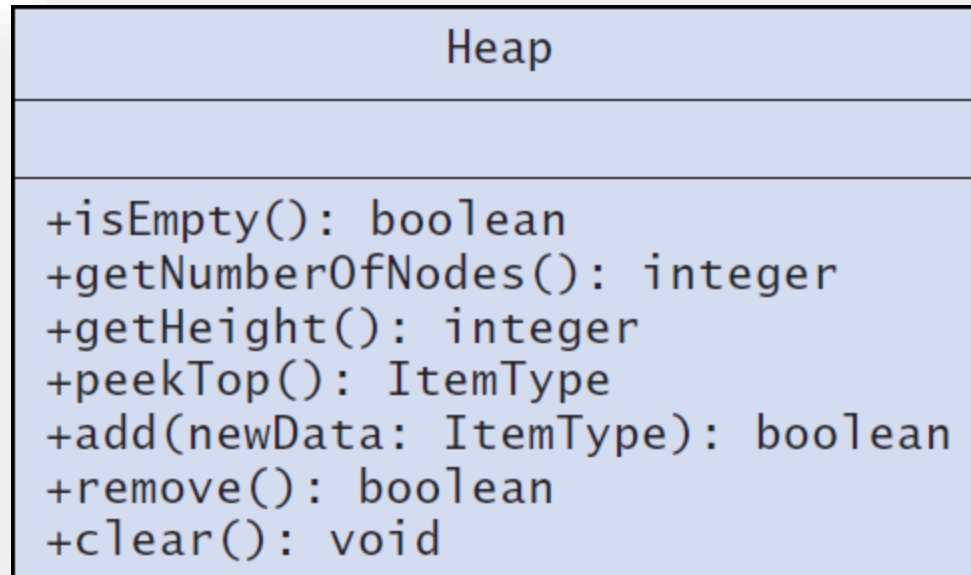


- For any node `tree.nodes[index]`
 - its left child is in
 - `tree.nodes[index*2 + 1]`
 - right child is in
 - `tree.nodes[index*2 + 2]`
 - its parent is in
 - `tree.nodes[(index - 1)/2]`

Leaf nodes:

`tree.nodes[numElements/2]` to
`tree.nodes[numElements - 1]`

- UML diagram for the class **Heap**



- An interface for the ADT **heap**

```
1  /** Interface for the ADT heap.
2   @file HeapInterface.h */
3
4  #ifndef HEAP_INTERFACE_
5  #define HEAP_INTERFACE_
6
7  template<class ItemType>
8  class HeapInterface
9  {
10 public:
11     /** Sees whether this heap is empty.
12     @return True if the heap is empty, or false if not. */
13     virtual bool isEmpty() const = 0;
14
15     /** Gets the number of nodes in this heap.
16     @return The number of nodes in the heap. */
17     virtual int getNumberOfNodes() const = 0;
```

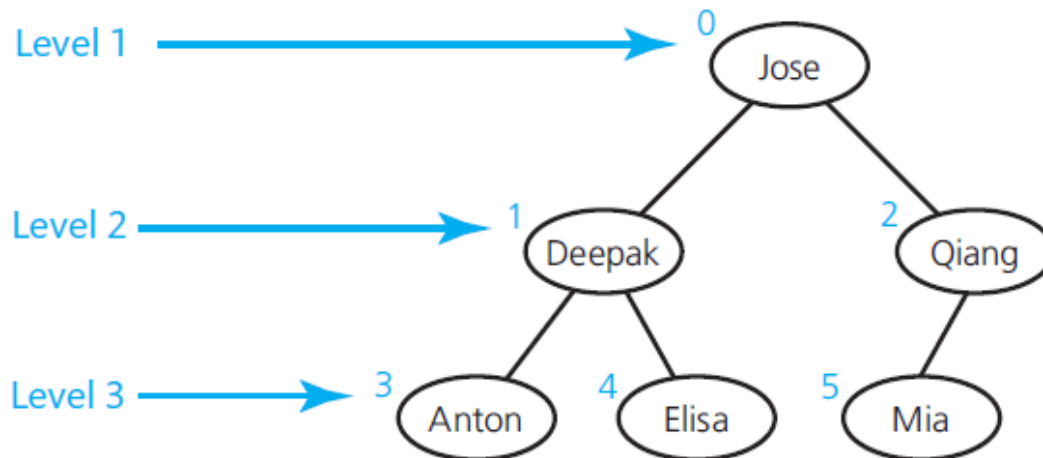
- An interface for the ADT **heap**

```
18
19     /** Gets the height of this heap.
20         @return The height of the heap. */
21     virtual int getHeight() const = 0;
22
23     /** Gets the data that is in the root (top) of this heap.
24         For a maxheap, the data is the largest value in the heap;
25         for a minheap, the data is the smallest value in the heap.
26         @pre The heap is not empty.
27         @post The root's data has been returned, and the heap is unchanged.
28         @return The data in the root of the heap. */
29     virtual ItemType peekTop() const = 0;
30
31     /** Adds a new data item to this heap
```

- An interface for the ADT **heap**

```
30
31  /** Adds a new data item to this heap.
32   @param newData  The data to be added.
33   @post  The heap has a new node that contains newData.
34   @return True if the addition is successful, or false if not. */
35  virtual bool add(const ItemType& newData) = 0;
36
37  /** Removes the data that is in the root (top) of this heap.
38   @return True if the removal is successful, or false if not. */
39  virtual bool remove() = 0;
40
41  /** Removes all data from this heap. */
42  virtual void clear() = 0;
43
44  /** Destroys this heap and frees its assigned memory. */
45  virtual ~HeapInterface() { }
46 }; // end HeapInterface
47 #endif
```


(a) Level-by-level numbering



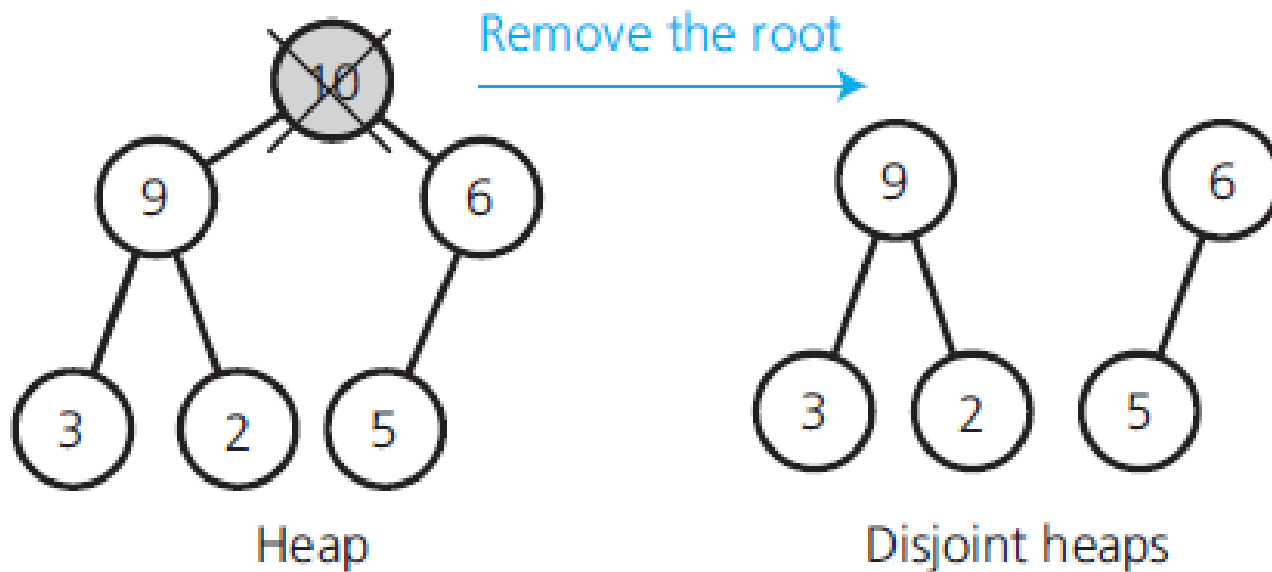
(b) Array-based implementation
items

0	Jose
1	Deepak
2	Qiang
3	Anton
4	Elisa
5	Mia
6	
7	

- A complete binary tree and its array-based implementation

- Assume following private data members
 - **items**: an array of heap items
 - **itemCount**: an integer equal to the number of items in the heap
 - **maxItems**: an integer equal to the maximum capacity of the heap

- Disjoint heaps after removing the heap's root



- Recursive algorithm to transform semiheap to heap.

// Converts a semiheap rooted at index nodeIndex into a heap.

```
heapRebuild(nodeIndex: integer, items: ArrayType, itemCount: integer): void
{
```

*// Recursively trickle the item at index nodeIndex down to its proper position by
// swapping it with its larger child, if the child is larger than the item.*

// If the item is at a leaf, nothing needs to be done.

if (the root is not a leaf)

```
{
```

// The root must have a left child; find larger child

```
leftChildIndex = 2 * rootIndex + 1
```

```
rightChildIndex = leftChildIndex + 1
```

```
largerChildIndex = rightChildIndex // Assume right child exists and is the larger
```

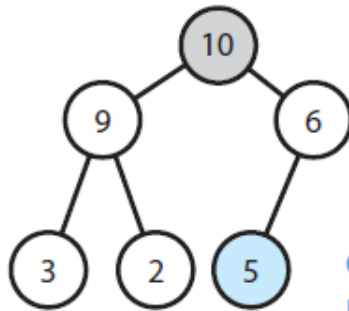
// Check whether right child is larger than left child, if so, swap them

- Recursive algorithm to transform semiheap to heap.

```
// Check whether right child exists; if so, is left child larger?  
// If no right child, left one is larger  
if ((largerChildIndex >= itemCount) ||  
    (items[leftChildIndex] > items[rightChildIndex]))  
    largerChildIndex = leftChildIndex; // Assumption was wrong  
if (items[nodeIndex] < items[largerChildIndex])  
{  
    Swap items[nodeIndex] and items[largerChildIndex]  
  
    // Transform the semiheap rooted at largerChildIndex into a heap  
    heapRebuild(largerChildIndex, items, itemCount)  
}  
// Else root is a leaf, so you are done  
}
```

