	KJSCE/IT/SYBTech/SEM IV/AA/2021-22
Experiment No. : 3	
Title: Virtual lab on Kruskal's algorithm Minimum spanning trees	n to construct
(Autonomous College Affiliated to Universi	ty of Mumbai)
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Batch: B1 Roll No.: 16010420133 Experiment No.: 3

**Aim:** Explore the virtual lab on Kruskal's algorithm to construct minimum spanning trees.

#### **Resource Needed:**

https://ds2-iiith.vlabs.ac.in/exp/min-spanning-trees/index.html

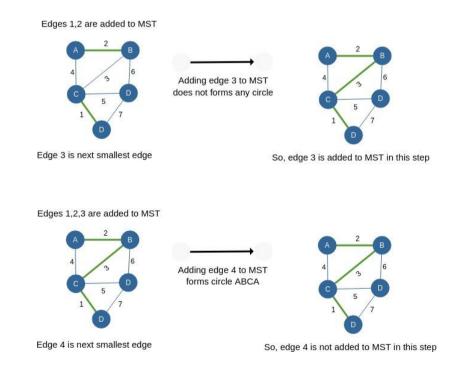
### **Kruskal Algorithm**

#### **Introduction:**

In Kruskal's Algorithm, we take the fastest possible approach to create MST.

- We first sort the edges of graph in increasing order.
- Check the edges in sorted order if they should be in MST or not.
- Whenever we see the newly taken edge making circle (checked by Union-find) (refer picture below), we do not keep it in MST or else we will proceed further.
- We keep performing the above steps over the array again and again till all the edges are checked

### When should we add edge to mst?



### **Algorithm of Kruskal Algorithm:**

```
MST-KRUSKAL(G, w)
   A = \emptyset
1
2
   for each vertex \nu \in G.V
3
        Make-Set(\nu)
   sort the edges of G.E into nondecreasing order by weight w
4
5
   for each edge (u, v) \in G.E, taken in nondecreasing order by weight
        if FIND-SET(u) \neq FIND-SET(v)
6
            A = A \cup \{(u, v)\}\
7
8
            Union(u, v)
   return A
```

### **Explanation and Working of Kruskal Algorithm:**

Let's have a final look at the consolidated algorithm to find MST of given graph:

- **STEP 1:** Sort the given edges.
- **STEP 2:** Check each edge in sorted order if it forms a cycle with already selected edges. Ifnot Add it to list of MST. If it does move to next edge.
- **STEP 3:** Run steps 1 and 2 till v-1 edges are selected (v= no. of vertices).

#### Time complexity and derivation of Kruskal Algorithm:

Running Time of Kruskal's

Lets assume that we are finding MST of a N vertices graph using Kruskal's.

- To check edges we need to sort the given edges based on weights of edges. The best way to sort has a order of  $O(N \log(N))$ .
- To Check one edge if it needs to be in MST or not, we apply Union-find to check if it forms a circle with edges present and add to MST exactly once and apply Union-Find algorithm of order log(E).
- Since we perform at most N checks for a graph total complexity is O(Nlog(E)) for the checkings.
- So, total complexity is O(Nlog(E)+Nlog(N))

Best and Worst Cases for Kruskal's

For regular Kruskal's, time complexity will be O(Nlog(E)+Nlog(N)) in all cases. For Kruskal's:

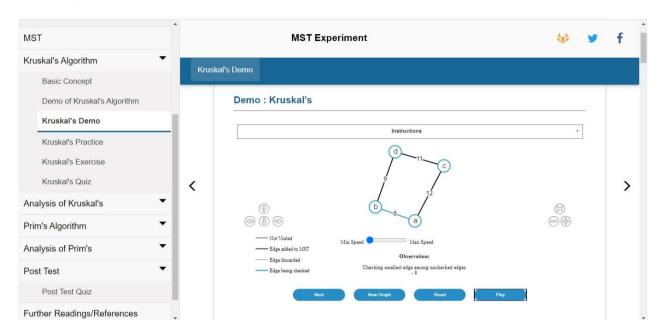
- In best case scenario, we have N no cycles and we have to run N-1 iterations to determine MST.
- Time complexity will be O((N-1)\*log(E)+Elog(E)) in this case.
- In worst case we will have to check all E edges. Time complexity in such case would be O(Elog(E) + Nlog(N))

