



Energy Efficient Resource Management and Task Scheduling for IoT Services in Edge Computing Paradigm

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Outline

- ❖ Background
- ❖ System Model
- ❖ Scheduling Scheme
- ❖ Real Data Based Evaluation
- ❖ Conclusion and Future Work



Background

❖ Internet of Things (IoT)

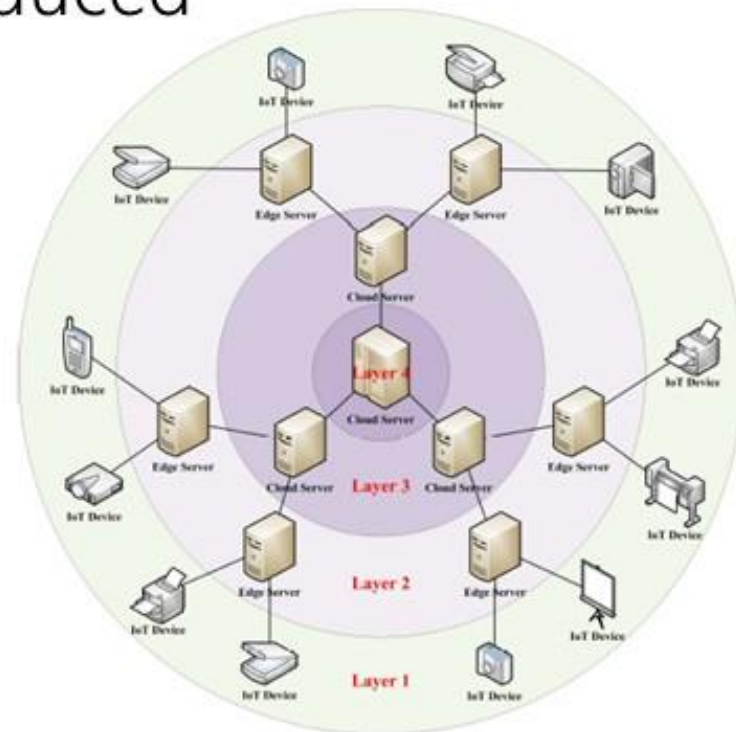
- Connect interrelated devices over the network
- Cooperate with each other to reach common goals
- Wide application in several aspects
 - ✓ Internet of Vehicles
 - ✓ Smart home
 - ✓ Transportation
 - ✓ ...



Background

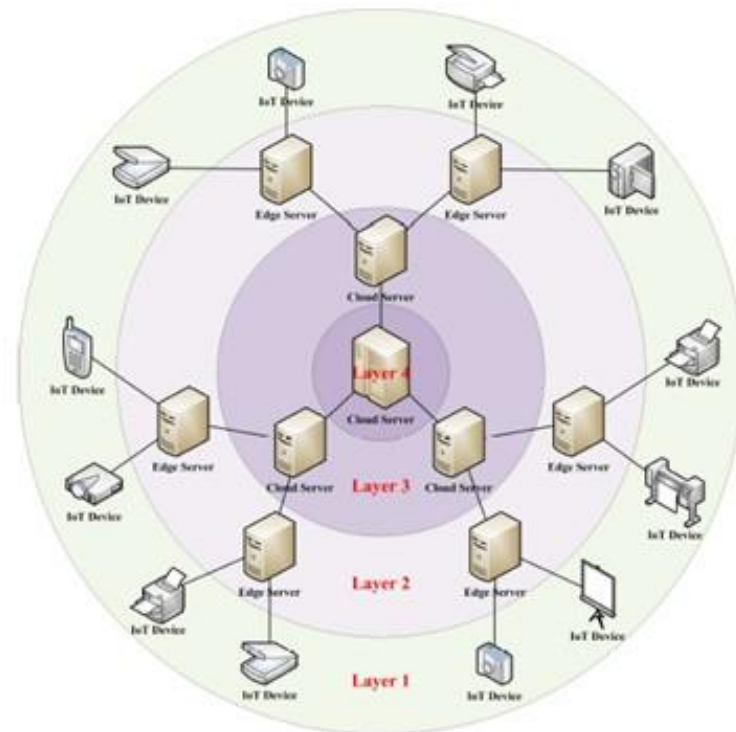
❖ Edge Computing

- Novel computing paradigm for IoT service implementation
- Computation overload distributedly balanced
- Communication delay reduced



Background

- ❖ Energy Efficiency Issues in Edge Computing
 - Resource management according to dynamic workload
 - Task assignment between layers



Background

❖ Our contribution

- Dynamic Formulation for IoT Services Under the Edge Computing Paradigm
- Quantitative Analyses for Performance and Energy Costs of IoT Services
- Joint Decision Process of Resource Management and Task Scheduling with High Efficiency
- Empirical Experiments Based on Real-World Data