Algorithms for Array-Based Heap Operations

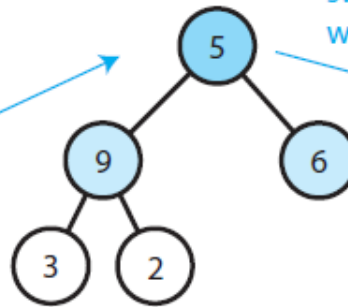
(a) Heap



Copy entry in last node to root

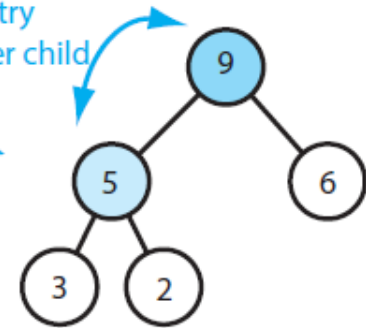
(b) Semiheap

Entry in root is smaller than entries in its children



Trickle down by swapping root entry with entry in larger child

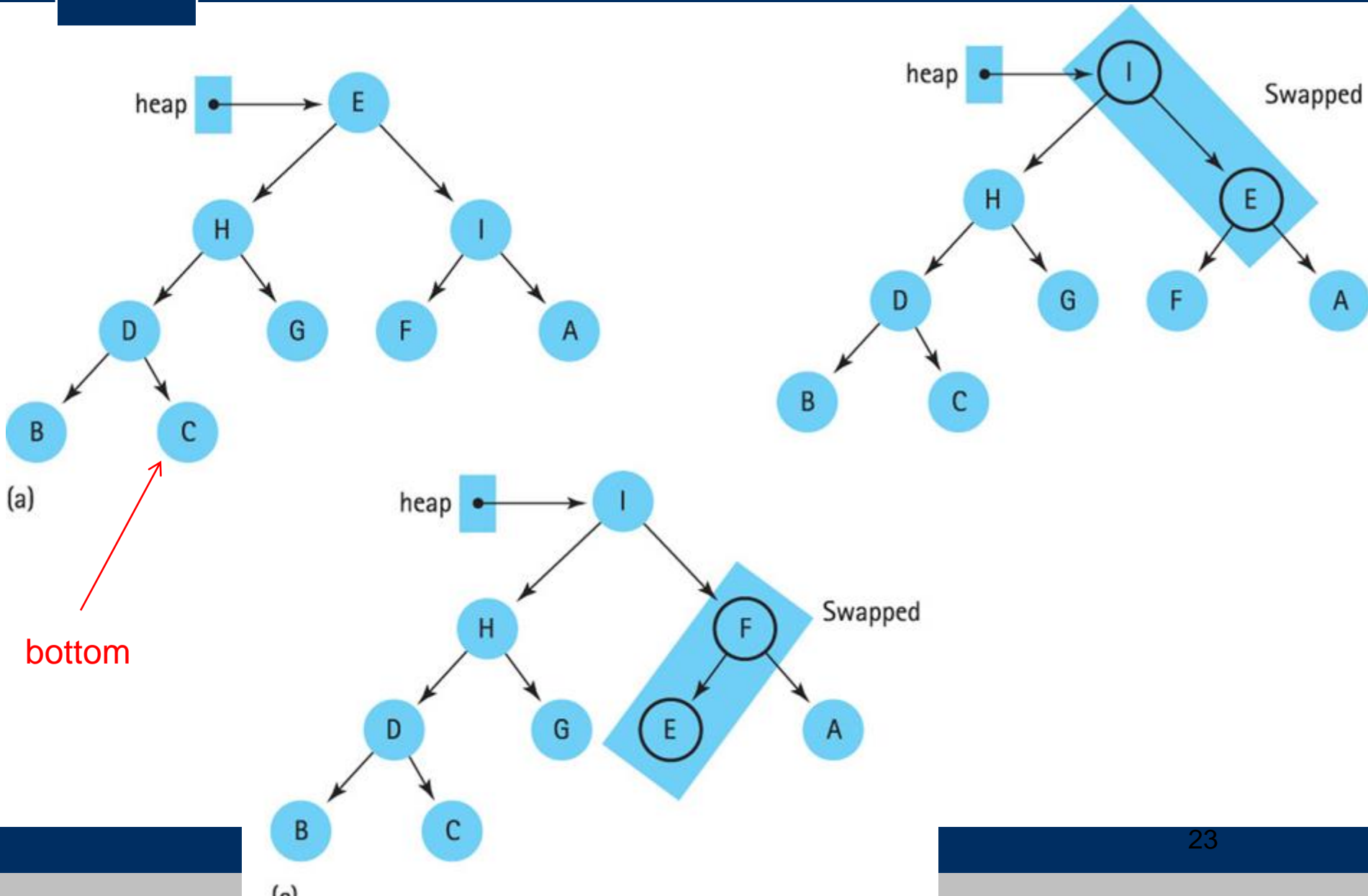
(c) Restored heap



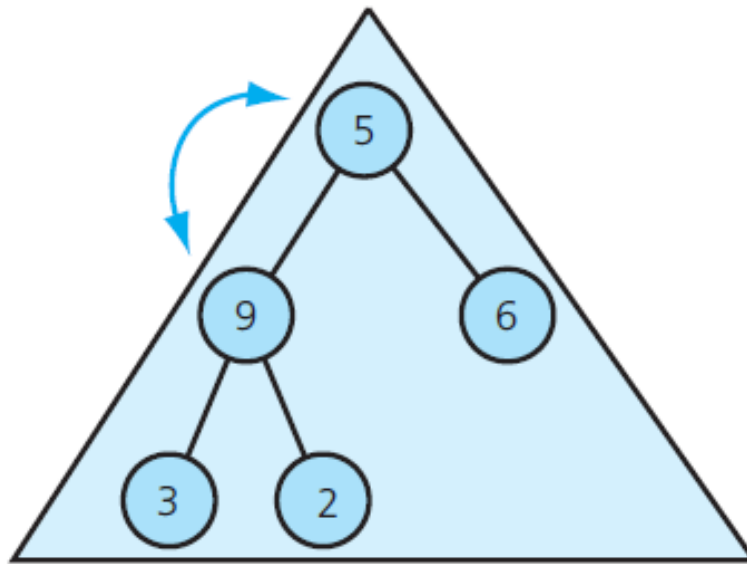
10	9	6	3	2	5
0	1	2	3	4	5

5	9	6	3	2	
0	1	2	3	4	5

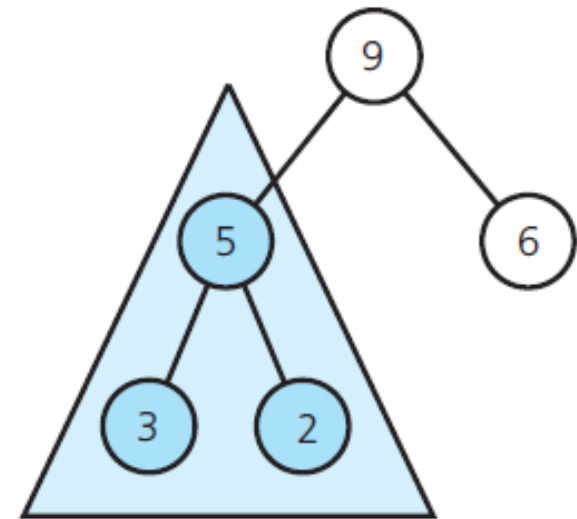
9	5	6	3	2	
0	1	2	3	4	5



- Recursive calls to `heapRebuild`

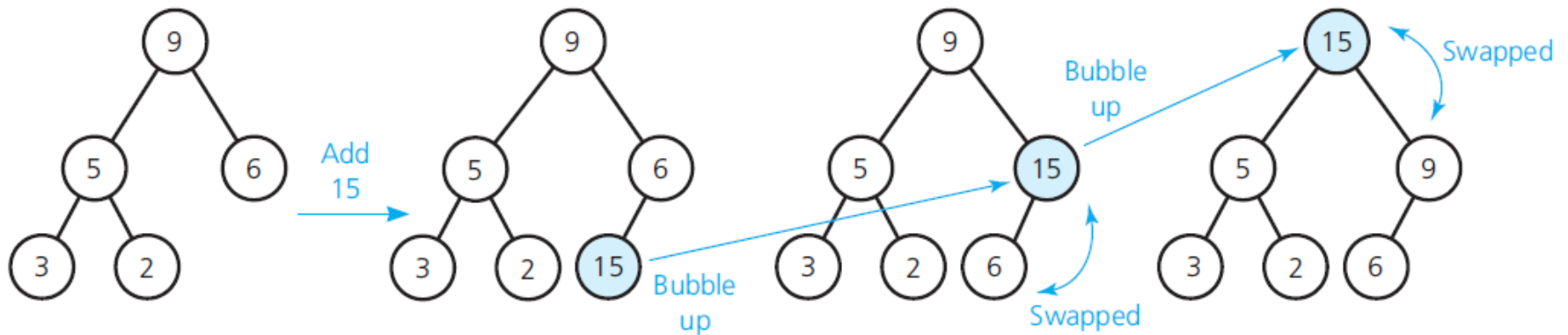


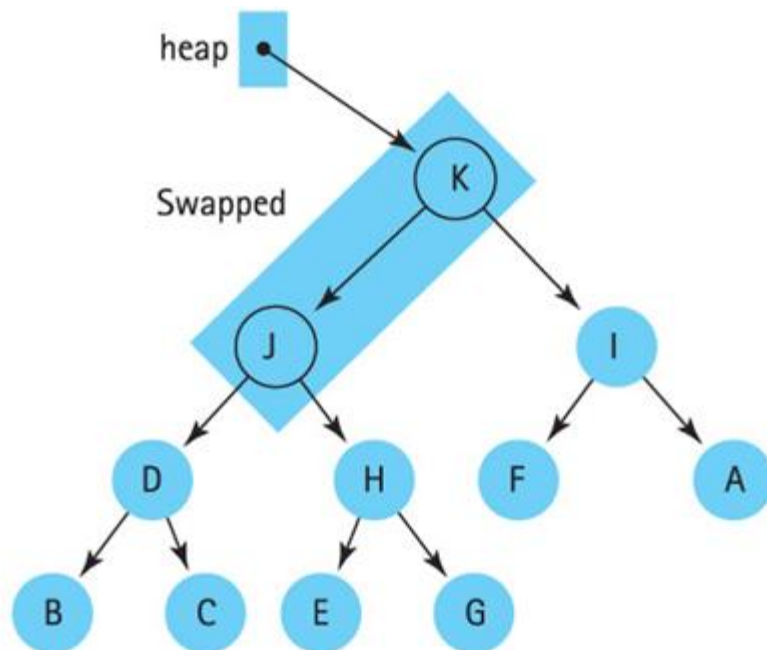
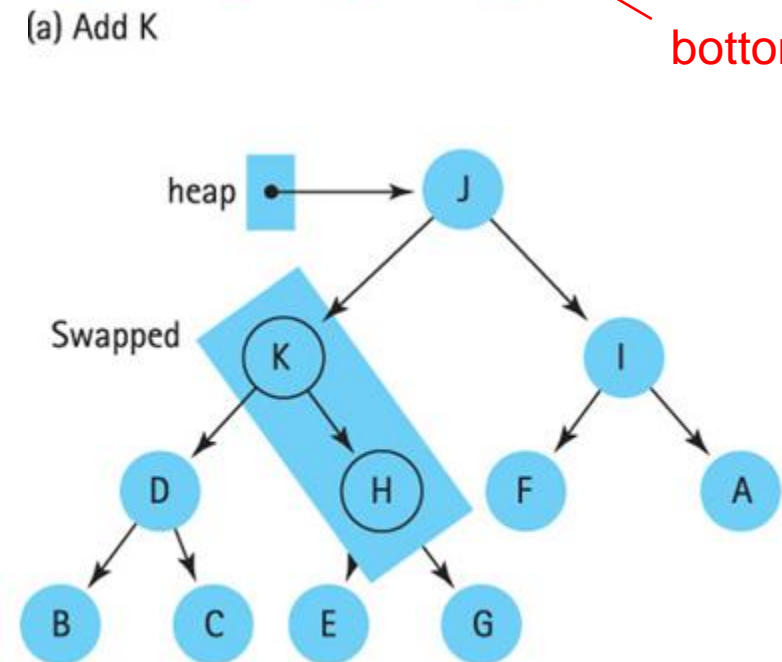
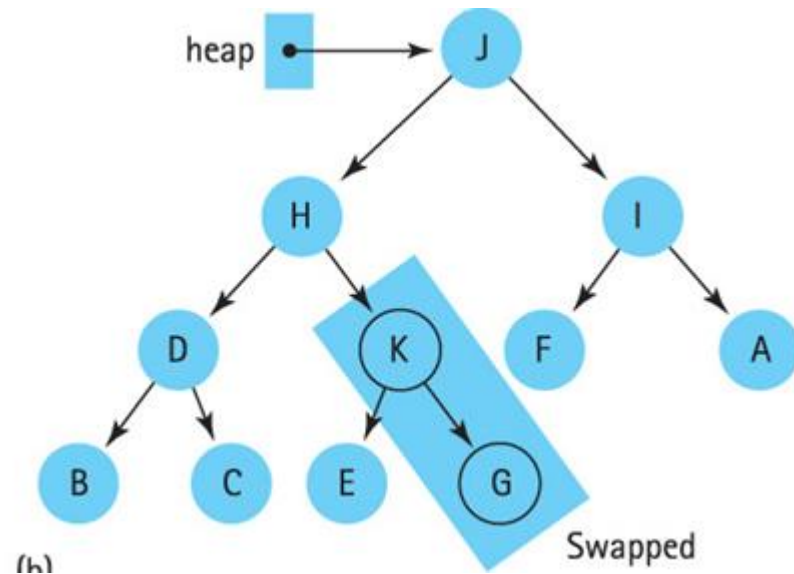
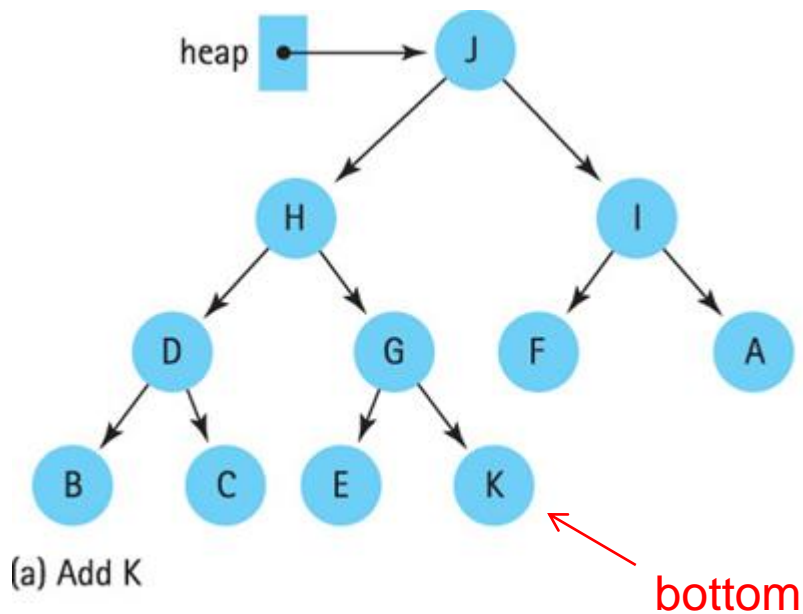
First semiheap passed
to `heapRebuild`



