

# An STT-MRAM Based Strong PUF

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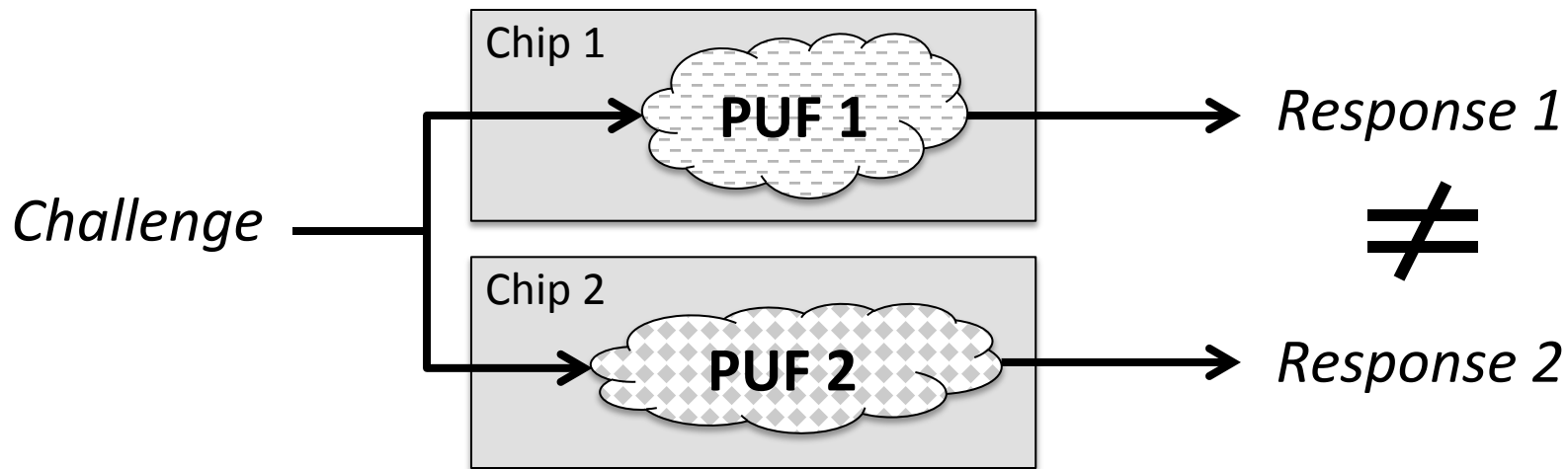
# What is a PUF?

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- No two chips are exactly the same
  - Due to manufacturing *process variations*

## ✧ **PUF:** Physically Unclonable Function

- A device, based on physical disorders of chips



- **CRP:** Challenge-Response Pair

# PUFs at a Glance

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- Manufactured with the same layout
- A unique function per chip
  - *Challenge-Response Pairs (**CRPs**)*
- Physically unclonable
  - Impossible to avoid random process variations
- Unpredictable behavior
  - Unless by testing all CRPs
  - Known CRPs cannot be used to predict responses to new challenges

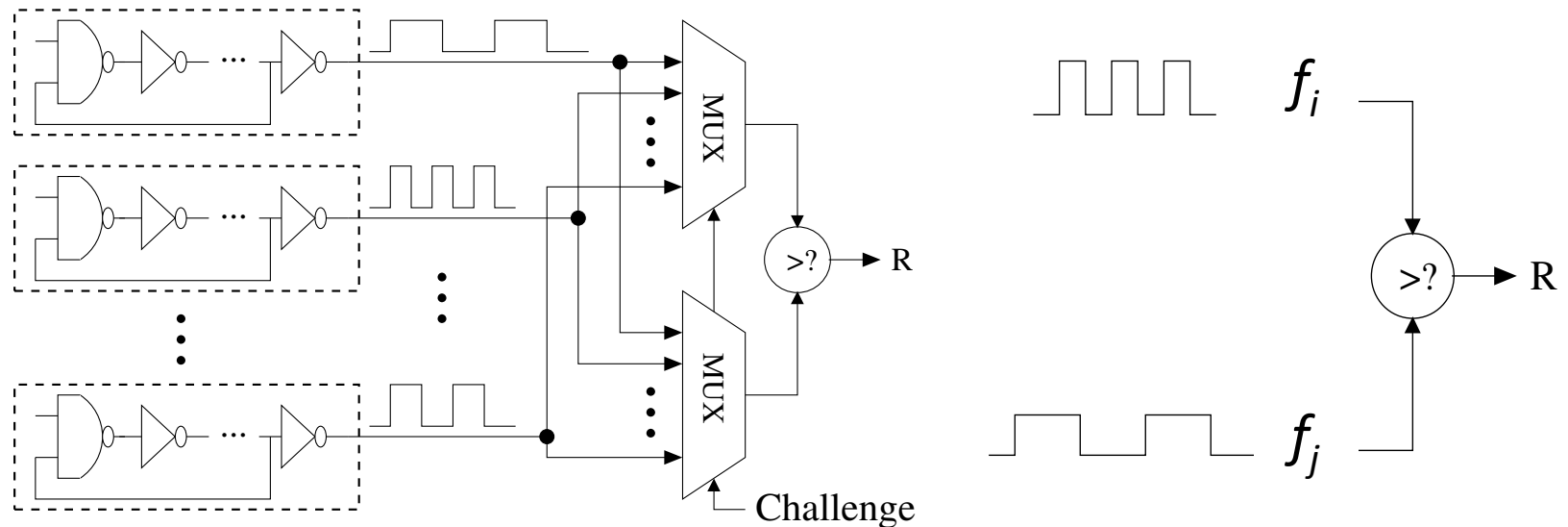
# Security Advantages

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- Keys are generated on demand
  - No need to program the key
  - No non-volatile memory required
- No need to store information on chip
  - Unlike conventional digital storage
  - Security achieved through a *Challenge-Response* mechanism
- No attack when the chip is OFF
- Resilient against *invasive* attacks
  - The PUF would be changed/destroyed

