Project 3: Dijkstra Sequence

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Chapter 1: Introduction

Dijkstra's algorithm, a renowned greedy algorithm devised by computer scientist Edsger W. Dijkstra, stands as a pivotal solution for the single-source shortest path problem. It efficiently computes the shortest paths from a designated source vertex to all other vertices within a given graph. The algorithm maintains a set of vertices included in the shortest path tree, systematically adding vertices at each step based on their minimum distance from the source. This sequential addition produces what we term a "Dijkstra sequence."

However, a critical aspect arises when considering that a given graph might have more than one valid Dijkstra sequence. This project tackles the challenge of determining whether a provided sequence qualifies as a Dijkstra sequence for a given connected graph.

This project will provide specific code for checking whether a given sequence is Dijkstra sequence or not and also a comprehensive analysis for the algorithm.

Chapter 2: Algorithm Specification

• Full code is presented at the end of this PDF document and also submitted in the code directory.

we will first construct a **weighted undirected graph** using the input data(**NOTE**: We can get from the project requirements that the given graph is undirected), then assess whether the given sequence is a *Dijkstra sequence*.

Basic Data Structure: Weighted Undirected Graph

In the project, i use the following implementation to represent the given graph

Define "struct nei" to represent the neighbours of a specified node in the graph.

- the variable "distance" with type "int" represents the distance between this neighbouring vertex and the specified vertex
- the variable "index" with type "int" represents the index of this neighbouring vertex

```
// Struct definition: Adjacent node
struct nei {
    int distance; // Weight of the edge
    int index; // Index of the target vertex
};
```

Define "struct ver" to represent a vertex in the graph.

- the variable "distance" with type "int" represents the shortest distance updated form the source vertex to this specified vertex.
- the variable "list" with type "struct nei*" points to an array with type "struct nei", containing all the neighbouring vertices of this specified vertex.
- the variable "neinum" with type "int" represents the number of vertices adjacent to this specified vertex.
- the variable "known" with type "int" represents whether this vertex has been processed(Which is a key step of the Dijkstra's algorithm)

```
struct ver {

int distance; // Shortest path length from the source to this node

struct nei *list; // Adjacency list, storing information about adjacent vertices

int neinum; // Number of vertices adjacent to this node

int known; // Whether this node has been processed};
```

Use an array with type "struct ver" to represent the Graph ,each element in the array represents a vertex in the graph.

```
int vertice, edge;
scanf("%d %d", &vertice, &edge);
struct ver Graph[vertice];
```

Initialization of the Graph

- The first loop initializes each vertex, allocating memory for an empty array **Graph[i].list** (whose type is **struct nei**) for every vertice in the given Graph with size: vertices*sizeof(struct nei).
- Through the second loop reading edge information, the code updates the adjacency lists for each vertex based on the edges.

As it is an undirected graph, each edge updates the adjacency lists of two vertices.

1 Graph Initialization

```
1: for i = 0 to vertice -1 do
      Graph[i].neinum = 0
 2:
      Graph[i].list \leftarrow allocate memory for an empty list of neighbors with size
 3:
      vertice \times sizeof(struct nei)
 4: end for
 5: for i = 0 to edge -1 do
      int from, to, distance;
 6:
      \operatorname{scanf}("\%d \%d \%d", \&from, \&to, \&distance);
 7:
      int index \leftarrow Graph[from - 1].neinum;
 8:
      Graph[from - 1].neinum + +;
 9:
      Graph[from - 1].list[index].distance \leftarrow distance;
10:
      Graph[from - 1].list[index].index \leftarrow to - 1;
11:
      index \leftarrow Graph[to - 1].neinum;
12:
      Graph[to-1].neinum++;
13:
      Graph[to-1].list[index].distance \leftarrow distance;
14:
      Graph[to-1].list[index].index \leftarrow from-1;
15:
16: end for
```

Sketch of the whole program

- In the main function, we first read in "cnt", which is the number of the test cases.
- Then, for each test case, we read in a sequence, initialize the Graph by setting the distance of each edge to the source edge to -1 (which represents infinity), and each vertex to be "unknown" at the very beginning.
- After that, we call the function Dijkstra to check whether the input sequence is a dijsktra sequence.
- Also, we have to free the memory we malloced previously.

```
Algorithm 0: Program Sketch
   Input: cnt
 1 while cnt > 0 do
       Input: vertice, Graph
       int a[vertice];
 2
       for i \leftarrow 0 to vertice - 1 do
 3
        | Input: a[i]
       for i \leftarrow 0 to vertice - 1 do
 4
            Graph[i].distance \leftarrow -1;
 5
           Graph[i].known \leftarrow 0;
 6
       Dijkstra(Graph, vertice, a);
 7
       cnt \leftarrow cnt - 1;
 9 for i \leftarrow 0 to vertice - 1 do
       free(Graph[i].list);
11 return 0;
```