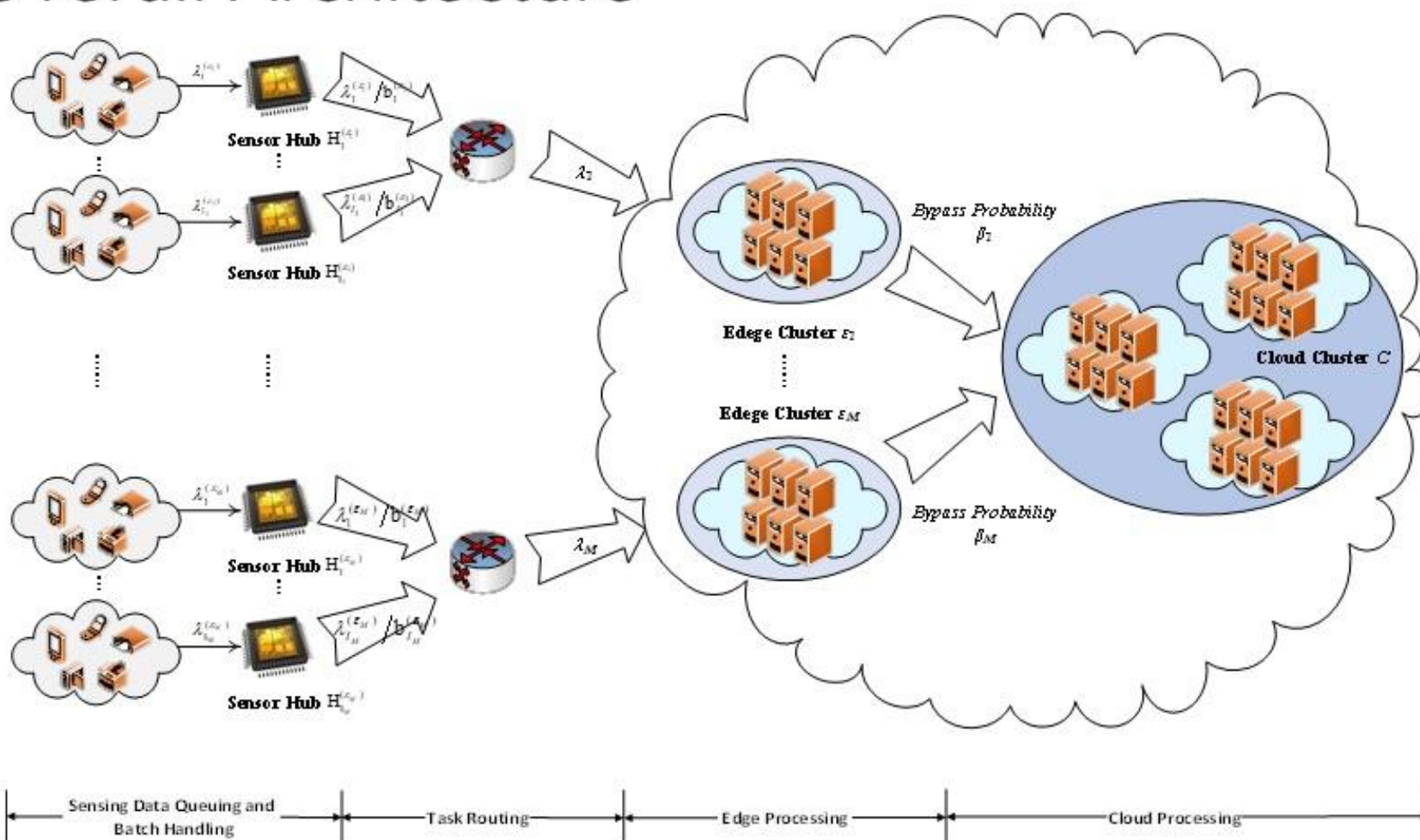
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System Model

❖ Overall Architecture



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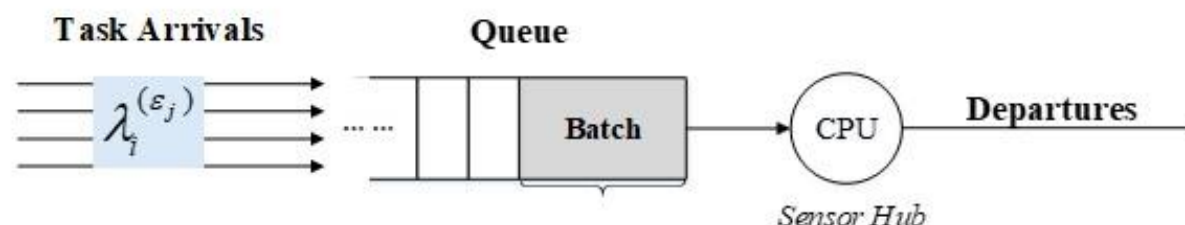


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System Model

❖ Sensor Hub $H_i^{\varepsilon_j}$

- M/M/1 Queue Model with Batch Service
- Average Queue Length $E[q_i] = \frac{\rho_i}{1 - \rho_i}$
- Average Response Time $T_i^{RS} = \frac{b_i \cdot E[q_i]}{\lambda_i}$
- Power Consumption $P_i^{\varepsilon_j} = \sigma_i \cdot (p_i^{(static)} + \rho_i \cdot p_i^{(dyn)})$



System Model

❖ Sensor Hub $H_i^{\varepsilon_j}$

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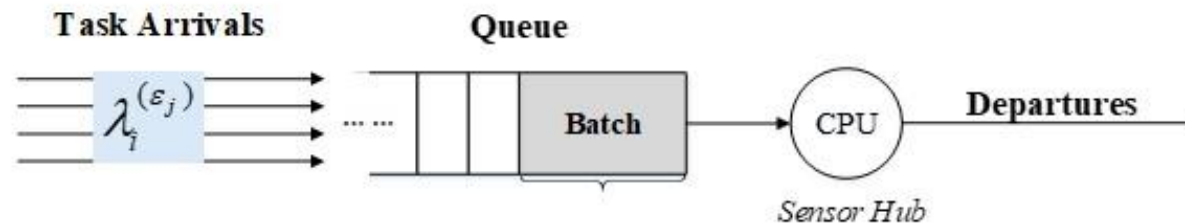
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Batch Arrival Rate

$$\hat{\lambda}_i^{(\varepsilon_j)} = \frac{\lambda_i^{(\varepsilon_j)}}{b_i^{(\varepsilon_j)}}$$