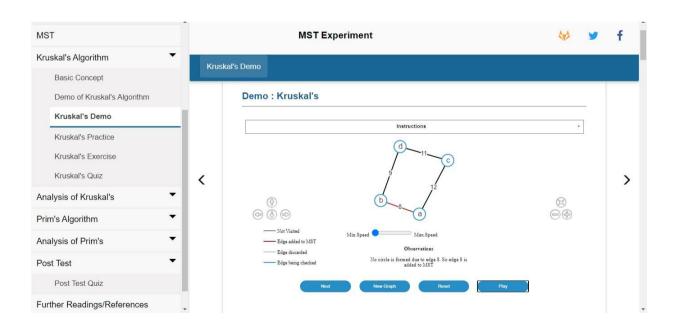
Space Complexity of Kruskal's

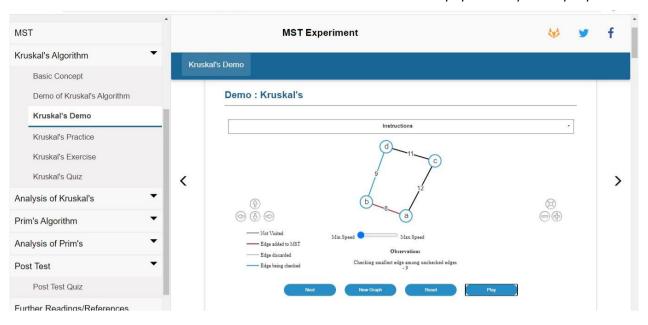
While sorting, we need an extra array to store sorted array of edges (Space comlexity of O(E)), Another array for Union-Find of size O(E). So, total Space Complexity would be  $O(\log(E))$ .

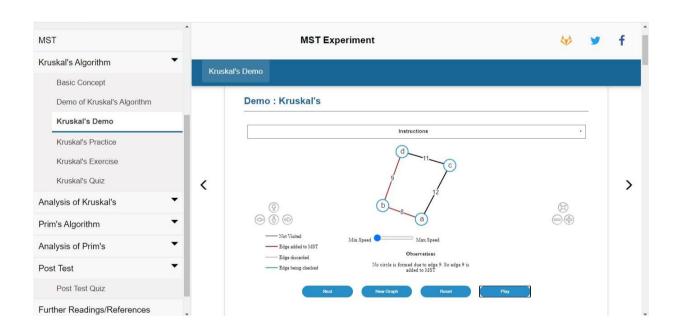
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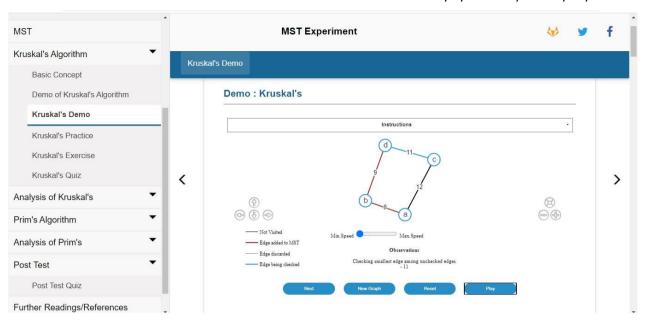
## Observations from Simulation: Kruskal's Demo:

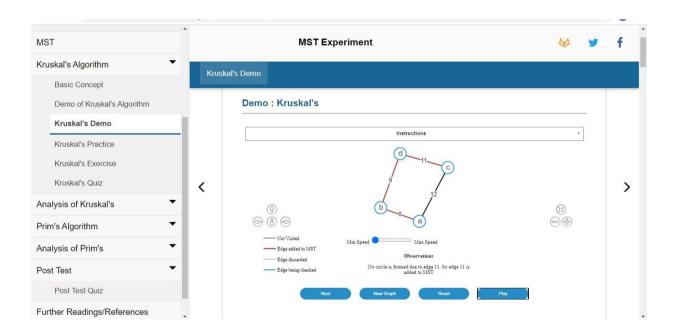


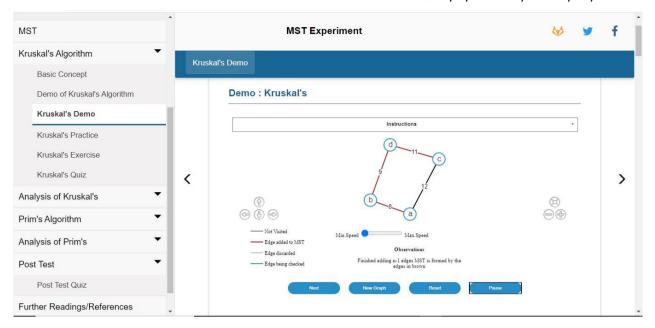




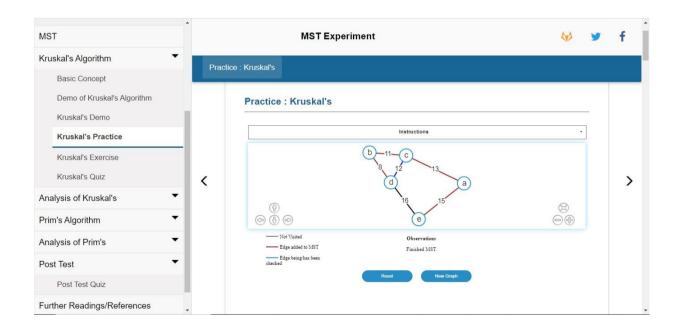




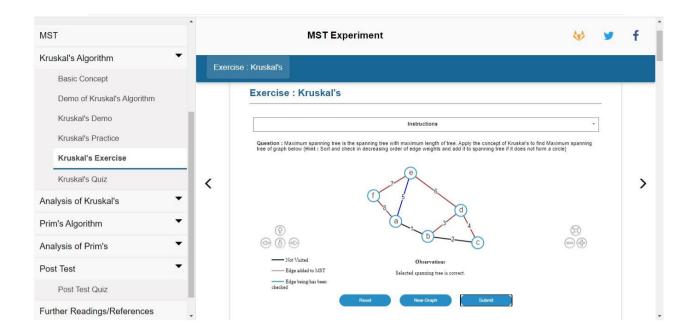




#### Kruskal's Practice:

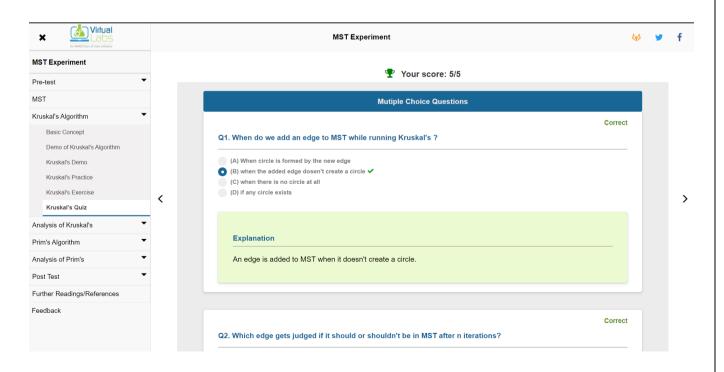


#### Kruskal's Exercise:

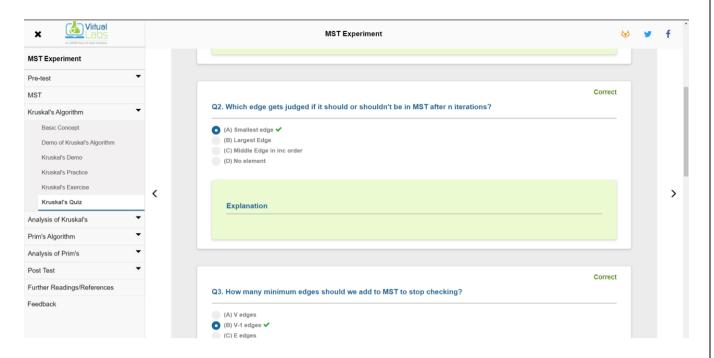


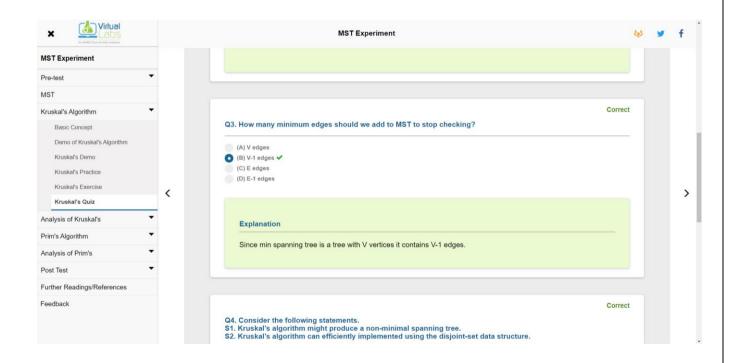