# Weak PUFs

## ❑ Example: *Ring-Oscillator (RO) based PUF\**



❖  $n$  ROs  $\Rightarrow \binom{n}{2}$  CRPs :  $O(n^2)$

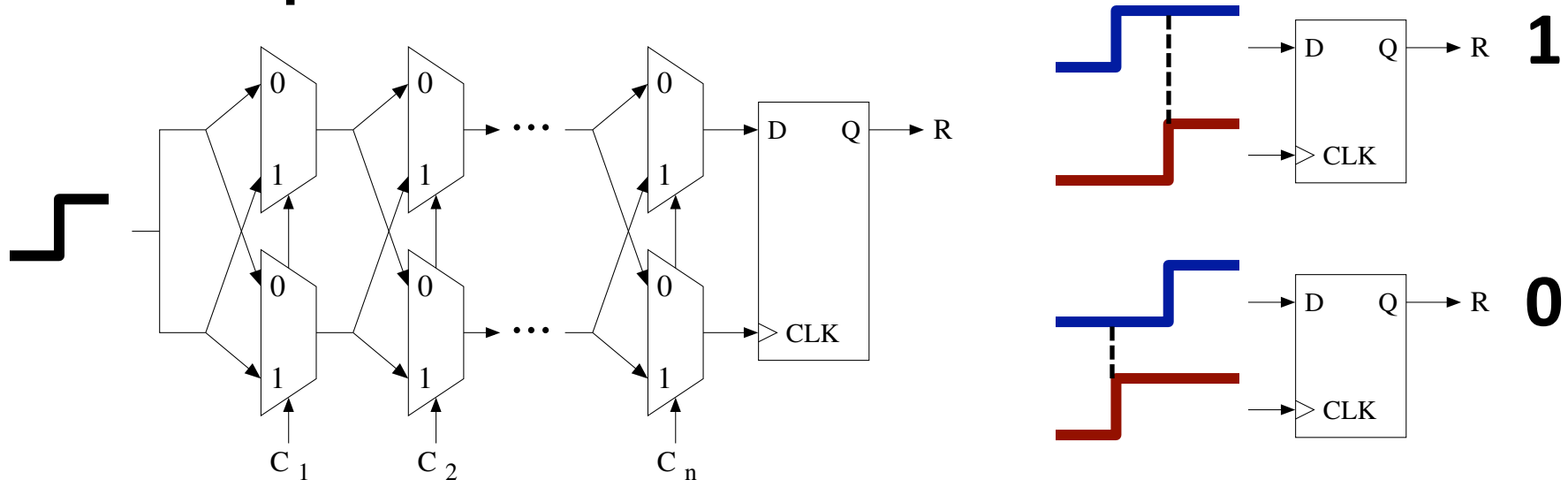
## ➤ **Weak PUF:** Limited number of independent CRPs

- Polynomial w.r.t. the number of components

\* G. Sush and S. Devadas, *IEEE/ACM Design Automation Conference*, 2007

# Strong PUFs

## ❑ Example: *Arbiter* PUF\*



❖  $2n$  Multiplexers  $\Rightarrow 2^n$  CRPs :  $O(2^n)$

➤ **Strong PUF:** Huge number of independent CRPs

- Exponential w.r.t. the number of components

# Weak PUF vs. Strong PUF

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## Weak PUF

- **Limited** number of CRPs
  - Polynomial
- CRPs must be kept **secret**
  - Otherwise, the attacker can fully characterize the PUF
- Applications
  - Key generation, fingerprint

## Strong PUF

- **Huge** number of CRPs
  - Exponential
- CRPs are **not secret**
  - Infeasible to fully recover the truth-table
- Applications
  - Device authentication, logic obfuscation, etc.