The Checking algorithm Dijsktra

(1) Then 'FindShortest' Function

First implement a function to find the minimum distance among unknown vertices (from the specified source vertex)

The project requirements specified that "the positive integer weight (≤ 100)", the number of vertices $Nv (\leq 10^3)$ and the number of edges $Ne (\leq 10^5)$.

Thus, the use of INT_MAX for initialization ensures that any valid distance in the graph will be smaller during the
first iteration.

Note: In the function Dijsktra(to be analysed below), the distance of the source vertex is already set to be zero, so there is no need to worry that we cannot get a real **min** distance (say, all the vertices' distance is -1)

The function checks all the vertices of the given Graph, and returns the minimum distance between the source vertex and the unknown vertices.

- Initialize min to the maximum possible integer value (INT_MAX).
- For each vertex, check if it's unprocessed (!Graph[i].known) and if its distance is smaller than the current minimum (Graph[i].distance <= min). If true, update min with the vertex's distance.
- After the loop, min will contain the minimum distance among unprocessed nodes.
- Return the minimum distance

Algorithm 1: Find the Minimum Distance among Unknown Nodes

(2) The 'Dijsktra' function

• Initialization:

• The function starts by initializing the distance of the source node (given by a[0]) to itself as 0.

• Node Traversal:

- The function iterates over all vertices in the graph using a loop.
- For each loop, it finds the minimum distance among unprocessed nodes by calling the findshortest function.

• Validity Check:

- It checks whether the distance of the current node is invalid (if the distance is less than zero, it means this vertex is not reachable at the current state) or greater than the minimum distance found in the previous step.
- If so, it prints "No" and returns, indicating that the given sequence is not a valid Dijsktra sequence.

• Processing Mark:

• It marks the current node as processed by setting Graph[vertex].known to 1.

• Update Neighbor Distances:

- o For each neighbor of the current node, it checks if the distance from the source node to the neighbor through the current node is smaller than the known distance to the neighbor.
- If true, it updates the distance.

Result Output:

O After processing all nodes without any violations, it prints "Yes," indicating that the shortest paths have been successfully computed.

Algorithm 2: Dijkstra's Algorithm

```
Data: Graph: struct ver*, vertice: int, a: int*
 1 Function Dijkstra(Graph, vertice, a):
 2
       source \leftarrow a[0];
 3
       Graph[source-1].distance \leftarrow 0; // Set the distance from the
        source node to itself as 0
       for i \leftarrow 0 to vertice - 1 do
 4
           \min \leftarrow \text{findshortest(Graph, vertice)};
 5
 6
           vertex \leftarrow a[i] - 1;
          if Graph/vertex. distance < 0 or Graph/vertex. distance > min
 7
            then
              // If the distance of the current node is invalid
                  or greater than the minimum distance
 8
              Output: No;
 9
              return;
10
           Graph[vertex].known \leftarrow 1; // Mark the current node as
            processed
           cnt \leftarrow Graph[vertex].neinum;
11
12
           for j \leftarrow 0 to cnt - 1 do
              neighbour \leftarrow Graph[vertex].list[j].index;
13
              distance \leftarrow Graph[vertex].list[j].distance;
14
              if \neg Graph/neighbour/.known then
15
                  // If the distance from the current node to the
                      adjacent node is smaller, update the distance
                  if Graph/vertex. distance + distance <
16
                    Graph/neighbour/.distance or Graph/neighbour/.distance <
                   0 then
                      Graph[neighbour].distance \leftarrow
17
                       Graph[vertex].distance + distance;
       Output: Yes;
18
```

Chapter 3: Testing Results

• Basic tests given in the question

Sample Input

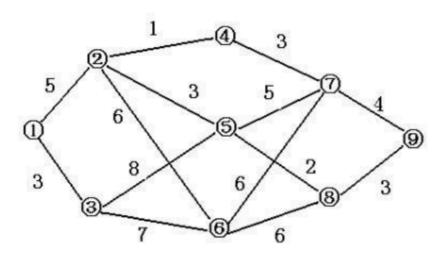


OUTPUT

1	Yes				
2	Yes Yes Yes				
3	Yes				
4	No				

• Other Test Cases

Graph:



Input Graph

1	9 14			
2	1 2 5			
3	1 3 3			
4	2 4 1			
5	253			
6	266			
7	3 5 8			
8	3 6 7			
9	473			
10	5 7 5			
11	582			
12	676			
13	686			
14	794			
15	893			

Input Sequence	Output
132457689	Yes
1 3 2 4 5 7 8 6 9	Yes
123456789	No
132547869	No
245781639	Yes
2 4 5 7 8 1 6 9 3	Yes
2 4 5 7 1 8 6 3 9	Yes
2 4 5 7 1 8 6 9 3	Yes
245871639	No
254367891	No
427518963	Yes
427518693	Yes
427581963	Yes
427581693	Yes

The Program has past all the above test cases,indicating that it is to a degree complete and efficient to judge a red-black tree

.

$Chapter 4: Analysis \ and \ Comments$

1.The 'FindShortest' Function

Time complexcity

• Loop through vertices once O(vertice) with constant time operations inside the loop.

Therefore, the overall time complexity is O(vertice).

Space complexity

Constant space used, thus space complexity is O(1)

2.The "Dijsktra" Function

NOTE: Use variable edge to represent the count of edges, and the variable vertice for the number of vertices.

time complexity

- The main loop runs for each vertex O(vertice).
 - Inside each loop, call findshortest function with whose time complexity is O(vertice), contributes to a total complexity $O(vertice^2)$
 - o Inside each loop, update distances with complexity O(neinum), contributes to a total complexity O(2*edgenum) = O(edge).
- So the total time complexity for the function is $O(vertice^2 + edge)$, consider that the number of edges is $O(vertice^2)$, so the total time complexity for the function is $O(vertice^2)$

space complexity

- *Memory for variables*: Constant space for variables in Dijkstra.
- findshortest Function:
 - Called once in each iteration of the outer loop using constant space
- Overall Space Complexity:

Constant space O(1)

3.The "main"Function

time complexity

- Initialization of Graph Array:
 - \circ Loop through each vertex O(vertice).
 - Inside the loop, allocate memory for the adjacency list O(vertice), contributes to an overall time complexity $O(vertice^2)$
 - \circ Overall, $O(vertice^2)$.

• Input Edges:

- Loop through each edge O(edge).
- Inside the loop, update adjacency lists (constant time).
- \circ Overall, $O(edge) = O(vertice^2)$.

• Multiple Queries:

- Loop through each query O(cnt).
- Inside the loop, call Dijkstra's algorithm with time complexity $O(vertice^2)$ (see analysed above).
- \circ Overall, $O(cnt * vertice^2)$.

• Free Allocated Memory:

- \circ Loop through each vertex to free memory O(vertice).
- Overall Time Complexity: $O(cnt * vertice^2)$.

Space Complexity

• Graph Array:

- \circ Array of vertices O(vertice).
- \circ Each vertex has an adjacency list O(vertice).
- \circ Overall, $O(vertice^2)$.

• Queries:

- Array for each query : O(vertice) for each query.
- Stack
 - For cnt queries, the main function calls the Dijsktra function for cnt times, each with constance space of stack memory.
- Overall O(vertice)
- Overall Space Complexity: O(vertice²)

Summary:

- The overall time complexity is $O(cnt * vertice^2)$
- The space complexity is $O(vertice^2)$, mainly determined by the size of the graph array.

Keep in mind that these are upper bounds, and the actual performance could vary based on the specific ,input and graph characteristics.