**Self evaluation:** Solve both Kruskal Quiz and Analysis quiz and Display the result of your first attempt.

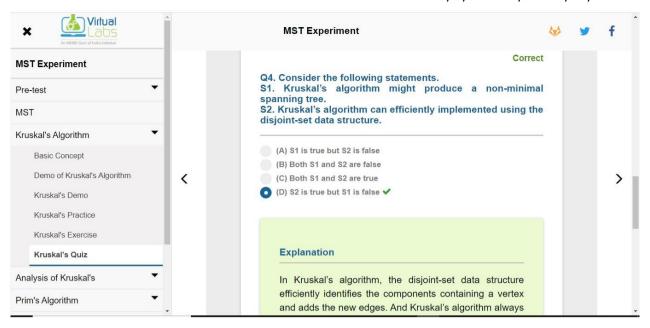
## Kruskal's Quiz:

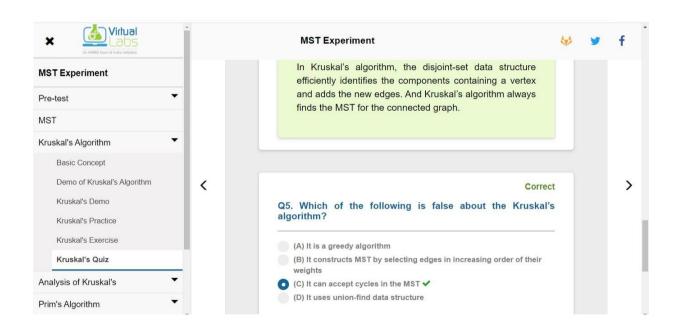


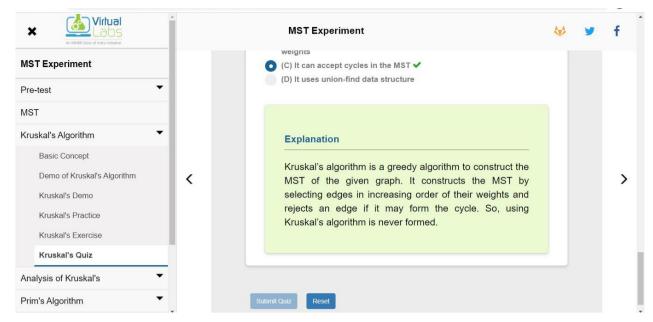
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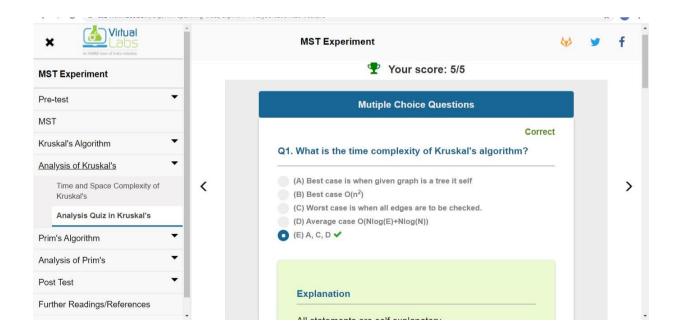