# Overview of the Proposed Work

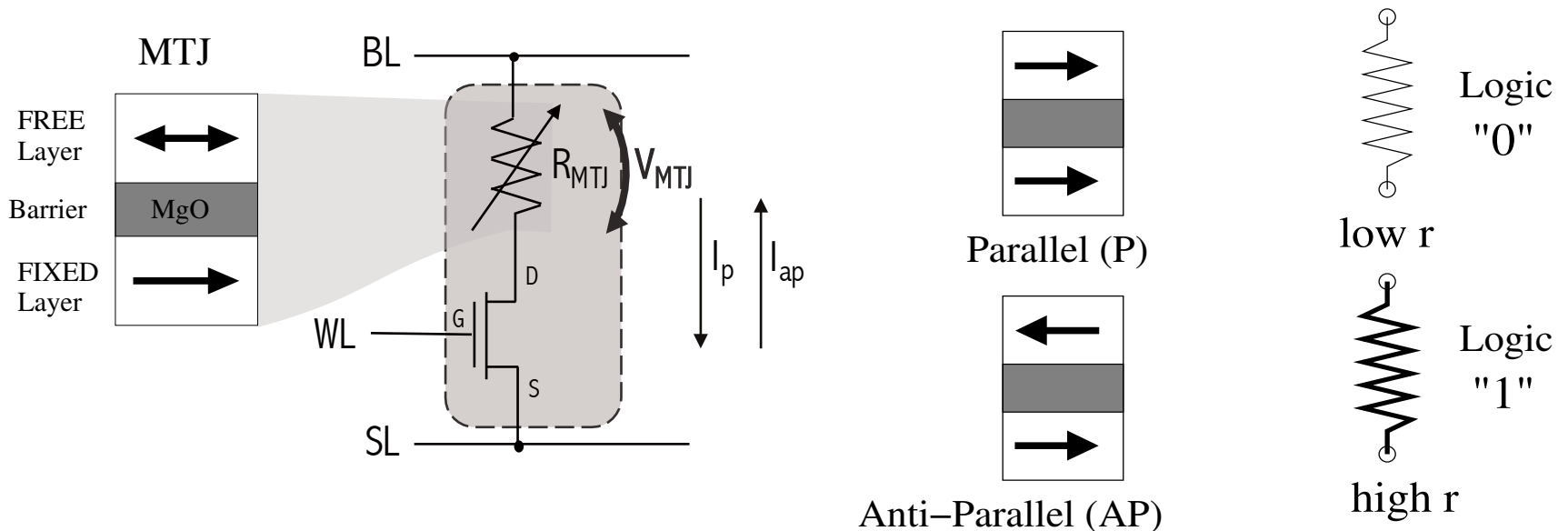
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- 1) Proposing a strong PUF based on STT-MRAM devices
  - ❖ Previous work: an STT-MRAM based **weak PUF**
  - ❖ Introducing the idea of *Group Formation*
  - ❖ The proposed **strong PUF**
  
- 2) Generalizing the idea of ***Group Formation***
  - ❖ How to make a strong PUF in general?
  - ❖ Applying the idea of group formation to RO-based weak PUF



# Introduction to STT-MRAM

- Spin-Transfer Torque Magnetic RAM (STT-MRAM)
  - A nano device, with Magnetic Tunnel Junction (MTJ) as its storage



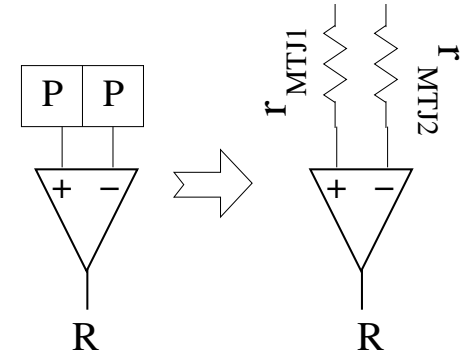
- Two states:
  - 1) *Parallel (P)*: **low** resistance, associated with logic 0
  - 2) *Anti-Parallel (AP)*: **high** resistance, associated with logic 1

# An STT-MRAM based Weak PUF\*

## ➤ Main idea

- Two identical cells with the same magnetization

$$R = \begin{cases} 0 & \text{if } r_{\text{MTJ1}} > r_{\text{MTJ2}} \\ 1 & \text{if } r_{\text{MTJ1}} < r_{\text{MTJ2}} \end{cases}$$

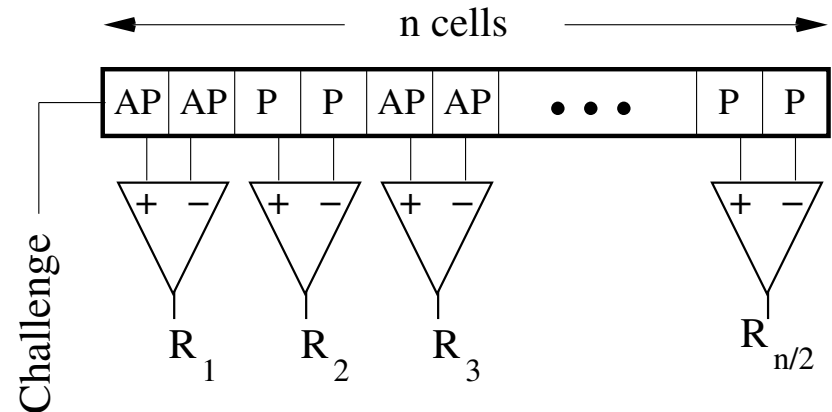


## ➤ Architecture

- $n/2$  pairs  $\Rightarrow 2^{n/2}$  CRPs

❖ Only 2 independent CRPs

- 1) All pairs set to state  $P$
- 2) All pairs set to state  $AP$



## ✓ Weak PUF

# Weak PUF vs. Strong PUF

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- PUFs are based on noisy analog features
  - STT-MRAM: resistances of the MTJ devices
  - Arbiter, RO-based: gate delays
- Analog features “collapsing” into digital bits ==> PUF
  - STT-MRAM: Resistance comparison of two cells
- Analog features can offer an unlimited number of CRPs
  - Infinite precision
- Weak PUF vs. Strong PUF
  - How much of such infinite precision is exploited before collapsing into digital?