Second semiheap passed
to `heapRebuild`

- Adding 15 to a heap





```
add(newData: itemType): boolean
{
    // Place newData at the bottom of the tree
    items[itemCount] = newData

    // Make new item bubble up to the appropriate spot in the tree
    newDataIndex = itemCount
    inPlace = false
    while ((newDataIndex >= 0) and !inPlace)
    {
        parentIndex = (newDataIndex - 1) / 2
        if (items[newDataIndex] <= items[parentIndex])
            inPlace = true
        else
        {
            Swap items[newDataIndex] and items[parentIndex]
            newDataIndex = parentIndex
        }
    }
    itemCount++
    return inPlace
}
```

- Pseudocode for **add**

- The header file for the class `ArrayMaxHeap`

```
1  /** Array-based implementation of the ADT heap.
2   * @file ArrayMaxHeap.h */
3  #ifndef ARRAY_MAX_HEAP_
4  #define ARRAY_MAX_HEAP_
5
6  #include "HeapInterface.h"
7  #include "PrecondViolatedExcept.h"
8
9  template<class ItemType>
10 class ArrayMaxHeap : public HeapInterface<ItemType>
11 {
12 private:
13     static const int ROOT_INDEX = 0;           // Helps with readability
14     static const int DEFAULT_CAPACITY = 21;    // Small capacity for testing
15     std::unique_ptr<ItemType[]> items;         // Array of heap items
16     int itemCount;                             // Current count of heap items
17     int maxItems;                             // Maximum capacity of the heap
18 }
```

- The header file for the class `ArrayMaxHeap`

```
18
19 // -----
20 // Most of the private utility methods use an array index as a parameter
21 // and in calculations. This should be safe, even though the array is an
22 // implementation detail, since the methods are private.
23 // -----
24
25 // Returns the array index of the left child (if it exists).
26 int getLeftChildIndex(const int nodeIndex) const;
27
28 // Returns the array index of the right child (if it exists).
29 int getRightChildIndex(int nodeIndex) const;
30
31 // Returns the array index of the parent node.
32 int getParentIndex(int nodeIndex) const;
33
34 // Tests whether this node is a leaf.
35 bool isLeaf(int nodeIndex) const;
36
```

- The header file for the class `ArrayMaxHeap`

```
37 // Converts a semiheap to a heap.
38 void heapRebuild(int nodeIndex);
39
40 // Creates a heap from an unordered array.
41 void heapCreate();
42
43 public:
44     ArrayMaxHeap();
45     ArrayMaxHeap(const ItemType someArray[], const int arraySize);
46     virtual ~ArrayMaxHeap();
47
48     // HeapInterface Public Methods:
49     bool isEmpty() const;
50     int getNumberOfNodes() const;
51     int getHeight() const;
52     ItemType peekTop() const throw(PrecondViolatedExcept);
53     bool add(const ItemType& newData);
54     bool remove();
55     void clear();
56 }; // end ArrayMaxHeap
57 #include "ArrayMaxHeap.cpp"
58 #endif
```



The Implementation

- Definition of method `getLeftChildIndex`

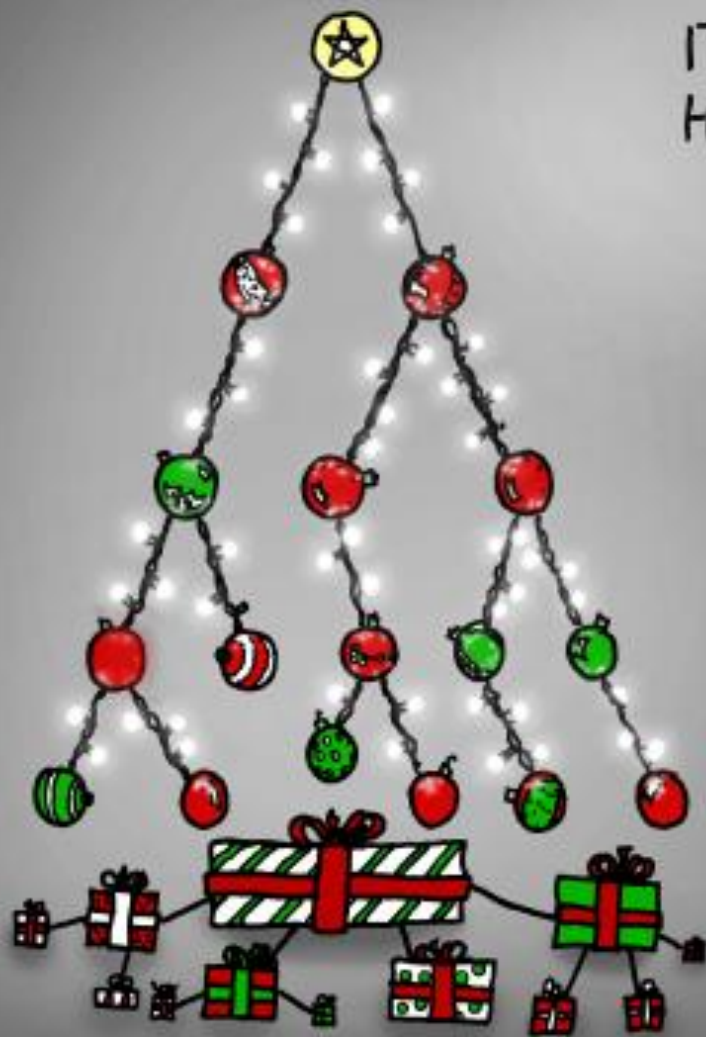
```
template<class ItemType>
int ArrayMaxHeap<ItemType>::getLeftChildIndex(const int nodeIndex) const
{
    return (2 * nodeIndex) + 1;
} // end getLeftChildIndex
```

- Definition of the constructor

```
template<class ItemType>
ArrayMaxHeap<ItemType>::
ArrayMaxHeap(const ItemType someArray[], const int arraySize):
    itemCount(arraySize), maxItems(2 * arraySize)
{
    // Allocate the array
    items = std::make_unique<ItemType[]>(maxItems);

    // Copy given values into the array
    for (int i = 0; i < itemCount; i++)
        items[i] = someArray[i];

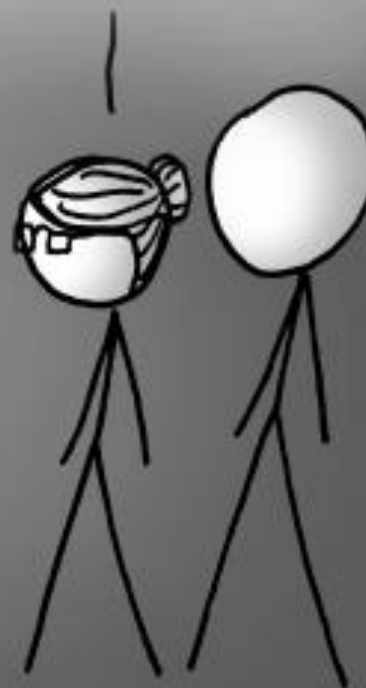
    // Reorganize the array into a heap
    heapCreate();
} // end constructor
```

IT'S A CHRISTMAS TREE WITH A
HEAP OF PRESENTS UNDERNEATH!

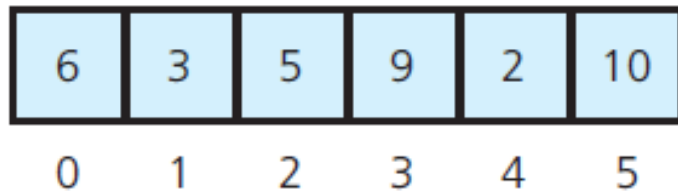


... WE'RE NOT INVITING
YOU HOME NEXT YEAR.

