x86 Assembly

CS350 — Undergraduate Operating Systems 02/10/2021

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Today's agenda

- Hints about the Primer
- x86 Assembly
 - Architecture Overview
 - The AT&T Operand Format
 - Data Movement Instructions
 - Arithmetic Instructions
 - Bitwise Logical Instructions
 - Control Flow Instructions
 - GNU Inline Assembly

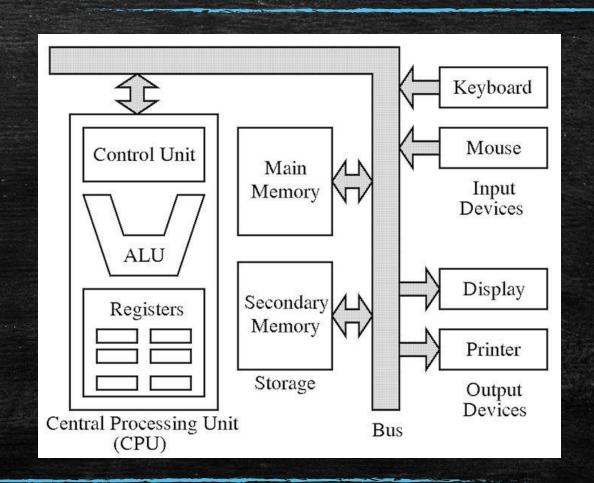
How to go about the Keyboard Driver?

- Start as early as you can!
- Start by extending the provided IOCTL module/test code
- Polling method: You need to disable the native driver (i8o42).
- Interrupt Based: Look for Linux kernel functions to define and install your own ISR. Google it!
- At this point, you will get double characters as you are co-existing with the native i8042! Look for a way to bypass that.
- Distinguish key-press from key-release events. Look at the bit definition of the data port (ox6o). Read this and then this!
- Work out how to manipulate wait queues in Linux. You need to make the caller task block until there's a keyboard event!

x86 Architecture Overview

Von Neumann Architecture

- Unified memory for Instructions and Data, i.e., the Main Memory
- Move data between Main Memory and Registers
- System Bus:
 - Data Bus: 32-bit wide
 - Address Bus: 32-bit wide
 - Control Bus: R/W, Memory/IO



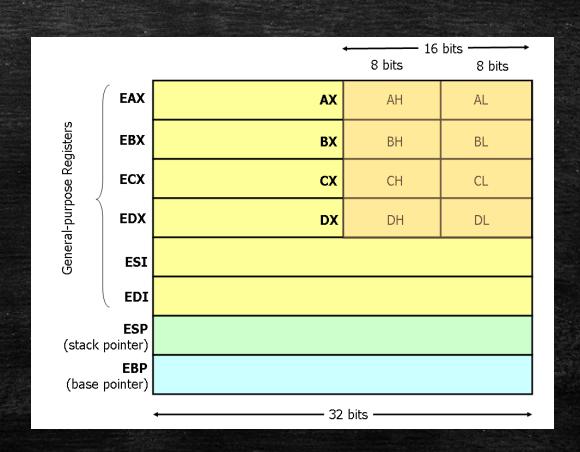
x86 Registers

General Purpose Registers

- EAX, EBX, ECX, EDX
- AX, BX, CX, DX
- AH, AL, BH, BL, ...
- ESI, EDI

Special Purpose Register

- ESP: Stack pointer
- EBP: Base pointer
- EIP: Instruction Pointer
- And a lot more: Control Registers, Model-Specific Registers and so on....



Program Organization (1)

- A program is a collection of Data and Code spread over different sections.
- We use GNU-AS to compile our x86 assembly code.
- Some useful GNU-AS directives to organizing your program:
 - Definition of sections:
 - .data, .rodata, .bss, .text, .section [name-of-custom-section]
 - Definition of labels: Labels are name of constants, variables, functions and anything that can be addressed in the program!
 - [name-of-label]:
 - Data definition directives:
 - .byte, .short, .long, .zero, .string, .space, .float, .double

Program Organization (2)

Examples data section of an assembly program

```
.data /* initialized global variables */
my_byte_arr: /* define an array of bytes labeled my_byte_arr. */
    .byte 64, 0x10, 0xFF
x: /* define a 2-byte integer variable labeled x initialized to 42 */
    .short 42
y: /* define a 4-byte integer variable labeled y = 0x1234ABCD */
    .long 0x1234ABCD
s1: /* define a null-terminated string initialized to "Hello World" */
    .string "Hello World"
.bss /* Uninitialzed global variables */
buf: /* Reserve 256 bytes in a buffer labeled buf */
    .space 256
.text
/* This is where our code (x86 instructions) goes! */
```