# Strong PUF Component 1: Group Formation

## ➤ Previous *Weak* PUF:

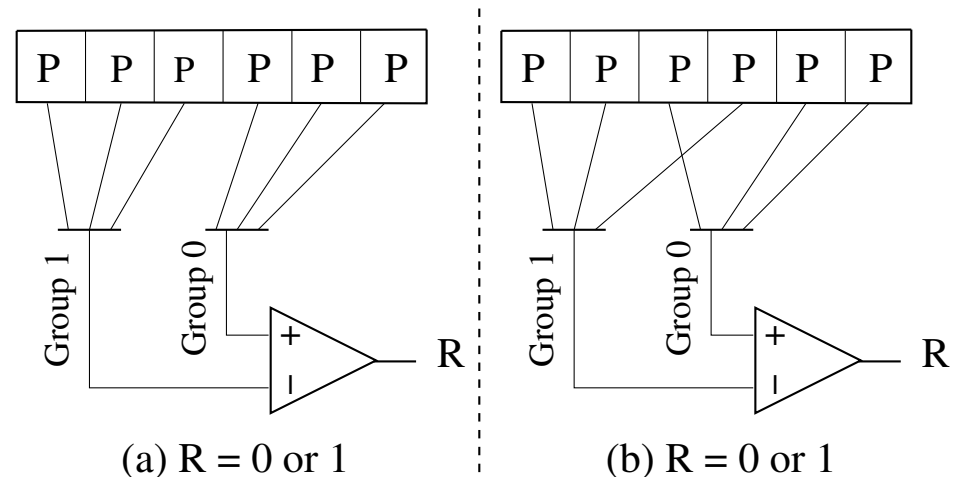
- Comparing the resistances of two **single** cells

## ➤ **Group Formation:**

- Comparing the resistances of two **groups** of cells

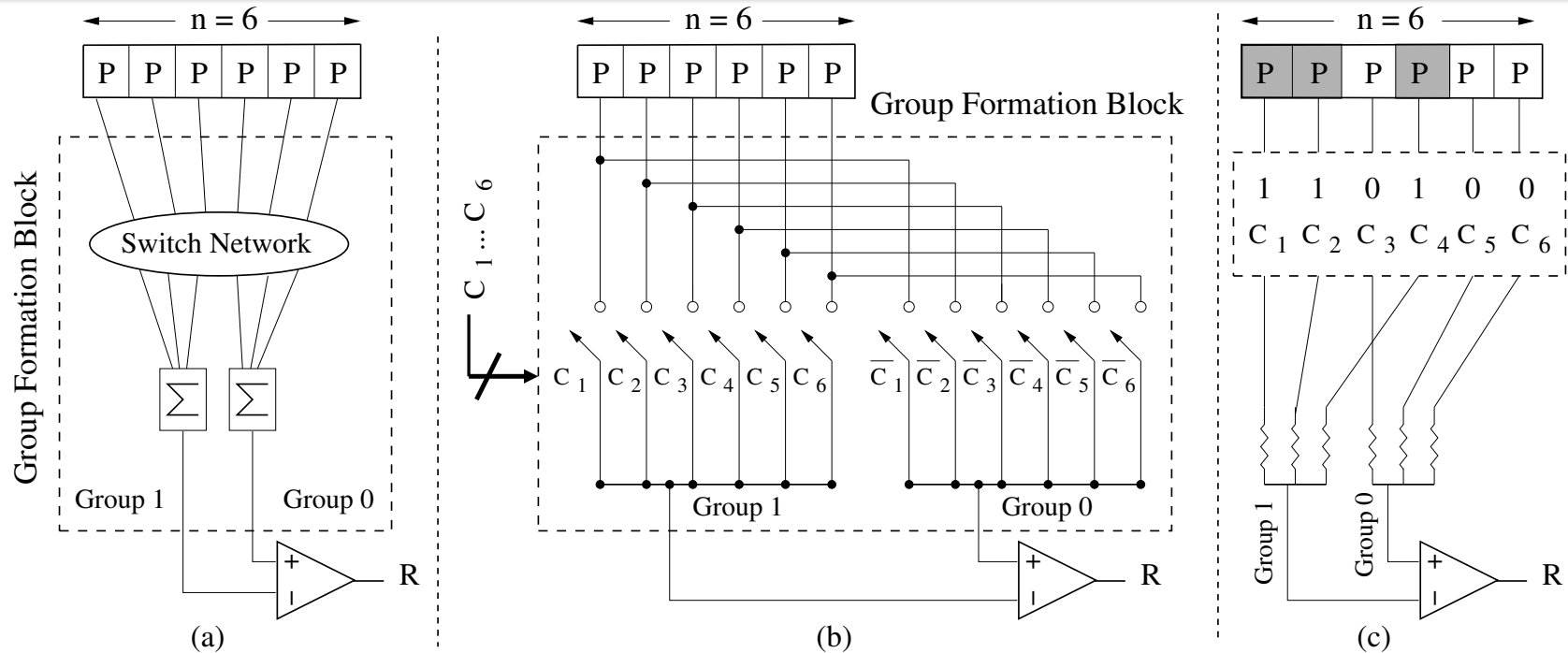
### □ Example:

- 6 STT-MRAM cells
- 3 cells per group
- Fixed magnetization
- Group 1 vs. Group 0



❖ Configurable grouping: **10 independent CRPs**

# Group Formation Block



- Two groups (*Group 1* and *Group 0*) with 3 cells
- Each cell can join either of groups, but not both simultaneously
- Challenge bits:  $C_1$  to  $C_6$  , controlling 12 switches

## ❑ Example:

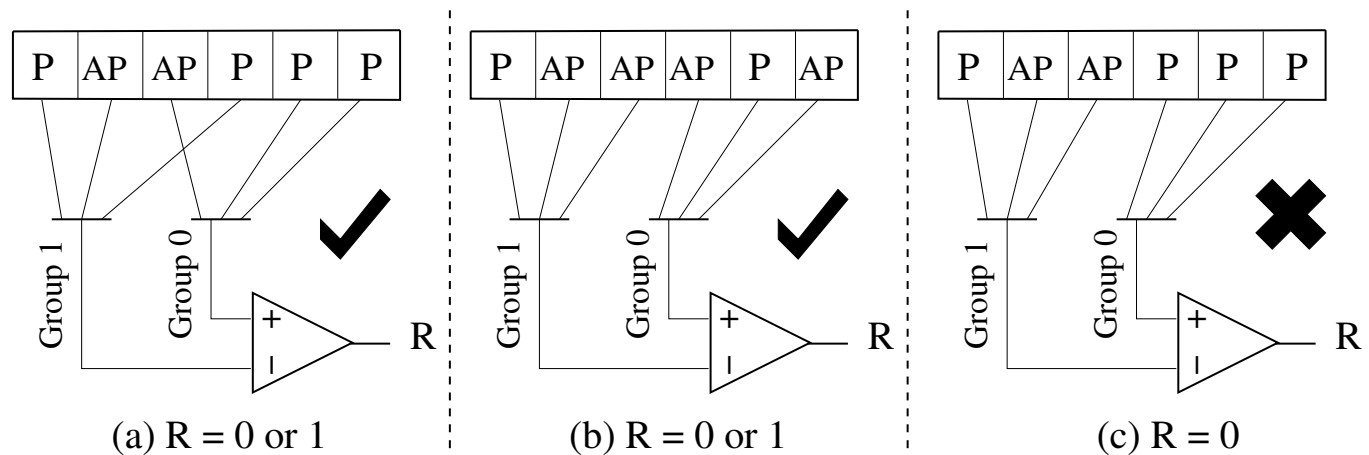
- *Group 1*: Cells in Grey      *Group 0*: cells in white

# Strong PUF Component 2: Bit Pattern

## ➤ Bit Pattern:

- Changing the magnetization of cells in each group  
(assumed to be fixed so far)

## □ Example:



- Same number of cells in states *P* and *AP* in both groups
- ❖ **20 independent CRPs** (per 6 cells) for every group choice

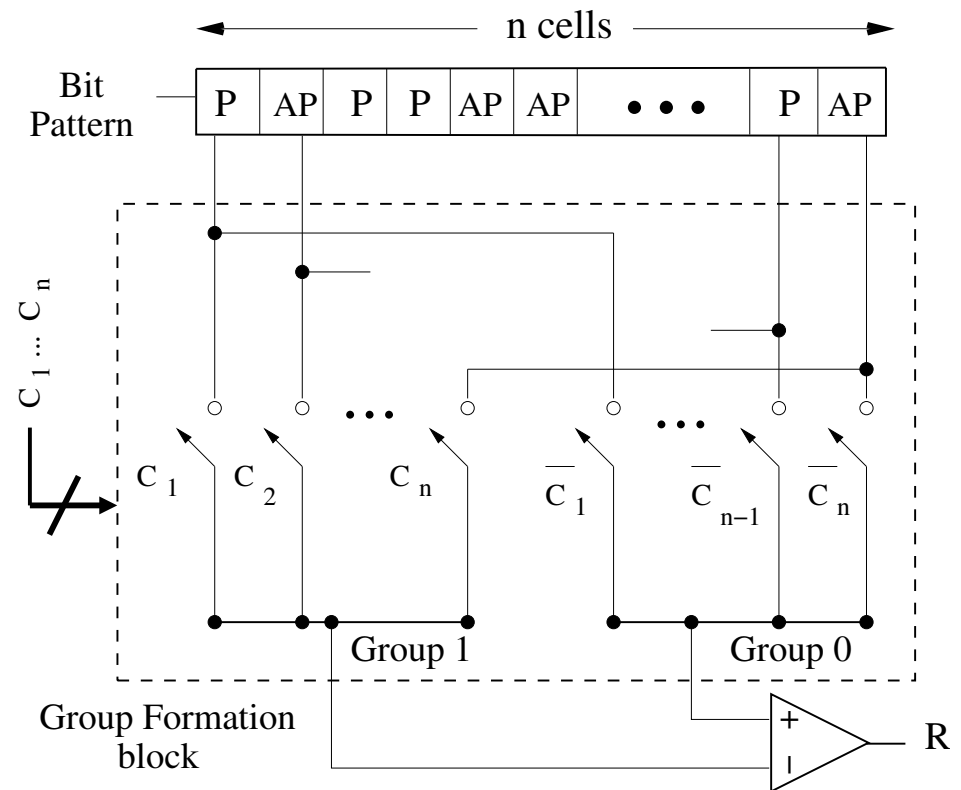
# Proposed STT-MRAM Based Strong PUF

- $n/2$  cells per group
- Each cell can join either of groups, but not both simultaneously
- Groups with less than  $n/2$  cells: not allowed by the architecture
  - To avoid using the info. of smaller groups to determine the larger ones