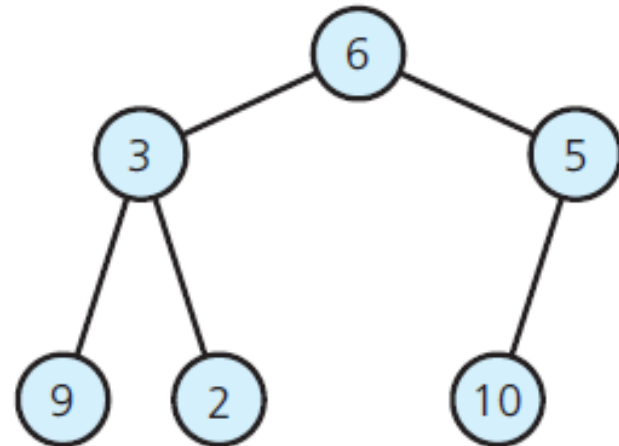


- Array and its corresponding complete binary tree

(a) An array of given values



(b) The complete binary tree represented by the array



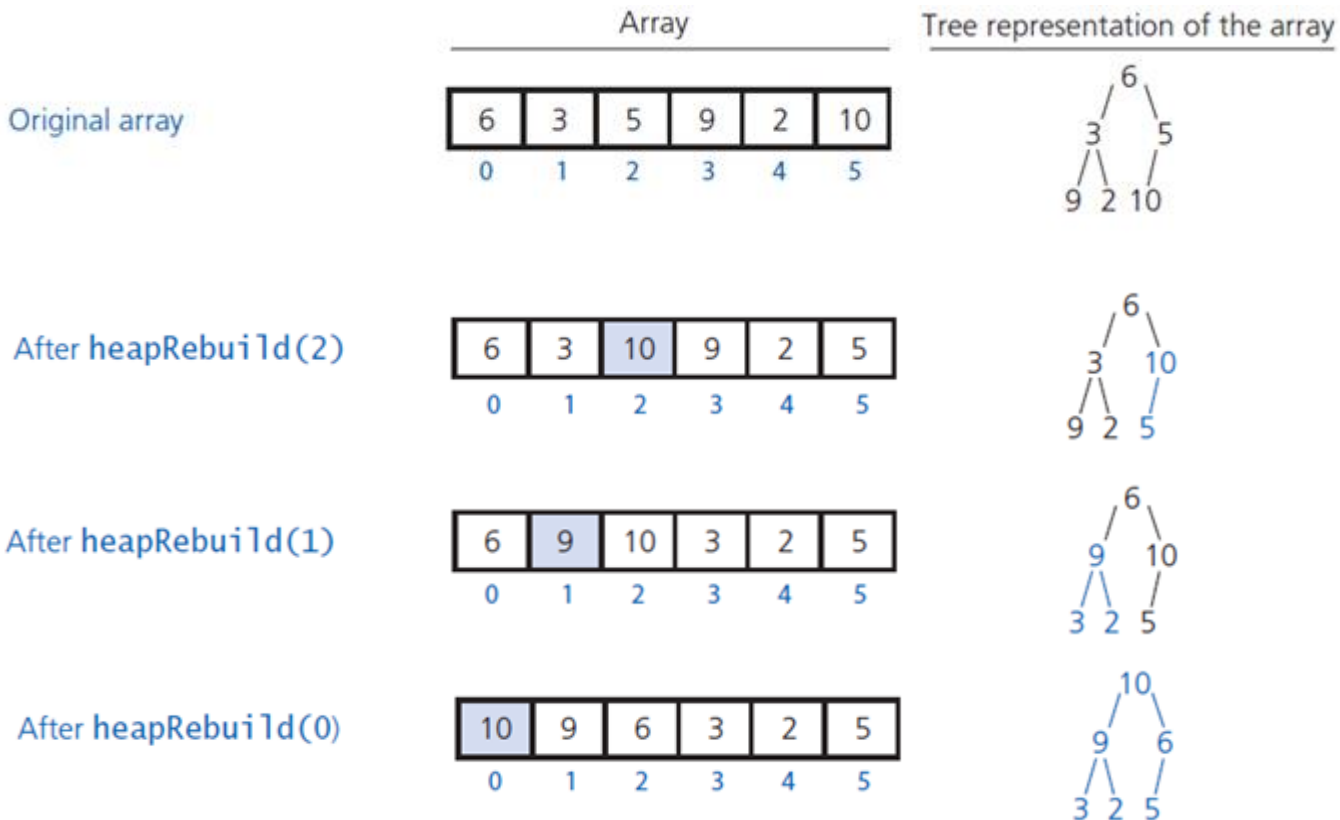


The Implementation

- Building a heap from an array of data

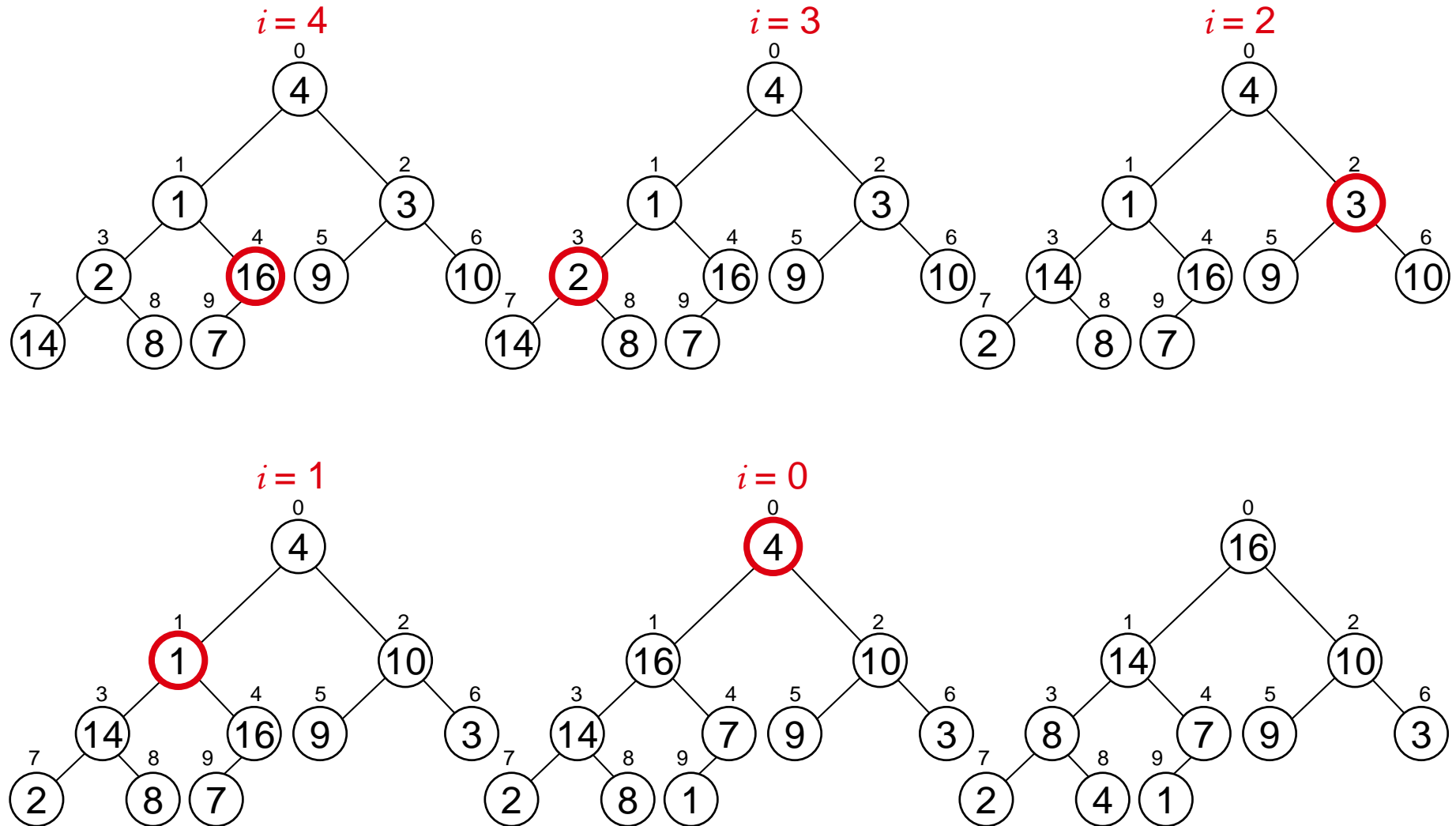
```
for (index = itemCount - 1 down to 0)
{
    // Assertion: The tree rooted at index is a semiheap
    heapRebuild(index)
    // Assertion: The tree rooted at index is a heap
}
```

- Transforming an array into a heap



Building a Heap:

4	1	3	2	16	9	10	14	8	7
---	---	---	---	----	---	----	----	---	---



- C++ method `heapCreate`

```
template<class ItemType>
void ArrayMaxHeap<ItemType>::heapCreate()
{
    for (int index = itemCount / 2; index >= 0; index--)
        heapRebuild(index);
} // end heapCreate
```

- C++ method `peekTop` which tests for an empty heap

```
template<class ItemType>
ItemType ArrayMaxHeap<ItemType>::peekTop() const throw(PrecondViolatedExcept)
{
    if (isEmpty())
        throw PrecondViolatedExcept("Attempted peek into an empty heap.");

    return items[0];
} // end peekTop
```

```
1  /** ADT priority queue: Heap-based implementation.
2   @file HeapPriorityQueue.h */
3  #ifndef HEAP_PRIORITY_QUEUE_
4  #define HEAP_PRIORITY_QUEUE_
5  #include "ArrayMaxHeap.h"
6  #include "PriorityQueueInterface.h"
7
8  template<class ItemType>
9  class HeapPriorityQueue : public PriorityQueueInterface<ItemType>,
10                          private ArrayMaxHeap<ItemType>
11  {
12  public:
13      HeapPriorityQueue();
14      bool isEmpty() const;
15      bool enqueue(const ItemType& newEntry);
16      bool dequeue();
17
18      /** @pre The priority queue is not empty. */
19      ItemType peekFront() const throw(PrecondViolatedExcept);
20  }; // end HeapPriorityQueue
21
22  #include "HeapPriorityQueue.cpp"
23  #endif
```

- A header file for the class **HeapPriorityQueue**


```
1  /** Heap-based implementation of the ADT priority queue.
2   * @file HeapPriorityQueue.cpp */
3
4  #include "HeapPriorityQueue.h"
5
6  template<class ItemType>
7  HeapPriorityQueue<ItemType>::HeapPriorityQueue()
8  {
9      ArrayMaxHeap<ItemType>();
10 } // end constructor
11
12 template<class ItemType>
13 bool HeapPriorityQueue<ItemType>::isEmpty() const
14 {
15     return ArrayMaxHeap<ItemType>::isEmpty();
16 } // end isEmpty
17
18 template<class ItemType>
19 bool HeapPriorityQueue<ItemType>::enqueue(const ItemType& newEntry)
20 {
21     return ArrayMaxHeap<ItemType>::add(newEntry);
22 } // end add
```

- An implementation of the class **HeapPriorityQueue**

```
23
24 template<class ItemType>
25 bool HeapPriorityQueue<ItemType>::dequeue()
26 {
27     return ArrayMaxHeap<ItemType>::remove();
28 } // end remove
29
30 template<class ItemType>
31 ItemType HeapPriorityQueue<ItemType>::peekFront() const throw(PrecondViolatedExcept)
32 {
33     try
34     {
35         return ArrayMaxHeap<ItemType>::peekTop();
36     }
37     catch (PrecondViolatedExcept e)
38     {
39         throw PrecondViolatedExcept("Attempted peek into an empty priority queue.");
40     } // end try/catch
41 } // end peekFront
```

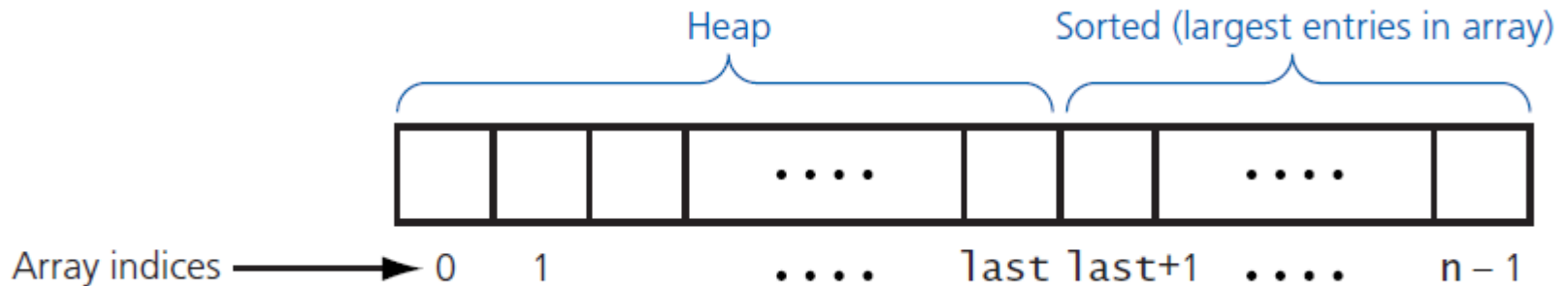
- An implementation of the class **HeapPriorityQueue**



Heap Implementation of the ADT Priority Queue

- Heap versus a binary search tree
 - If you know maximum number of items in the priority queue, heap is the better implementation
- Finite, distinct priority values
 - Many items likely have same priority value
 - Place in same order as encountered

- Heap sort partitions an array into two regions

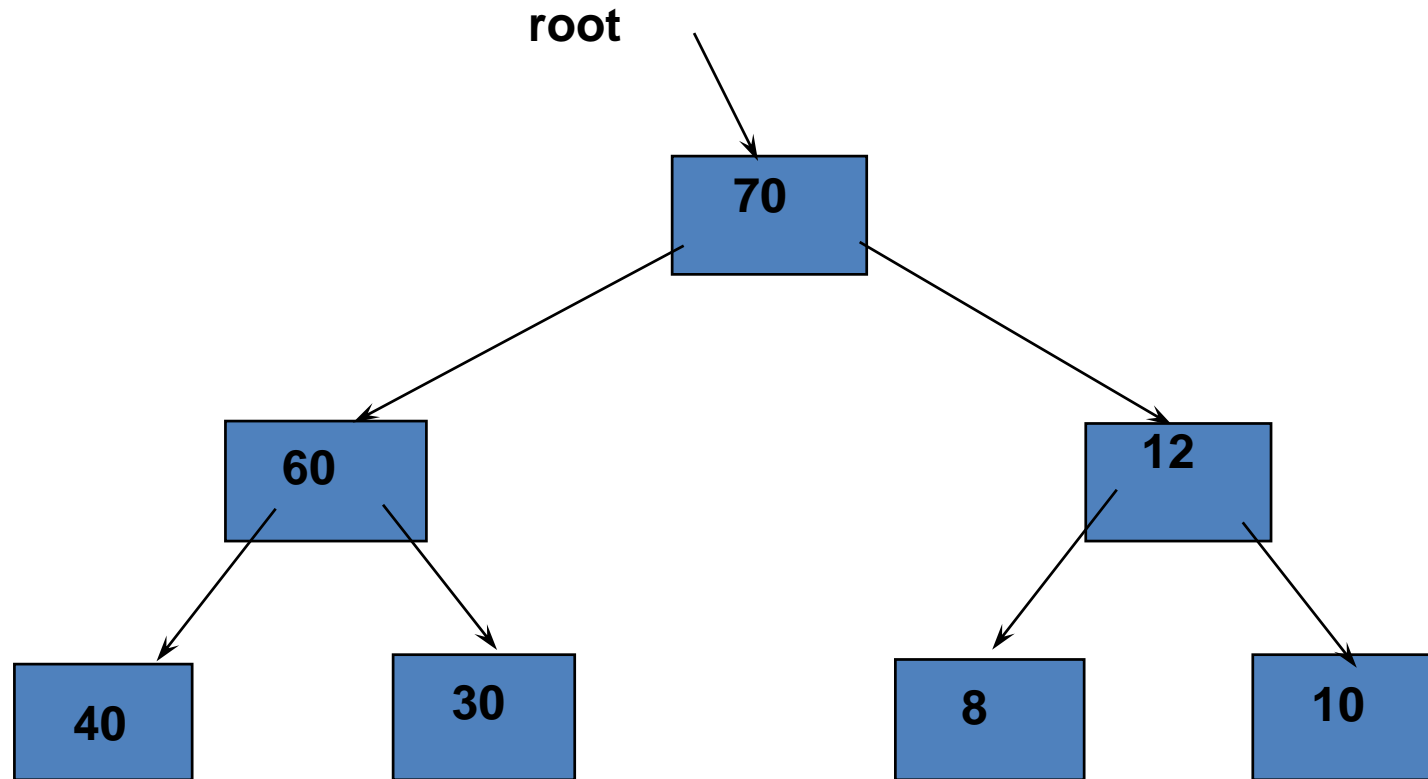


A heap is a binary tree that satisfies these special **SHAPE** and **ORDER** properties:

- Its shape must be a complete binary tree.
- For each node in the heap, the value stored in that node is greater than or equal to the value in each of its children.

