Miller-Rabin Primality Test

Miller-Rabin primality test is an optimization of Fermat primality test (based on Fermat's little theorem).

Full Implementation Codes:

```
import random
This function is called for all k trials. It returns False if n is composite
and returns True if n is probably prime.
d is an odd number such that d * 2 ^ r = n - 1 for some r >= 1
def millerTest(d, n):
    # Pick a random number in [2, n-2]
    \# Corner cases in isPrime function make sure that n > 4
    a = 2 + random.randint(1, n - 4)
    # Compute a ^ d % n
    x = pow(a, d, n)
    if (x == 1 \text{ or } x == n - 1):
        return True
    Keep squaring x while one of the following does not happen
    (1) d does not reach n - 1
    (2) (x ^ 2) % n is not 1
    (3) (x ^ 2) % n is not n - 1
    while (d != n - 1):
        x = (x * x) % n
        d *= 2
        if (x == 1):
            return False
        if (x == n - 1):
            return True
    # If no x satisfies, n is a composite
    return False
11 11 11
It returns False if n is composite and returns True if n is probably prime
(pseudoprime).
k is an input parameter that determines accuracy level. Higher level of k
indicates more accuracy.
def isPrime(n, k):
    # Corner cases
    if (n \le 1 \text{ or } n == 4):
        return False
```

```
if (n \le 3):
        return True
    \# Find r such that n = 2 ^ s * d + 1 for some d >= 1
    # d is an odd number
    d = n - 1
    while (d % 2 == 0):
        d //= 2
    # Iterate given number of "k" times
    for i in range(k):
        if (millerTest(d, n) == False):
            return False
    return True
11 11 11
Main programme
def main():
    # Number of iterations
    k = 4
    upperBound = int(input("Find all primes below: "))
    print(f"All primes smaller than {upperBound}: ")
    print()
    counter = 0
    for n in range(1, upperBound):
        if (isPrime(n, k)):
            print(n, end=" ")
            counter += 1
    print("\n")
    print(f"{counter} primes in total")
    print("\n" * 3)
main()
```

Mathematical Theories:

Fermat's Little Theorem: $a^{n-1} \equiv 1 \mod n$ (n is a prime, a is an integer that 1 < a < n-1) $\therefore n$ is a prime $\therefore n-1$ is an even number let $n-1=2^r \times d$ [1] (d is an odd number, which chould not be further divided by 2) $\therefore a^{n-1} \equiv 1 \mod n$ $\Rightarrow a^{n-1} - 1 \equiv 0 \mod n$ $\Rightarrow a^{2^s \times d} - 1 \equiv 0 \mod n$ $\Rightarrow (a^{2^{s-1} \times d} - 1)(a^{2^{s-1} \times d} + 1) \equiv 0 \mod n$ $\Rightarrow (a^{2^{s-1} \times d} + 1)(a^{2^{s-2} \times d} + 1) \cdots (a^d + 1)(a^d - 1) \equiv 0 \mod n$

 $\Rightarrow (a^{2^{s-1} \times d} + 1 \equiv 0 \mod n) \vee (a^{2^{s-2} \times d} + 1 \equiv 0 \mod n) \vee \dots \vee (a^d + 1 \equiv 0 \mod n) \vee (a^d - 1 \equiv 0 \mod n)$ (if any of these is true, n is a pseudoprime, if none of these is true, n is a composit)

Back to the Python Codes:

1. Firstly, let d = n - 1. Keep taking out 2 from d, until d is an odd number (to fulfill the format of [1]).

Specify the number of iterations. The Miller-Rabin primality test can find pseudoprimes, which means there is a possibility that the number found is a composite. Repeating the algorithm with different random value of a will increase the accuracy, but take longer time.

```
# Number of iterations
k = 4
```

3. Generate random number a according to the requirement in Fermat's little theorem

```
# Pick a random number in [2, n-2]
# Corner cases in isPrime function make sure that n > 4
a = 2 + random.randint(1, n - 4)
```

```
4. Verify if a^d + 1 \equiv 0 \mod n or a^d - 1 \equiv 0 \mod n.
    If a^d \mod n = 1, (a^d - 1) \mod n = 0
    If a^d \mod n = n - 1, (a^d + 1) \mod n = 0
# Compute a ^ d % n
x = pow(a, d, n)
if (x == 1 \text{ or } x == n - 1):
    return True
5. If first 2 cases not satisfied, keep trying the rest, until d = n - 1 (the case a^{2^s \times d} + 1 \equiv 0 \mod n).
    x = a^d \mod n
    x' = x^2 \mod n = (a^d \mod n)^2 \mod n = a^{2d} \mod n
    (see Diffie-Hellman Key Exchange file for proof)
    If a^{2^k \times d} \mod n = n - 1, (a^{2^k \times d} + 1) \mod n = 0
while (d != n - 1):
    x = (x * x) % n
    d *= 2
    if (x == 1):
            return False
    if (x == n - 1):
            return True
# If no x satisfies, n is a composite
return False
```

Running Outcome:

Although "pseudoprime" sounds not very accurate, the result is actually very reliable. In the case above, k is set to 4.