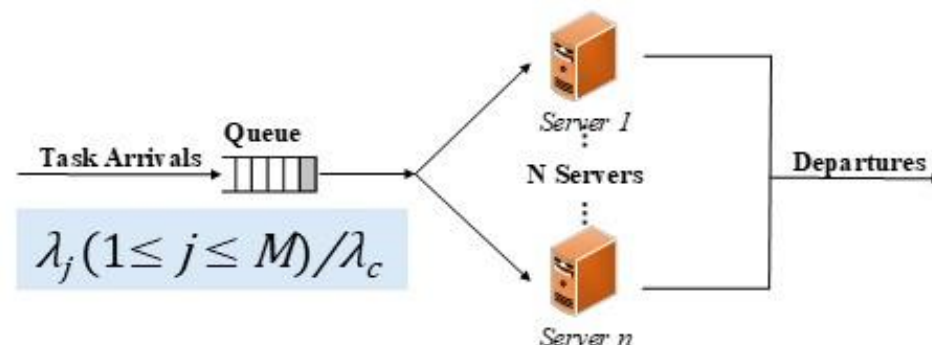
System Model

❖ Edge (ε_j) and Cloud (c) Server Cluster

- M/M/n Queue Model
- Average Queue Length

$$E[q_j] = n_j \rho_j + \frac{\rho_j (n_j \rho_j)^{n_j}}{n_j! (1 - \rho_j)^2} \left[\sum_{k=0}^{n_j-1} \frac{(n_j \rho_j)^k}{k!} + \frac{(n_j \rho_j)^{n_j}}{n_j! (1 - \rho_j)} \right]^{-1}$$

- Average Response Time $T_j^{RS} = \frac{E[q_j]}{\lambda_j}$



System Model

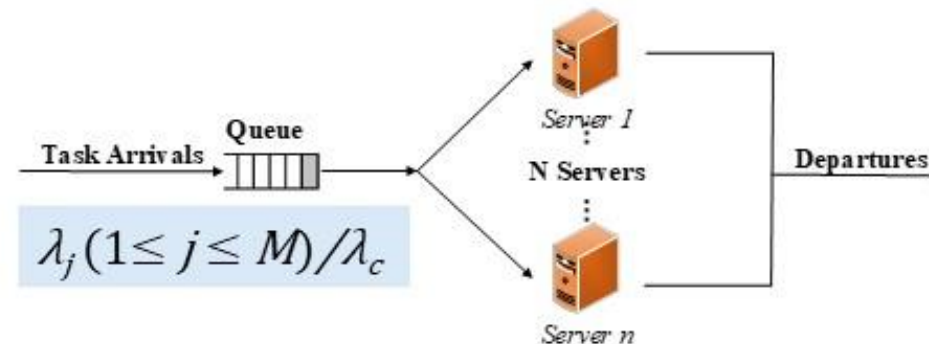
❖ Edge and Cloud Server Cluster ε_j

- M/M/n Queue Model

- Average Queue Length $E[q_j] = n_j \rho_j + \frac{\rho_j (n_j \rho_j)^{n_j}}{n_j! (1 - \rho_j)^2} \left[\sum_{k=0}^{n_j-1} \frac{(n_j \rho_j)^k}{k!} + \frac{(n_j \rho_j)^{n_j}}{n_j! (1 - \rho_j)} \right]^{-1}$

- Average Response Time $T_j^{RS} = \frac{E[q_j]}{\lambda_j}$

- Power Consumption $P_j = \sum_{k=1}^{N_j} \left(\sigma_k \cdot \left(p_k^{(static)} + \rho_k \cdot p_k^{(dyn)} \right) \right)$



System Model

❖ Cross-Layer Transmission

- Bypass Transmission Probability $\beta_j (1 \leq j \leq M)$
- Power Consumption $T_j^{TR} = D_j \cdot t$
- Communication Delay $P_j^{TR} = p^{(dyn)} \cdot \lambda_j \cdot \beta_j + p^{(static)}$



System Model

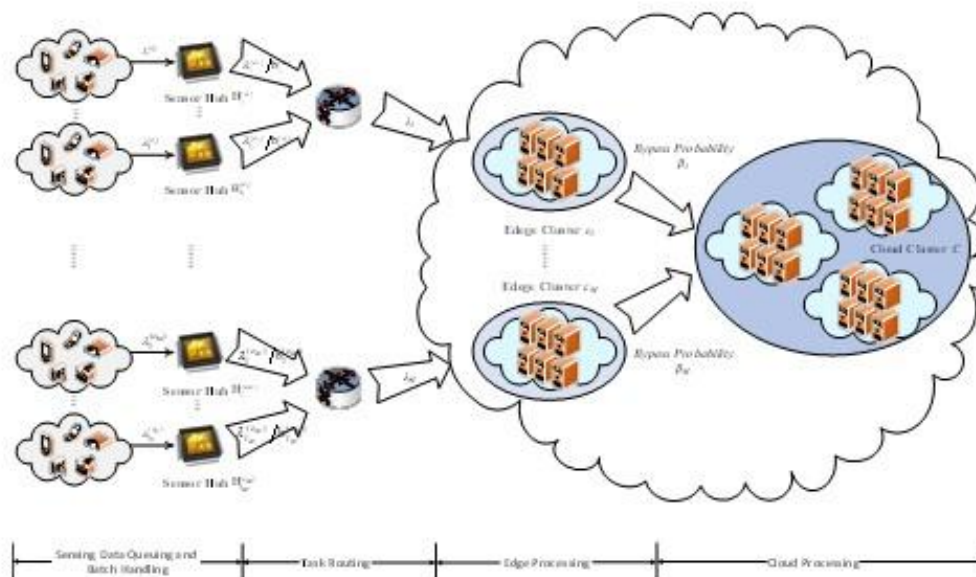
❖ Overall Architecture

- Central Cloud Cluster c
- Edge Clusters $\varepsilon_j (1 \leq j \leq M)$
- Sensor Hubs Subordinating to ε_j $H_i^{\varepsilon_j} (1 \leq i \leq I_j)$

