

Experiment No.: 2

Title: Divide and Conquer Strategy

Batch:B1 Roll No.: 16010420133 Experiment No.:2

Aim: To implement and analyse time complexity of Quick-sort and Merge sort and compare both.

Explanation and Working of Quick sort & Merge sort:

Algorithm of Quick sort & Merge sort:

Merge sort:

```
MERGE-SORT(A, p, r)

1 if p < r

2 q = \lfloor (p+r)/2 \rfloor

3 MERGE-SORT(A, p, q)

4 MERGE-SORT(A, q+1, r)

5 MERGE(A, p, q, r)
```

```
MERGE(A, p, q, r)
 1 \quad n_1 = q - p + 1
 2 n_2 = r - q
 3 let L[1...n_1 + 1] and R[1...n_2 + 1] be new arrays
 4 for i = 1 to n_1
        L[i] = A[p+i-1]
 5
 6 for j = 1 to n_2
       R[j] = A[q+j]
7
 8 L[n_1 + 1] = \infty
 9 R[n_2 + 1] = \infty
10 i = 1
11 \quad j = 1
12 for k = p to r
13
        if L[i] \leq R[j]
            A[k] = L[i]
14
            i = i + 1
15
        else A[k] = R[j]
16
            j = j + 1
17
```

Quick sort:

```
QUICKSORT(A, p, r)
```

```
PARTITION (A, p, r)
1 if p < r
        q = PARTITION(A, p, r)
2
                                       1 \quad x = A[r]
        QUICKSORT(A, p, q - 1)
3

\begin{array}{ccc}
2 & i = p - 1 \\
3 & \text{for } j = p
\end{array}

        Ouicksort(A, q + 1, r)
                                          for j = p to r - 1
                                      4
                                             if A[j] \leq x
                                       5
                                                      i = i + 1
                                                      exchange A[i] with A[j]
                                       7 exchange A[i+1] with A[r]
                                           return i + 1
```

Derivation of Analysis Quick sort:

Worst Case Analysis

$$T(n) = O(n2)$$
.

Best Case Analysis

$$T(n) = O(n(\log(n))).$$

Derivation of Analysis Merge sort:

Worst Case Analysis

$$T(n) = O(n2)$$
.

Best Case Analysis

$$T(n) = O(n(\log(n))).$$

Program(s) of Quick sort:

#include <iostream>
using namespace std;

i++;



```
swap(&arr[i], &arr[j]);
   c++;
  }
 }
 swap(&arr[i+1], &arr[r]);
 c++;
 return (i + 1);
void quickSort(int arr[], int p, int r) {
 if (p < r) {
  int pi = partition(arr, p, r);
  quickSort(arr, p, pi - 1);
  quickSort(arr, pi + 1, r);
}
int main() {
 int size = 100;
 int data[size];
 for(int i=0;i < size;i++){
  data[i] = i+1;
 int n = sizeof(data) / sizeof(data[0]);
 cout << "Unsorted Array: \n";</pre>
 printArray(data, n);
 quickSort(data, 0, n - 1);
 cout << "Sorted arr in ascending order: \n";</pre>
 printArray(data, n);
 cout << "No of passes \n"<<c;</pre>
Merge sort:
#include <iostream>
using namespace std;
int c=0;
void merge(int arr[], int p, int q, int r) {
```

```
int n1 = q - p + 1;
 int n2 = r - q;
 int L[n1], M[n2];
 for (int i = 0; i < n1; i++)
  L[i] = arr[p + i];
 for (int j = 0; j < n2; j++)
  M[i] = arr[q + 1 + i];
 int i, j, k;
 i = 0;
 j = 0;
 k = p;
 while (i < n1 \&\& j < n2) {
  if (L[i] \le M[j]) {
   arr[k] = L[i];
   i++;
   c++;
   } else {
   arr[k] = M[j];
   j++;
   c++;
  k++;
 while (i < n1) {
  arr[k] = L[i];
  i++;
  k++;
  c++;
 while (j < n2) {
  arr[k] = M[j];
  j++;
  k++;
  c++;
}
void mergeSort(int arr[], int l, int r) {
 if (1 < r) {
  int m = 1 + (r - 1) / 2;
  mergeSort(arr, l, m);
  mergeSort(arr, m + 1, r);
```



```
merge(arr, l, m, r);
}
void printArray(int arr[], int size) {
 for (int i = 0; i < size; i++)
  cout << arr[i] << " ";
 cout << endl;
int main() {
 int n = 100;
 int arr[n];
 for(int i=0;i<n;i++){
  arr[i] = i+1;
 int size = sizeof(arr[0]);
 mergeSort(arr, 0, size - 1);
 /* cout << "Sorted array: \n";
 printArray(arr, size); */
 cout << "No of passes \n" << c;
 return 0;
```