N

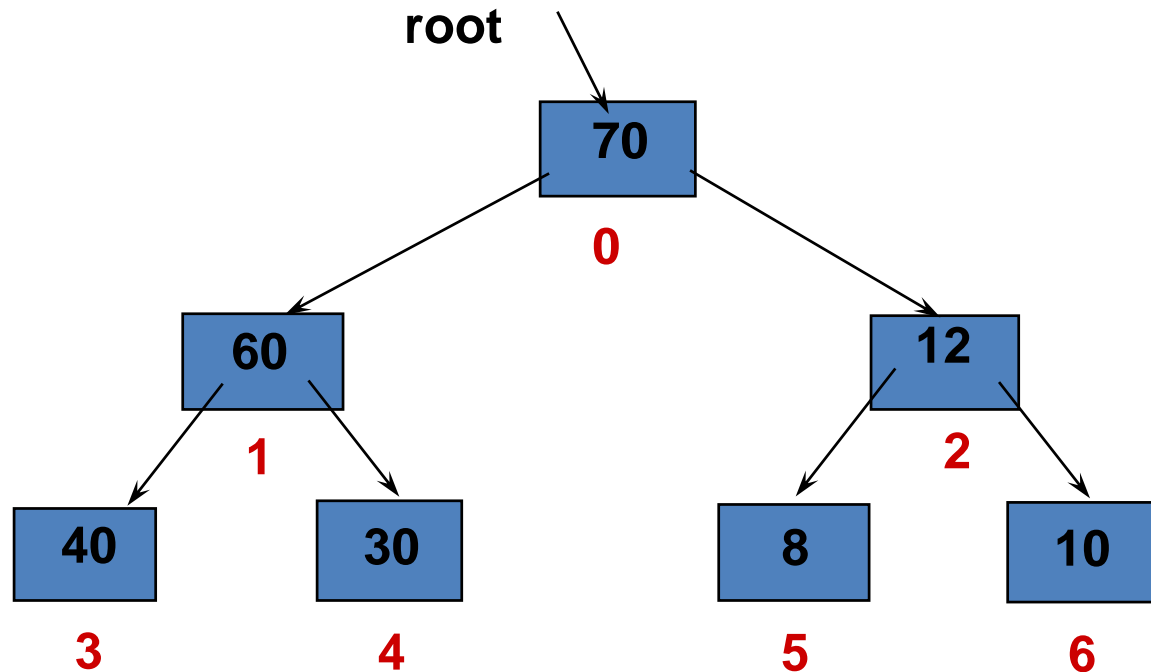
The largest element in a heap is always found in the root node





The heap can be stored

	values
[0]	70
[1]	60
[2]	12
[3]	40
[4]	30
[5]	8
[6]	10





Heap Sort Approach

First, make the unsorted array into a heap by satisfying the order property.

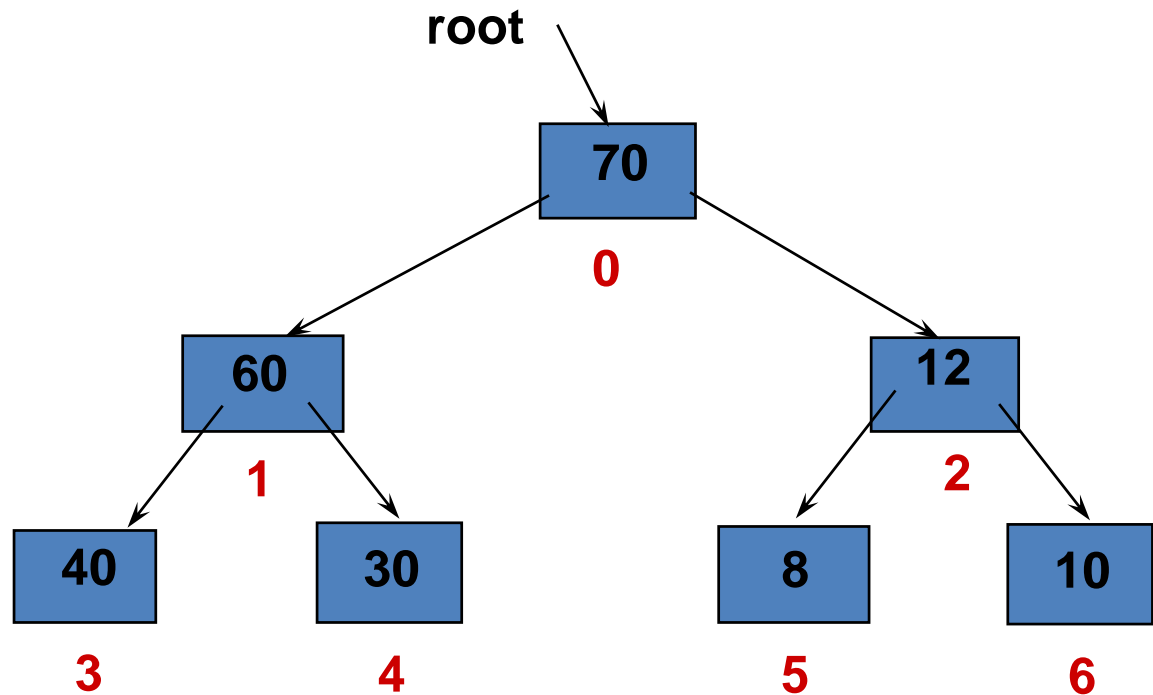
Then repeat the steps below until there are no more unsorted elements.

- Take the root (maximum) element off the heap by swapping it into its correct place in the array at the end of the unsorted elements.
- Reheap the remaining unsorted elements.
 - This puts the next-largest element into the root position.



After creating the original heap

	values
[0]	70
[1]	60
[2]	12
[3]	40
[4]	30
[5]	8
[6]	10

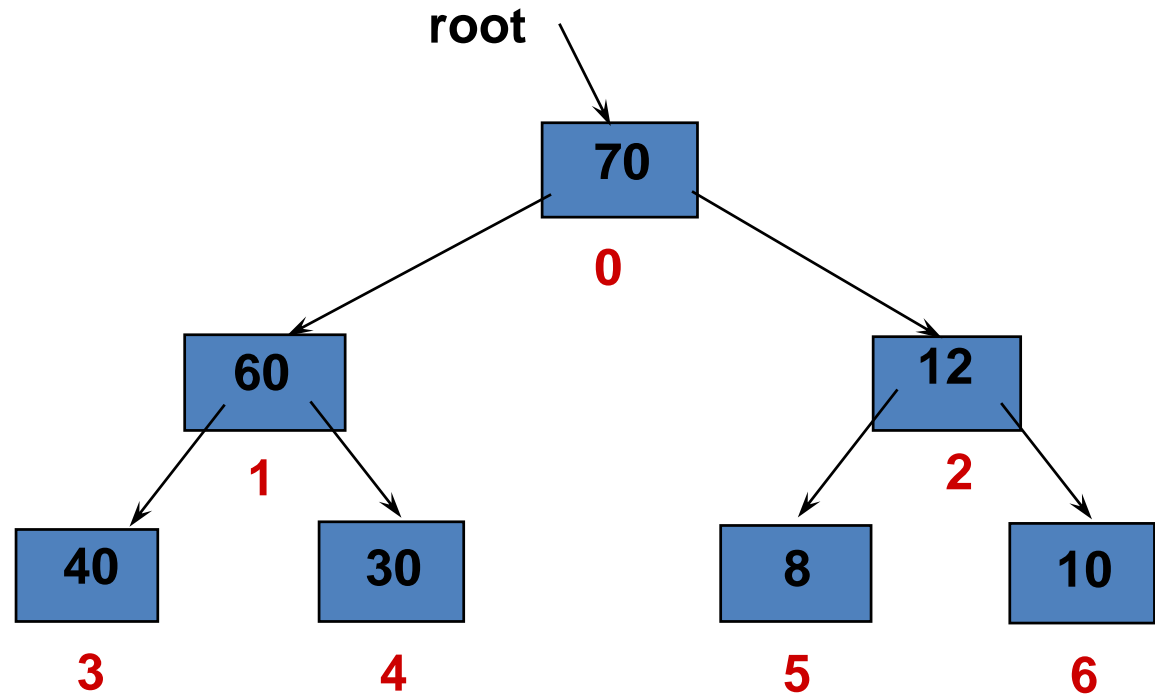




Swap root element into last place in unsorted array

values

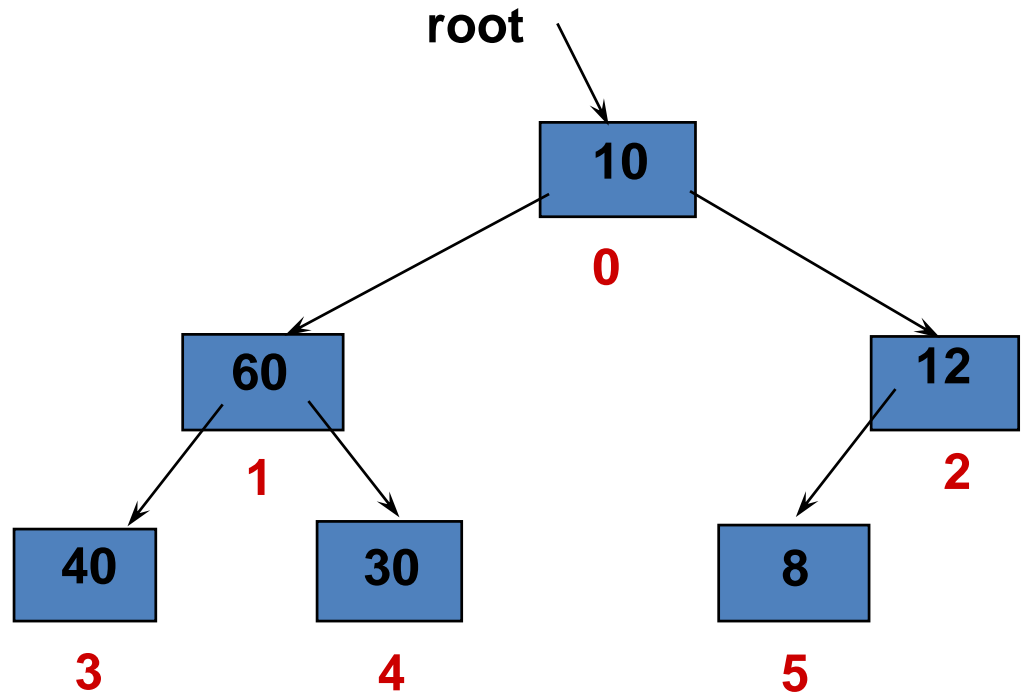
[0]	70
[1]	60
[2]	12
[3]	40
[4]	30
[5]	8
[6]	10