Arrival Rate λ_c

Arrival Rate λ_j

Batch Arrival Rate
 $\hat{\lambda}_i^{(\varepsilon_j)}$



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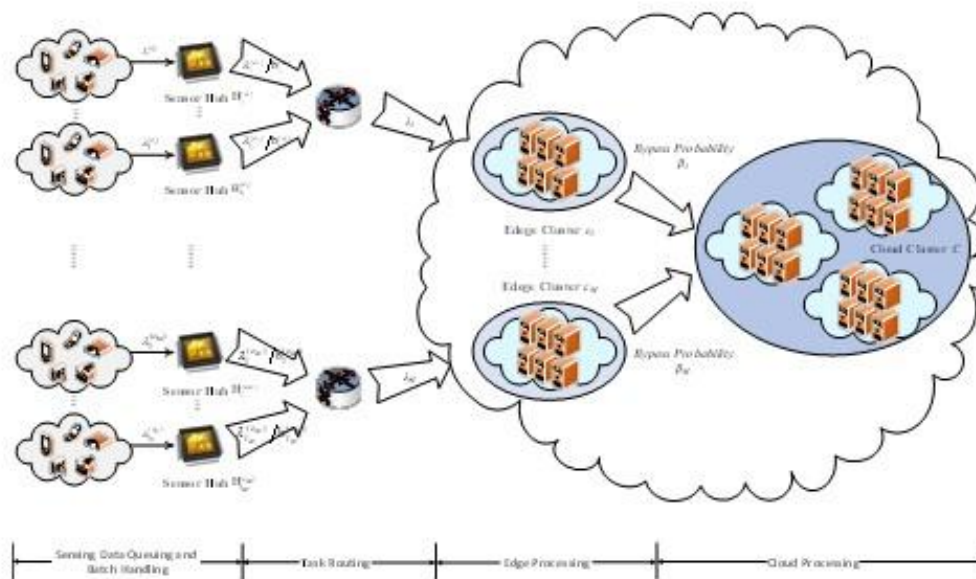


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System Model

❖ Overall Architecture

- Central Cloud Cluster c
- Edge Clusters $\varepsilon_j (1 \leq j \leq M)$
- Sensor Hubs Subordinating to ε_j $H_i^{\varepsilon_j} (1 \leq i \leq I_j)$



$$\lambda_j = (1 - \beta_j) \cdot \sum_{i=1}^{I_j} \hat{\lambda}_i^{(\varepsilon_j)}$$

$$\lambda_c = \beta_j \cdot \sum_{i=1}^{I_j} \hat{\lambda}_i^{(\varepsilon_j)}$$



System Model

❖ Total System Power Consumption

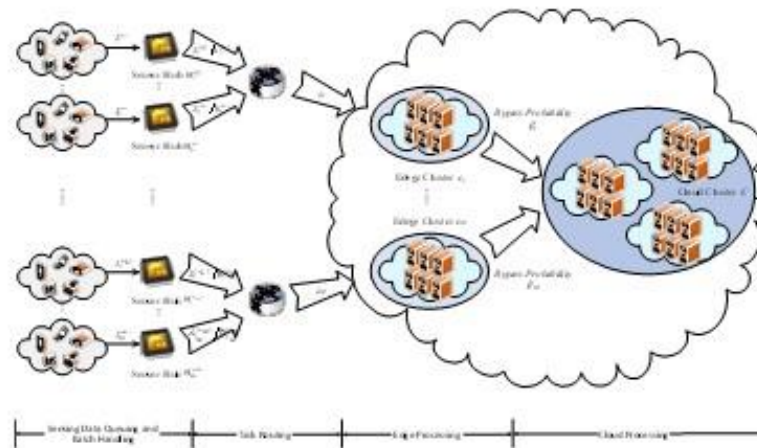
$$P^{sys} = \sum_{j=1}^M (P_j + P_j^{TR}) + P_c$$

❖ Total System Response Time

$$T^{sys} = \sum_{j=1}^M \lambda_j \cdot \left(\frac{q_j \cdot (1 - \beta_j)}{n_j \cdot \mu_j} + \beta_j \cdot (T_j^{TR} + T_c) \right) \bigg/ \sum_{j=1}^M \lambda_j$$

$$\lambda_j = (1 - \beta_j) \cdot \sum_{i=1}^{I_j} \hat{\lambda}_i^{(\varepsilon_j)}$$

$$\lambda_c = \beta_j \cdot \sum_{i=1}^{I_j} \hat{\lambda}_i^{(\varepsilon_j)}$$



System Model

❖ Total System Power Consumption

$$P^{sys} = \sum_{j=1}^M (P_j + P_j^{TR}) + P_c$$

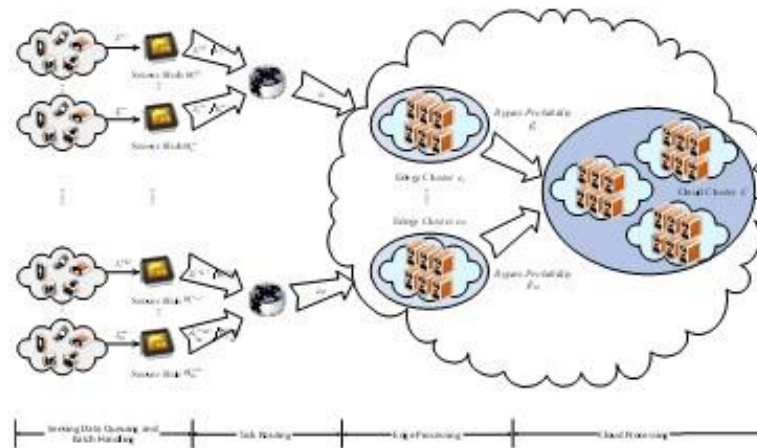
❖ Total System Response Time

Stationary Response Time at the Cloud Site

$$T^{sys} = \sum_{j=1}^M \lambda_j \cdot \left(\frac{q_j \cdot (1 - \beta_j)}{n_j \cdot \mu_j} + \beta_j \cdot (T_j^{TR} + T_c) \right) / \sum_{j=1}^M \lambda_j$$