## ❖ Challenges:

- 1) Group Formation
- 2) Bit Pattern Selection

## ❖ Responses: 1 or 0 (1-bit)



# Analysis of the Proposed PUF (1/2)

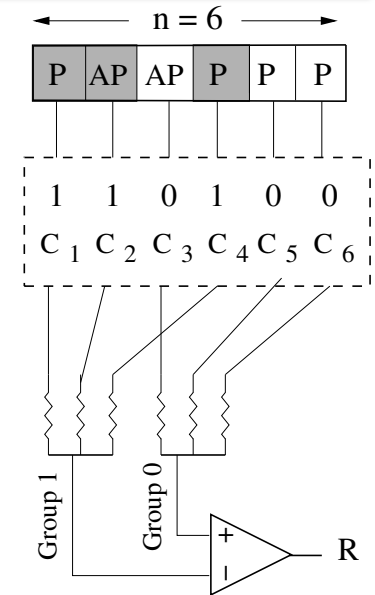
## ➤ Total resistance of each group

- Parallel equivalence of all cells

## ❑ Example:

- 3 cells per group (two  $P$ 's and one  $AP$ )

$$\frac{1}{r_{G1}} = \frac{1}{r_1^P} + \frac{1}{r_2^{AP}} + \frac{1}{r_4^P} \quad \text{and} \quad \frac{1}{r_{G0}} = \frac{1}{r_3^{AP}} + \frac{1}{r_5^P} + \frac{1}{r_6^P}$$



## ➤ Hardware overhead

- Previous weak PUF:  $n$  sense-amplifiers
- Proposed PUF: 1 sense-amplifier +  $2n$  switches

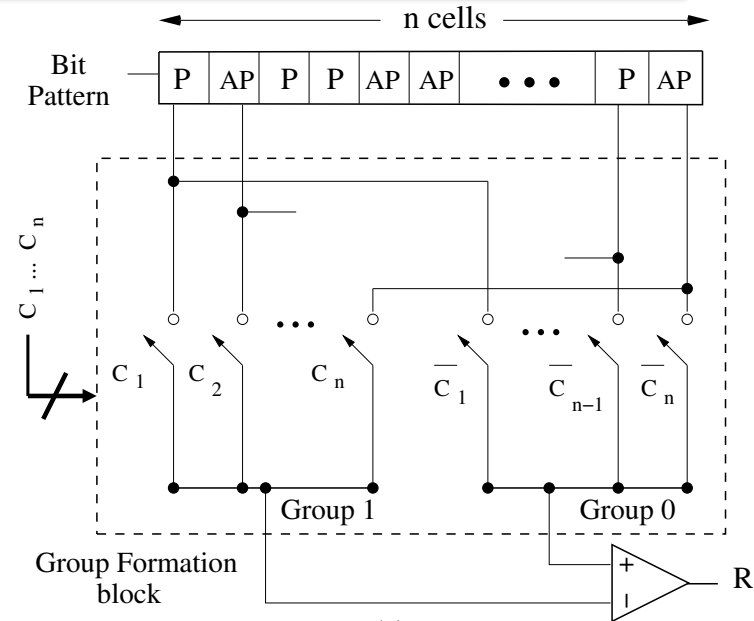


# Analysis of the Proposed PUF (2/2)

## ➤ Number of independent CRPs

1) Group Formation:  $\frac{1}{2} \binom{n}{n/2}$

2) Bit Pattern Selection:  $\sum_{i=0}^{n/2} \binom{\frac{n}{2}}{i}^2$



## ➤ Total number of independent CRPs

$$\frac{1}{2} \times \binom{n}{\frac{n}{2}} \times \sum_{i=0}^{n/2} \binom{\frac{n}{2}}{i}^2 = \frac{1}{2} \binom{n}{\frac{n}{2}}$$