An Introduction to x86 Instruction Set

- Data Movement Instructions
 - mov, push, pop, lea, in, out
- Arithmetic Instructions
 - add, sub, inc, dec, imul, idiv
- Logical Instructions
 - and, or, xor, not, neg, shl, shr, sar
- Control-Flow (Branching) Instructions
 - jmp, je, jne, jz, jnz, jg, jge, jl, jle
 - call, ret
 - int
- Many more that we can't possibly cover in this lab.

The AT&T Syntax - Operands

- The AT&T syntax for instructions w/ more than one operand: INSTR src, dst or INSTR src1, src2, dst
- Let's consider the MOV instruction. It is used to move data between the registers and the main memory.
- The source and destination operands can be one of the following:
 - 1. Registers: %[register name] e.g., mov %eax, %ebx
 - 2. Immediate: **\$[constant value]** e.g., mov **\$0x10**, %eax
 - 3. Memory Location: Follows the format Offset(Base, Index, Scale)
 - Offset is an immediate (Constant number or a label)
 - Base and Index are x86 registers
 - Scale is an immediate (Constant number or a label)
 - Memory Location = Offset + (Base + Index*Scale)

The AT&T Syntax - Operands

- Sometimes the instruction operands cannot unequivocally specify how many bytes should the instruction operate on. E.g.,:
- Copying some content from address X in the memory into EAX: Should the CPU copy 1, 2 or 4 bytes from X?
- The following suffixes are attached to the name of the instruction to clarify the size of the operand:
 - **b**: One byte
 - w: A word, i.e., Two bytes
 - 1: A long word, i.e., Four bytes
- Then the instruction mov has four forms: mov, movb, movw, mov1

Data Movement Instructions (1)

Move!

- Syntax
 - mov <reg>, <reg>
 - mov <reg>, <mem>
 - mov <mem>, <reg>
 - mov <imm>, <reg>
 - mov <imm>, <mem>
- Remember: The first operand is the source and the last one is the destination.
- Also remember that the <mem> operands follow the format: Offset(Base, Index, Scale)

MOV Examples

- 1. Direct Memory: MOVb var(,1), %ecx
 - Write the 1-byte value at memory address of the label var into ECX. The suffix b means "a byte"! Label are translated into address values by the linker.
- 2. Indirect Memory: MOVw (%ebx), %eax
 - Write the 2-byte value at memory address stored in EBX into EAX i.e., Dereference EBX into EAX. The suffix w means "a word" that is 2 bytes.
- 3. Indexed Memory: MOV1 8(%ebx, %esi, 4), %edx
 - Move the 4-bytes value at address 8+(EBX+ESI*4) into EDX. The suffix I means "a long word" that is 4 bytes
 - Assume you want to address the 5th element of an array of 4-byte integers located at offset 8 of a data-structure whose base address is at 0x100000

Data Movement Instructions (2)

LEA (Load Effective Address): Computes the absolute value of a memory location specified in the **Offset(Base, Index, Scale)** format.

Think of MOV as dereferencing a pointer and LEA as reading the address in a pointer.

Syntax

- lea <mem>, <reg32>
- Examples

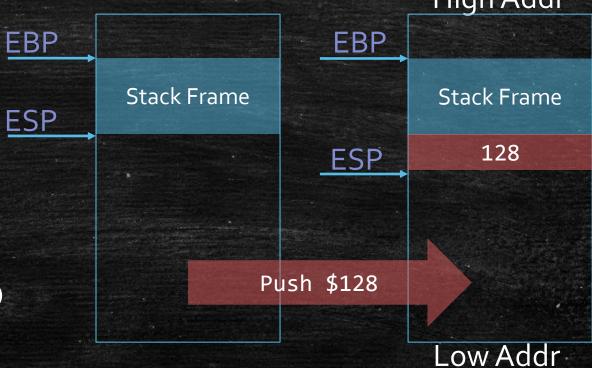
```
- lea (%ebx,%esi,8), %edi /* EDI <- EBX+8*ESI */
- lea val(,1), %eax /* EAX <- val */</pre>
```

Data Movement Instructions (3)

PUSH: Places its operand onto the top of the hardware supported stack in memory, where ESP points.