After swapping root element into its place

values	
[0]	10
[1]	60
[2]	12
[3]	40
[4]	30
[5]	8
[6]	70

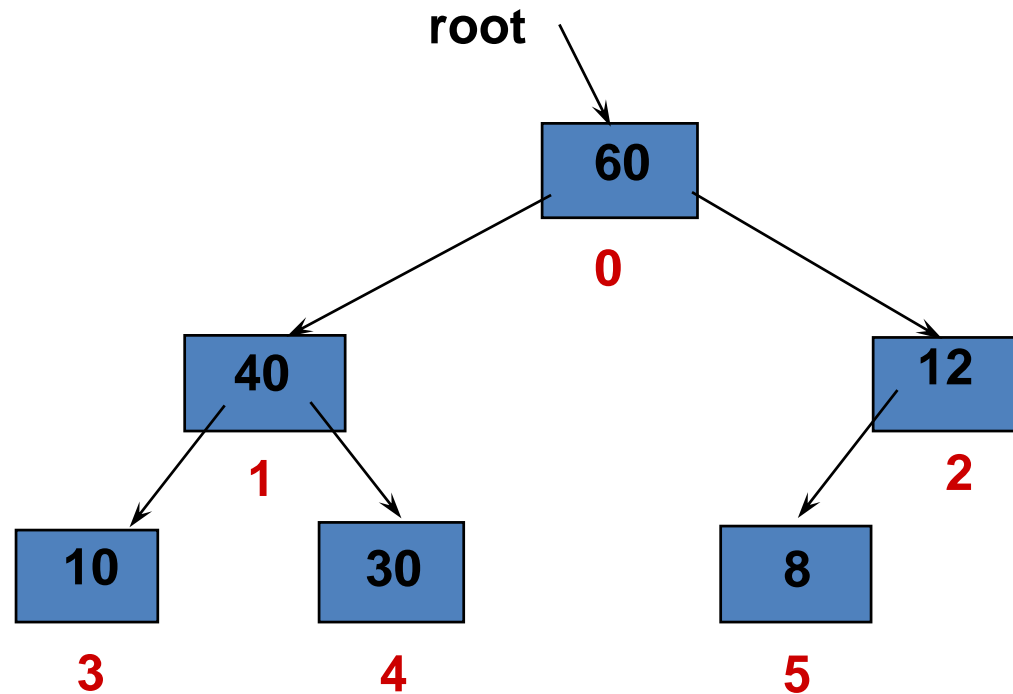


NO NEED TO CONSIDER AGAIN



After reheaping remaining unsorted elements

values	
[0]	60
[1]	40
[2]	12
[3]	10
[4]	30
[5]	8
[6]	70

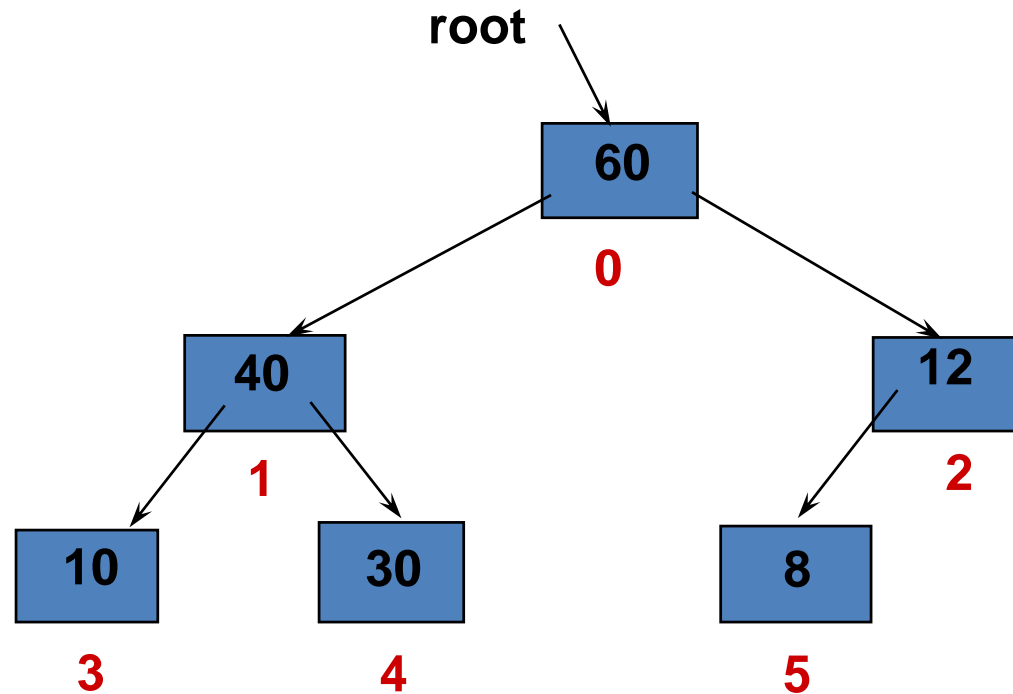




Swap root element into last place in unsorted array

values

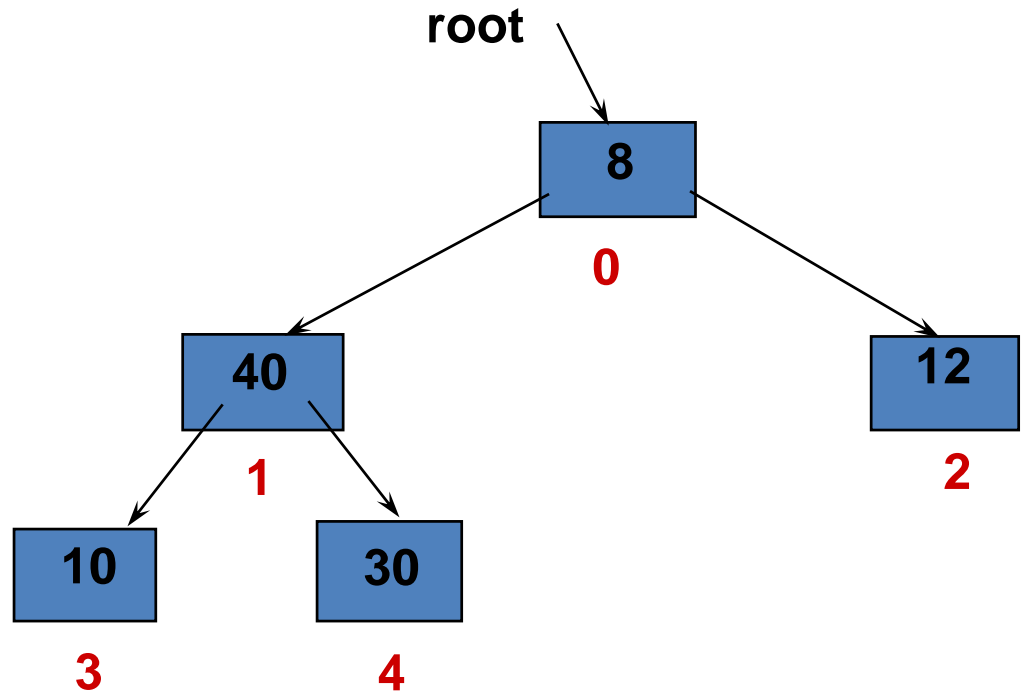
[0]	60
[1]	40
[2]	12
[3]	10
[4]	30
[5]	8
[6]	70





After swapping root element into its place

values	
[0]	8
[1]	40
[2]	12
[3]	10
[4]	30
[5]	60
[6]	70

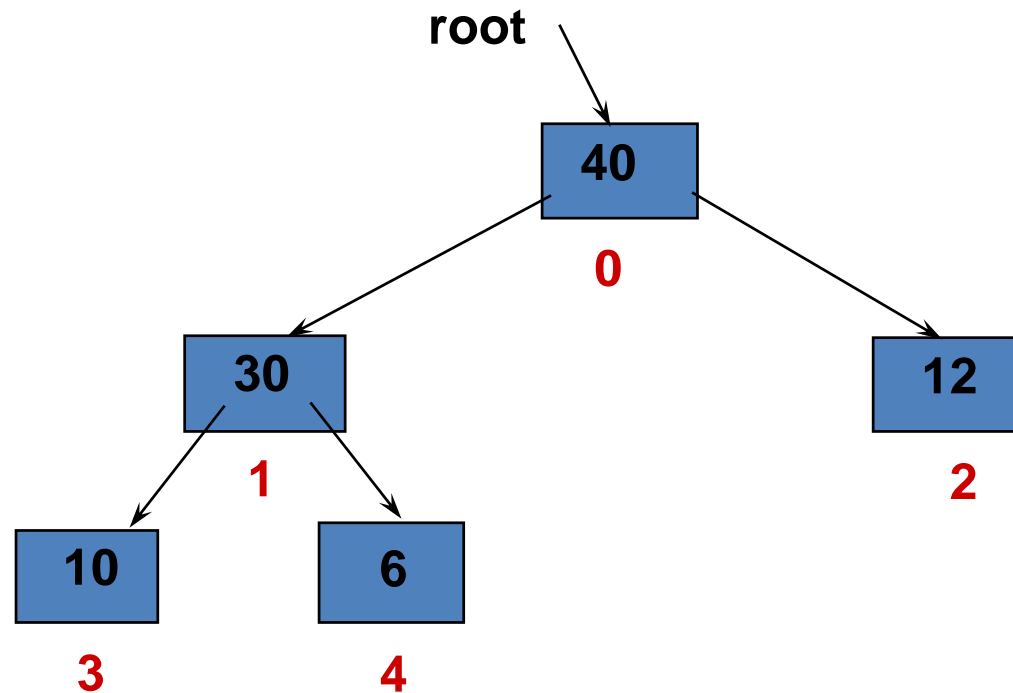


NO NEED TO CONSIDER AGAIN



After reheaping remaining unsorted elements

values	
[0]	40
[1]	30
[2]	12
[3]	10
[4]	6
[5]	60
[6]	70

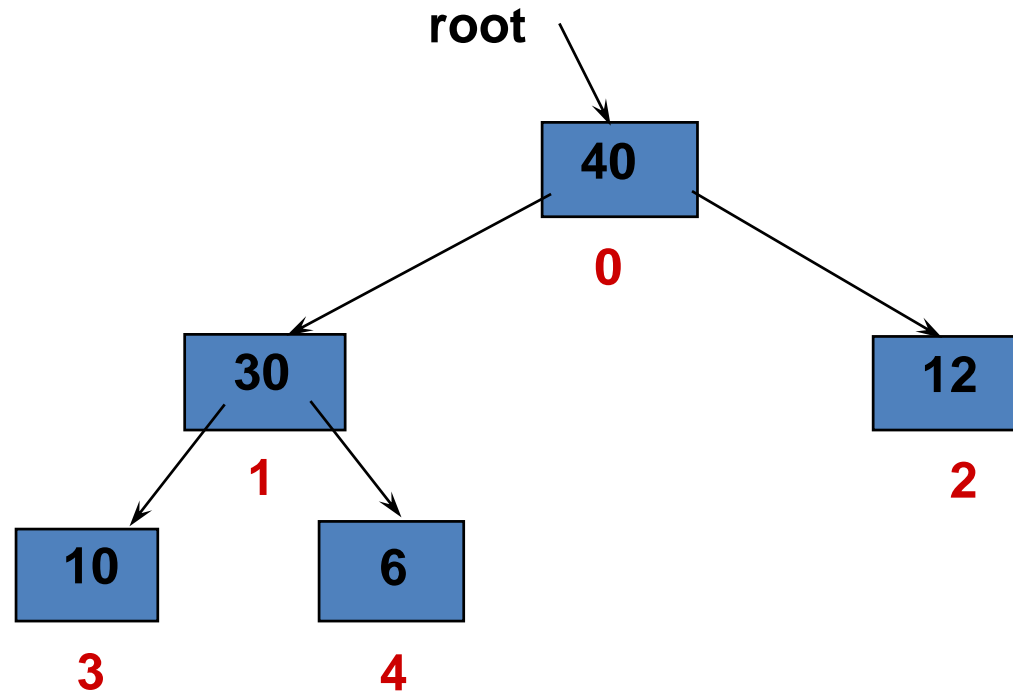




Swap root element into last place in unsorted array

values

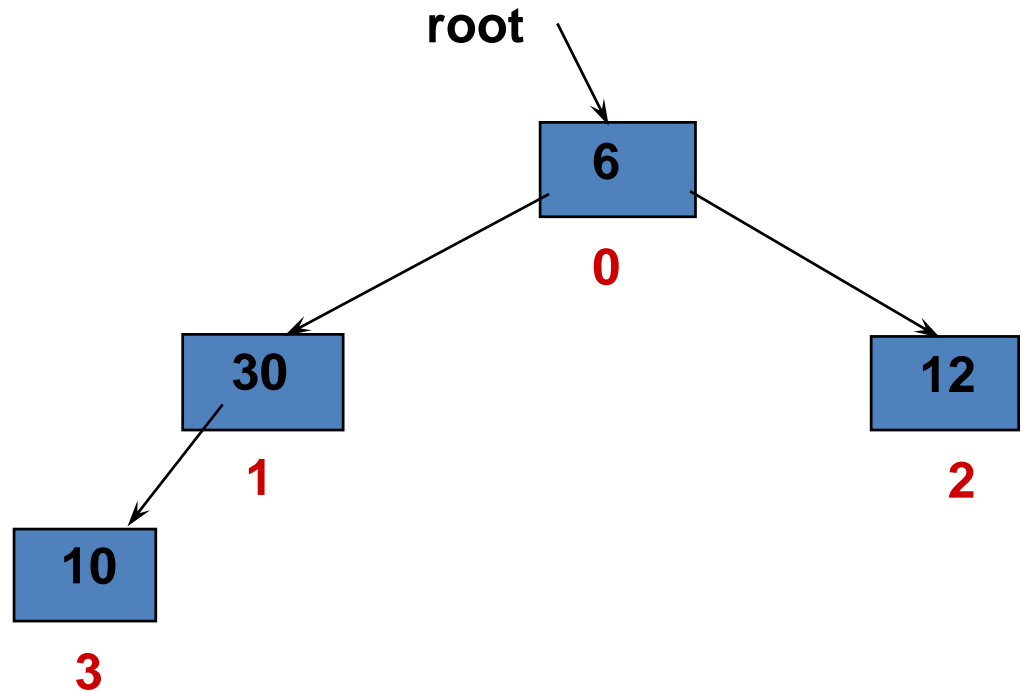
[0]	40
[1]	30
[2]	12
[3]	10
[4]	6
[5]	60
[6]	70