$$\lambda_j = (1 - \beta_j) \cdot \sum_{i=1}^{I_j} \hat{\lambda}_i^{(\varepsilon_j)}$$

$$\lambda_c = \beta_j \cdot \sum_{i=1}^{I_j} \hat{\lambda}_i^{(\varepsilon_j)}$$



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Scheduling Scheme

❖ Reward Model

- Total System Power Consumption

$$P^{sys} = \sum_{j=1}^M (P_j + P_j^{TR}) + P_c$$

- Total System Response Time

$$T^{sys} = \sum_{j=1}^M \lambda_j \cdot \left(\frac{q_j \cdot (1 - \beta_j)}{n_j \cdot \mu_j} + \beta_j \cdot (T_j^{TR} + T_c) \right) / \sum_{j=1}^M \lambda_j$$

Stationary Response
Time at the Cloud Site

- Service Level Agreement (SLA)

$$T_{SLA}$$

Upper Bound of the
Response Time



Scheduling Scheme

❖ Reward Model

- Total System Power Consumption

$$P^{sys} = \sum_{j=1}^M (P_j + P_j^{TR}) + P_c$$

- Total System Response Time

$$T^{sys} = \sum_{j=1}^M \lambda_j \cdot \left(\frac{q_j \cdot (1 - \beta_j)}{n_j \cdot \mu_j} + \beta_j \cdot (T_j^{TR} + T_c) \right) / \sum_{j=1}^M \lambda_j$$

Stationary Response
Time at the Cloud Site

- Reward Function

$$R = \frac{T_{SLA} - T^{sys}}{T_{SLA}} \bigg/ \frac{P^{sys}}{P_{max}}$$



Scheduling Scheme

❖ MDP Formulation

- Decision Epoch $t = \tau, 2\tau, 3\tau, 4\tau, \dots$

- State Space

$(q_1, \dots, q_j, \dots, q_M)$

Workload in M edge clusters

- Action Space
- Reward Function
- State Probability



Scheduling Scheme

❖ MDP Formulation

- Decision Epoch $t = \tau, 2\tau, 3\tau, 4\tau, \dots$

- State Space $(q_1, \dots, q_j, \dots, q_M)$

- Action Space

$$(\beta_1, \dots, \beta_M) \times (n_1, \dots, n_M, n_c)$$

Cartesian Product for bypass transmission probability n_j and #of power-on machines β_j

- Reward Function

- State Probability



Scheduling Scheme

❖ MDP Formulation

- Decision Epoch $t = \tau, 2\tau, 3\tau, 4\tau, \dots$
- State Space $(q_1, \dots, q_j, \dots, q_M)$
- Action Space $(\beta_1, \dots, \beta_M) \times (n_1, \dots, n_M, n_c)$
- Reward Function

$$R = \frac{T_{SLA} - T^{sys}}{T_{SLA}} \bigg/ \frac{P^{sys}}{P_{max}}$$

- State Transition Probability



MDP Formulation

❖ State Transition Probability

■ Discrete Time Markov Chain

$$p_{q_j \rightarrow q_{j+1}} = \lambda_j (1 - \beta_j) \left(1 - e^{-\sum_{k=1}^M (\lambda_k + q_k \mu_k) \cdot \tau} \right) / \sum_{k=1}^M (\lambda_k + q_k \mu_k)$$

$$p_{q_j \rightarrow q_{j+1}} = q_j \mu_j \left(1 - e^{-\sum_{k=1}^M (\lambda_k + q_k \mu_k) \cdot \tau} \right) / \sum_{k=1}^M (\lambda_k + q_k \mu_k)$$

$$p_{loop} = e^{-\sum_{k=1}^M (\lambda_k + q_k \mu_k) \cdot \tau} + \frac{\lambda_j \beta_j \left(1 - e^{-\sum_{k=1}^M (\lambda_k + q_k \mu_k) \cdot \tau} \right)}{\sum_{k=1}^M (\lambda_k + q_k \mu_k)}$$



Scheduling Scheme

❖ OO-Based Solution

- Ordinal Optimization (Edward et al.1997)
- Coarse Model
- Selection Rule



Scheduling Scheme

❖ OO-Based Solution

- Ordinal Optimization
- Coarse Model - Ignore the future earnings and simply focus on the current reward
- Selection Rule

$$R = \frac{T_{SLA} - T^{sys}}{T_{SLA}} \bigg/ \frac{P^{sys}}{P_{max}}$$



Scheduling Scheme

❖ OO-Based Solution

- Ordinal Optimization
- Coarse Model - Ignore the future earnings and simply focus on the current reward
- Selection Rule
 - Ordered Performance Curve (OPC)
 - Error Level