High Addr

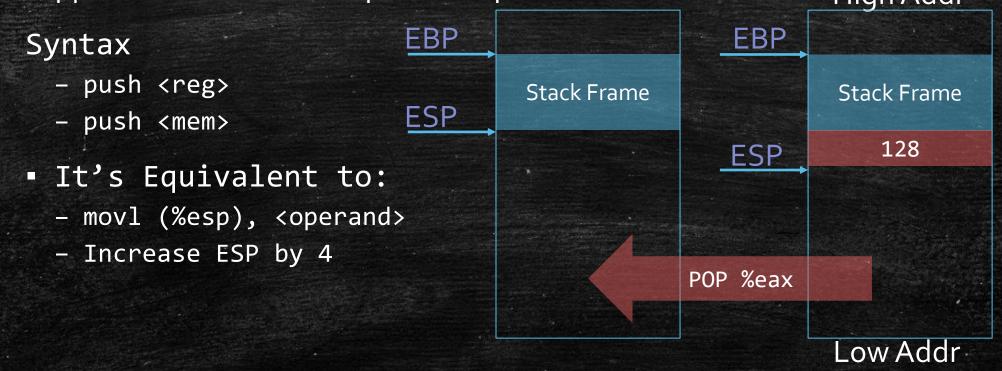
- Syntax
 - push <reg>
 - push <mem>
 - push <imm>
- It's Equivalent to:
 - Decrement ESP by 4
 - movl <operand>, (%esp)



Data Movement Instructions (4)

POP: Removes the 4-byte data element from the top of the hardwaresupported stack into the specified operand.

High Addr



Data Movement Instructions (5)

Reading and Write Hardware I/O Ports

- Syntax
 - in <imm8>, <AL, AX, or EAX>
 - in DX , <AL, AX, or EAX >
 - out <AL, AX, or EAX >, <imm8>
 - out <AL, AX, or EAX >, DX
- Examples
 - inb \$0x64, %al /* Read one byte from port# 0x64 into AL */
 - inw %dx, %ax /* Read two bytes from port# in DX into AX */
 - outb %al, %dx /* Write the byte in AL into the port in DX */

Arithmetic Instructions (1)

Integer Addition

- Syntax
 - add <reg>, <reg>
 - add <mem>, <reg>
 - add <reg>, <mem>
 - add <imm>, <reg>
 - add <imm>, <mem>
- Examples
 - add \$10, %eax /* EAX is set to EAX + 10 */
 - addb \$10, (%eax) /* add 10 to the single byte stored at memory address stored in EAX */

Arithmetic Instructions (2)

Integer Subtraction

- Syntax
 - sub <reg>, <reg>
 - sub <mem>, <reg>
 - sub <reg>, <mem>
 - sub <imm>, <reg>
 - sub <imm>, <mem>
- Examples
 - sub %ah, %al /* AL <- AL AH */</pre>
 - sub \$55, %eax /* EAX <- EAX 55 */</pre>

Arithmetic Instructions (3)

Increment, Decrement

Syntax
- inc <reg>
- inc <mem>
- dec <reg>

- dec <mem>

- Examples
 - dec %eax /* subtract 1 from the contents of EAX */
 incl var(,1) /* add 1 to the 32-bit int. at location var */,

Arithmetic Instructions (4)

Integer Multiplication

- Syntax
 - imul <reg32>, <reg32>
 - imul <mem>, <reg32>
 - imul <con>, <reg32>, <reg32>
 - imul <con>, <mem>, <reg32>
- Examples
 - imul (%ebx), %eax /* EAX <- EAX * 32-bit value @ mem[EBX] */</pre>
 - imul \$25, %edi, %esi /* ESI <- EDI * 25 */</pre>

Arithmetic Instructions (5)

Integer Division

- Syntax
 - idiv <reg32>
 - idiv <mem>
- Divides the contents of the 64-bit integer EDX:EAX by the specified operand value. The quotient result of the division is stored into EAX, while the remainder is placed in EDX.
- Examples

```
- idiv %ebx /* EDX:EAX / EBX */
```

- idivw (%ebx) /* EDX:EAX / <16-bit value at mem[EBX]> */

Logical Instructions (1)