Output(0) of Quick sort:

```
Input array
Input array
1 23 3 43 51 35 19 45
Array sorted with quick sort
3 12 19 23 35 43 45 51
...Program finished with exit code 0
Press ENTER to exit console.
```

Merge sort:

```
input
Enter number of elements to be sorted:5
Kater 5 elements to be sorted: 5 3 4 2 1
Sorted array
1
2
3
4
5
. Program finished with exit code 0
Press ENTER to exit console.
```

Results:

Time Complexity of Quick sort:

Worst Case Analysis:

Sr. No.	Input size	No: of steps from Algorithm analysis	No: of steps from Theoretical analysis
1	100	5030	10000
2	500	125247	250000
3	10000	50003998	100000000
4	500000	~125 * 109	250000000000

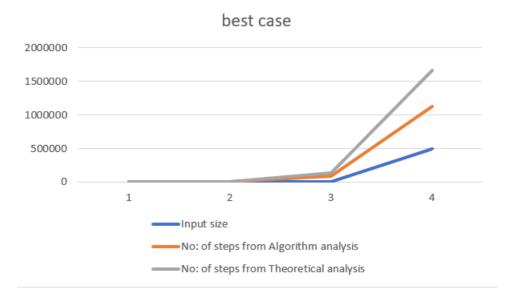
GRAPH:



Best Case Analysis:

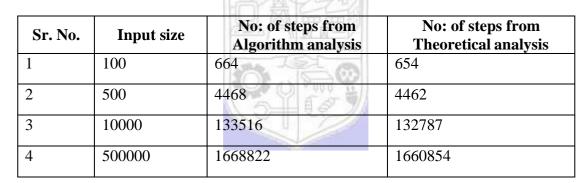
Sr. No.	Input size	No: of steps from Algorithm analysis	No: of steps from Theoretical analysis
1	100	320	664
2	500	2784	4482
3	10000	85031	132877
4	500000	1130033	1660964

GRAPH

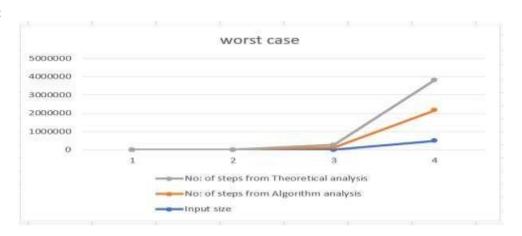


Time Complexity of Merge sort:

Worst Case Analysis:



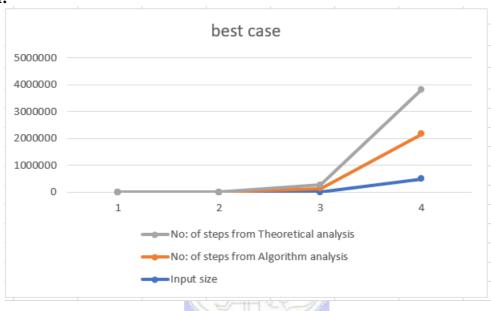
GRAPH:



Best Case Analysis:

Sr. No.	Input size	No: of steps from Algorithm analysis	No: of steps from Theoretical analysis
1	100	664	654
2	500	4468	4462
3	10000	133516	132787
4	500000	1668822	1660854

GRAPH:



Conclusion: (Based on the observations): learned to implement and analyse time complexity of Quick-sort and Merge sort and compare both.

Outcome: CO1: Analyze time and space complexity of basic algorithms.

References:

- 1. Richard E. Neapolitan, "Foundation of Algorithms", 5th Edition 2016, Jones & Bartlett Students Edition
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- 3. T.H. Coreman ,C.E. Leiserson,R.L. Rivest, and C. Stein, "Introduction to algorithms", 3rd Edition 2009, Prentice Hall India Publication
- 4. Jon Kleinberg, Eva Tardos, "Algorithm Design", 10th Edition 2013, Pearson India Education Services Pvt. Ltd.