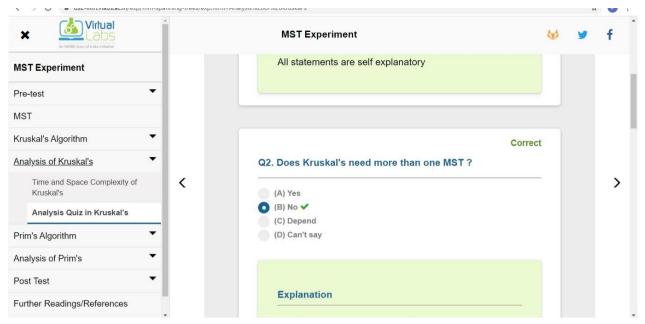


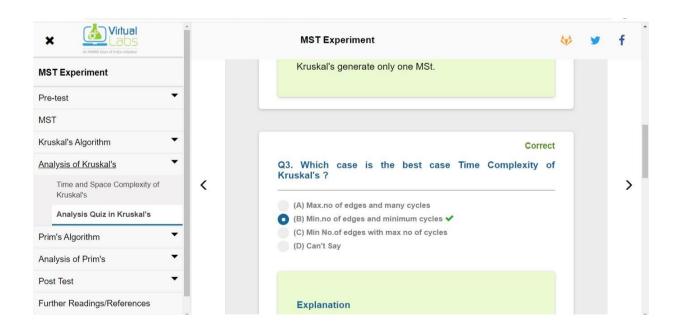


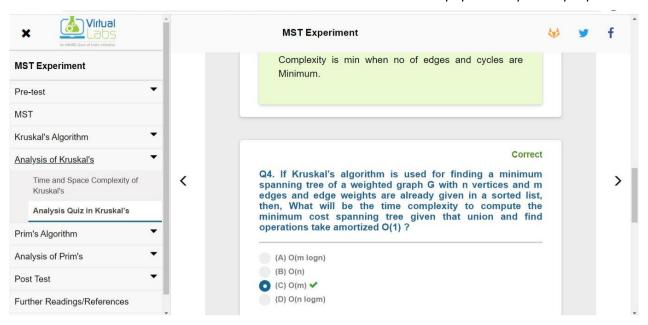


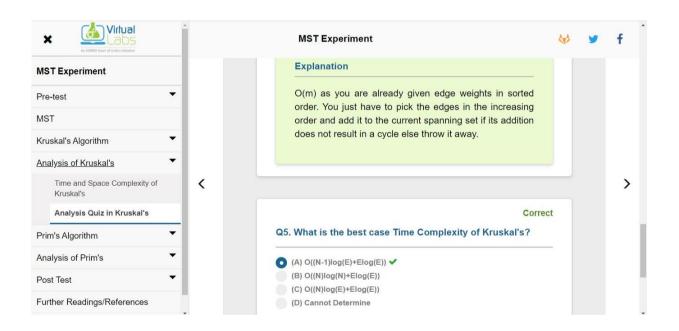
# Analysis Quiz in Kruskal's:

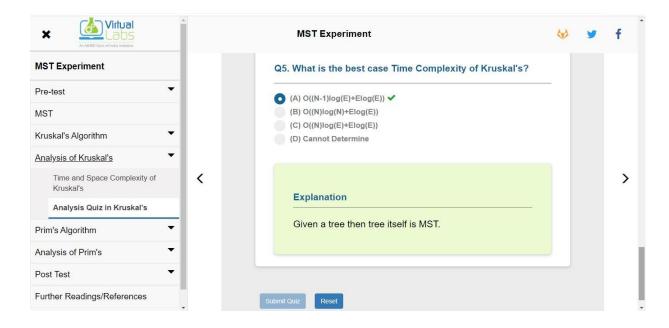












**Self-evaluation:** Solve both Kruskal Quiz and Analysis quiz and Display the result of your first attempt.

## Conclusion: (Based on the observations): MAIYA COLLEGE OF ENGG

Performed the experiment using virtual lab to construct minimum spanning trees using Kruskal's algorithm.

Outcome: CO 3: Implement Greedy and Dynamic Programming algorithms.

#### **References:**

- 1. Richard E. Neapolitan, "Foundation of Algorithms", 5th Edition 2016, Jones & Bartlett Students Edition
- 2. Harsh Bhasin, "Algorithms: Design & Analysis", 1st Edition 2013, Oxford Higher education, India
- 3. T.H. Coreman ,C.E. Leiserson,R.L. Rivest, and C. Stein, "Introduction to algorithms", 3rd Edition 2009, Prentice Hall India Publication
- 4. Jon Kleinberg, Eva Tardos, "Algorithm Design", 10th Edition 2013, Pearson India Education Services Pvt. Ltd.