*Factorial* growth of **independent** CRPs ==> Strong PUF

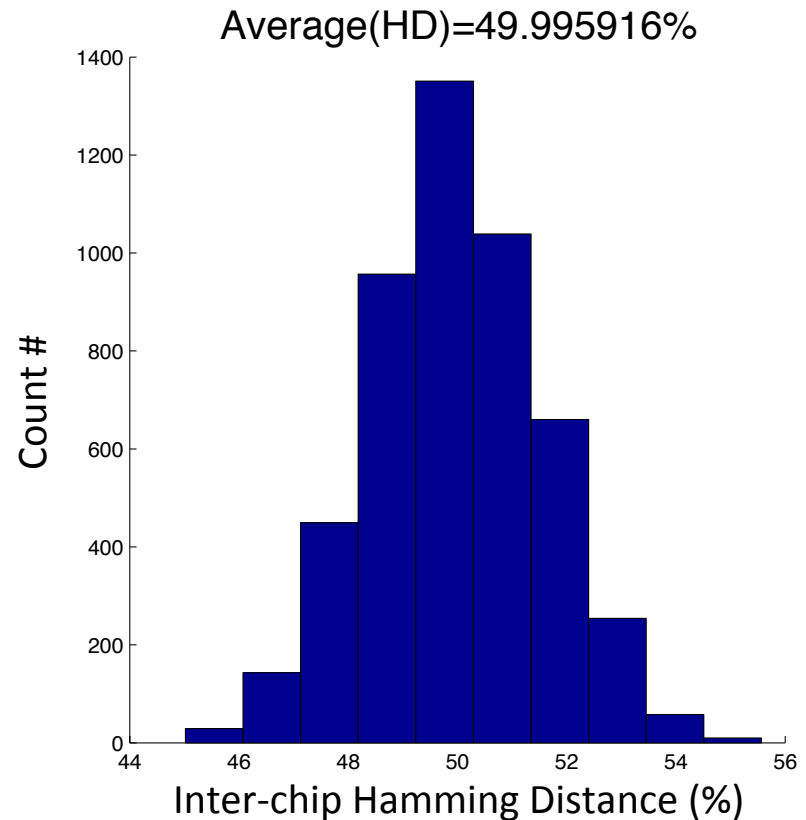
# Simulation Results: Inter-chip

## ❑ Experiment Setup:

- 1024-bit responses
- 100 Chips
- MATLAB model\*

## ➤ Inter-chip randomness

- Hamming Distance



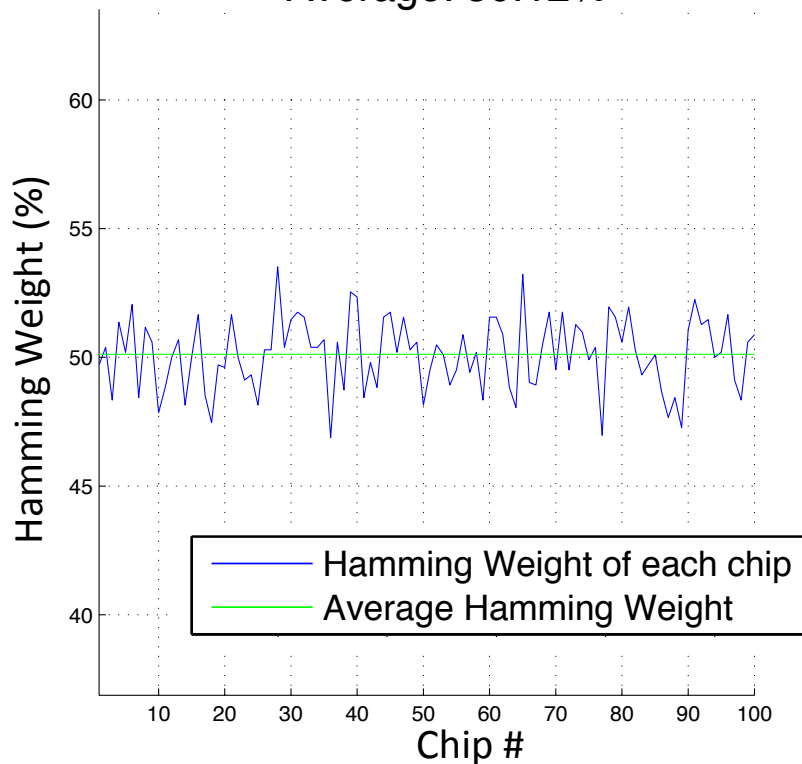
\* L. Zhang et al, *IEEE International Symposium on Circuits and Systems*, 2014

# Simulation Results: Intra-chip

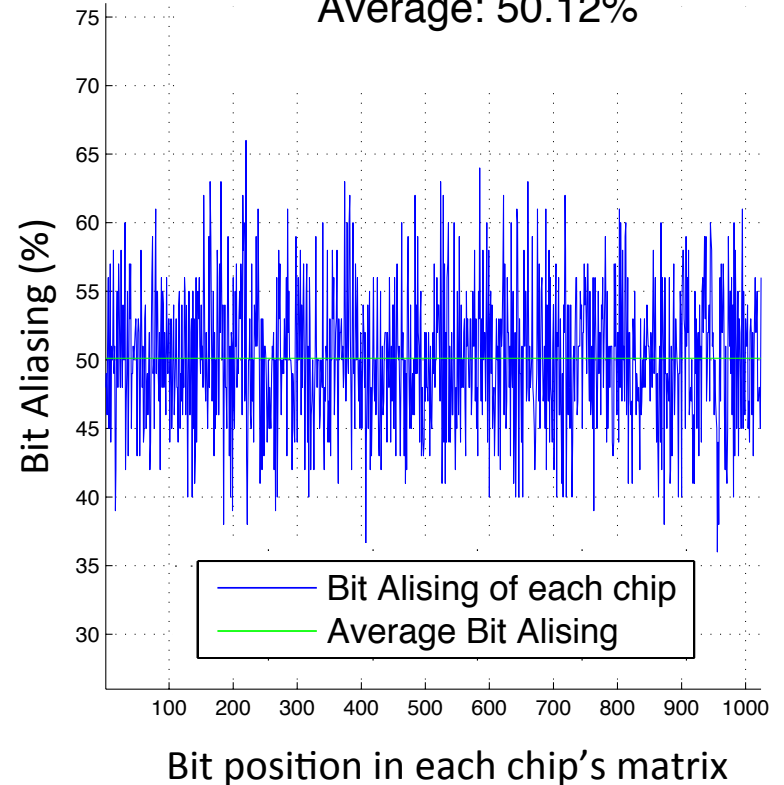
## ➤ Intra-chip randomness

- 1) *Hamming Weight*: randomness of **bits** within the *same* response
- 2) *Bit Aliasing*: **single bit** randomness among *different* responses