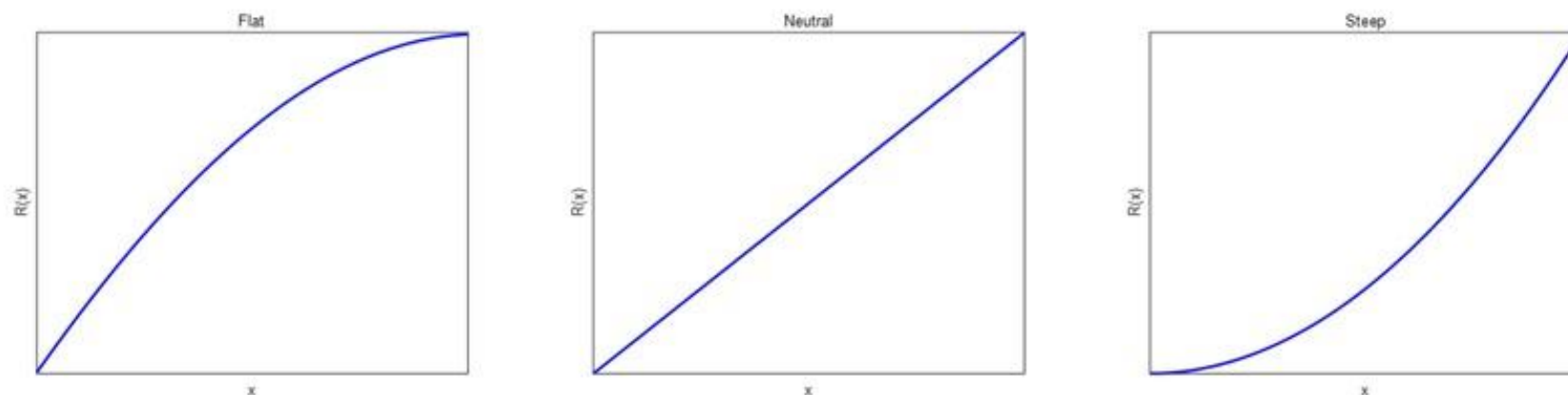
$$R = \frac{T_{SLA} - T^{sys}}{T_{SLA}} \bigg/ \frac{P^{sys}}{P_{max}}$$



Selection Rule

❖ Ordered Performance Curve (OPC)

$$r(x_i) = \frac{r_e^{[i]} - r_e^{[1]}}{r_e^{[N]} - r_e^{[1]}} \quad x_i = \frac{i-1}{N-1}$$



- Non-decreasing sequence for estimated actions

$$\vec{R}_e = (r_e^{[1]}, r_e^{[2]}, \dots, r_e^{[N]})$$



Selection Rule

❖ Error Level

$$error = \max_{1 \leq k \leq N} \left\{ 1 - \frac{\vec{R}_e^{(k)T} \cdot \vec{R}^{(k)}}{\|\vec{R}_e^{(k)}\|_2 \times \|\vec{R}^{(k)}\|_2} \right\}$$

- $0 \leq error < 0.5$ (Small Level)
- $0.5 \leq error < 1$ (Medium Level)
- $error \geq 1$ (Large Level)

Non-decreasing sequence for estimated actions

$$\vec{R}_e^{(k)} = (r_e^{[1]}, r_e^{[2]}, \dots, r_e^{[N]})$$

Non-decreasing sequence for evaluated actions

$$\vec{R}^{(k)} = (r^{[1]}, r^{[2]}, \dots, r^{[N]})$$



Scheduling Scheme

❖ TSECS Algorithm

Algorithm 1 Task Scheduling of Edge-Cloud System (TSECS)

Input: Decision-making space of scheduling scenarios \bar{S} , the number of good enough solutions g , alignment level k

Output: Determined action a_n .

- 1: Calculate the reward values of all the scheduling scenarios using coarse model
 - 2: Estimate the OPC type based
 - 3: Estimate the normalized error level
 - 4: Calculate the number s of selected scenarios and the theory of OO ensures that s scheduling scenarios contains at least k good enough scheduling scenarios with probability no less than 0.95
 - 5: Use the iterative algorithm for MDP optimization to obtain the determined action a_n within the selected s scheduling scenarios
 - 6: Return the determined action a_n .
-



Outline

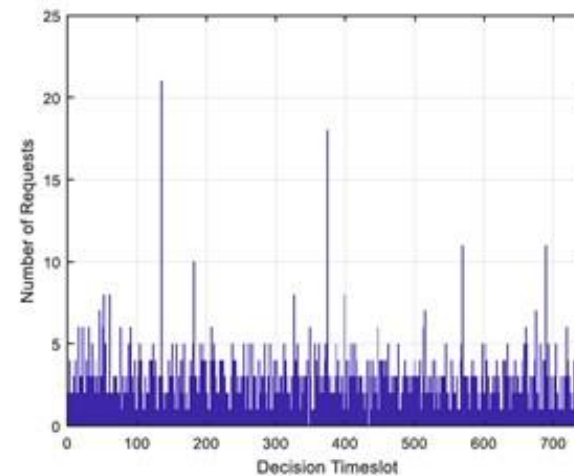
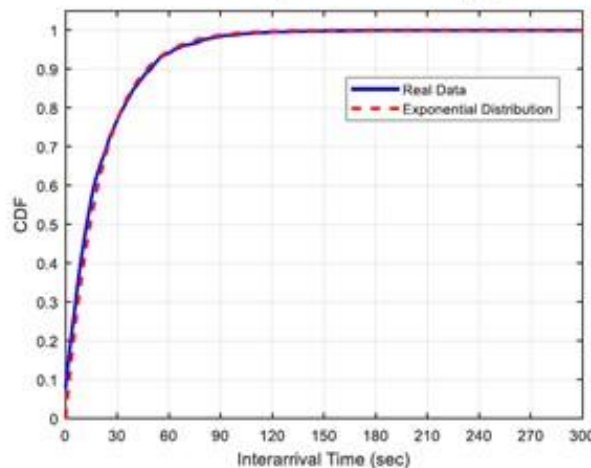
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Evaluation