After swapping root element into its place

values	
[0]	6
[1]	30
[2]	12
[3]	10
[4]	40
[5]	60
[6]	70

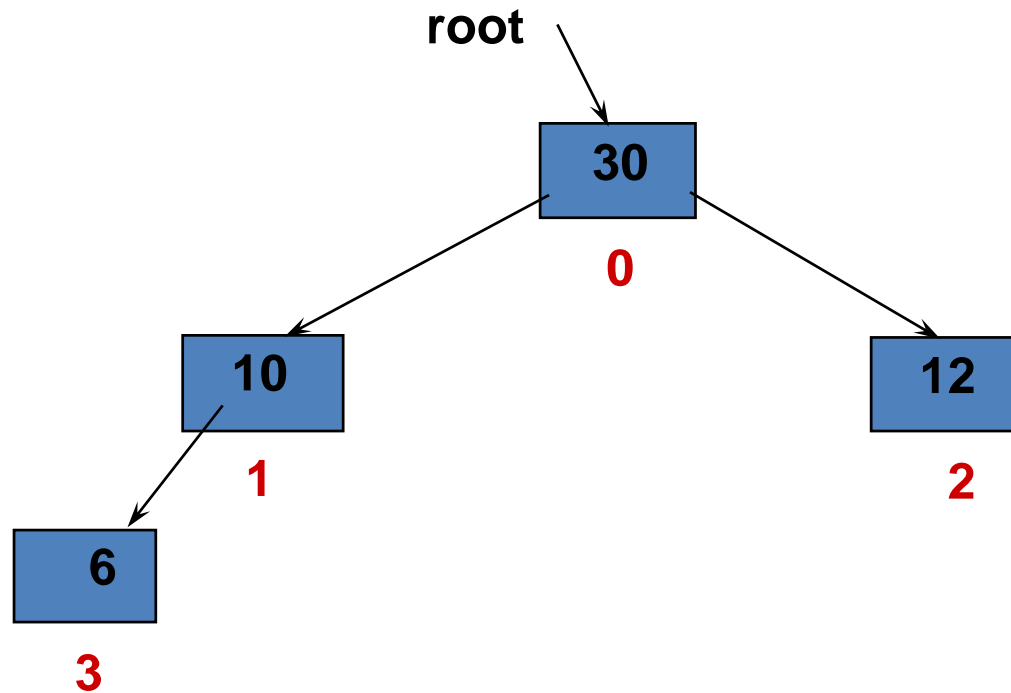


NO NEED TO CONSIDER AGAIN



After reheaping remaining unsorted elements

values	
[0]	30
[1]	10
[2]	12
[3]	6
[4]	40
[5]	60
[6]	70

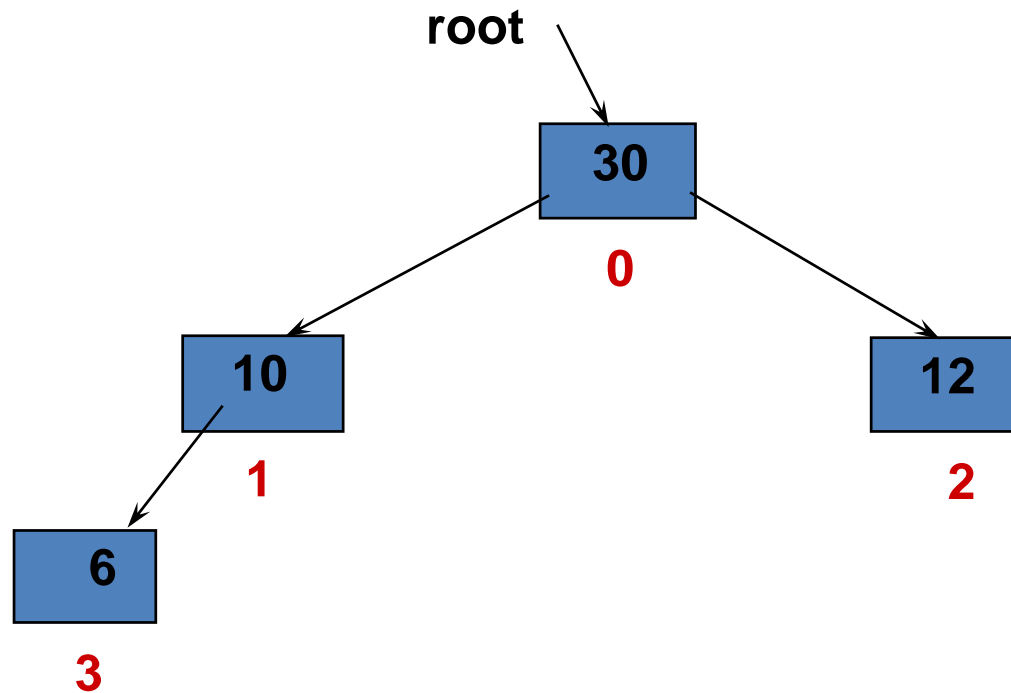




Swap root element into last place in unsorted array

values

[0]	30
[1]	10
[2]	12
[3]	6
[4]	40
[5]	60
[6]	70

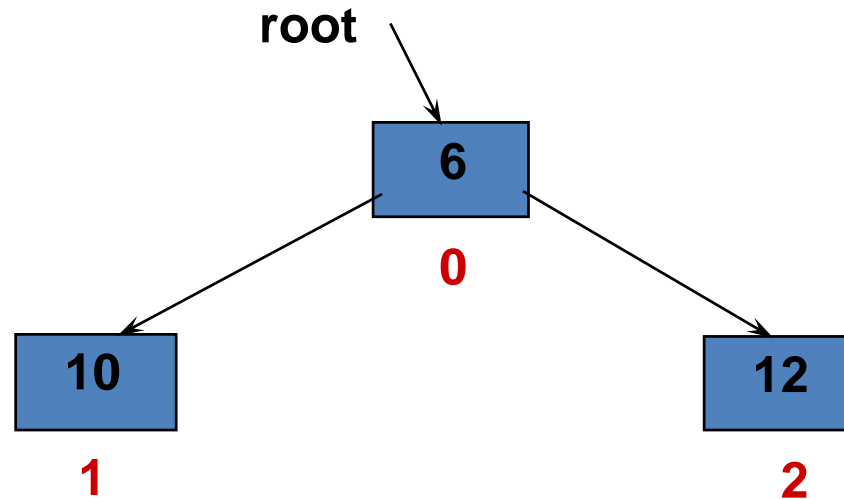




After swapping root element into its place

values

[0]	6
[1]	10
[2]	12
[3]	30
[4]	40
[5]	60
[6]	70

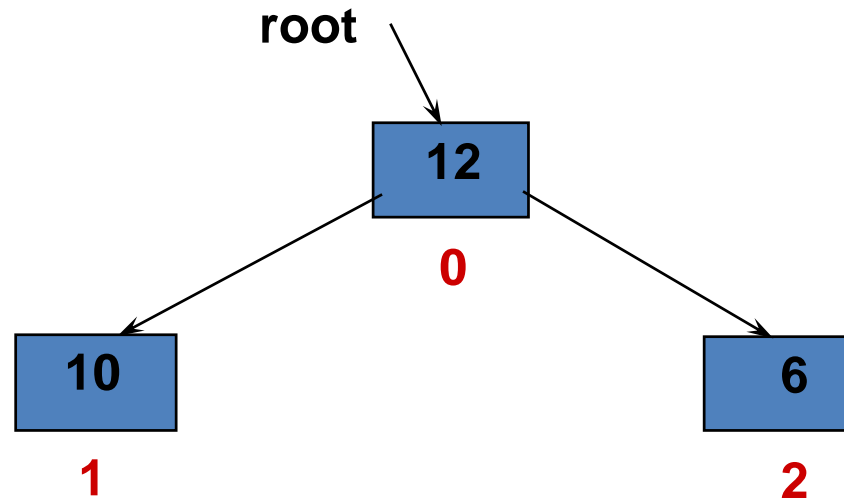


NO NEED TO CONSIDER AGAIN



After reheaping remaining unsorted elements

values	
[0]	12
[1]	10
[2]	6
[3]	30
[4]	40
[5]	60
[6]	70

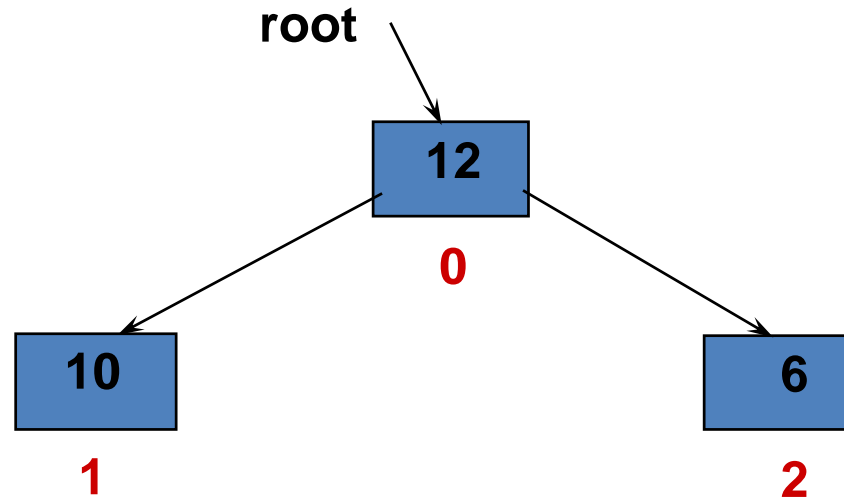




Swap root element into last place in unsorted array

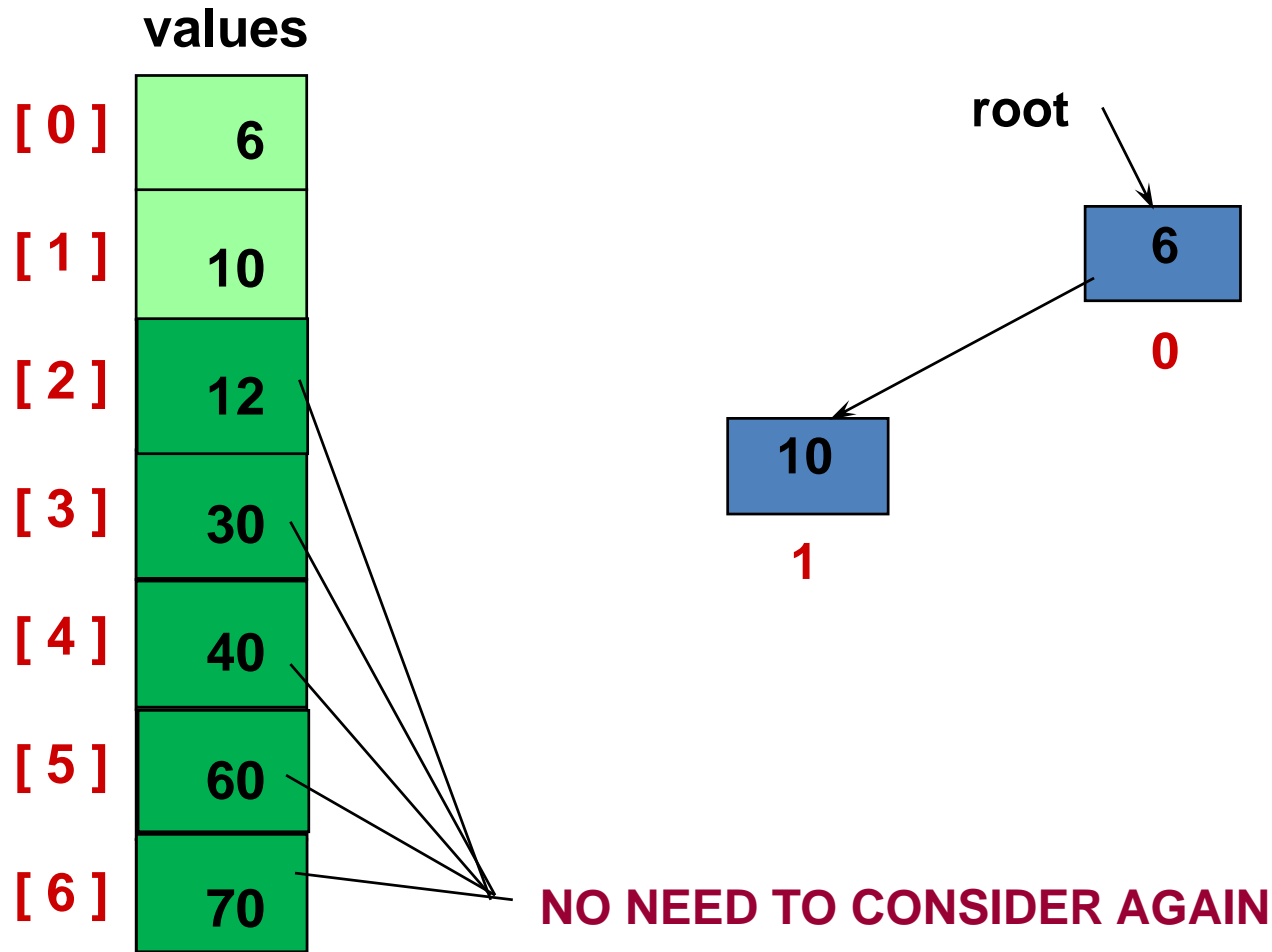
values

[0]	12
[1]	10
[2]	6
[3]	30
[4]	40
[5]	60
[6]	70





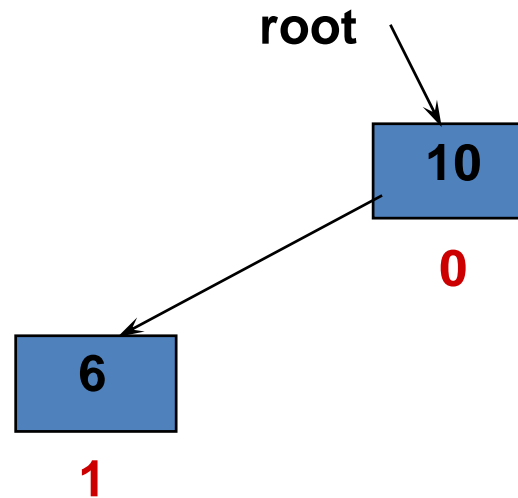
After swapping root element into its place