Average: 50.12%



Average: 50.12%



# Strong PUF Beyond STT-MRAM

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## ➤ Necessary conditions for making a strong PUF

### 1) **Device-level** compatibility for group formation

- STT-MRAM: parallel combination of MTJ cells' resistances
- Arbiter PUF: gate delays
- Not supported in “SRAM-based power-up PUF” \*

### 2) **Architecture-level** support to achieve independent CRPs

- STT-MRAM: group formation flexibility
- Arbiter PUF: path selection

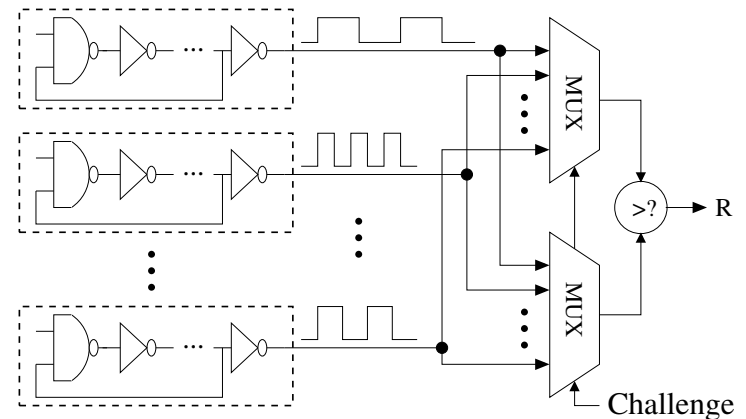
\* D. Holcomb et al, *IEEE Transactions on Computers*, 2009

# RO-based Strong PUF

## ➤ RO-based weak PUF

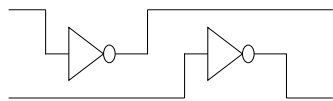
- Two ROs compared at a time

❖  $n$  ROs  $\Rightarrow \binom{n}{2}$  CRPs :  $O(n^2)$

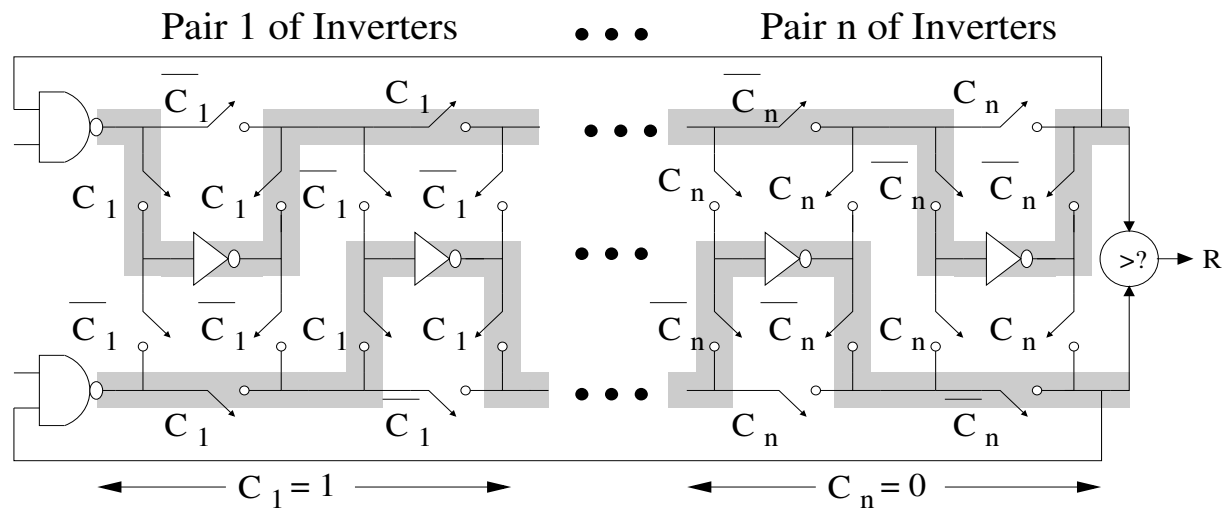
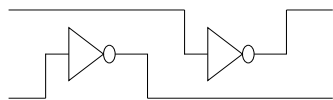


## ➤ RO-based **strong** PUF

1st option:  $C_i = 1$



2nd option:  $C_i = 0$



❖  $2n$  Inverters  $\Rightarrow 2^n$  CRPs :  $O(2^n)$

# Conclusions

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## ✧ Weak PUF vs. Strong PUF

- How much to exploit the noisy analog feature?

## ✧ An STT-MRAM based strong PUF

- Based on the idea of Group Formation
  - Fixing the group size
  - Allowing each device to join either of groups, but not both
- Simulation results confirmed high quality of randomness

## ✧ Making a strong PUF in general

- 1) Device-level compatibility for group formation
- 2) Architecture-level support for huge CRPs