❖ T-Drive Dataset

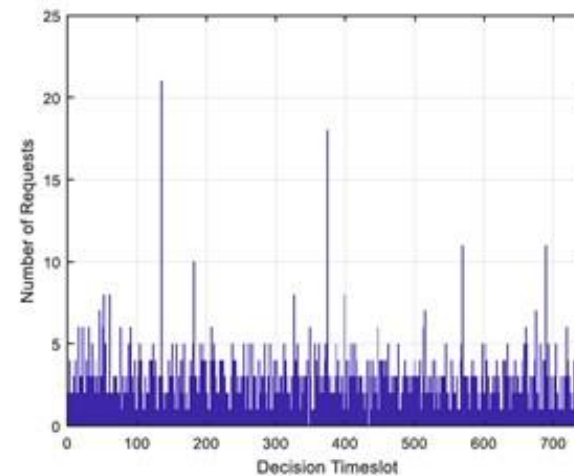
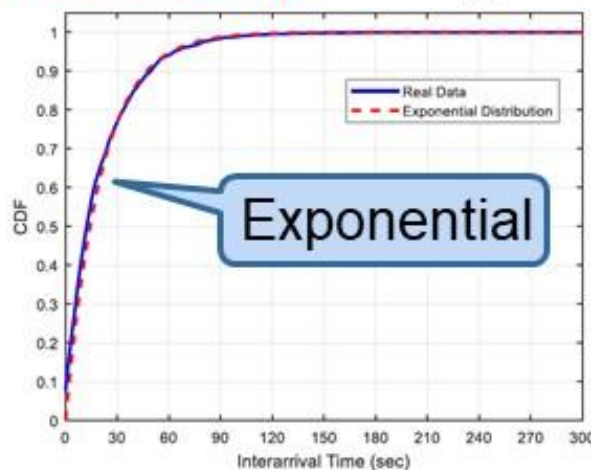
- Released by Microsoft Research
- GPS trajectories of 10,357 taxis within Beijing during a period of one week in 2008
- 4 fields for each piece of data, including taxi id, timestamp, longitude and latitude



Evaluation

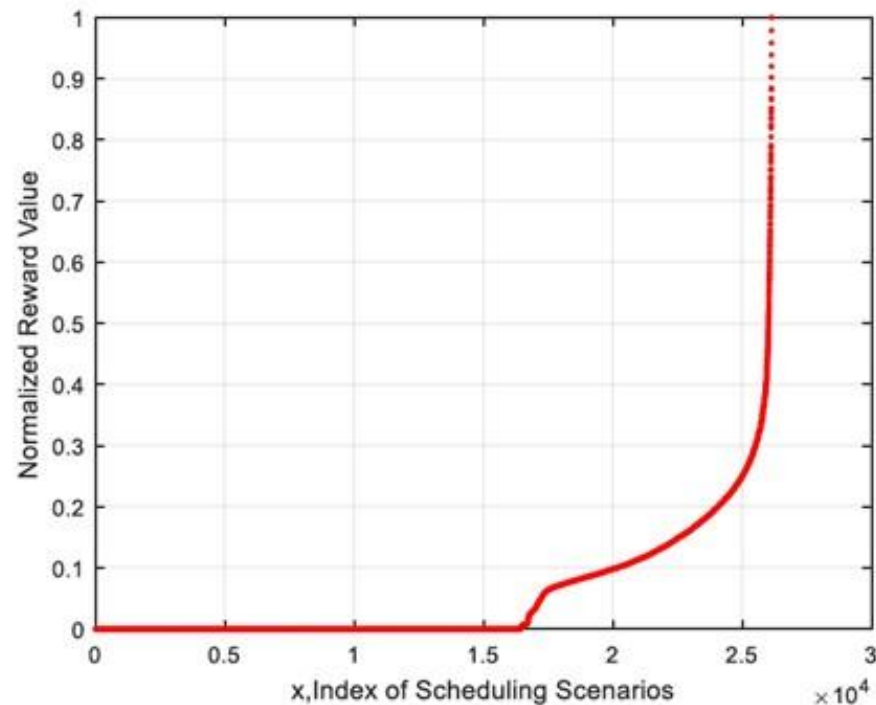
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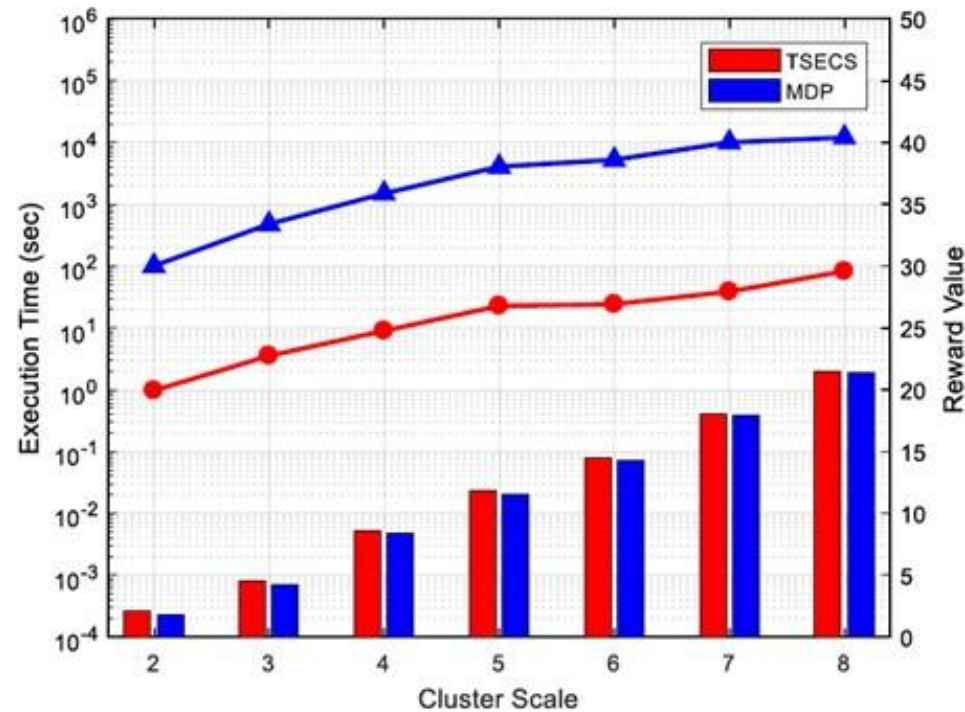
Evaluation