$Appendix: Source\ Code$

- 1 #include<stdio.h>
- 2 #include<stdlib.h>
- 3 #include imits.h>
- 4 // Struct definition: Adjacent node

```
5
      struct nei {
 6
        int distance; // Weight of the edge
 7
        int index; // Index of the target vertex
 8
     };
 9
     // Struct definition: Vertex of the graph
     struct ver {
11
        int distance; // Shortest path length from the source to this node
12
        struct nei *list; // Adjacency list, storing information about adjacent vertices
13
                         // Number of vertices adjacent to this node
        int neinum;
14
        int known:
                         // Whether this node has been processed
15
     };
16
     // Function definition: Find the minimum distance among unknown nodes
17
     int findshortest(struct ver* Graph, int vertice) {
        int min = INT\_MAX;
18
19
        for(int i = 0; i < vertice; i++) {
20
           // Unprocessed node with a smaller distance
21
           if(Graph[i].distance >= 0 && Graph[i].distance <= min && (!Graph[i].known)) {
22
             min = Graph[i].distance;
23
           }
24
        }
25
        return min; // return the minimum distance
26
27
     // Function definition: Dijkstra's algorithm to find the shortest path
28
     void Dijkstra(struct ver* Graph, int vertice, int *a) {
29
        int source = a[0];
                                   // Set the source vertex of the shortest path problem to be first element of the input sequence.
30
        Graph[source - 1].distance = 0; // Set the distance from the source node to itself as 0
31
32
        // Traverse all nodes
33
        for(int i = 0; i < vertice; i++) {
34
           int min = findshortest(Graph, vertice); //find the shortest distance at the current state between the source vertex and the
     unknown vertices.
35
           int vertex = a[i] - 1;
36
           // If the distance of the current node is invalid or greater than the minimum distance, output No and return
37
           if(Graph[vertex].distance < 0 | Graph[vertex].distance > min) {
38
39
             printf("No\n");
40
             return;
41
           }
42
43
           Graph[vertex].known = 1; // Mark the current node as processed
44
           int cnt = Graph[vertex].neinum; // get the count of the neighbours of the processing vertex
45
           // Update the distance of adjacent nodes
46
47
           for(int j = 0; j < cnt; j++) {
             int neighbour = Graph[vertex].list[j].index;
48
49
             int distance = Graph[vertex].list[j].distance;
50
51
             if(!Graph[neighbour].known) {
52
                // If the distance from the current node to the adjacent node is smaller, update the distance
```

```
53
                 if(Graph[vertex].distance + distance < Graph[neighbour].distance || Graph[neighbour].distance < 0) {
 54
                   Graph[neighbour].distance = Graph[vertex].distance + distance;
 55
                 }
 56
 57
            }
 58
         }
 59
 60
         printf("Yes\n"); // No violations OUTPUT Yes
 61
 62
       }
 63
       int main() {
 64
         int vertice, edge;
 65
         scanf("%d %d", &vertice, &edge); // get the number of vertices and edges in the weighted undirected graph
 66
 67
         // Initialization of the array of graph vertices
 68
         struct ver Graph[vertice];
 69
         for(int i = 0; i < vertice; i++) {
 70
            Graph[i].neinum = 0;
                                      //initialize the number of neighbours to be zero
 71
            Graph[i].list = (struct nei*)malloc(vertice * sizeof(struct nei));
 72
         }
 73
         // Input information about edges and build the adjacency list
 74
         for(int i = 0; i < edge; i++) {
 75
            int from, to, distance;
 76
            scanf("%d %d %d", &from, &to, &distance); //get the information of an edge
 77
            // initialize edge <from,to>
 78
            int index = Graph[from - 1].neinum;
 79
            Graph[from - 1].neinum++;
 80
            Graph[from - 1].list[index].distance = distance;
 81
            Graph[from - 1].list[index].index = to - 1;
 82
            //initialize edge <to,from>
 83
            index = Graph[to - 1].neinum;
 84
            Graph[to - 1].neinum++;
 85
            Graph[to - 1].list[index].distance = distance;
 86
            Graph[to - 1].list[index].index = from - 1;
 87
         }
 88
         int cnt;
 89
         scanf("%d", &cnt);
 90
         // Process multiple queries
 91
         while(cnt--) {
 92
            int a[vertice];
 93
            for(int i = 0; i < vertice; i++) {
 94
              scanf("%d", &a[i]);
 95
            }
 96
            // Initialize the distances and processing status of the vertices
 97
 98
            for(int i = 0; i < vertice; i++) {
 99
              Graph[i].distance = -1;
100
              Graph[i].known = 0;
101
```

```
102
            // Call Dijkstra's algorithm
103
            Dijkstra(Graph, vertice, a);\\
104
         // Free allocated memory
105
106
         for(int i = 0; i < vertice; i++) {
107
            free(Graph[i].list);\\
108
         }
109
         return 0;
110
```

Declaration

I hereby declare that all the work done in this project titled "Dijkstra Sequence" is of my independent effort.