After reheaping remaining unsorted elements

	values
[0]	10
[1]	6
[2]	12
[3]	30
[4]	40
[5]	60
[6]	70

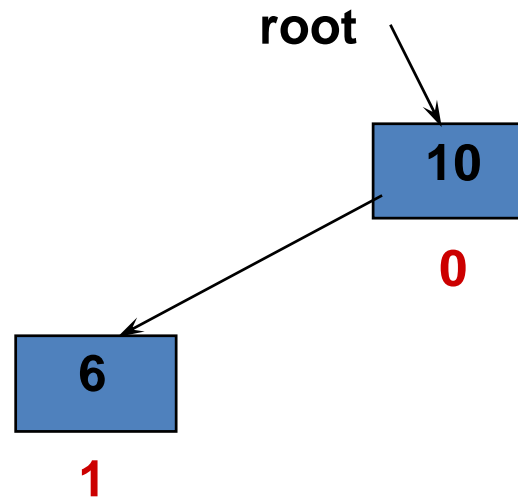
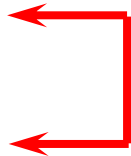




Swap root element into last place in unsorted array

values

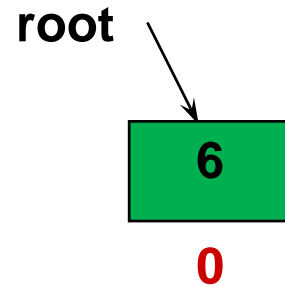
[0]	10
[1]	6
[2]	12
[3]	30
[4]	40
[5]	60
[6]	70





After swapping root element into its place

	values
[0]	6
[1]	10
[2]	12
[3]	30
[4]	40
[5]	60
[6]	70



ALL ELEMENTS ARE SORTED

```
template <class ItemType >
void HeapSort ( ItemType values[], int numValues )

// Post: Sorts array values[ 0 . . numValues-1 ] into
// ascending order by key
{
    int index ;

    // Convert array values[0..numValues-1] into a heap
    for (index = numValues/2 - 1; index >= 0; index--)
        ReheapDown ( values , index , numValues-1 ) ;

    // Sort the array.
    for (index = numValues-1; index >= 1; index--)
    {
        Swap (values[0], values[index]);
        ReheapDown (values , 0 , index - 1);
    }
}
```

N

ReheapDown

```
template< class  ItemType >
void  ReheapDown ( ItemType  values[], int root, int bottom )

//  Pre:  root is the index of a node that may violate the
//         heap order property
//  Post:  Heap order property is restored between root and
//         bottom

{
    int  maxChild ;
    int  rightChild ;
    int  leftChild ;

    leftChild  =  root * 2 + 1 ;
    rightChild =  root * 2 + 2 ;
```

```

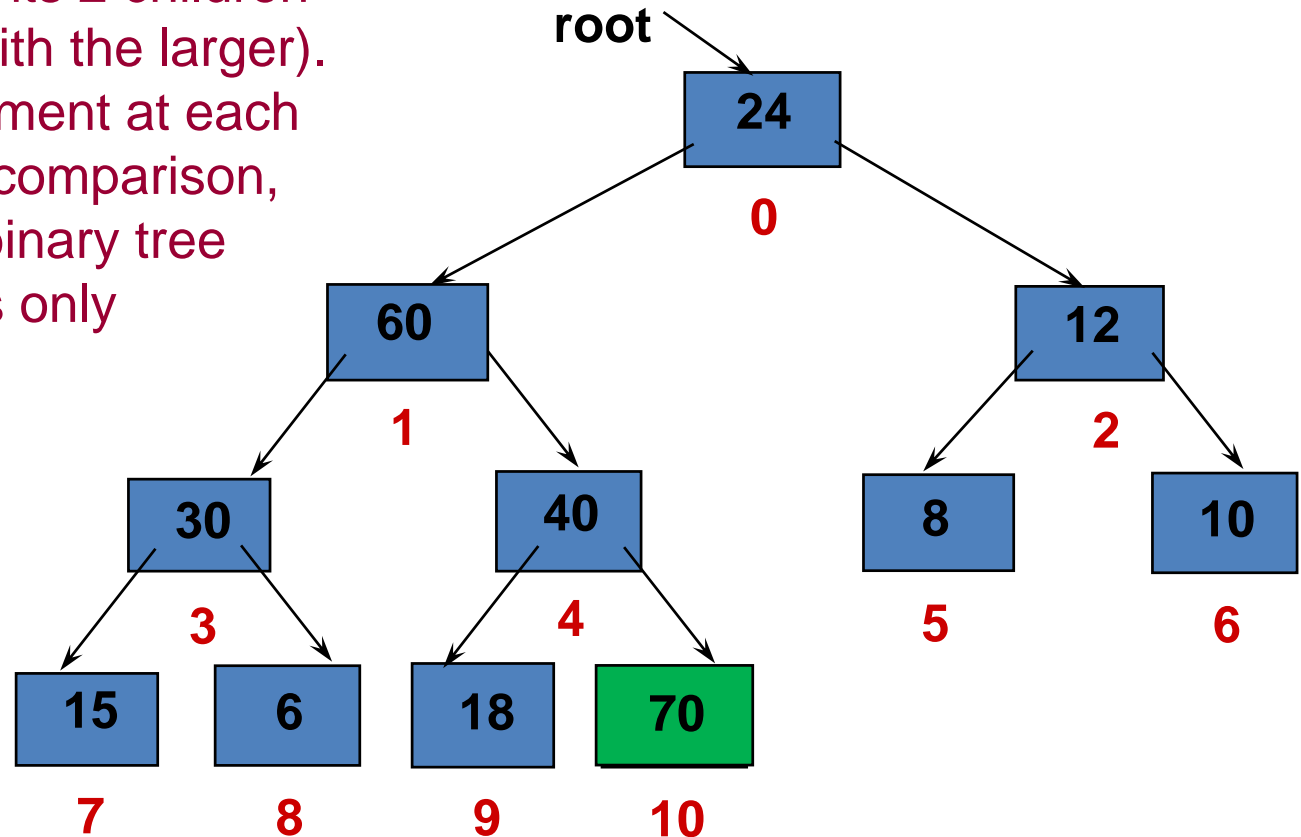
if (leftChild <= bottom) // ReheapDown continued
{
    if (leftChild == bottom)
        maxChild = leftChild;
    else
    {
        if (values[leftChild] <= values [rightChild])
            maxChild = rightChild ;
        else
            maxChild = leftChild ;
    }
    if (values[ root ] < values[maxChild])
    {
        Swap (values[root], values[maxChild]);
        ReheapDown ( maxChild, bottom ) ;
    }
}
}

```



Heap Sort: How many comparisons?

In reheap down, an element is compared with its 2 children (and swapped with the larger). But only one element at each level makes this comparison, and a complete binary tree with N nodes has only $O(\log_2 N)$ levels.





Heap Sort of N elements: How many comparisons?

$(N/2) * O(\log N)$ compares to create original heap

$(N-1) * O(\log N)$ compares for the sorting loop

= $O(N * \log N)$ compares total

(reheap takes $O(\log N)$)

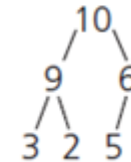


Heap Sort

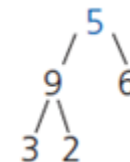
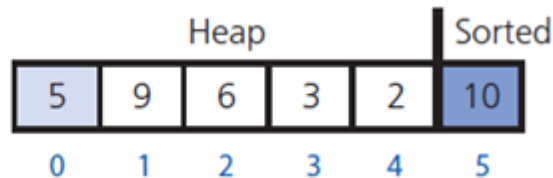
Array anArray

Tree representation of Heap region

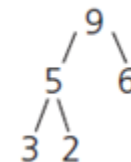
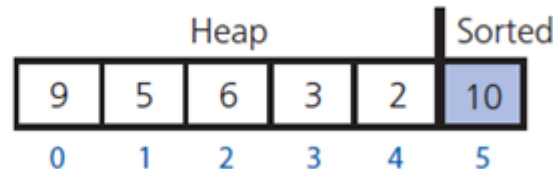
After making anArray a heap



After swapping anArray[0] with anArray[5] and decreasing the size of the Heap region



After heapRebuild(0, anArray, 4)

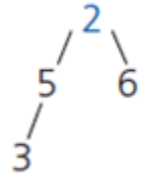


- A trace of heap sort, beginning with the heap

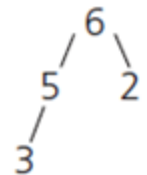
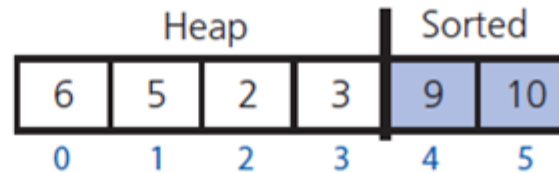


Heap Sort

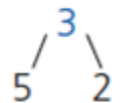
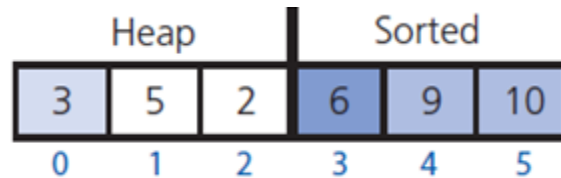
After swapping `anArray[0]` with `anArray[4]` and decreasing the size of the Heap region



After `heapRebuild(0, anArray, 3)`



After swapping `anArray[0]` with `anArray[3]` and decreasing the size of the Heap region

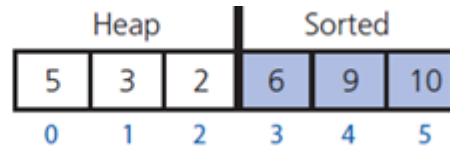


- A trace of heap sort, beginning with the heap

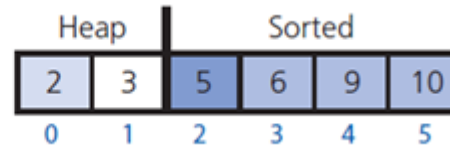


Heap Sort

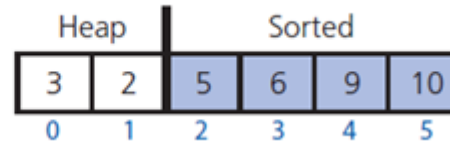
After `rebuildHeap(0, anArray, 2)`



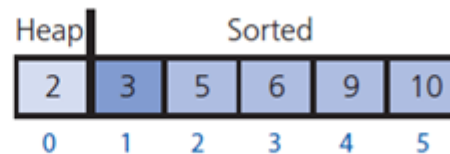
After swapping `anArray[0]` with `anArray[2]` and decreasing the size of the Heap region



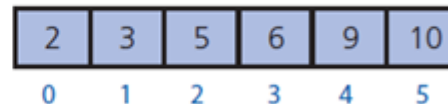
After `heapRebuild(0, anArray, 1)`



After swapping `anArray[0]` with `anArray[1]` and decreasing the size of the Heap region



Array is sorted



- A trace of heap sort, beginning with the heap