- ❖ Normalized Error= 0.5789 (Medium Level)
- ❖ Ordered Performance Curves (Steep)



Evaluation

❖ TSECS(OO-Assisted) V.S. MDP(Classic)



Evaluation

❖ State-of-the-Art

- Load-Balanced Algorithm
 - Requests are dispatched to edge or cloud servers based on the serving capacity
- Best Effort Algorithm
 - All of servers are switched on and keep constantly running



Evaluation

❖ State-of-the-Art

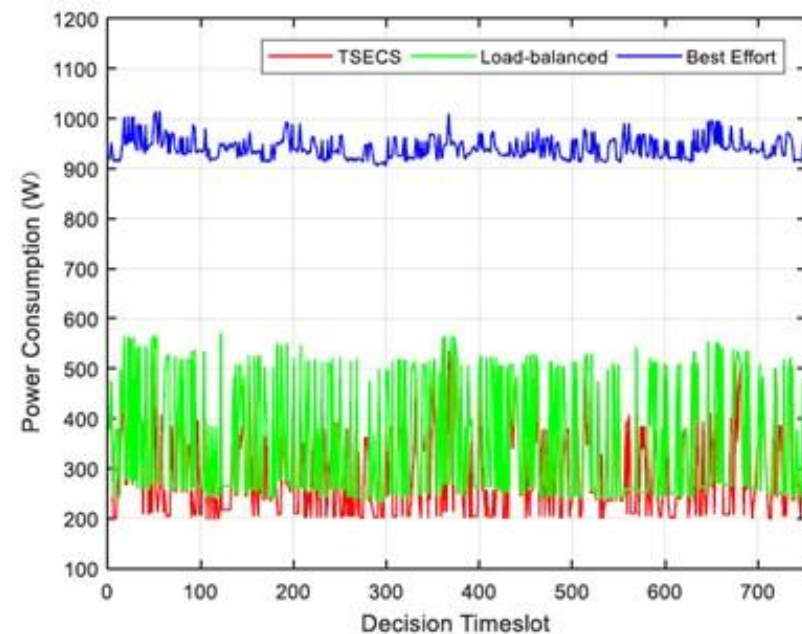
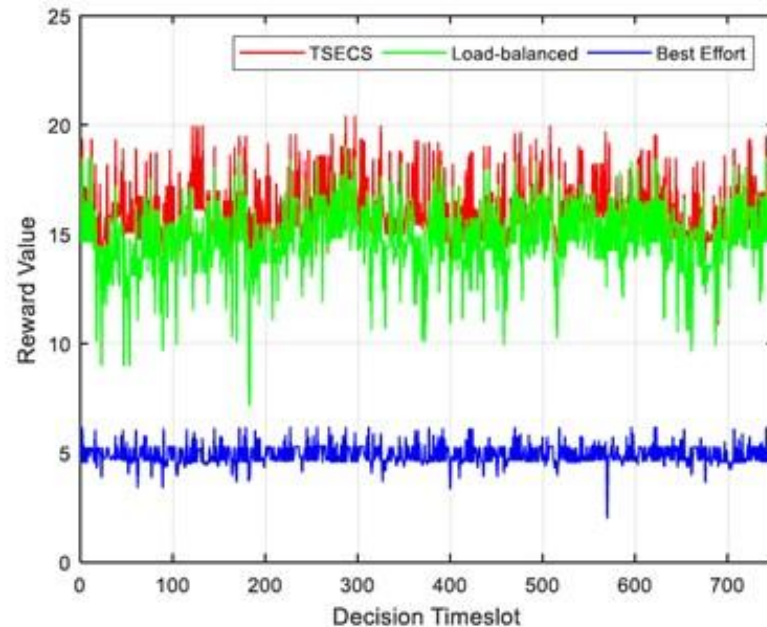


TABLE. AVERAGE RESPONSE TIME OF DIFFERENT ALGORITHMS

T_{SLA} (sec)	Algorithms Response Time (sec)		
	TSECS	Load-balanced	Best Effort
8.00	7.67	7.16	7.59



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Conclusion

- ❖ Modeling approach of energy-aware performance evaluation for an IoT system with edge computing paradigm
- ❖ Joint scheme for task scheduling and resource allocations with high efficiency
- ❖ Scheduling efficiency improvement with OO techniques applied
- ❖ Simulation experiments based on real-world data from IoT and cloud systems



Future Work

- ❖ Detailed specifications on models
 - More variety of real-life IoT systems
 - More adaptive to various task arrivals
- ❖ Elaborate algorithm design
 - Further theoretical validation of efficiency



Thank you!

Q&A

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