Bitwise AND, OR, XOR

- Syntax
 - and <reg>, <reg>
 - and <mem>, <reg>
 - and <reg>, <mem>
 - and <imm>, <reg>
 - and <imm>, <mem>
 - Similar syntaxes for OR and XOR
- Examples
 - and \$0x0F, %eax /* clear all but the last 4 bits of EAX. */.
 - xor %edx, %edx /* set the contents of EDX to zero. */

Logical Instructions (2)

Bitwise Logical NOT, 2's Complement Negation

- Syntax
 - not <reg>
 - not <mem>
 - neg <reg>
 - neg <mem>
- Examples
 - not %eax /* flip all the bits of EAX */
 - neg %eax /* EAX is set to (- EAX) */

Logical Instructions (3)

Shift Left (SHL) and Right (SHR, SAR)

- Syntax
 - shl <imm8>, <reg>
 - shl <imm8>, <mem>
 - shl %cl, <reg>
 - shl %cl, <mem>
 - Similar syntaxes for SHR and SAR
- Examples

```
- shl $1, %eax  /* EAX = EAX << 2 = EAX *= 2 (if the most significant bit is 0) */
- shr %cl, %ebx /* EBX = EBX >> CL (LOGICAL SHIFT TO RIGHT) */
- sar %cl, %ebx /* EBX = EBX >> CL = EBX <- floor(EBX/(2^CL)) */</pre>
```

Control Flow Instructions (1)

Modify the Instruction Pointer (EIP) based on result of the last arithmetic state

```
Syntax
```

- cmp <reg>, <reg> (Equivalent to sub <reg>, <reg> but does not change the destination value)
- cmp <mem>, <reg>
- cmp <reg>, <mem>
- cmp <imm>, <reg>
- je <label> (jump when equal)
- jne <label> (jump when not equal)
- jz <label> (jump when last result was zero)
- jnz <label> (jump when last result was non-zero)
- jg <label> (jump when greater than)
- jge <label> (jump when greater than or equal to)
- jl <label> (jump when less than)
- jle <label> (jump when less than or equal to)
- jmp <label> (unconditional jump)

Control Flow Instructions (2)

Example - A count-up FOR loop:
begin:
 xor %ecx, %ecx /* Zero out the counter register */
 mov (%esi), %eax /* Store the final count in EAX */
loop:
 /* DO SOMETHING HERE */
 inc %ecx
 cmp %eax, %ecx
 jl loop /* if counter < final_count then jump to loop */</pre>

Example - A count-down FOR loop:
begin:
 mov (%esi), %ecx /* Store the count in ECX */
1:
 /* DO SOMETHING HERE */
 dec %ecx
 jnz 1b /* if ECX != 0 then jump to label 1 before this line */

Control Flow Instructions (3)

- Call/Ret : Call a function/Return from a function
- Syntax
 - call <label>
 - ret
- call is equivalent to

```
push %eip /* Save the address of the next instruction on top of the stack */
jmp foo /* jump to the label foo, i.e., name of the destination function */
```

- ret is equivalent to
 pop %eip /* Retrieve the return address from the stack-top */
- More on x86 Calling Convention: <u>Here</u>

- Can't do everythin in C! E.g., How to make a system call (INT 0x80) in C? :/
- So, what to do now?
 - 1. Make object files from Assembly and link them with your C code!
 - 2. Use inline Assembly!

 Basic syntax: You can have something like the following in your C functions:

Example

- Extended syntax: We can let the compiler decide what registers to use and transfer data between registers and our C variables!
 - asm ("statements" : outputs : inputs : clobbered);
- Format of outputs and inputs:
 - "flags"(variable_name), "flags"(variable_name), ...
- Flags:
 - "r" or "q": Use a register for as an input operand
 - "=r" or "=q": Use a register as an output operand
 - "m", "=m": Memory input/output operand
 - "a", "b", "c", "d", "S", "D", "N": Use registers EAX, EBX, ECX, EDX, ESI, EDI, and 0-255 immediate value, respectively

- Example 1: x_times_5 = x + (4*x)
 asm ("leal (%1,%1,4), %0" : "=r"(x_times_5) : "r"(x));
 - asm ("leal (%%ebx, %%ebx, 4), %%ebx" : "=b"(x) : "b"(x));
- Example 2: Read one byte from an I/O port
 - asm ("inb %1,%0" : "=a"(value) : "Nd"(port_no));
- More info? Here
- Still need more info? Here

References

- X86 Assembly Guide by the Flint Group @ Yale
- Brennan's Guide to Inline Assembly
- GCC-Inline-Assembly-HOWTO by Sandeep.S